CSCI-4320/6360 - Assignment 1: C Program for a Carry-Lookahead Adder

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1 Overview

You are to construct a C program that specifically "simulates" a 256 bit Carry Lookahead Adder using 4 bit blocks. We use the term "simulate" here because unlike in real hardware, your CLA adder will not run faster than the serial ripple-carry adder. Here, you will be doing the steps. However, keep in mind that if each level where done completely in parallel via hardware, then it clear executes much, much faster.

For a list of bitwise operators and the associated C language syntax, please see: `https://en.wikipedia.org/wiki/Operators_in_C_and_C%2B%2B`. Make sure you used the older ANSI-C operators and not the newer operator forms/syntax (e.g., XOR is ˆ in C).

Recall, that a CLA adder can be constructed from the follow:

1. Calculate g_i and p_i for all 256 bits i.

2. Calculate gg_j and gp_j for all 64 groups j using g_i and p_i.

3. Calculate sg_k and sp_k for all 16 sections k using gg_j and gp_j.

4. Calculate ssg_l and sp_l for all 4 super sections l using sg_k and sp_k. *Note, it is at this point, we can shift to computing the top-level sectional carries. This is because the number of sections is less than or equal the block size which is 4 bits.*

5. Calculate ssc_l using ssg_l and ssp_l for all l super sections and 0 for ssc_{-1}.

6. Calculate sc_k using sg_k and sp_k and correct $ssc_l, l = k \ div \ 4$ as super sectional carry-in for all sections k.

7. Calculate gc_j using gg_j, gp_j and correct $sc_k, k = j \ div \ 4$ as sectional carry-in for all groups j.

8. Calculate c_i using g_i, p_i and correct gc_j, $j = i \ div \ 4$ as group carry-in for all bits i.

9. Calculate sum_i using $a_i \oplus b_i \oplus c_{i-1}$ for all i.

You need to construct a C-program that executes on `submitty.cs.rpi.edu` that reproduces this algorithm. Each step in the above algorithm will be implemented as a separate function. Also, the arrays can be globally/statically allocated in memory. No need for dynamic memory allocation.
More specifically, you will do the following:

1. Read two 64 digit hex numbers as input.

2. Convert the "hex" (base-16) string to a "binary" (base-2) string. NOTE: That hex[0] is the most significant digit, so when you convert to the "binary" string, you need to make sure that binary[255] is the most significant digit. This can be done by translating each hex digit into it's binary format (e.g., $B = 1011$) and then reversing the whole binary array to make sure that binary[255] is the most signficant digit and in the right order with repect to how the hex digits where translated.

3. Each bit can be a full integer and so you'll have arrays of unsigned ints for each of the components parts of the adder.

4. Write functions for each step in the above algorithm. You can do it using `for` loops and do not have to perform the equation substitutions by hand as we did in class.

5. A master "cla" routine will run thru all the steps.

6. You main routine will invoke your input reading function, followed by your master "cla" function and then reconvert the binary sum of the two numbers into hex output.

7. Last, you can check your answer by creating a simple ripple carry adder based on computing the $c_i = g_i + p_i * c_{i-1}$ for all i and then computing the final sum, based on $sum_i = a_i \oplus b_i \oplus c_{i-1}$.

8. Convert the binary sum_i array to "hex" and output your result in 128 digit "hex" format. Note, you'll need to re-invert/reverse the answer so that it prints correctly.

2 HAND-IN and GRADING INSTRUCTIONS

Please submit your C-code to the `submitty.cs.rpi.edu` grading system. The grading rubic will be available to test your solution.