

TDT4205 Compiler Construction (2019)

Assignment 2

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1 LL(1) Parsing table construction

1.1 Starting grammar

```
/* Specified in bison notation */  
S: c L E;  
L: L E | l E;  
E: E s | U;  
U: b | c | n | %empty;
```

1.2 Modified grammar, without left recursion

```
/* Specified in bison notation */  
S: c L E;  
L: l E L1; /* alpha=E, beta=lE */  
L1: E L1 | %empty;  
E: U E1; /* alpha=s, beta=U */  
E1: S E1 | %empty;  
U: b | c | n | %empty;
```

1.3 Table of $\text{First}(\alpha)$, $\text{Follow}(\alpha)$ and $\text{Nullable}(\alpha)$

	First	Follow	Nullable
S	c	\$	No
L	l	{b, c, n, s}	No
L1	{b, c, n, s}	\$	Yes
E	{b, c, n, s}	{b, c, n, s}	Yes
E1	s	\$	Yes
U	{b, c, n}	\$	Yes

The grammar is ambiguous, as is seen in the table above, because $First(E) = Follow(E)$.

1.4 Parsing table of the ambiguous LL(1) grammar

	c	l	b	n	s
S	$S \rightarrow cLE$				
L		$L \rightarrow lEL'$			
L'	$L' \rightarrow EL'$		$L' \rightarrow EL'$	$L' \rightarrow EL'$	$L \rightarrow EL'$
E	$E \rightarrow UE' / E \rightarrow \epsilon$		$E \rightarrow UE' / E \rightarrow \epsilon$	$E \rightarrow UE' / E \rightarrow \epsilon$	$E \rightarrow UE'$
E'					$E' \rightarrow sE'$
U	$U \rightarrow c$		$U \rightarrow b$	$U \rightarrow n$	