

Quantum Computing and Shor's Algorithm

Applying Quantum Mechanics to Computation

A Brief Introduction

- What is a Quantum Computer?

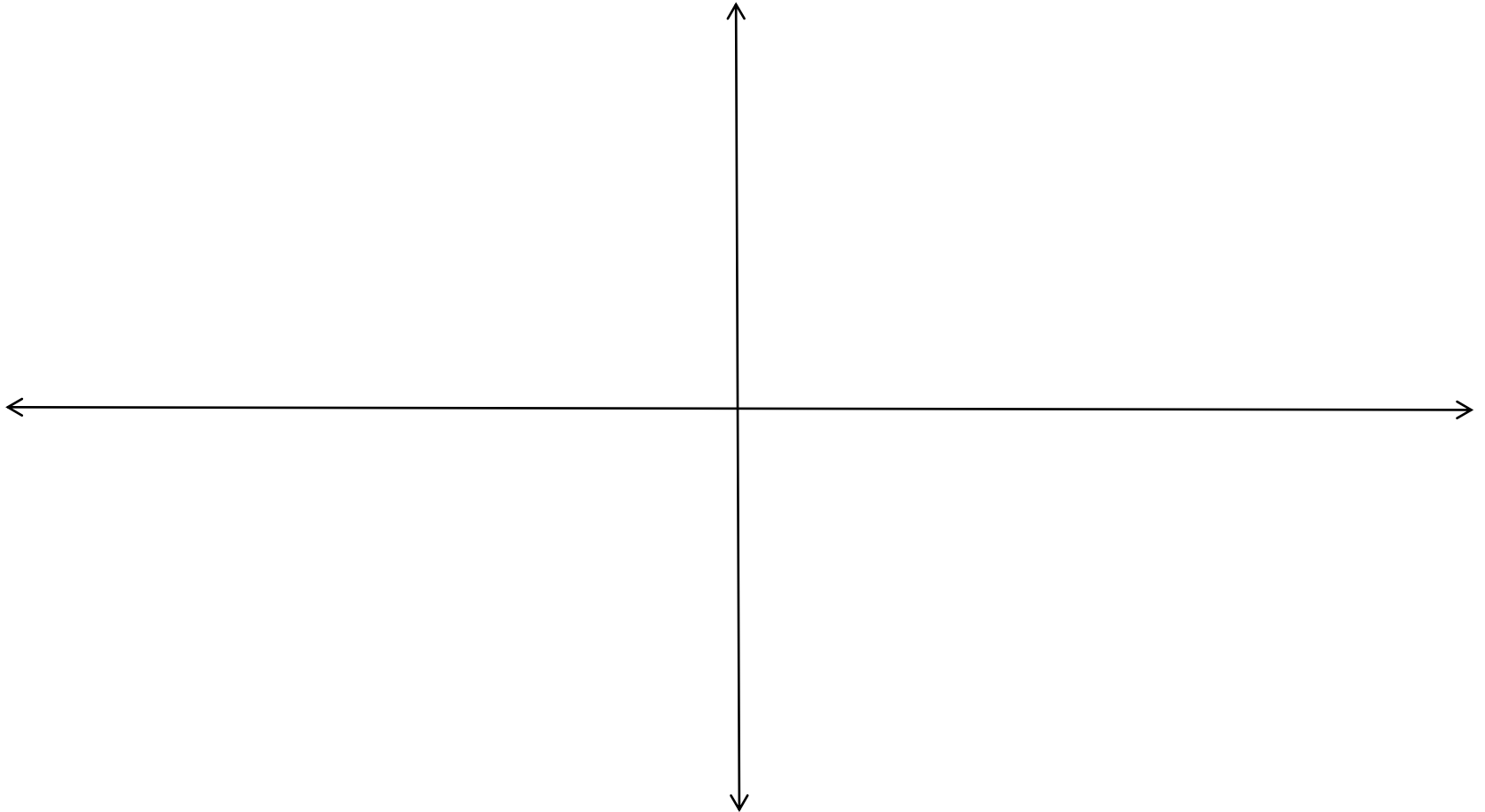
A Brief Introduction

- What is a Quantum Computer?
- We don't know where an electron is (until we observe it)

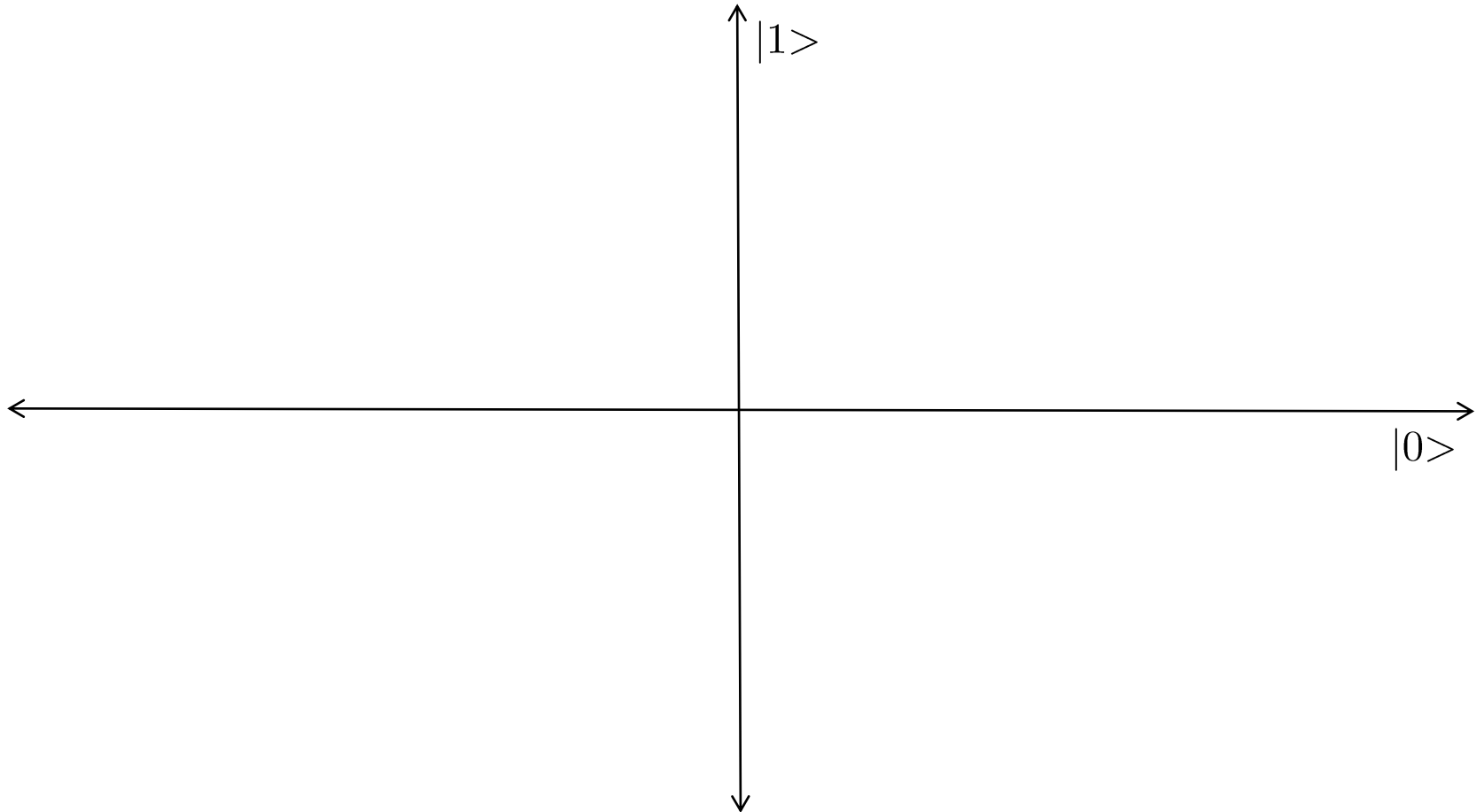
A Brief Introduction

- What is a Quantum Computer?
- We don't know where an electron is (until we observe it)
- Amplitudes – square to find probabilities

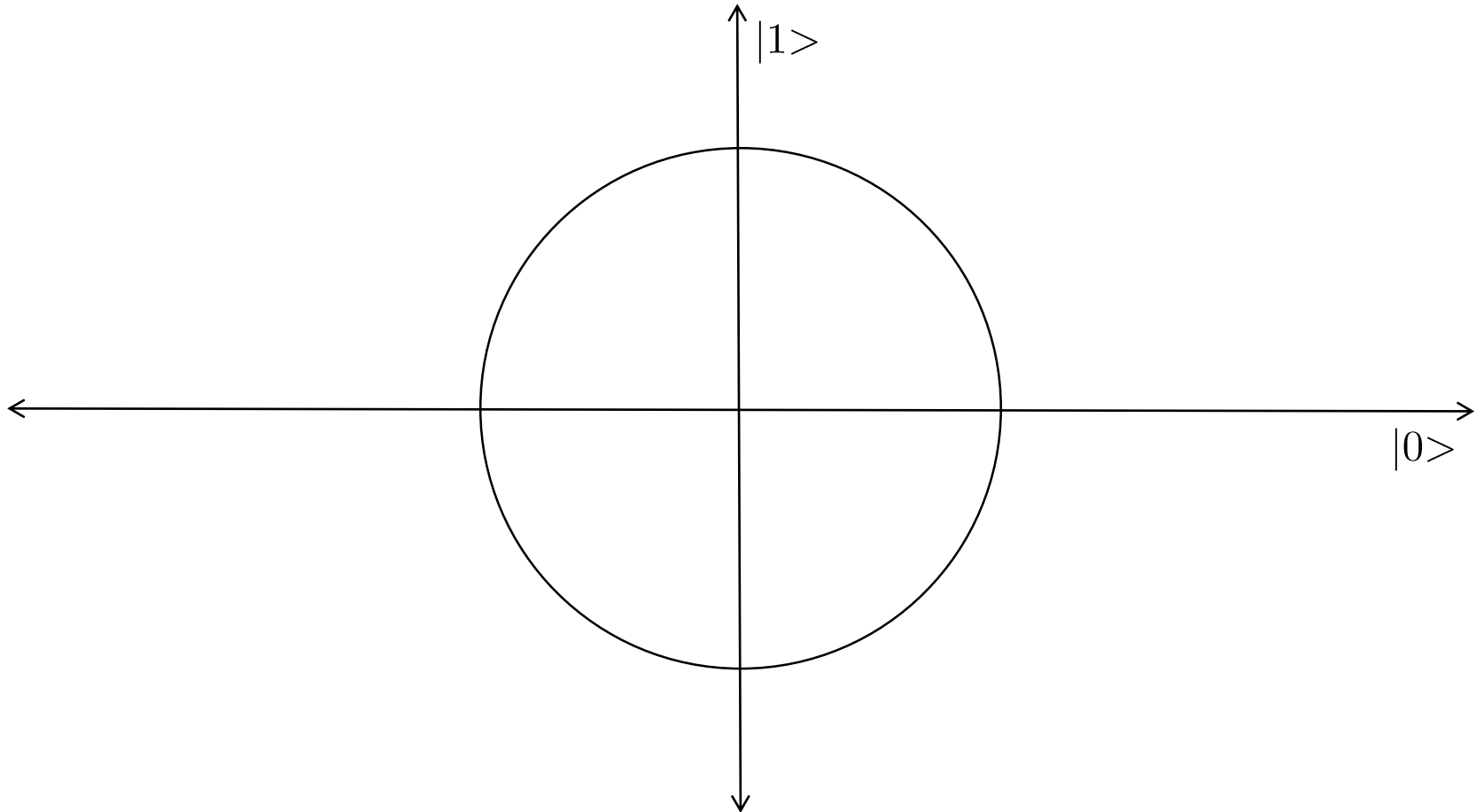
Quantum States and Dirac Notation



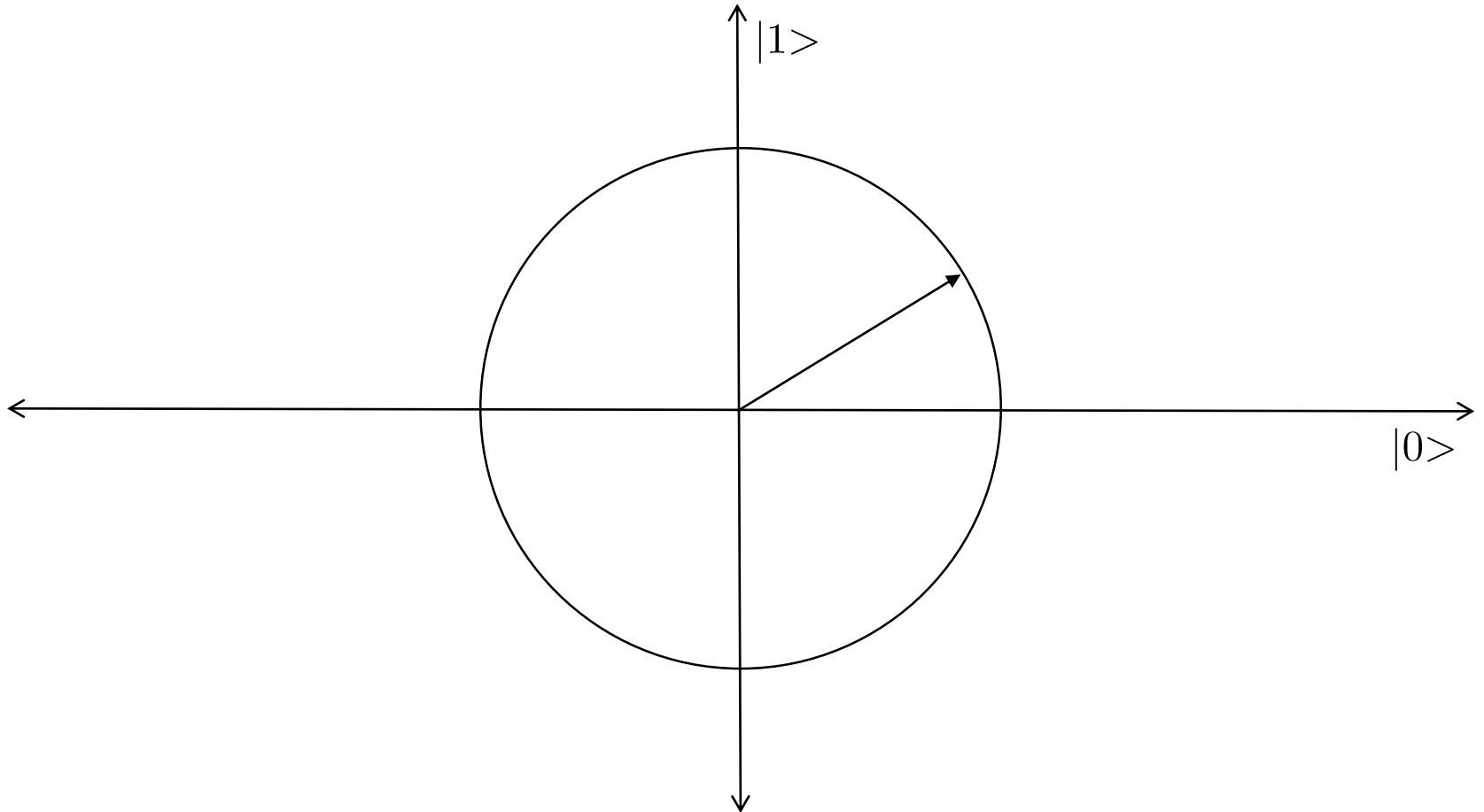
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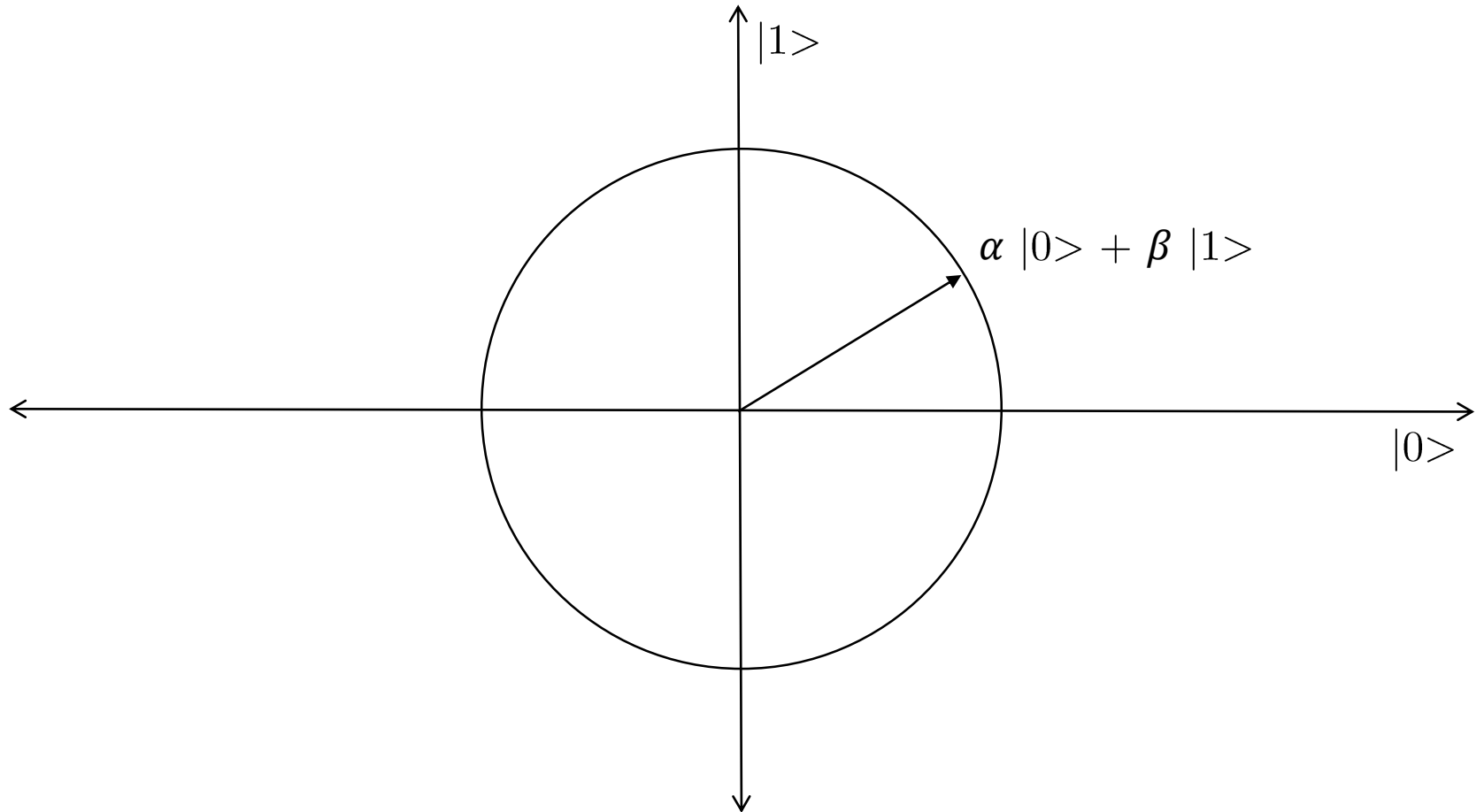
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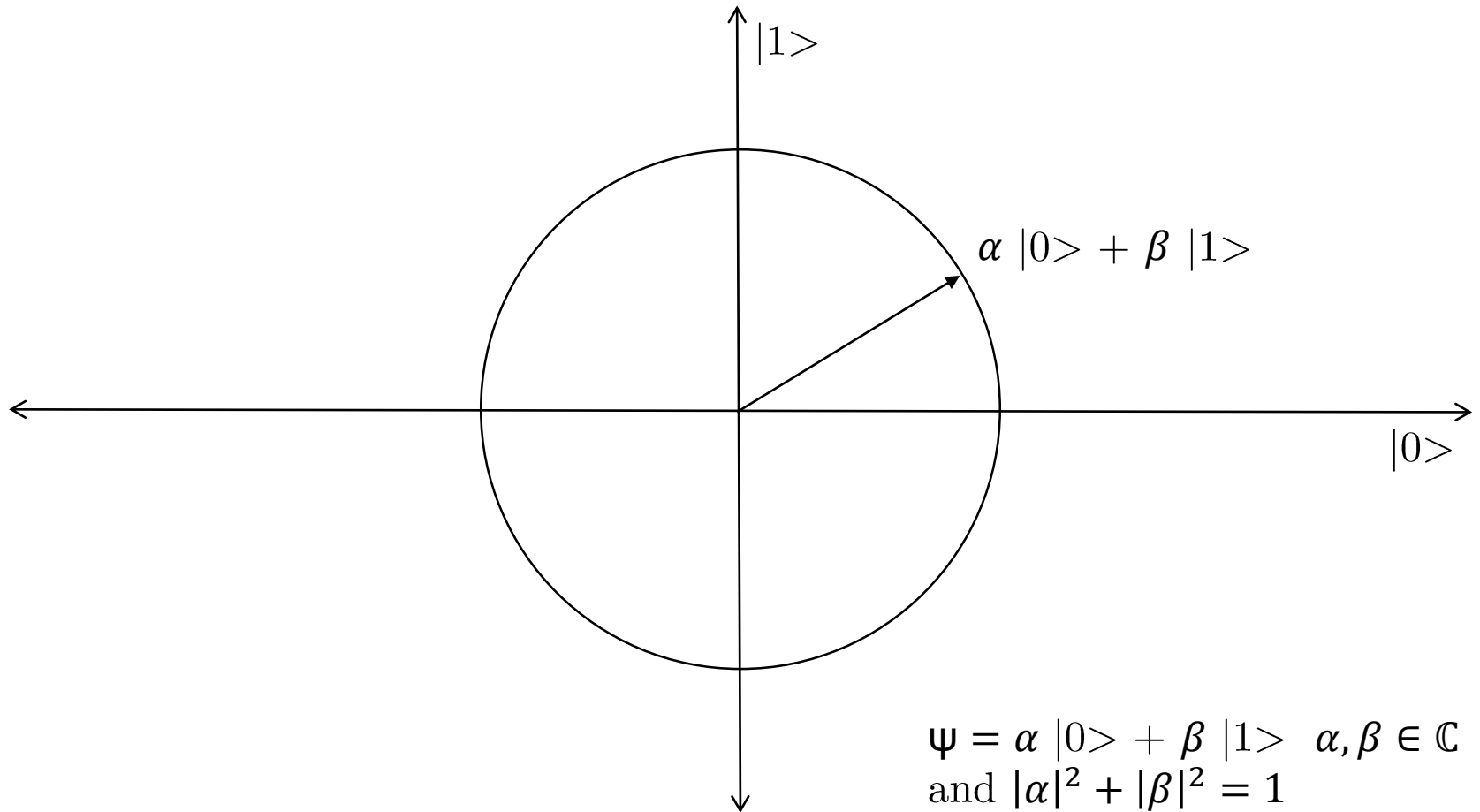
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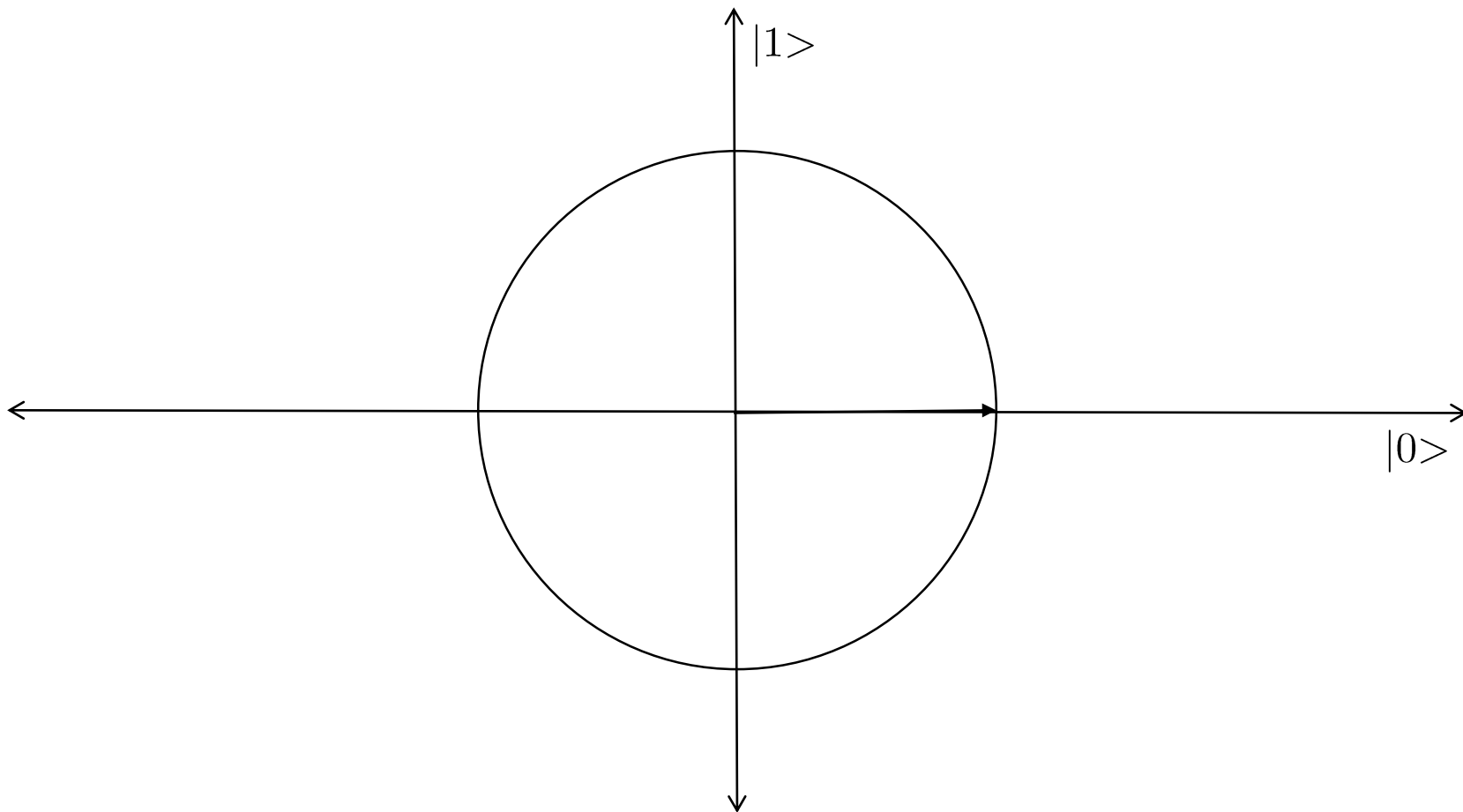


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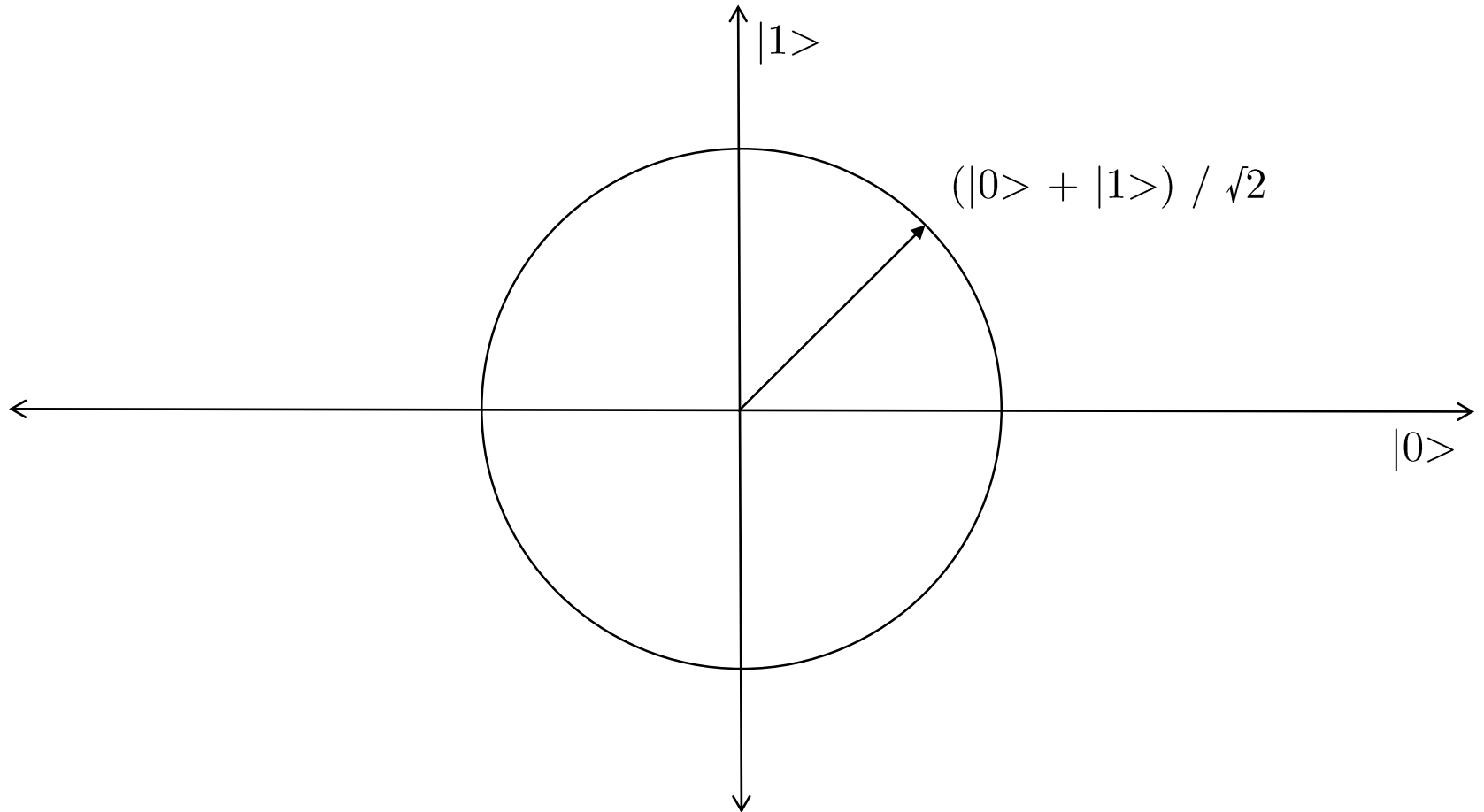


Intrinsic Randomness

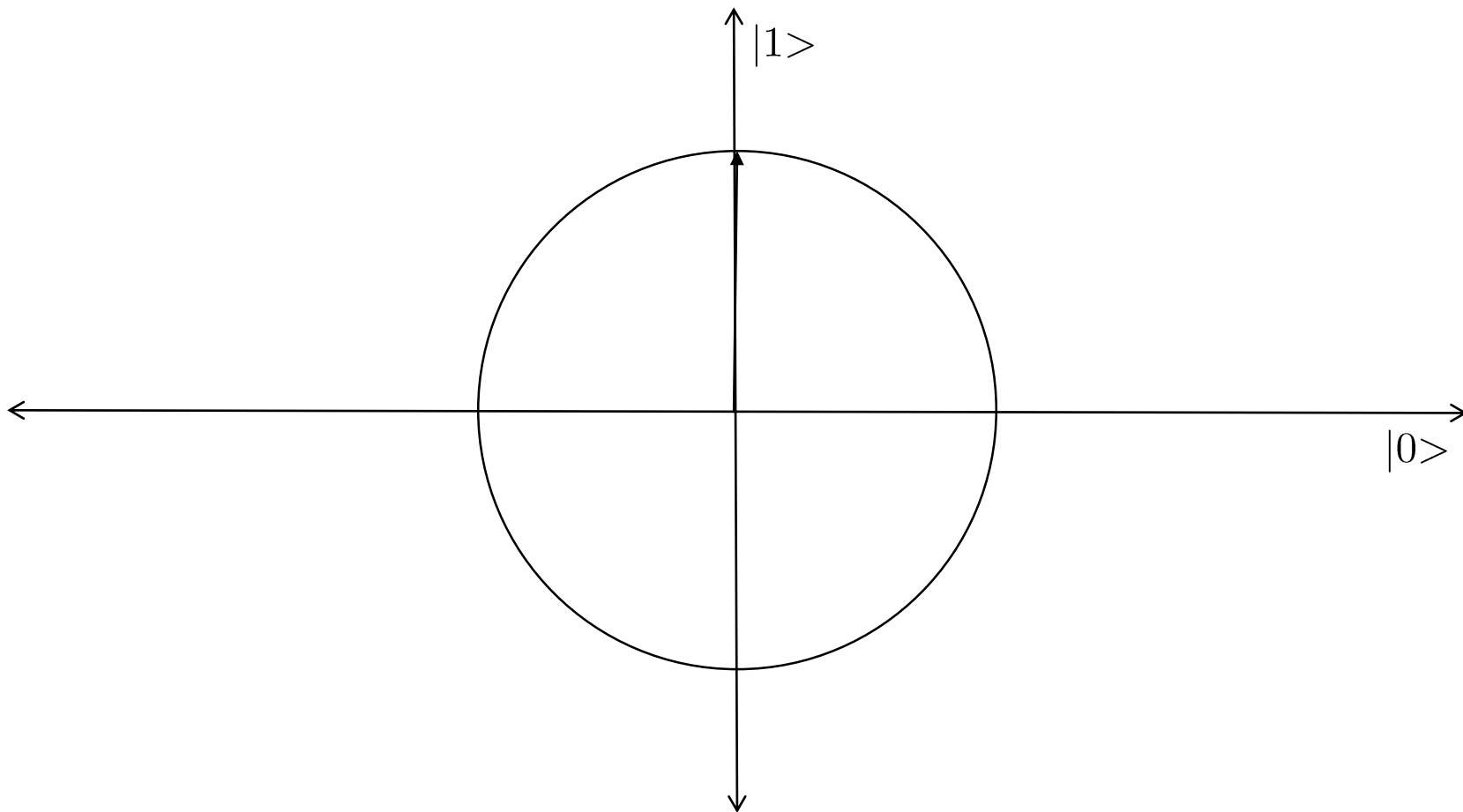
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Interference

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$$((|0\rangle + |1\rangle - |0\rangle + |1\rangle) / \sqrt{2}) / \sqrt{2}$$

$$= ((2|1\rangle) / \sqrt{2}) / \sqrt{2} = |1\rangle \text{ as required}$$

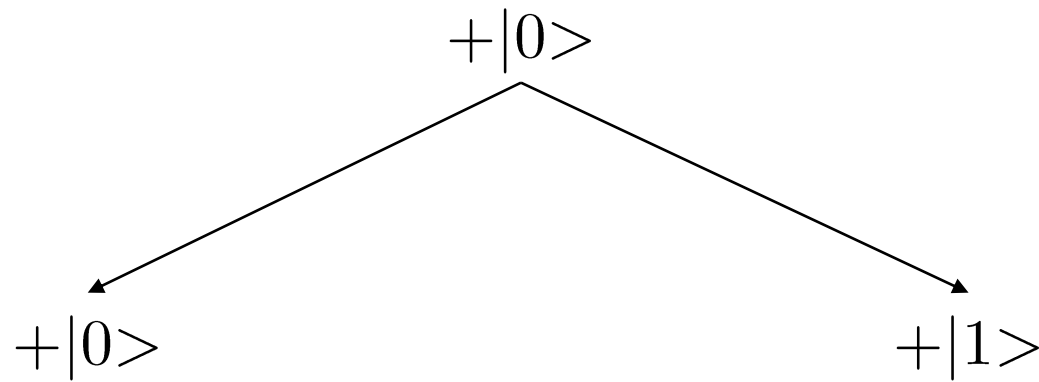
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$$+|0\rangle$$

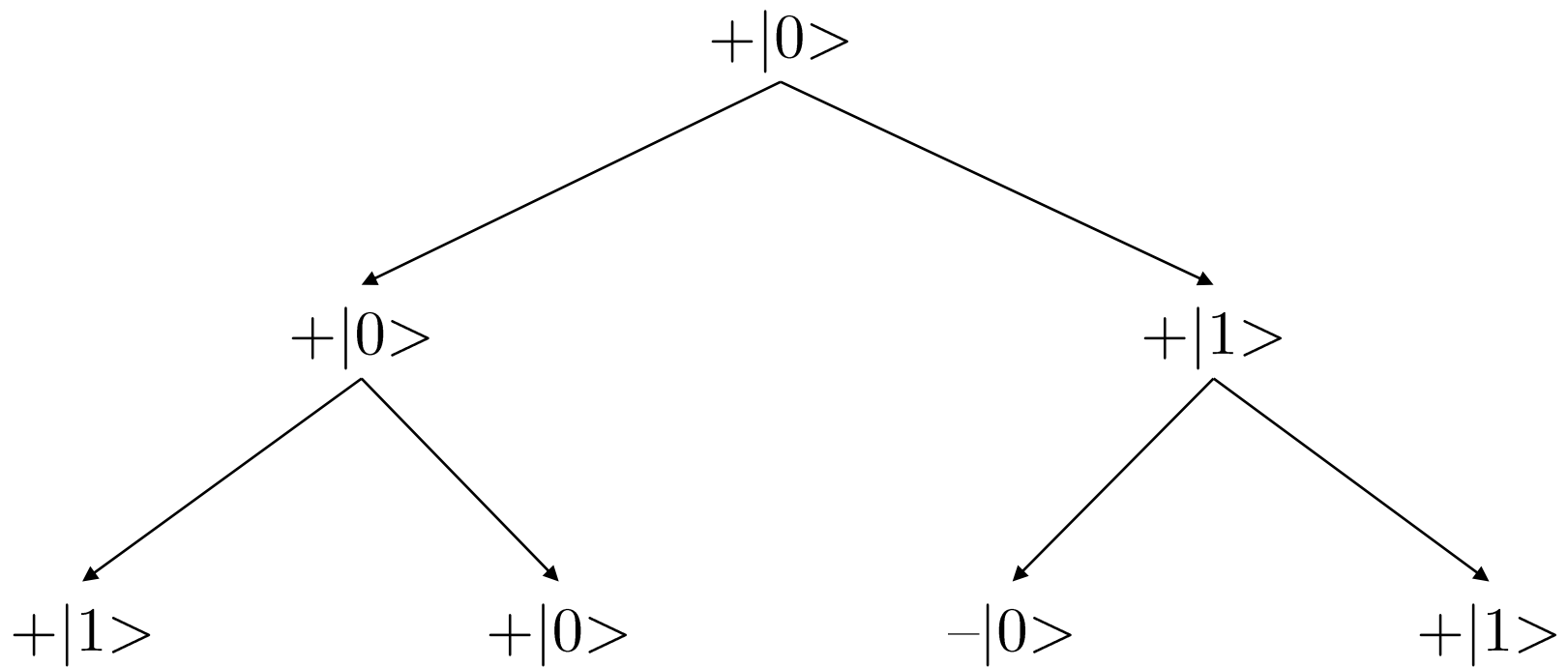
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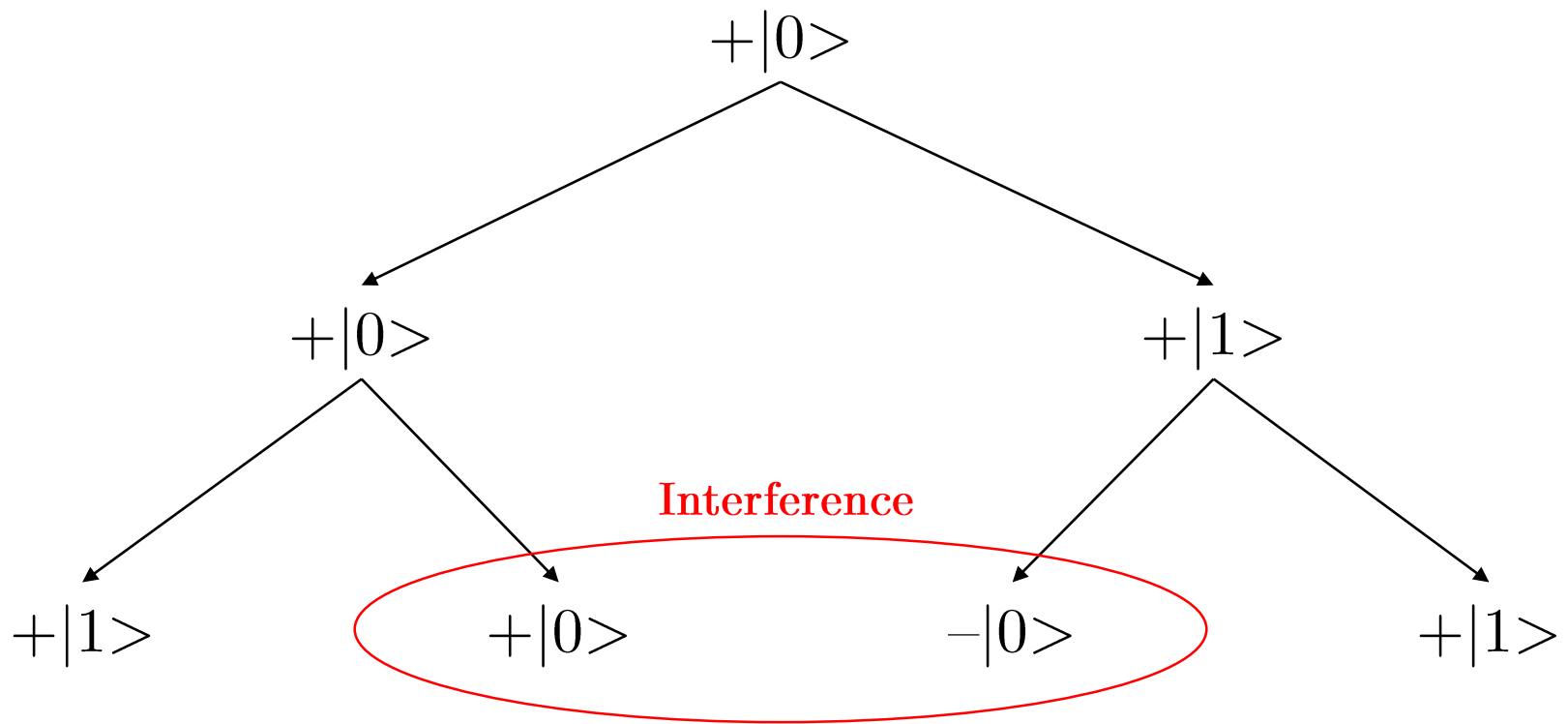
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Applying to Computation

Taking 2 bits / qubits:

CLASSICAL COMPUTER (Bits)	QUANTUM COMPUTER (Qubits)
00	$ 0, 0\rangle$
01	$ s\rangle = 0, 1 - 1, 0\rangle / \sqrt{2}$
10	$ T_0\rangle = 0, 1 + 1, 0\rangle / \sqrt{2}$
11	$ 1, 1\rangle$

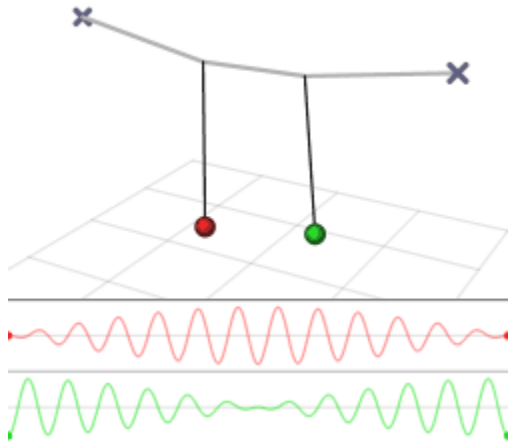
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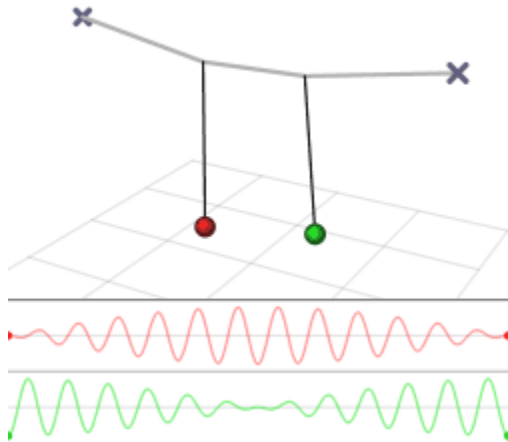
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Entangled States

Explaining the Entangled States

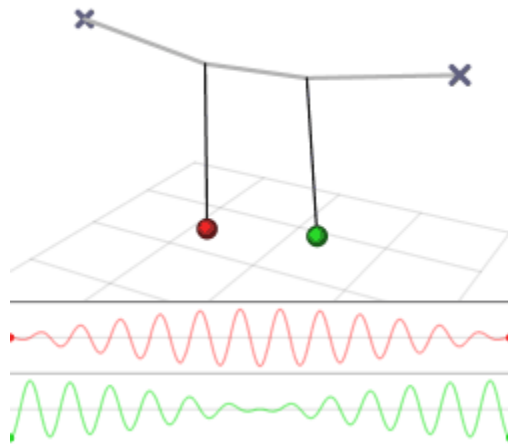


Explaining the Entangled States



$$\begin{array}{l} |0,1\rangle \\ |1,0\rangle \end{array}$$

Explaining the Entangled States



~~$|0,1\rangle$
 $|1,0\rangle$~~

$$\frac{1}{\sqrt{2}} |0, 1 \rangle - \frac{1}{\sqrt{2}} |1, 0 \rangle \text{ (Singlet State)}$$

$$\frac{1}{\sqrt{2}} |0, 1 \rangle + \frac{1}{\sqrt{2}} |1, 0 \rangle \text{ (T}_0 \text{ State)}$$

(Qubits are not independent of each other!)

Superposition of States

$$\alpha |0, 0\rangle$$

$$\beta |T_0\rangle = |0, 1 - 1, 0\rangle / \sqrt{2}$$

$$\gamma |s\rangle = |0, 1 + 1, 0\rangle / \sqrt{2}$$

$$\delta |1, 1\rangle$$

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No. of bits	No. of values held by classical bits	No. of values held by qubits
1	1	2
2	2	4
3	3	8
n	n	2^n

Problems

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- Classical Algorithms

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- Superpositions

The Prime Factorisation Problem

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$O(e^{(\log n \log \log n)^{\frac{1}{2}}})$ problem using classical algorithms

The Period Finding Problem

2, 4, 8, 16, 32, 64, 128, 256, 512, 1024, ...

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2, 4, 8, 1, 2, 4, 8, 1, 2, 4, ...

2, 4, 8, 1, 16, 11, 2, 4, 8, 16, ...

Euler's Totient Function

$x \bmod N, x^2 \bmod N, x^3 \bmod N, x^4 \bmod N, \dots$

Period evenly divides $(p-1)(q-1)$

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$x=2, N=21; p=3, q=7$

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Hence $p=3, q=5$

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$x=2, N=21; p=3, q=7$

$(p-1)(q-1) = 12$ and $12|6$ as required

Shor's Algorithm: The Classical Part

$x^r \bmod N$ (r may be very large)

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$$r = 2^3 + 2^2 + 2^1$$

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$x^r \bmod N$ (r may be very large)

Let $N=17$, $x=3$, $r=14$

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$$x^r = 3^{14} = 3^{2^3+2^2+2^1} = 3^{2^3} \cdot 3^{2^2} \cdot 3^{2^1} = ((3^2)^2)^2 \cdot (3^2)^2 \cdot 3^2$$

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$$6561 \bmod 17 * 81 \bmod 17 * 9 \bmod 17 = 16 * 13 * 9$$

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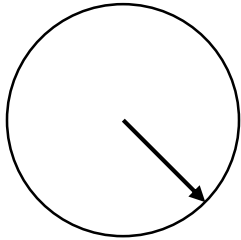
$$6561 \bmod 17 * 81 \bmod 17 * 9 \bmod 17 = 16 * 13 * 9$$

$$\text{and } (16 * 13 * 9) \bmod 17 = 2$$

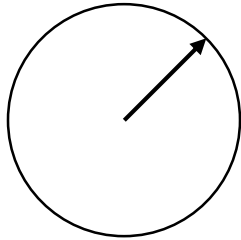
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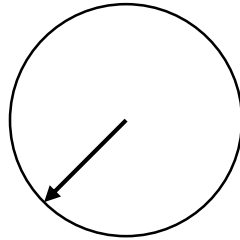
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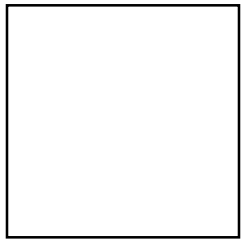
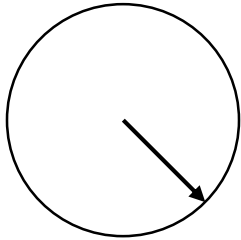


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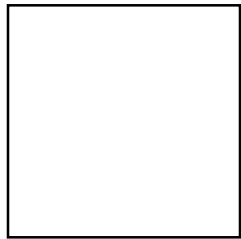
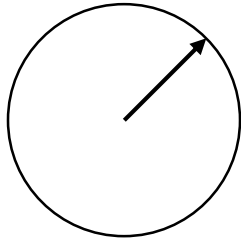


Shor's Algorithm: The Quantum Part

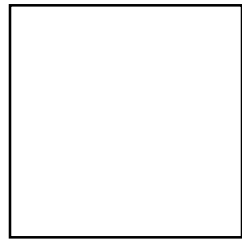
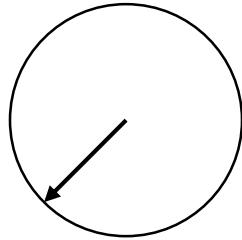
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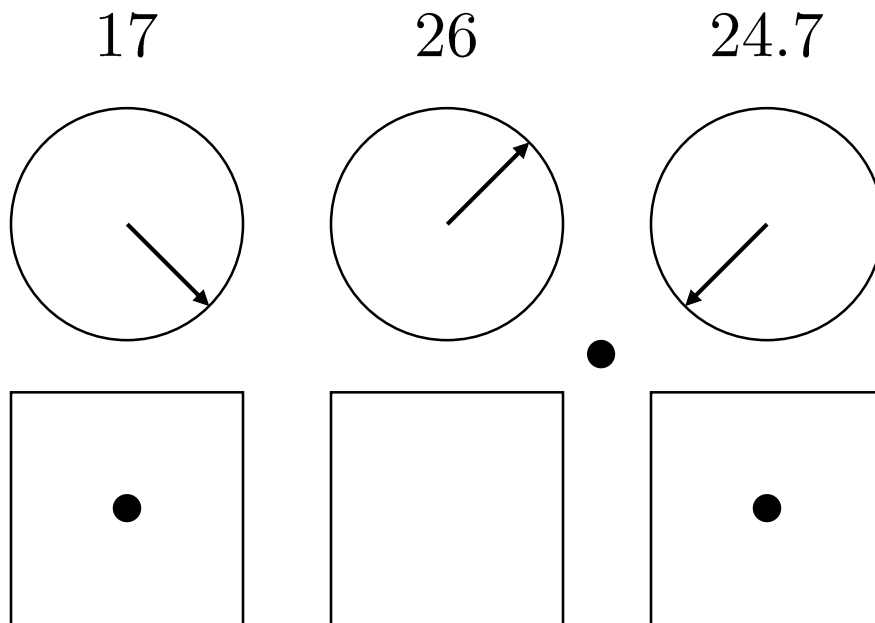
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24.7



Shor's Algorithm: The Quantum Part



The Quantum Fourier Transform

The Quantum Fourier Transform

$$X_k = \frac{1}{N} \sum_{n=0}^{N-1} x_n e^{i2\pi k \frac{n}{N}}$$

To find the energy at a particular frequency, spin your signal around a circle at that frequency, and average a bunch of points along that path.

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<http://betterexplained.com/articles/an-interactive-guide-to-the-fourier-transform/>

Back to Interference