Quantum-mechanical systems in traps and density functional theory

by

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Preface

blah blah

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Chapter 1

Introduction

blah blah

Part I

Theory

Chapter 2

Scientific Programming

The introduction of the computer around 1945 had a major impact on the mathematical fields of science. Previously unsolvable problems were now easily solvable. The question was no longer whether or not it was possible, but rather to what precision and with which method. The computer spawned a new branch of physics, computational physics, breaching barriers no one could even imagine existed. The first major result of this synergy between science and computers came with the atomic bombs as a result of the Manhattan Project at the end of the second world war citation needed.

2.1 Programming languages

Writing a program, or a code, is a list of instructions for the computer. It is in many ways similar to writing human-to-human instructions. You may use different programming languages, such as C++, Python, Java, as long as the reader is able to translate it. The translator, called *compiler*, translates your program from e.g. C++ code into machine code. Other languages as Python are interpreted real-time and therefore require no compilation. Although the latter seems like a better solution, it comes at the price of efficiency, a key concept in programming.

As a rule of thumb, efficiency is inverse proportional to the complexity of the programming language. It is therefore natural to sort languages into different subgroups depending on where they are at the efficiency-complexity scale.

2.1.1 High-level languages

This subgroup of languages are often referred to as *scripting languages*. A script is a short code, often with a specific purpose such as analyzing output by e.g. generating tables and figures from raw data, or gluing together different programs which are meant to be run in sequential order.

For simple jobs as these, taking only seconds to run on modern computers, we do not need an optimized code, but rather an easily read, clutter-free code. The languages which prefer simplicity over efficiency are referred to as *High-level* ¹. Examples of high-level languages are Python, Ruby, Perl, Visual Basic and UNIX shells. In this thesis I will emphasize the use of Python as a scripting language.

¹There are different definitions of high-level vs. low-level. You have languages such as *assembly*, which is extremely complex and close to machine code, leaving all machine-independent languages as high-level ones. However, for the purpose of this thesis I will not go into assembly languages, and keep the distinction at a higher level.

Python

Python is an open source programming language with a focus on simplicity over complexity. To mention a few of the entries in the Zen of $Python^2$, "Beautiful is better than ugly. Simple is better than complex. Readability counts. If the implementation is hard to explain, it's a bad idea."

To demonstrate the simplicity of Python, let us have a look at a simple implementation and execution of the following expression

$$S = \sum_{i=1}^{100} i = 5050.$$

```
1 #Sum100Python.py
2 print sum(range(101))
```

```
~$ python Sum100Python.py 5050
```

2.1.2 Low-level languages

A huge part of scientific programming is solving complex equations. Complexity does not necessarily imply that the equations themselves are hard to understand. Frankly, this is often not the case. In most cases of linear algebra, the problem can be boiled down to solving $A\vec{x} = B$, however, the complexity lies in the dimensionality of the problem at hand. Matrix dimensions range as high as millions. With each element being a double precision number, it is crucial that we have full control of the memory, and execute operations as efficiently as possible.

This is where lower level languages excel. Hiding few to none of the details, the power is in the hand of the programmer. This comes at a price. More technical concepts such as memory pointers, declarations etc. makes the development process slower than that of a higher-level language. If you e.g. try to access an element outside the bounds of an array, Python would tell you exactly this, whereas C++ would crash runtime leaving nothing but a "segmentation fault" for the user to interpret. However, when the optimized program ends up running for days, the extra time spent developing it pays off. As a rule of thumb, higher-level languages (without focus on *vectorization* or other optimization techniques) run 100 times slower than lower level ones.

C++

C++ is a language developed by Bjarne Stroustrup in 1979 at Bell Labs. It serves as an extension to the original C language, adding e.g. object oriented features such as classes. To calculate the sum in Eq. 2.1.1 with C++, we would need something like this:

²Retrieved by typing "import this" in your Python shell.

```
9     }
10
11     std::cout << S << std::endl;
12
13     return 0;
14 }</pre>
```

```
~$ g++ Sum100C++.cpp -o sum100C++.x
~$ ./sum100C++.x
5050
```

As we can see in lines five and six, we need to declare S and i as integer variables, exactly as described in section 2.1.2. In comparison with the Python version, it is clear that lower level languages are more complicated, and not designed for simple jobs as calculating a single sum.

Even though this is an extremely simply example, it illustrates the difference in coding styles between high- and low-level languages: Complexity vs. simplicity, efficiency vs. readability. I will not go through all the basic details of C++, but rather focus on the more complicated parts involving object orientation in scientific programming.

2.2 Object orientation

Object orientated programming was introduced in the language Simula, developed by the Norwegian scientists Ole-Johan Dahl and Kristen NygaardNeed specifics and citations. It quickly became the state-of-the-art in programming, and is today used throughout the world in all branches of programming. It is brilliant in the way that it ties our everyday intuition into the programming language - our brain is object oriented. It is focused around the concept of classes, an extension to standard structs. A class is nothing but a collection of variables and functions under a sensible name, however, they provide a great deal of functionality like inheritance and accessibility control. To better illustrate the concept of classes and inheritance, let us first look at an example.

2.2.1 Inheritance

All keyboards have two things in common: A board and keys. In object orientation we would say that the *superclass* of keyboards describe a board with keys. It is *abstract* in the sense that you do not need to know what the keys look like, or what function they possess, in order to define the concept of a keyboard.

However, we can have different types of keyboards, for example a computer keyboard, or a musical keyboard. They are different in design and function, but they both relate to the same concept of a keyboard described previously. They are both *subclasses* of the same superclass, inheriting the basic concepts, but expands upon them defining their own specific case. Let us look at how this concept can be converted into Python code

```
def pressKey(self, key):
11
12
13
14
15
   class ComputerKeyboard (Keyboard):
16
17
            language
18
19
            def __init__(self, language):
20
                     self.language = language
21
22
                     self.setupKeys()
23
24
            def setupKeys(self):
25
                     if self.language == "Norwegian":
26
                              keys [0] = "I"
27
28
                              keys[1] = "1"
29
                              #...
30
31
32
                    #...
33
34
            def pressKey(self, key):
35
                    return self.keys[key]
36
37
38
39
40
41
   class MusicalKeyboard (Keyboard):
42
43
            def __init__(self , nKeys, noteLength , amplitude):
44
                     self.amplitude = amplitude
45
                     self.noteLenght = noteLength
46
47
                     keys = zeros(nKeys)
48
                     self.setupKeys()
49
50
            def setupKeys(self):
51
                     keys[0] = lowestNote
52
53
54
                     keys[30] = 440 \text{ #Hz}
55
56
                    #...
57
            def pressKey(self, key):
58
59
60
                     t = linspace(0, self.noteLength, 100)
                     pi = 3.141592
61
62
                     return self.amplitude*sin(2*pi*keys[key]*t)
63
```

As we can see, the only thing differentiating the two keyboard types are how the keys are set up, and what happens when we press one of them. A superclass function designed to be overridden is referred to as *virtual*.

2.2.2 Virtual Functions, pointers and types

It a virtual function is overridden, the latest implementation will be called. setupKeys and pressKey are examples of this, however, they are in a sense more than virtual, since they are not even implemented. They have to be overridden in order to work. These functions are referred to as *pure virtual* functions. In python, all class functions, or member functions, are virtual. In C++ however, we have to specify whether or not a function is virtual in the declaration.

Higher level languages like Python handles all the pointers by itself. In low-level languages like C++, however, you need to control these yourself. A pointer is a hexadecimal number representing a memory address. If you pass a pointer to an object, e.g. Someclass* someobject, as an argument to a function, whenever that function changes a class variable, the value is changed globally, since the memory address is directly accessed. If you instead choose to send the object without a pointer declaration, e.g. Someclass someobject, changing the value will not change the object globally. What happens instead is that you change a local copy of the object. It is a misconception that pointers are your enemies, they are, quite frankly, making your codes much easier.

A pointer holds a type, that is, int, double, or any class that you have access to. In the following example we will study the interplay between virtual functions, pointers and types:

```
1 #include <iostream>
  using namespace std;
3
   class VirtualTest {
  public:
6
            virtual void virtualFunc(){
                      \mathrm{cout} << "superclass virtualFunc called" << \mathrm{endl};
9
            }
10
            void notVirtualFunc(){
11
                      {\tt cout} <\!< \verb"superclass" not Virtual Func called "<< endl;
12
13
14
15
   };
16
   class subclass : public VirtualTest{
17
            virtual void virtualFunc(){
19
                      cout << "subclass virtualFunc called" << endl;</pre>
20
22
23
            void notVirtualFunc(){
                      {\tt cout} <\!< \verb"subclass notVirtualFunc called" << endl;
24
            }
25
26
   };
27
28
   void function(VirtualTest* object){
            object -> virtualFunc();
30
31
            object->notVirtualFunc();
32
33
  int main(){
35
            \mathrm{cout} << "-Calling subclass object of type VirtualTest*" << \mathrm{endl};
36
            VirtualTest* object = new subclass();
37
            function (object);
38
39
            \mathrm{cout} << \mathrm{endl} << \text{"-Calling subclass object of type subclass*"} << \mathrm{endl};
40
            subclass* object2 = new subclass();
41
            function (object2);
42
43
            cout << endl << "-Directly calling object of type subclass*" << endl;</pre>
44
            object2->virtualFunc();
45
            object2->notVirtualFunc();
46
47
48
            return 0;
49
```

~\$./virtualFunctionsC++.x

⁻Calling subclass object of type VirtualTest* subclass virtualFunc called

```
superclass notVirtualFunc called
-Calling subclass object of type subclass*
subclass virtualFunc called
```

superclass notVirtualFunc called

```
-Directly calling object of type subclass*
subclass virtualFunc called
subclass notVirtualFunc called
```

In the first call, the pointer is declared as a VirtualTest* type, however, it is still initialized to be a subclass pointer in the sense that the subclass' functions are loaded into the object. This results in the virtual function being overridden (as mentioned previously), but since the object type is VirtualTest*, C++ does not dig deeper than the superclass when it looks for the non-virtual function implementation. In other words: virtual induce a search for deeper implementations of the same function, given that the function is loaded through e.g. new subclass().

In the second call, the same thing happens, even though it is set as a subclass* type. This is because the function is instructed to receive a superclass object. If it receives anything else, it simply attempts to convert it, or *cast* it to a different type; *typecasting*³. In this case it works just fine. The third call, outside the function, demonstrates that if we declare it as a subclass type, both functions are overridden, since we do not have to go through the superclass at all.

The strength of using virtual functions in a class hierarchy is that you can easily expand or implement new functionality without completely rewriting the code. As an example, in the QMC code, changing potentials, or switching between closed form and numerical expressions for the dell and laplacian, is just a matter of which object's functions are called by the energy function.

```
double QMC::calculate_local_energy(Walker* walker) const {
2
      return kinetics ->get_KE(walker) + system->get_potential_energy(walker);
3
  }
1 class Walker {
2
  . . .
3
4
      mat r:
5
6 };
1 class Kinetics {
2
3
      virtual double get_KE(const Walker* walker) = 0;
5
6 };
8 class Numerical : public Kinetics {
10
      virtual double get_KE(const Walker* walker);
11
12
13 };
1 class Potential {
3
       virtual double get_pot_E(const Walker* walker) const = 0;
4
5
6 };
```

³The standard example of typecasting is converting a double to an integer, resulting in the stripping of all the decimal bits (flooring).

```
class Harmonic_osc : public Potential {
    ...

virtual double get_pot_E(const Walker* walker) const;
    ...
};
```

```
double Harmonic_osc::get_pot_E(const Walker* walker) const {

    double e_potential = 0;

    for (int i = 0; i < n_p; i++) {
        e_potential += 0.5 * w * w * walker->get_r_i2(i);
    }

    return e_potential;
}
```

The QMC member function receives a Walker* object, representing a set of positions for all particles. The kinetics object is of type Kinetics* and holds implementations of either a numerical calculation or a method for extracting closed form expressions from the orbitals. The get_potential_energy method is not virtual. It's purpose is to iterate over all potentials given (i.e. a Harmonic Oscillator and Coulomb), extracting their value at the walker's position. Adding a new potential to this list is extremely simple.

The point is that the local energy function is written completely independent of how the potential actually looks. The implementation of a new potential would mean a new subclass of Potential with a new implementation of the virtual function get_local_E. This makes the code extremely readable (given it is properly commented), since the reader would not have to care about the details if the algorithm is that of interest (and vice versa).

2.2.3 Const Correctness

In the QMC code example above, function declarations with const are used. If an object is declared with const on input, e.g. void f(const x), the function itself cannot alter the value of x. It is a safeguard that nothing will happen to x as it passes through f. This is practical in situations where major bugs will arise if anything happens to an object.

If you declare a member function itself with const on the right hand side, it safeguards the function from changing any of the class variables. If you e.g. have a variable representing the electron charge, you do not want this changed by the Coulomb class member function. This should only happen through specific functions whose sole purpose is changing the charge, and taking care of any following consequences.

In other words: **const** works as a safeguard for changing values which should remain unchanged. A change in such a variable is then followed by a compiler error instead of infecting your code with bugs, resulting in unforeseen consequences.

2.2.4 Accessibility levels and Friend classes

const is a direct way to avoid any change what so ever. However, sometimes we want to keep the ability to alter variables, but only in certain situations, as e.g. internally in the class. As an example, from the main file, you should not have access to QMC member functions such as dump_output, since it does not make sense to do out of a context. However, you obviously want access to the run_method function.

The solution to this problem is to set accessibility levels. Declaring a variable under the public part of a class sets its accessibility level to public, meaning that anything, anywhere can access it as long as it

has access to the object. All public variables are inherited to subclasses when inherited as in the virtual function example⁴. Declarations beneath the private part stops all other classes than instances of it self from reaching it, even subclass instances. If you want private variables inherited, the protected accessibility level should be used.

There is one exception to the rule of protected and private variables, namely *friend* classes. In the QMC library, there is a output class called OutputHandler. This class needs access to protected variables, since the user should be able to output anything he wants. If we friend the output class with QMC, we get exactly this behavior:

```
class QMC {
  protected:
3
       int n_c;
4
5
       int n_p;
6
       int n2;
7
       int dim;
9
10
       int cycle;
11
12
13
       Walker* original_walker;
14
15
       Walker* trial_walker;
16
17
18
  public:
19
20
21
       friend class Distribution;
22
23
  };
```

```
void Distribution::dump() {
          if ((qmc->cycle > qmc->n_c / 2) && (qmc->cycle % 100 == 0)) { for (int i = 0; i < qmc->n_p; i++) {
3
4
                      for (int j = 0; j < qmc -> dim; j++) {
                            \quad \textbf{if} \ (\texttt{j} = \texttt{qmc-} \texttt{>} \texttt{dim} \ - \ 1) \ \{
6
                                   file \ll qmc->original-walker->r(i, j);
                                   \label{eq:file_scale} file << qmc-> original_walker-> r(i, j) << " ";
9
10
11
                      file << endl;
12
13
14
15
16
```

Without going into details, we can see that Distribution has full access to protected members of QMC* qmc. Friend classes is a savior in those very specific cases when you really need full access to protected members of another class, but setting full public access would ruin the code. It is true that you could code your entire code without const and with solely public members, but in that case, it is very easy to put together a very disorganized code, with pointers flying everywhere and functions being called in all sorts of contexts. Clever use of accessibility levels will make your code easier to develop in an organized, intuitive way - you will be forced to implement things in an organized fashion.

⁴You could inherit with protected option as well, but it is rarely used and messes up the functionality.

2.2.5 Example: PotionGame

To end the section I would like to demonstrate the versatile power of object orientation. Consider the following codes (libraries)

```
1 #potionClass.py
  class Potion:
           def __init__(self , amount):
5
                    self.amount = amount;
6
                    self.setName()
8
           {f def} applyPotion(self, player):
9
                   #Pure virtual function
                    return None
11
12
           def setName(self):
13
                   #Virtual function
14
15
                    self.name = "Undefined"
16
17
  class HealthPotion(Potion):
19
           #Constructor is inherited
20
21
           def applyPotion(self, player):
22
23
                    player.changeHealth(self.amount)
24
25
           def setName(self):
                    self.name = "Health Potion (" + str(self.amount) + ")"
26
27
  class EnergyPotion(Potion):
30
31
           def applyPotion(self, player):
                    player.changeEnergy(self.amount)
32
33
34
           def setName(self):
                    self.name = "Energy Potion (" + str(self.amount) + ")"
35
```

```
1 #playerClass.py
  class Player:
3
           def __init__(self , name):
                    self.health = 100
                    self.energy = 100
6
7
                    self.name = name
8
                    self.dead = False
9
10
                    self.potions = []
11
12
           def addPotion(self, potion):
13
                    self.potions.append(potion)
14
15
           def usePotion(self, potionIndex):
16
                    self.potions[potionIndex].applyPotion(self)
17
                    self.potions.pop(potionIndex)
18
19
           def changeHealth(self, amount):
20
                    self.health += amount
21
22
                    \#Cap health at [0,100].
23
                    if self.health > 100:
24
                            self.health = 100
25
                    elif self.health <= 0:
26
                            self.health = 0
27
```

```
self.dead = True;
28
29
30
           def changeEnergy(self, amount):
31
                    self.energy += amount
32
33
                    \#Cap energy at [0,100].
34
                    if self.energy > 100:
35
36
                             self.energy = 100
                    elif self.energy < 0:
37
                             self.energy = 0
38
39
40
41
           def superPowers(self):
                    #Pure virtual function
42
                    return None
43
44
           def displayPotions(self):
45
                    if len(self.potions) = 0:
46
47
                             print "No potions aviable"
                    for potion in self.potions:
48
49
                             print potion.name
50
  class Superman (Player):
51
52
           def superPowers(self):
                    self.canFly = True;
53
```

We have a Player class keeping track of a players energy and health level, and which potions the player is carrying. The Potion class described all potions, that is, an object of type Potion with the ability to affect a player in some way. The subclasses define specifically which effect is to be applied, e.g. HealthPotion changes the health level of the player by a certain amount. This code does nothing by itself, but let us use it in an example where two players fight each other:

```
#potionGameMain.py
  from potionClass import *
3
  from playerClass import *
  def roundOutput(players, n):
6
           header= "\nRound %d: " % n
           print header.replace('0','start')
8
9
           for player in players:
                   print "%s (hp/e=%d/%d):" % (player.name, player.health, player.energy)
10
                   player.displayPotions()
11
12
                   print
13
14
  player1 = Player('john')
  player2 = Player('james')
16
17
  player1.addPotion(HealthPotion(10))
  player1.addPotion(EnergyPotion(30))
19
  player2.addPotion(EnergyPotion(20))
21
  player2.addPotion(EnergyPotion(20))
22
23
  roundOutput([player1, player2], 0)
24
  print "\nPlayer2 hit Player1 for 50 using 100 energy"
  player1.changeHealth(-50)
27
28
  player2.changeEnergy(-100)
30 print "Players used all potions"
  player1.usePotion(1)
32 player1.usePotion(0)
33 player2.usePotion(1)
  player2.usePotion(0)
35
```

36 print "Player1 hit Player2 for 40 using 90 energy\n"

```
player1.changeEnergy(-90); player2.changeHealth(-40)

roundOutput([player1, player2], 1)

** python potionGameMain.py

Round start:
john (hp/e=100/100):
Health Potion (10)
Energy Potion (30)

james (hp/e=100/100):
Energy Potion (20)
Energy Potion (20)

Player2 hit Player1 for 50 using 100 energy
```

Round 1: john (hp/e=60/10): No potions aviable james (hp/e=60/40): No potions aviable

Players used all potions

Player1 hit Player2 for 40 using 90 energy

The readability of this code is pretty good. Imagine if we had no objects, but just a lot of parameters per player juggled around in variables such as Player1health etc. Increasing the number of players then requires a total rewriting of the entire program, where as in this object oriented style, it is just a matter of adding another player object. Object orientation is truly brilliant when it comes to developing codes.

In this section I have not focused too much on scientific computing, but rather on the use of object orientation in general. When the physical methods are discussed in section **Insert physics section**, I will get back to a more specific description of scientific programming.

As an introduction to the next section, take a look at the way the classes and the main file are separated in this section's example. In a small code like this there's really no point of doing so, but when the class structures span thousands of lines, having a good structure and the right editor is crucial to the development process and the code's readability.

2.3 Structuring the code

blah blah messy=bad, organized = good

2.3.1 File structures

blah blah src folder good yes

2.3.2 IDEs

Netbeans good yes, spider good yes

Part II

Results

Chapter 3

Results

- 3.1 Validating the code
- 3.1.1 Calculation for non-interacting particles

[1]

Appendix A

Dirac Notation

Due to the orthogonal nature of Hermitian operators' eigenfunctions¹, the inner product between two states constructed from them will result in a lot of integrals being zero, one, or eigenvalues for that matter. Writing the integrals in their full form then feels like a waste of space and time. Even specifying e.g. a position basis is obsolete. Abstracting the wave functions from a given parameter space (e.g. \mathbb{C}^n) into a *Hilbert space*² is what is called the *Dirac notation*, or the *Bra-ket notation*.

The basic idea is that since the coordinate representation of a wave function is the projection of an abstract state on the position basis through an inner product, we can separate the different pieces of the inner product:

$$\psi(\vec{r}) = \langle r, \psi_i \rangle \equiv \langle r | \psi_i \rangle = \langle r | \times | \psi_i \rangle.$$

The notation is designed to be simple. The right hand side of the inner product is called a ket, while the left hand side is called a bra. Combining both of them leaves you with an inner product bracket, hence the names. Let us look at an example where this notation is extremely powerful. Imagine a coupled two-particle spin- $\frac{1}{2}$ system in the following state

$$|\psi\rangle = N \left[|\uparrow\downarrow\rangle - i |\downarrow\uparrow\rangle \right]$$
 (A.1)

$$\langle \psi | = N \left[\langle \uparrow \downarrow | + i \langle \downarrow \uparrow | \right]$$
 (A.2)

Using the fact that both the full two-particle state and the two-level spin states should be orthonormal, we can with this notation calculate the normalization factor without explicitly calculating anything.

$$\begin{split} \langle \psi | \psi \rangle &= N^2 \Big[\left< \uparrow \downarrow \right| + i \left< \downarrow \uparrow \right| \Big] \Big[\left| \uparrow \downarrow \right> - i \left| \downarrow \uparrow \right> \Big] \\ &= N^2 \Big[\left< \uparrow \downarrow \right| \uparrow \downarrow \right> + i \left< \downarrow \uparrow \right| \uparrow \downarrow \right> - i \left< \uparrow \downarrow \right| \downarrow \uparrow \right> + \left< \downarrow \uparrow \right| \downarrow \uparrow \right> \Big] \\ &= N^2 \Big[1 + 0 - 0 + 1 \Big] \\ &= 2N^2 \end{split}$$

 $^{^{1}\}mathrm{Eigenfunctions}$ of a Hermitian operator always make up a complete orthogonal set.

²A Hilbert space is an inner product space spanned by the different states. For every state, there exists a complementary state which is the Hermitian conjugate of the original[2].

This implies as we expected $N=1/\sqrt{2}$. With this powerful notation at hand, we can easily show properties such as the *completeness relation* of a set. We start by expanding one state $|\phi\rangle$ in a complete set of different states $|\psi_i\rangle$:

$$\begin{split} |\phi\rangle &=& \sum_{i} c_{i} \, |\psi_{i}\rangle \\ \langle \psi_{k} | \phi\rangle &=& \sum_{i} c_{i} \, \langle \psi_{k} | \psi_{i}\rangle \\ &=& c_{k} \\ |\phi\rangle &=& \sum_{i} \langle \psi_{i} | \phi\rangle \, |\psi_{i}\rangle \\ &=& \left[\sum_{i} |\psi_{i}\rangle \, \langle \psi_{i} | \right] |\phi\rangle \,, \end{split}$$

which implies that

$$\sum_{i} |\psi_{i}\rangle \langle \psi_{i}| = \mathbb{1} \tag{A.3}$$

for any complete set of orthonormal states $|\psi_i\rangle$. For a continuous basis like e.g. the position basis we have a similar relation:

$$\int |\psi(x)|^2 dx = 1$$

$$\int |\psi(x)|^2 dx = \int \psi^*(x)\psi(x)dx$$

$$= \int \langle \psi|x\rangle \langle x|\psi\rangle dx$$

$$= \langle \psi| \left[\int |x\rangle \langle x| dx\right] |\psi\rangle.$$
(A.4)

Combining eq. A.4 and eq. A.5 with the fact that $\langle \psi | \psi \rangle = 1$ yields the identity

$$\int |x\rangle \langle x| \, dx = \mathbb{1}. \tag{A.6}$$

Appendix B

Matrix representation of states and operators

One of the most common ways to represent states and operators, at least in computational quantum mechanics, is using vectors and matrices. It is crucial to note, however, that we are discussing a *representation* of states and operators; the theory itself is general, and independent of whatever convenient choice of representation we make.

The matrix representation of an operator $\widehat{\mathbf{A}}$ in necessarily dependent of our choice of basis. To illustrate this we look at the matrix representation of the Hamiltonian. It satisfies the time independent $\operatorname{Schr} \widehat{\mathbf{A}} \P$ dinger equation

$$\widehat{\mathbf{H}} \ket{\psi_{E_i}} = E_i \ket{\psi_{E_i}}.$$

Using spectral decomposition on $\widehat{\mathbf{H}}$ we get

$$\widehat{\mathbf{H}} = \sum_{k} E_k |\psi_{E_k}\rangle \langle \psi_{E_k}|, \qquad (B.1)$$

which by definition of $\hat{\mathbf{H}}$ is diagonal in the energy eigenstates:

$$H = \begin{pmatrix} E_0 & 0 & 0 & \cdots & 0 \\ 0 & E_1 & 0 & & \vdots \\ 0 & 0 & \ddots & & \\ \vdots & & & & \\ 0 & \cdots & & E_N \end{pmatrix}. \tag{B.2}$$

However, if we perform a change of basis from $|\psi_{E_i}\rangle$ to an arbitrary complete set of orthogonal states $|\phi_i\rangle$ by using the completeness relation from eq. A.3, we get the following relation

$$\widehat{\mathbf{H}} = \sum_{k} E_{k} |\psi_{E_{k}}\rangle \langle \psi_{E_{k}}|
= \sum_{k} \sum_{i,j} E_{k} |\phi_{i}\rangle \langle \phi_{i}| \psi_{E_{k}}\rangle \langle \psi_{E_{k}}| \phi_{j}\rangle \langle \phi_{j}|
= \sum_{k} \sum_{i,j} |\phi_{i}\rangle \langle \phi_{i}| \widehat{\mathbf{H}} |\psi_{E_{k}}\rangle \langle \psi_{E_{k}}| \phi_{j}\rangle \langle \phi_{j}|
= \sum_{i,j} |\phi_{i}\rangle \langle \phi_{i}| \widehat{\mathbf{H}} \Big[\sum_{k} |\psi_{E_{k}}\rangle \langle \psi_{E_{k}}| \Big] |\phi_{j}\rangle \langle \phi_{j}|
= \sum_{i,j} |\phi_{i}\rangle \langle \phi_{i}| \widehat{\mathbf{H}} |\phi_{j}\rangle \langle \phi_{j}|
= \sum_{i,j} |\phi_{i}\rangle \langle \phi_{i}| \widehat{\mathbf{H}} |\phi_{j}\rangle \langle \phi_{j}|
= \sum_{i,j} |\psi_{i}\rangle \langle \phi_{j}|,$$
(B.3)

which is not diagonal in the new basis. This is usually the starting point when we do physics, since the goal of the computation is to obtain the true eigenstates and eigenvectors of a Hamiltonian. If we choose an initial complete orthonormal basis, we can always set up the matrix and diagonalize it¹.

Doing this basis change, we have also deduced the general form of the matrix elements in a given basis:

$$A_{ij} = \langle \psi_i | \widehat{\mathbf{A}} | \psi_i \rangle, \tag{B.4}$$

$$A = \begin{pmatrix} \langle \psi_{0} | \widehat{\mathbf{A}} | \psi_{0} \rangle & \langle \psi_{0} | \widehat{\mathbf{A}} | \psi_{1} \rangle & \langle \psi_{0} | \widehat{\mathbf{A}} | \psi_{2} \rangle & \cdots & \langle \psi_{0} | \widehat{\mathbf{A}} | \psi_{N} \rangle \\ \langle \psi_{1} | \widehat{\mathbf{A}} | \psi_{0} \rangle & \langle \psi_{1} | \widehat{\mathbf{A}} | \psi_{1} \rangle & \langle \psi_{1} | \widehat{\mathbf{A}} | \psi_{2} \rangle & & \vdots \\ \langle \psi_{2} | \widehat{\mathbf{A}} | \psi_{0} \rangle & \langle \psi_{2} | \widehat{\mathbf{A}} | \psi_{1} \rangle & \ddots & & \\ \vdots & & & & & \\ \langle \psi_{N} | \widehat{\mathbf{A}} | \psi_{0} \rangle & \cdots & & \langle \psi_{N} | \widehat{\mathbf{A}} | \psi_{N} \rangle \end{pmatrix}.$$
(B.5)

The matrix elements are calculated as integrals, e.g. the expectation value of the energy in an interacting quantum dot using single particle harmonic oscillator wave functions.

¹The brute force method of doing this (up to a given truncation in the infinite basis) is called *full configuration interaction* (FCI) or *full scale diagonalization*.

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