

# Smarter dating moves and faster proposals: revisiting the phylogenetic relaxed clock model

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## Abstract

## Author summary

## Introduction

## Models

## Preliminaries

Let  $\mathcal{T}$  be a binary rooted time tree with  $N$  taxa (and  $2N - 2$  branches). Let  $L$  be the number of sites within the multiple sequence alignment  $D$ , and let  $L_{\text{eff}}$  be the *effective* number of sites in the alignment (ie. the number of site patterns). The posterior density of a phylogenetic model is described by

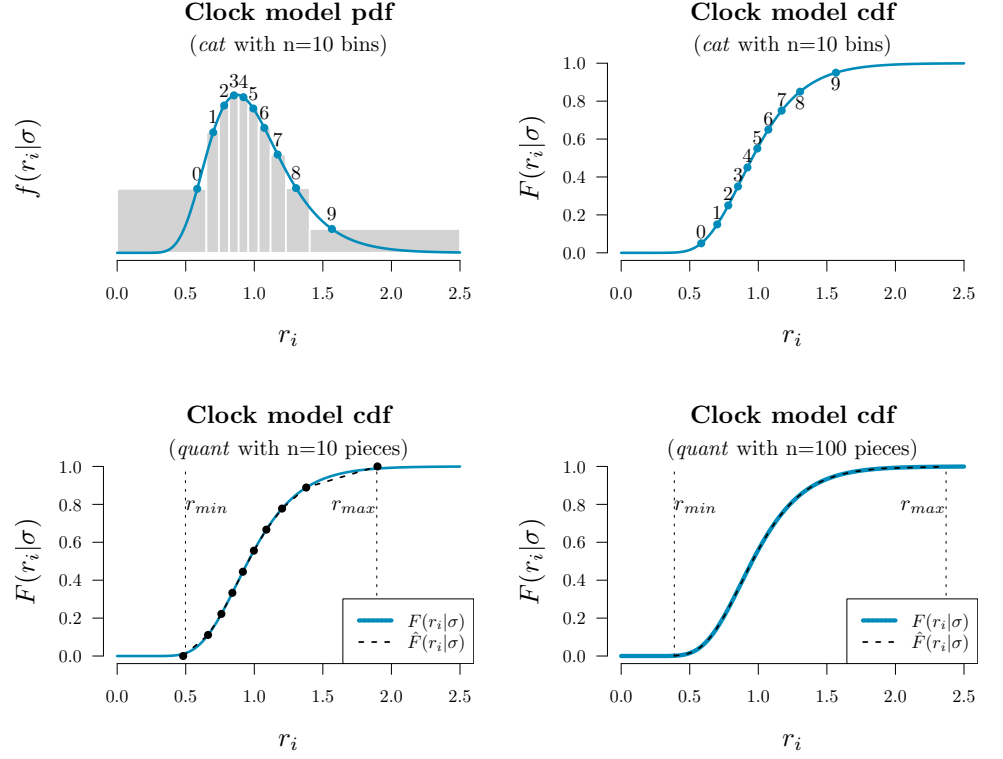
$$p(\mathcal{T}, \vec{\mathcal{R}}, \sigma, \theta | D) \propto p(D | \mathcal{T}, r(\vec{\mathcal{R}}), \theta) p(\mathcal{T} | \theta) p(\vec{\mathcal{R}} | \sigma) p(\sigma) p(\theta), \quad (1)$$

for rate standard deviation  $\sigma$  and other model parameters  $\theta$ .  $\vec{\mathcal{R}}$  is a vector of abstracted substitution rates, which is transformed into real rates by  $r(\vec{\mathcal{R}})$ . Three methods of representing rates as  $\vec{\mathcal{R}}$  are presented in **Rate parameterisations**.

Under the *relaxed clock model*, each internal and leaf node is assigned a substitution rate  $r_i = r(\mathcal{R}_i)$ , which corresponds to its parent branch. There are a total of  $|\vec{\mathcal{R}}| = 2N - 2$  rates, which are independently distributed under the relaxed clock model prior [1].

## Rate parameterisations

In Bayesian inference, the way parameters are represented in the model can affect the mixing ability of the model and the meaning of the model itself. Three methods for parameterising substitution rates are described below, and are later evaluated in **Results and Discussion**. Each parameterisation technique is associated with i) an abstraction of the rates  $\vec{\mathcal{R}}$ , ii) some function for transforming this parameter into real rates  $r(\vec{\mathcal{R}})$ , and iii) a prior density function of the abstraction  $p(\vec{\mathcal{R}} | \sigma)$ . The three methods are summarised in **Fig 1**.



**Fig 1. Methods of rate parameterisation.** The *cat* and *quant* approximations are plotted on top of the true underlying rate prior distribution (*real*). In this example, rates are drawn from a  $\text{LogNormal}(\mu = -0.045, \sigma = 0.3)$  distribution. The probability density function (pdf) and cumulative density function (cdf) of this distribution are shown.

## 1. Real rates

The natural (and unabridged) parameterisation of a substitution rate is a real number  $\mathcal{R}_i \in \mathbb{R}, \mathcal{R}_i > 0$  which is equal to the rate itself. Thus, under the *real* parameterisation:

$$r(\vec{\mathcal{R}}) = \vec{\mathcal{R}}. \quad (2)$$

Under the prior distribution  $p(\vec{\mathcal{R}}|\sigma)$ , rates are often log-normally or exponentially distributed with a mean of 1:

$$p(\mathcal{R}_i|\sigma) = \frac{1}{\mathcal{R}_i\sigma\sqrt{2\pi}} \exp\left(-\frac{(\ln \mathcal{R}_i - \mu)^2}{2\sigma^2}\right) \quad (\text{LogNormal}(\mu, \sigma)), \text{ or} \quad (3)$$

$$p(\mathcal{R}_i|\sigma) = p(\mathcal{R}_i) = e^{-\mathcal{R}_i} \quad (\text{Exponential}(\lambda = 1)) \quad (4)$$

where  $\mu = -0.5\sigma^2$  is set such that the expected value of the log-normal distribution is 1.

Zhang and Drummond 2020 present a series of tree operators which propose internal/root node heights, and then recompute the rates of incident branches such that their genetic distances ( $r_i \times \tau_i$ ) remain constant after the proposal. By maintaining genetic distances the likelihood can also be maintained. These operators account for the

correlation which exists between branch rates and branch times – a correlation which is induced by the likelihood function.

## 2. Categories

The category parameterisation (*cat*) is an abstraction of the *real* parameterisation. Each branch is assigned an integer from 0 to  $n - 1$ :

$$\vec{\mathcal{R}} \in \{0, 1, \dots, n - 1\}^{2N-2}. \quad (5)$$

The domain of  $\vec{\mathcal{R}}$  is uniformly distributed:

$$p(\mathcal{R}_i|\sigma) = p(\mathcal{R}_i) = \frac{1}{n}. \quad (6)$$

Let  $f(x|\sigma)$  be the probability density function (pdf) and let  $F(x|\sigma) = \int_0^x f(t|\sigma) dt$  be the cumulative distribution function (cdf) of the prior distribution used by the underlying *real* clock model. Then, in the *cat* parameterisation,  $f(x|\sigma)$  is discretised into  $n$  bins and the elements of  $\vec{\mathcal{R}}$  each point to one of these bins. Each bin contains uniform probability density  $1/n$ . The rate of each bin is equal to the median value within the bin

$$r(\mathcal{R}_i) = F^{-1}\left(\frac{\mathcal{R}_i + 0.5}{n}\right), \quad (7)$$

where  $F^{-1}$  is the inverse cumulative distribution function (i-cdf).

The key advantage of the *cat* parameterisation is the removal of a term from the posterior density (Equation 1), or more accurately the replacement of a non-trivial  $p(\vec{\mathcal{R}}|\sigma)$  term with that of a uniform prior. Thus, one fewer term needs to be estimated per rate.

This method was suggested in the original BEAST2 relaxed clock paper [1] and has been widely used. However, the constant distance operators since introduced by Zhang and Drummond 2020 – which are incompatible with the *cat* parameterisation – yield an increase in mixing rate under *real* by up to an order of magnitude over that of *cat*.

## 3. Quantiles

Finally, rates can be parameterised as real numbers  $0 < \mathcal{R}_i < 1$  which describe the rate's quantile with respect to some underlying clock model distribution. Under the *quant* parameterisation, each element in  $\vec{\mathcal{R}}$  is uniformly distributed.

$$\vec{\mathcal{R}} \in \mathbb{R}^{2N-2}, 0 < \mathcal{R}_i < 1 \quad (8)$$

$$p(\mathcal{R}_i|\sigma) = p(\mathcal{R}_i) = 1 \quad (9)$$

Transforming these quantiles into rates invokes the i-cdf of the underlying *real* clock model distribution. Thus, while this approach has clear similarities with *cat*, the domain of rates here is continuous (as opposed to being confined to a discrete number of bins) and is therefore compatible with the class of operators described by Zhang and Drummond 2020.

A potential disadvantage of the *quant* method would be the computational requirements of continuously evaluating the i-cdf, especially for trees with large  $N$ .

Hence, rather than evaluating the exact i-cdf  $F^{-1}$ , an approximation  $\hat{F}^{-1}$  will be used instead:

$$r(\mathcal{R}_i) = \hat{F}^{-1}(\mathcal{R}_i). \quad (10)$$

In this article we have extended *quant* through a linear piecewise approximation of the i-cdf. As the piecewise approximation is linear, evaluating the derivatives  $\frac{\partial}{\partial \mathcal{R}_i} \hat{F}^{-1}(\mathcal{R}_i) = D\hat{F}^{-1}(\mathcal{R}_i)$  and  $\frac{\partial}{\partial r_i} \hat{F}(r_i) = D\hat{F}(r_i)$  – which are required for computing Hastings ratios – is trivial. The approximation is comprised of  $n$  pieces (where  $n$  is fixed) and upper and lower rate boundaries  $r_{\min}$  and  $r_{\max}$ . The approximation is displayed in **Fig 1** and further detailed in **S1 Appendix**.

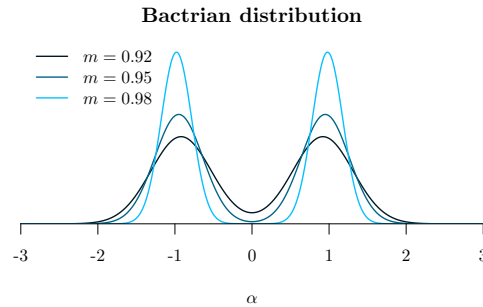
Zhang and Drummond 2020 introduced several tree operators for the *real* parameterisation – including **Constant Distance**, **Simple Distance**, and **Small Pulley**. In this project, we extended these three operators so that they are compatible with the *quant* parameterisation. These are presented in **S1 Appendix**.

## Bactrian proposal kernel

The step size of a proposal kernel  $q(x'|x)$  should be such that the proposed state  $x'$  is sufficiently far from the current state  $x$  to explore vast areas of parameter space, but not so large that the proposal is rejected too often [2]. Yang et al. have challenged the widely used uniform proposal kernel in place of the Bactrian kernel [3, 4]. The Bactrian( $m$ ) distribution is defined as the sum of two Normal distributions:

$$\Sigma \sim \text{Bactrian}(m) \equiv \frac{1}{2}\text{Normal}(-m, 1 - m^2) + \frac{1}{2}\text{Normal}(m, 1 - m^2) \quad (11)$$

where  $0 \leq m < 1$  describes the modality of the Bactrian distribution. When  $m = 0$ , the Bactrian distribution is equivalent to a Normal(0, 1) distribution. As  $m \rightarrow 1$ , the distribution becomes increasingly bimodal (**Fig. 2**). Yang et al. 2013 [3] suggest that Bactrian( $m = 0.95$ ) yields a proposal kernel superior to the uniform kernel, by placing minimal probability on steps which are too small or too large.



**Fig 2. The Bactrian proposal kernel.** Y-axis corresponds to probability density  $f(\Sigma|m)$ .

In this article we compare the performance of uniform and Bactrian proposal kernels in the clock model. Two Bactrian distributions are compared ( $m = 0.95$  and  $m = 0.98$ ). The clock model operators which these proposal kernels apply to are described in **Table 1**.

	Operator	Proposal	Parameters
1	<b>Random walk operator</b>	$x' \leftarrow x + s\Sigma$	$\sigma, r, q$
2	<b>Scale operator</b>	$x' \leftarrow x \times e^{s\Sigma}$	$\sigma, r$
3	<b>Interval operator</b>	$y \leftarrow \frac{u-x}{x-l} \times e^{s\Sigma}$ $x' \leftarrow \frac{u+l*y}{y+1}$	$q \ (l = 0, u = 1)$
4	<b>Constant distance operators</b>	$x' \leftarrow x + s\Sigma$	$t$

**Table 1.** Summary of proposal kernels  $q(x'|x)$  of clock model operators. In each operator,  $\Sigma$  is drawn from either a Bactrian( $m$ ) or Uniform distribution (distributions are normalised so that they have a mean of 0 and a variance of 1). The scale size  $s$  is tunable. The proposal kernel may apply to node heights  $t$ , clock standard deviation  $\sigma$ , clock rates  $r$  (*real* only), and clock rate quantiles  $q$  (*quant* only). The Scale operator acts on parameters with non-negative domains. The Interval operator proposes values which respect its domain ie.  $l < x' < u$ .

## Narrow Exchange Rate

The **Narrow Exchange** operator [5], widely used in BEAST [6] and BEAST2 [7], is similar to NNI, and works as follows (**Fig. 3**):

*Step 1.* Sample an internal/root node  $E$  from tree  $\mathcal{T}$ , where  $E$  has grandchildren.

*Step 2.* Identify the child of  $E$  with the greater height. Denote this child as  $D$  and its sibling as  $C$  (ie.  $t_D > t_C$ ).

*Step 3.* Randomly identify the two children of  $D$  as  $A$  and  $B$ .

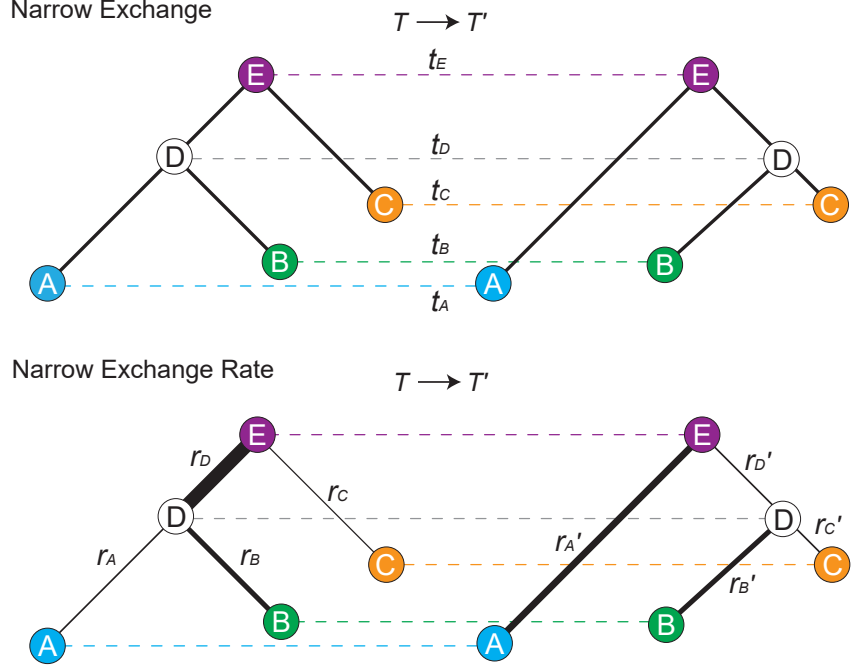
*Step 4.* Relocate the  $B - D$  branch onto the  $C - E$  branch, so that  $B$  and  $C$  become siblings and their parent is  $D$ . All node heights are unchanged.

Lakner et al. 2008 [8] found that tree operators which perturb topology (such as NNI and SPR) consistently perform better than those which also change branch lengths (such as LOCAL [9] and Continuous Change [10]). If Narrow Exchange was adapted to the relaxed clock model by ensuring that genetic distances remain constant after the proposal, its performance may be improved even further. This may in turn permit proposing a new node height  $t_D$  and therefore changing branch (time) lengths.

Here, we present the **Narrow Exchange Rate** (NER) operator. Let  $r_A, r_B, r_C$ , and  $r_D$  be the clock rates of nodes  $A, B, C$ , and  $D$ , respectively. In addition to the modest topological change applied by Narrow Exchange, NER also proposes new clock rates  $r_A', r_B', r_C'$ , and  $r_D'$ . While NER does not alter  $t_D$  (ie.  $t_D' \leftarrow t_D$ ), we also consider NERw - a special case of the NER operator which embarks  $t_D$  on a random walk:

$$t_D' \leftarrow t_D + s\Sigma \quad (12)$$

for random walk step size  $s\Sigma$  where  $s$  is a tunable scalar parameter and  $\Sigma$  is drawn from a uniform or **Bactrian proposal kernel**. NER (and NERw) are compatible with both the *real* and *quant* parameterisations. Analogous to the Constant Distance operator, new rates are proposed such that genetic distances between nodes  $A, B, C$ , and  $E$  are maintained. Thus, there are  $\binom{4}{2} = 6$  pairwise distance constraints.



**Fig 3. Depiction of Narrow Exchange and Narrow Exchange Rate operators.** Proposals are denoted by  $\mathcal{T} \rightarrow \mathcal{T}'$ . The vertical axis corresponds to node height  $t$ . In the bottom figure, branch rates  $r$  are indicated by line thickness. In this example, the  $\mathcal{D}_{AE}$  and  $\mathcal{D}_{CE}$  constraints are satisfied.

$$\mathcal{D}_{AB} : \begin{aligned} r_A(t_D - t_A) + r_B(t_D - t_B) = \\ r_A'(t_E - t_A) + r_D'(t_E - t_D') + r_B'(t_D' - t_B) \end{aligned} \quad (13)$$

$$\mathcal{D}_{AC} : \begin{aligned} r_A(t_D - t_A) + r_D(t_E - t_D) + r_C(t_E - t_C) = \\ r_A'(t_E - t_A) + r_D'(t_E - t_D') + r_C'(t_D' - t_C) \end{aligned} \quad (14)$$

$$\mathcal{D}_{AE} : \begin{aligned} r_A(t_D - t_A) + r_D(t_E - t_D) = \\ r_A'(t_E - t_A) \end{aligned} \quad (15)$$

$$\mathcal{D}_{BC} : \begin{aligned} r_B(t_D - t_B) + r_D(t_E - t_D) + r_C(t_E - t_D) = \\ r_B'(t_D' - t_B) + r_C'(t_D' - t_C) \end{aligned} \quad (16)$$

$$\mathcal{D}_{BE} : \begin{aligned} r_B(t_D - t_B) + r_D(t_E - t_D) = \\ r_B'(t_D' - t_B) + r_D'(t_E - t_D') \end{aligned} \quad (17)$$

$$\mathcal{D}_{CE} : \begin{aligned} r_C(t_E - t_C) = \\ r_C'(t_D' - t_C) + r_D'(t_E - t_D') \end{aligned} \quad (18)$$

Further constraints are imposed by the model itself:

$$r_i > 0 \text{ and } r_i' > 0 \text{ for } i \in \{A, B, C, D\} \quad (19)$$

$$\max\{t_B, t_C\} < t_D' < t_E. \quad (20)$$

Unfortunately, it is not possible to solve all six  $\mathcal{D}_{ij}$  constraints without permitting non-positive rates or illegal trees. Therefore rather than conserving all six pairwise

distances, NER conserves a *subset* of distances. It is not immediately clear which distances should be conserved.

### Automated generation of operators and constraint satisfaction

The total space of NER operators is comprised of all possible subsets of distance constraints (ie.  $\{\}, \{\mathcal{D}_{AB}\}, \{\mathcal{D}_{AC}\}, \dots, \{\mathcal{D}_{AB}, \mathcal{D}_{AC}, \mathcal{D}_{AE}, \mathcal{D}_{BC}, \mathcal{D}_{BE}, \mathcal{D}_{CE}\}$ ) which are solvable. The simplest NER – the null operator denoted by  $\text{NER}\{\}$  – does not satisfy any distance constraints. This is equivalent to Narrow Exchange.

As it is unclear which NER variants would perform the best, we developed an automated pipeline for generating and testing these operators.

**1. Solution finding.** Using standard analytical linear-system solving libraries in MATLAB, the  $2^6 = 64$  subsets of distance constraints are solved. 54 out of the 64 subsets were found to be solvable, and the unsolvables were discarded.

**2. Solving Jacobian determinants.** The determinant of the Jacobian matrix  $J$  is required for computing the Hastings ratio of the proposal.  $J$  is defined as

$$J = \begin{bmatrix} \frac{\partial r_A'}{\partial r_A} & \frac{\partial r_A'}{\partial r_B} & \frac{\partial r_A'}{\partial r_C} & \frac{\partial r_A'}{\partial r_D} \\ \frac{\partial r_B'}{\partial r_A} & \frac{\partial r_B'}{\partial r_B} & \frac{\partial r_B'}{\partial r_C} & \frac{\partial r_B'}{\partial r_D} \\ \frac{\partial r_C'}{\partial r_A} & \frac{\partial r_C'}{\partial r_B} & \frac{\partial r_C'}{\partial r_C} & \frac{\partial r_C'}{\partial r_D} \\ \frac{\partial r_D'}{\partial r_A} & \frac{\partial r_D'}{\partial r_B} & \frac{\partial r_D'}{\partial r_C} & \frac{\partial r_D'}{\partial r_D} \end{bmatrix}. \quad (21)$$

Computing the determinant  $|J|$  invokes standard analytical differentiation and linear algebra libraries of MATLAB. 6 of the 54 operators were found to have  $|J| = 0$ , corresponding to irreversible proposals, and were discarded.

**3. Automated generation of BEAST2 operators.** Java class files are generated using string processing. Each class corresponds to a single operator, extends the class of a meta-NER-operator, and is comprised of the solutions found in **1** and the Jacobian determinant found in **2**.  $|J|$  is further augmented if the *quant* parameterisation is employed.

The 48 operators generated by this pipeline are evaluated and compared in **Results and Discussion**. Each operator is considered with and without a random walk on  $t_D$  and thus there are 96 total settings.

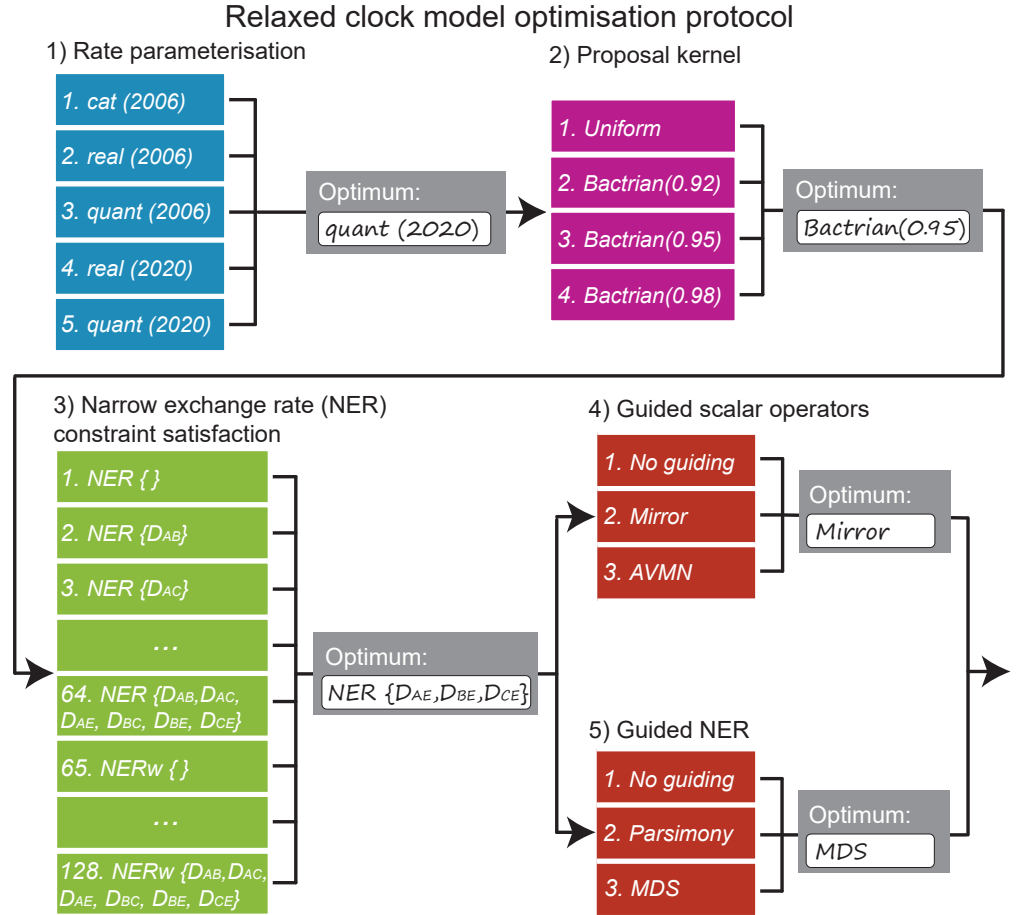
## Clock model averaging

## Results and Discussion

### Assessment criteria and datasets

To avoid a cross-product explosion, the three targets for clock model improvement (**Rate** parameterisations, **Bactrian proposal kernel**, and **Narrow Exchange Rate**) are evaluated sequentially, in the order presented in this paper. The setting(s) which are considered to be the best in each step are then incorporated into the following step. This protocol and its outcomes are summarised in **Fig. 4**.

Methodologies are assessed according to the following criteria.



**Fig 4. Protocol for optimising methodology settings.** The three areas (detailed in **Models**) are optimised sequentially, where the best setting from each step is used when optimising the following step.

**1. Validation.** This is assessed by measuring the coverage of all estimated parameters in a well-calibrated simulation study, using 100 simulated datasets (with  $N = 100$  taxa and  $L = 5000$  nucleotide alignments). These are presented in **S2 Appendix**.

**2. Time to convergence.** Two independent MCMC chains are run and the time is measured until: a) the absolute difference in clade posterior probability between the two chains is less than 0.05 for all clades, b) the Rubin-Gelman statistic  $\hat{R}$  [11] of every estimated parameter is less than 1.05, and c) the effective sample size [12] of every estimated parameter is greater than 100 in each chain.

**3. Mixing of parameters.** Key parameters are evaluated for the number of effective samples generated per hour (ESS/hr).

For the latter two criteria, methodologies are benchmarked using both simulated and empirical datasets – the latter were compiled [13] and partitioned [14] by Lanfear et al. as ‘benchmark alignments’. These are: the spider data by Richart et al. 2015 [15]; the turtle data by Crawford et al. 2012 [16]; the caterpillar data by Kawahara et al. 2013 [17]; and the songbird data by Moyle et al. 2016 [18]. Each setting is run  $10 \times$  per dataset and the average statistics across the  $10 \times 2 = 20$  chains are reported.

Methodologies are benchmarked using the Intel(R) Xeon(R) Gold 6138 CPU (2.00



	$N$	$P$	$L$ (kb)	$L_{\text{eff}}$ (kb)	Description
1	6	16	13.7	1.1	Spiders (Richart 2015 [15])
2	10	16	6.7	0.9	Turtles (Crawford 2012 [16])
3	40	1	5.0	2.2-2.6	Simulated data $\times 10$
4	70	3	2.2	0.9	Caterpillars (Kawahara 2013 [17])
5	106	1	0.6	0.3	Songbirds (Moyle 2016 [18])

**Table 2.** Datasets used during benchmarking, sorted in increasing order of taxa count  $N$ . Number of partitions  $P$ , total alignment length  $L$ , and number of patterns  $L_{\text{eff}}$  are also specified. Empirical datasets are benchmarked  $10\times$  on each setting, while the 10 simulated datasets are benchmarked once on each setting.

## Comparison of rate parameterisations

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We compared the three rate parameterisations described in **Rate parameterisations**. All three settings use the standard BEAST2 clock model operators from Drummond et al. 2006 [1]. *real* and *quant* additionally use the constant-distance tree operators described by Zhang and Drummond 2020. To determine whether the difference in performance between *real/quant* versus *cat* is because of the constant-distance tree operators or the parameterisation itself, we also included benchmarked two additional settings: *real 2006* and *quant 2006*, which do not use the constant-distance operators. These five settings are validated in **S2 Appendix**.

**Fig. 5** shows that the *real 2006* performs considerably worse than any of the other settings. This is due to the poor sampling of the prior under this setting (ie. low ESS of  $p(\theta)$ ). The failure of *real 2006* thus highlights the appeal of 1) the *cat* or *quant* parameterisations, both of which have trivial contributions to the prior density (ie. uniform priors), and 2) the smarter operators used by *real* (Zhang and Drummond 2020). Due to its computational burden, *real 2006* was not benchmarked for all of the datasets in **Table 2**.

Our results show that the *quant* parameterisation yields the best performance with respect to effective samples per hour. *quant* outperforms *quant 2006*, suggesting that the constant distance operators are effective. Furthermore, *quant* outperforms *real* especially at sampling from the posterior and prior distributions (ie. high ESS/hr for  $p(\theta|D)$  and  $p(\theta)$ ). This is most likely because of the uniform prior distribution of rate quantiles.

Overall, *quant* yields a median ESS/hr approximately 150 % faster with respect to sampling the posterior probability, and approximately 30 % faster with respect to sampling branch rates  $r$  and clock standard deviation  $\sigma$ , compared with *real*.

**Fig 5. Rate parameterisation evaluation.** Comparison of ESS/hr (averaged across two independent MCMC chains) with respect to relevant terms –  $L$ : likelihood,  $p$ : prior density,  $r$ : clock rate averaged across all leaves,  $\sigma$ : clock standard deviation. Each point is from one partition sample of the empirical data in **Table 2**.

## Comparison of Bactrian and uniform proposal kernels on the clock model

### Comparison of NER variants

We compared the Narrow Exchange Rate (NER) operators described in **Narrow Exchange Rate**. This protocol selects the best among 48 NER (no random walk) and 48 NERw (Bactrian(0.95) random walk) operators, and has two phases. First, the best of the 96 is selected by comparing operator acceptance rates on simulated data. Second, the selected operator is benchmarked with respect to convergence time and sampling rate on real data (**Table 2**). All the analyses in this section invoke the *quant* parameterisation and Bactrian(0.95) proposal kernels on clock model parameters.

#### Initial screening by acceptance rate

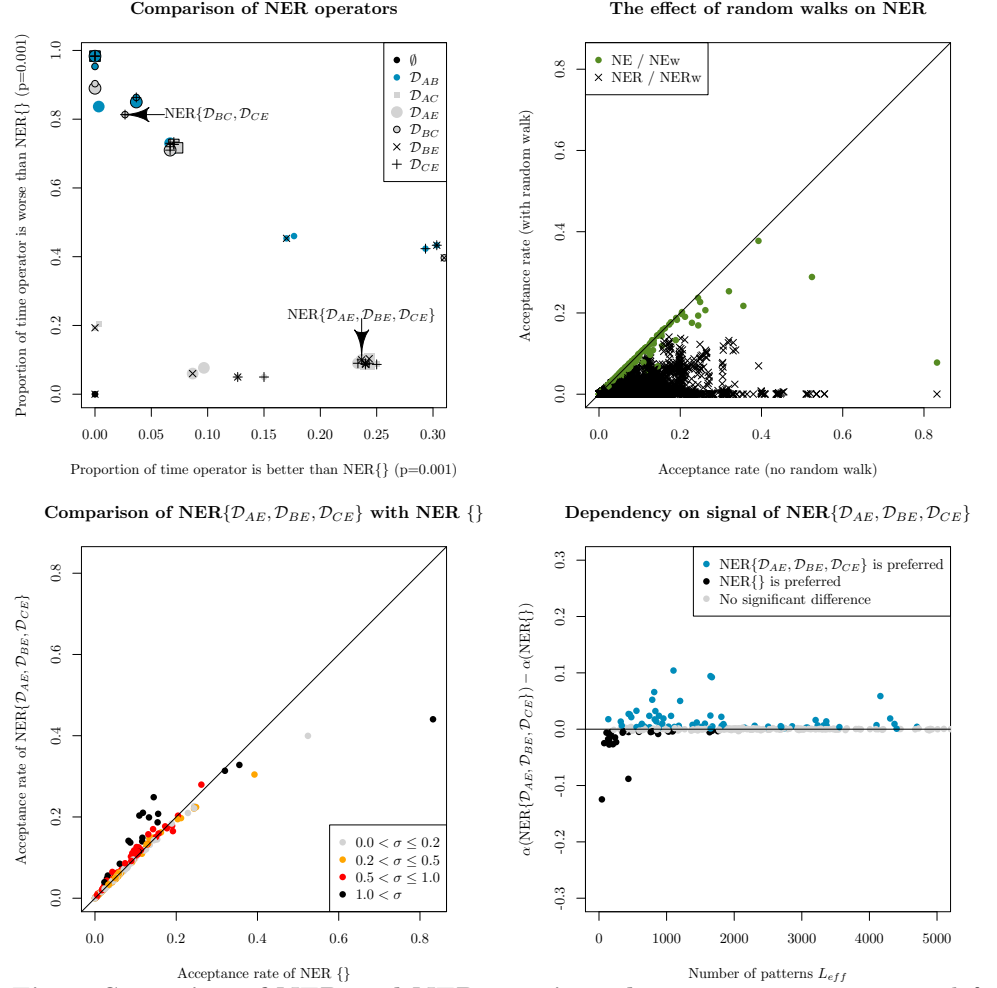
We selected the best operator variant by performing MCMC on 300 simulated datasets, where each MCMC employed all 96 NER/NERw variants. Simulated datasets have  $N = 30$  taxa and an alignment with  $L \sim \text{Uniform}(10^2, 10^4)$  sites. The acceptance rate of each operator is compared to that of the null operator NER{} (ie. Narrow Exchange).

**Fig. 6** shows that NER variants which satisfy the genetic distances between nodes  $B$  and  $A$  (ie.  $\mathcal{D}_{AB}$ ) or between  $B$  and  $C$  (ie.  $\mathcal{D}_{BC}$ ) usually perform worse than the standard Narrow Exchange operator, where  $B$  is the node being interchanged from the  $A$  branch to the  $C$  branch (**Fig. 3**). This is an intuitive result. If the posterior distribution is relatively flat, and the data presents high uncertainty in the positioning of  $B$ , with respect to  $A$  and  $C$ , then the topological rearrangement performed by Narrow Exchange will be favoured. However, this uncertainty in the *topology* is likely coupled with uncertainty in the *distance* between  $B$  and  $A$  or between  $B$  and  $C$ . Thus, in this case, respecting the  $\mathcal{D}_{AB}$  and  $\mathcal{D}_{BC}$  constraints (by proposing branch rates) makes too many unnecessary changes to the state and the operator performs worse.

**Fig. 6** also reveals a cluster of NER variants which – under the conditions of the simulation – performed better than the null operator NER{} around 25% of the time and performed worse around 10% of the time. One such operator is NER{ $\mathcal{D}_{AE}, \mathcal{D}_{BE}, \mathcal{D}_{CE}$ }. This variant conserves the genetic distance between all child nodes  $A$ ,  $B$ , and  $C$ , and the grandparent node  $E$ . This is performed by proposing rates for  $r_A$ ,  $r_B$ , and  $r_C$  while obeying the distance constraints imposed by the operator. Exploring this operator further, we can see that NER{ $\mathcal{D}_{AE}, \mathcal{D}_{BE}, \mathcal{D}_{CE}$ } is at its best when there is a large variance in branch rate ie. when clock standard deviation  $\sigma$  is high ( $\sigma \gtrsim 0.5$  for  $N = 30$ ), corresponding to data which is not clock-like. On the other hand, NER{} is much preferred when the operator’s acceptance rate is high ( $\gtrsim 0.15$ ) – corresponding with datasets with a small number of site patterns ( $L_{\text{eff}} \lesssim 500$  for  $N = 30$ ) and thus poor signal. Overall, NER{ $\mathcal{D}_{AE}, \mathcal{D}_{BE}, \mathcal{D}_{CE}$ } outperforms the standard Narrow Exchange operator when the data is not clock-like and contains enough signal.

Finally, **Fig. 6** shows that by applying a (Bactrian) random walk to  $t_D$  – the height of internal node  $D$  – the acceptance rate of NER plummets dramatically. This effect is most dominant for the NER variants which satisfy distance constraints (ie. the operators which are not NER{}). This result is unfortunate however not unexpected, and is consistent with Lakner et al. 2008 [8], who observed that tree operators perform best when they change either topology, or branch lengths, but not both.

Although there are several operators tying for first place, we selected the NER{ $\mathcal{D}_{AE}, \mathcal{D}_{BE}, \mathcal{D}_{CE}$ } operator to proceed to the next round of optimisation.



**Fig 6. Screening of NER and NERw variants by acceptance rate.** Top left: comparison of NER variants with the null operator NER{} (ie. Narrow Exchange). Each of the 48 operators are represented by a single point, uniquely encoded by the point stylings. The number of times each operator is proposed and accepted is compared with that of NER{}, and one-sided z-tests are performed to assess the statistical significance between the two acceptance rates ( $p = 0.001$ ). This process is repeated for each of 300 simulated datasets. The axes of each plot are the proportion of these 300 simulations for which there is evidence that the operator is better than NER{} (x-axis) or worse than NER{} (y-axis). Top right: comparison of NER and NERw acceptance rates. Each point is one NER/NERw variant from a single simulation. Bottom: relationship between the acceptance rates  $\alpha$  of NER $\{D_{AE}, D_{BE}, D_{CE}\}$  and NER{} with the clock model standard deviation  $\sigma$  and the number of patterns  $L_{eff}$ . Each point is a single simulation.

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