# Some Formulae Used in the Analysis of Algorithms

#### The Master Theorem

Suppose

- $r(n) = Cn^d + Ar(n/b)$  where A, B, and C are positive constants
- r(0) = B

Then

- If  $A < b^d$  then r(n) has  $O(n^d)$ ,  $\Theta(n^d)$ ,  $\Omega(n^d)$
- If  $A = b^d$  then r(n) has  $O(n^d \log n)$ ,  $\Theta(n^d \log n)$ ,  $\Omega(n^d \log n)$
- If  $A > b^d$  then r(n) has  $O(n^{\log_b A})$ ,  $\Theta(n^{\log_b A})$ ,  $\Omega(n^{\log_b A})$

### **Polynomial Series**

$$\sum_{i=1}^{n} i = \frac{n(n+1)}{2}$$

$$\sum_{i=1}^{n} i^2 = \frac{n(n+1)(2n+1)}{6}$$

$$\sum_{i=1}^{n} i^3 = \frac{n^2(n+1)^2}{4}$$

$$\sum_{i=1}^{n} i^4 = \frac{n(2n+1)(n+1)(3n^2+3n-1)}{30}$$

$$\sum_{i=1}^{n} i^{5} = \frac{n^{2}(2n^{2} + 2n - 1)(n+1)^{2}}{12}$$

#### Geometric Series

$$\sum_{i=0}^{n} a^{i} = \frac{a^{n+1} - 1}{a - 1} \text{ for } a \neq 1$$

## Logarithmic Series

$$\sum_{i=1}^{n} i^{k} \log i \text{ has } \Theta(n^{k+1} \log n) \text{ for } k = 0, 1, 2, \dots$$