

### Exercise 3.1

The hypothesis class  $\mathcal{H}$  being PAC learnable with sample complexity  $m_{\mathcal{H}}(\cdot, \cdot)$  means that there is a learning algorithm  $A$  such that when running  $A$  on  $m \geq m_{\mathcal{H}}(\epsilon, \delta)$  i.i.d. samples generated by  $\mathcal{D}$  and labeled by  $f$ , with probability at least  $1 - \delta$ ,  $A$  returns a hypothesis  $h \in \mathcal{H}$  with  $L_{D,f}(h) \leq \epsilon$ .

Given  $0 < \epsilon_1 \leq \epsilon_2 < 1$ , consider  $m \geq m_{\mathcal{H}}(\epsilon_1, \delta)$ , we have that with probability at least  $1 - \delta$ ,  $A$  returns a hypothesis  $h \in \mathcal{H}$  with  $L_{D,f}(h) \leq \epsilon_1 \leq \epsilon_2$ . This implies that  $m_{\mathcal{H}}(\epsilon_1, \delta)$  is a sufficient number of samples for accuracy  $\epsilon_2$ . Therefore,  $m_{\mathcal{H}}(\epsilon_1, \delta) \geq m_{\mathcal{H}}(\epsilon_2, \delta)$ .

The proof of  $m_{\mathcal{H}}(\epsilon, \delta_1) \geq m_{\mathcal{H}}(\epsilon, \delta_2)$  for  $0 < \delta_1 \leq \delta_2 < 1$  follows analogously from the definition.

### Exercise 3.3

The realizability assumption for  $\mathcal{H} = \{h_r : r \in \mathbb{R}_+\}$  implies that there is a circle such that any  $x$  inside it has label  $y = 1$ , and the learning task here is to distinguish this circle. Now consider an ERM algorithm which given a training sequence  $S = \{(x_i, y_i)\}_{i=1}^m$ , returns the hypothesis  $\hat{h}$  corresponding to the tightest circle which contains all the positive instances in  $S$  where  $y_i = 1$  and does not allow false negative predictions. With the realizability assumption let  $h^*$  be the circle with zero training error and  $r^*$  be the corresponding radius.

Let  $\bar{r} \leq r^*$  be a scalar such that  $\mathbb{P}_{x \sim \mathcal{D}}(x : \bar{r} \leq \|x\| \leq r^*) = \epsilon$  and  $E = \{x \in \mathbb{R}^2 : \bar{r} \leq \|x\| \leq r^*\}$ . We have

$$\begin{aligned} \mathbb{P}(L_{\mathcal{D}}(h_S) \geq \epsilon) &\leq \mathbb{P}(\text{no points in } S \text{ belongs to } E) \\ &= (1 - \epsilon)^m \\ &\leq e^{-\epsilon m} \end{aligned}$$

The desired bound on the sample complexity follows from requiring  $e^{-\epsilon m} \leq \delta$ .

### Exercise 3.7

Let  $g$  be any (potentially probabilistic) classifier from  $\mathcal{X}$  to  $\{0, 1\}$ . Note that for 0-1 loss

$$\begin{aligned} L_{\mathcal{D}}(g) &= \mathbb{E}_{(x,y) \sim \mathcal{D}}[\mathbb{1}_{g(x) \neq y}] = \mathbb{E}_{x \sim \mathcal{D}}[\mathbb{E}_{y \sim \mathcal{D}_{Y|x}}[\mathbb{1}_{g(x) \neq y}]] = \mathbb{E}_{x \sim \mathcal{D}}[\mathbb{P}_{y \sim \mathcal{D}_{Y|x}}(g(X) \neq Y | X = x)], \\ L_{\mathcal{D}}(f_{\mathcal{D}}) &= \mathbb{E}_{x \sim \mathcal{D}}[\mathbb{P}_{y \sim \mathcal{D}_{Y|x}}(f_{\mathcal{D}}(X) \neq Y | X = x)]. \end{aligned}$$

We should compare the two conditional probabilities inside the expectation. Let  $x \in \mathcal{X}$  and  $a_x = \mathbb{P}(Y = 1|X = x)$ . We have

$$\begin{aligned}
\mathbb{P}(g(X) \neq Y|X = x) &= \mathbb{P}(g(X) = 0|X = x) \cdot \mathbb{P}(Y = 1|X = x) \\
&\quad + \mathbb{P}(g(X) = 1|X = x) \cdot \mathbb{P}(Y = 0|X = x) \\
&= \mathbb{P}(g(X) = 0|X = x) \cdot a_x + \mathbb{P}(g(X) = 1|X = x) \cdot (1 - a_x) \\
&\geq \mathbb{P}(g(X) = 0|X = x) \cdot \min\{a_x, 1 - a_x\} \\
&\quad + \mathbb{P}(g(X) = 1|X = x) \cdot \min\{a_x, 1 - a_x\} \\
&= \min\{a_x, 1 - a_x\}.
\end{aligned}$$

When  $g = f_{\mathcal{D}}$  we should replace  $\mathbb{P}(g(X) = 0|X = x)$  by  $\mathbb{1}_{a_x < 1/2}$  and  $\mathbb{P}(g(X) = 1|X = x)$  by  $\mathbb{1}_{a_x \geq 1/2}$ . Then the above inequality is tight:

$$\mathbb{P}(f_{\mathcal{D}}(X) \neq Y|X = x) = \mathbb{1}_{a_x < 1/2} \cdot a_x + \mathbb{1}_{a_x \geq 1/2} \cdot (1 - a_x) = \min\{a_x, 1 - a_x\}.$$

Therefore, we have  $L_{\mathcal{D}}(f_{\mathcal{D}}) \leq L_{\mathcal{D}}(g)$ .

### Exercise 3.8

1. Solved already in Exercise 3.7.
2. We have shown in Exercise 3.7 that the Bayes optimal predictor  $f_{\mathcal{D}}$  is optimal w.r.t.  $\mathcal{D}$ ; in other words,  $f_{\mathcal{D}}$  is always better than any other learning algorithm w.r.t.  $\mathcal{D}$ .
3. Take  $\mathcal{D}$  to be any probability distribution and  $B = f_{\mathcal{D}}$ .

### Exercise 4.1

1  $\Rightarrow$  2: Assume for every  $\epsilon, \delta > 0$  there exists  $m(\epsilon, \delta)$  such that  $\forall m \geq m(\epsilon, \delta)$

$$\mathbb{P}_{S \sim \mathcal{D}^m}(L_{\mathcal{D}}(A(S)) > \epsilon) < \delta. \quad (1)$$

Then using the definition of expectation

$$\begin{aligned}
\mathbb{E}_{S \sim \mathcal{D}^m}[L_{\mathcal{D}}(A(S))] &\leq \mathbb{P}_{S \sim \mathcal{D}^m}(L_{\mathcal{D}}(A(S)) > \epsilon) \cdot 1 + \mathbb{P}_{S \sim \mathcal{D}^m}(L_{\mathcal{D}}(A(S)) \leq \epsilon) \cdot \epsilon \\
&\leq \mathbb{P}_{S \sim \mathcal{D}^m}(L_{\mathcal{D}}(A(S)) > \epsilon) + \epsilon \\
&\leq \delta + \epsilon,
\end{aligned}$$

where the last inequality follows from the assumption (1). Now set  $\delta = \epsilon$ . We have for every  $\epsilon > 0$  there exists  $m(\epsilon, \epsilon)$  such that  $\forall m \geq m(\epsilon, \epsilon)$

$$\mathbb{E}_{S \sim \mathcal{D}^m}[L_{\mathcal{D}}(A(S))] \leq 2\epsilon. \quad (2)$$

So it is valid to pass both sides of (2) to the limit  $\lim_{m \rightarrow \infty} \lim_{\epsilon \rightarrow 0}$ , which gives

$$\lim_{m \rightarrow \infty} \mathbb{E}_{S \sim \mathcal{D}^m}[L_{\mathcal{D}}(A(S))] \leq 0.$$

Also by definition  $\mathbb{E}_{S \sim \mathcal{D}^m}[L_{\mathcal{D}}(A(S))] \geq 0$ . Thus we conclude  $\lim_{m \rightarrow \infty} \mathbb{E}_{S \sim \mathcal{D}^m}[L_{\mathcal{D}}(A(S))] = 0$ .

2  $\Rightarrow$  1: Assume that  $\lim_{m \rightarrow \infty} \mathbb{E}_{S \sim \mathcal{D}^m} [L_{\mathcal{D}}(A(S))] = 0$ . For every  $\epsilon, \delta \in (0, 1)$  there exists some  $m_0 \in \mathbb{N}$  such that for every  $m \geq m_0$ ,  $\mathbb{E}_{S \sim \mathcal{D}^m} [L_{\mathcal{D}}(A(S))] \leq \epsilon\delta$ . By Markov's inequality,

$$\begin{aligned} \mathbb{P}_{S \sim \mathcal{D}^m} (L_{\mathcal{D}}(A(S)) > \epsilon) &\leq \frac{\mathbb{E}_{S \sim \mathcal{D}^m} [L_{\mathcal{D}}(A(S))]}{\epsilon} \\ &\leq \frac{\epsilon\delta}{\epsilon} \\ &= \delta. \end{aligned}$$

## Exercise 4.2

Using Hoeffding's inequality on  $L_{\mathcal{D}} \in [a, b]$  we have

$$\mathbb{P}_{S \sim \mathcal{D}^m} (|L_{\mathcal{D}}(h) - L_S(h)| > \epsilon) \leq 2 \exp \left( - \frac{2m\epsilon^2}{(b-a)^2} \right).$$

Then we substitute this into the step where the union bound is used:

$$\begin{aligned} \mathbb{P}_{S \sim \mathcal{D}^m} (\exists h \in \mathcal{H}, |L_S(h) - L_{\mathcal{D}}(h)| > \epsilon) &\leq \sum_{h \in \mathcal{H}} \mathbb{P}_{S \sim \mathcal{D}^m} (|L_{\mathcal{D}}(h) - L_S(h)| > \epsilon) \\ &\leq 2|\mathcal{H}| \exp \left( - \frac{2m\epsilon^2}{(b-a)^2} \right) \end{aligned}$$

The desired bound on the sample complexity follows from requiring  $2|\mathcal{H}| \exp \left( - \frac{2m\epsilon^2}{(b-a)^2} \right) \leq \delta$ .

1. For every  $\alpha \in [0, 1]$ , a convex function  $f$  satisfies

$$f(\alpha a + (1 - \alpha)b) \leq \alpha f(a) + (1 - \alpha)f(b).$$

Substituting  $f(X) = e^{\lambda X}$  and  $\alpha = \frac{b-X}{b-a} \in [0, 1]$  we get

$$e^{\lambda X} \leq \frac{b-X}{b-a} e^{\lambda a} + \frac{X-a}{b-a} e^{\lambda b}.$$

Taking the expectation on both sides and using  $\mathbb{E}[X] = 0$  we have

$$\mathbb{E}[e^{\lambda X}] \leq \frac{b}{b-a} e^{\lambda a} - \frac{a}{b-a} e^{\lambda b}.$$

2. With  $p = -a/(b-a)$  and  $h = \lambda(b-a)$ , we have

$$\begin{aligned} \log\left(\frac{b}{b-a} e^{\lambda a} - \frac{a}{b-a} e^{\lambda b}\right) &= \log(e^{\lambda a}) + \log\left(\frac{b}{b-a} - \frac{a}{b-a} e^{\lambda(b-a)}\right) \\ &= \lambda a + \log\left(1 + \frac{a}{b-a} - \frac{a}{b-a} e^{\lambda(b-a)}\right) \\ &= -hp + \log(1 - p + pe^h). \end{aligned}$$

3. Let  $\theta = \frac{pe^h}{1-p+pe^h}$ . One can compute

$$L'(h) = -p + \theta, \quad L''(h) = \theta(1 - \theta) = -\left(\theta - \frac{1}{2}\right)^2 + \frac{1}{4} \leq \frac{1}{4}.$$

One can also verify  $L(0) = L'(0) = 0$ . Using these remarks on the equation  $L(h) = L(0) + hL'(0) + (h^2/2)L''(\xi)$ , we obtain  $L(h) \leq h^2/8$ . Combining with the previous steps implies

$$\mathbb{E}[e^{\lambda X}] \leq e^{L(\lambda(b-a))} \leq e^{-\lambda^2(a-b)^2/8}.$$

4. Let  $X_i = Z_i - \mu$  and  $\bar{X} = \frac{1}{m} \sum_{i=1}^m X_i$ . Using the monotonicity of the exponent function and Markov's inequality, we have

$$\mathbb{P}(\bar{X} \geq \epsilon) = \mathbb{P}(e^{\lambda \bar{X}} \geq e^{\lambda \epsilon}) \leq e^{-\lambda \epsilon} \mathbb{E}[e^{\lambda \bar{X}}].$$

As  $X_i$  are independent, we have  $\mathbb{E}[e^{\lambda \bar{X}}] = \prod_{i=1}^m \mathbb{E}[e^{\lambda X_i/m}]$ . Also, the previous exercise provides  $\mathbb{E}[e^{\lambda X_i/m}] \leq e^{-\lambda^2(a-b)^2/(8m^2)}$ . So we conclude

$$\mathbb{P}(\bar{X} \geq \epsilon) \leq \exp\left(-\lambda \epsilon + \frac{\lambda^2(b-a)^2}{8m}\right).$$

5. The exponent  $-\lambda \epsilon + \frac{\lambda^2(b-a)^2}{8m}$  is a quadratic (convex) function of  $\lambda$ . It is minimized when  $\lambda = 4m\epsilon/(b-a)^2$ . This optimization gives the desired bound.

**5.1** We simply apply lemma from the hint to obtain

$$\begin{aligned}\mathbb{P}_{S \sim \mathcal{D}^m}(L_{\mathcal{D}}(A(S)) \geq 1/8) &= \mathbb{P}_{S \sim \mathcal{D}^m}(L_{\mathcal{D}}(A(S)) \geq 1 - 7/8) \\ &\geq \frac{\mathbb{E}[L_{\mathcal{D}}(A(S))] - (1 - 7/8)}{7/8} \\ &\geq \frac{1/8}{7/8} = 1/7.\end{aligned}$$

Alternatively, if you dislike Lemma B.1, you can also prove by contrapositive, i.e., showing that if  $\mathbb{P}_{S \sim \mathcal{D}^m}(L_{\mathcal{D}}(A(S)) \geq 1/8) < 1/7$  then  $\mathbb{E}[L_{\mathcal{D}}(A(S))] < 1/4$ . This is easily seen because

$$L_{\mathcal{D}}(A(S)) < 1 \cdot \mathbf{1}_{L_{\mathcal{D}}(A(S)) \geq 1/8} + \frac{1}{8} \cdot \mathbf{1}_{L_{\mathcal{D}}(A(S)) < 1/8}$$

and under the hypothesis

$$\mathbb{E}[L_{\mathcal{D}}(A(S))] < 1 \cdot \frac{1}{7} + \frac{1}{8} \cdot \frac{6}{7} = 1/4.$$

**6.2** (a) Consider a set of  $k+1$  elements. All-one labeling cannot be obtained, so  $\text{VCdim}(\mathcal{H}) \leq k$ . Analogously, for a set of  $|\mathcal{X}| - k + 1$  elements all-zero labeling cannot be obtained, so  $\text{VCdim}(\mathcal{H}_{=k}) \leq \min(k, |\mathcal{X}| - k)$ .

Take a set  $C$  of size  $m = \min(k, |\mathcal{X}| - k)$  and a labeling  $(y_1, \dots, y_m)$  with  $s$  ones,  $0 \leq s \leq m$ . We can pick a hypothesis  $h \in \mathcal{H}_{=k}$  such that  $h(x_i) = y_i$  for all  $x_i \in C$  and it has  $k - s$  ones at the set  $\mathcal{X} \setminus C$ . Therefore,  $C$  is shattered and  $\text{VCdim}(\mathcal{H}_{=k}) \geq \min(k, |\mathcal{X}| - k)$ .

(b) Consider set of  $2k + 2$  elements. It is clear that any labeling with  $k + 1$  ones and  $k + 1$  zeros cannot be obtained, so  $\text{VCdim}(\mathcal{H}_{at-most-k}) \leq 2k + 1$ . Note that it may happen that  $2k + 1 > |\mathcal{X}|$ , so the bound should be  $\text{VCdim}(\mathcal{H}_{at-most-k}) \leq \min(2k + 1, |\mathcal{X}|)$ .

Take a set of  $\min(2k + 1, |\mathcal{X}|)$  elements. Any labeling on this set has either  $\leq k$  zeros or  $\leq k$  ones, so it is shattered by  $\mathcal{H}_{at-most-k}$ . Therefore,  $\text{VCdim}(\mathcal{H}_{at-most-k}) = \min(2k + 1, |\mathcal{X}|)$ .

**6.5** We simply generalize the proof from the two-dimensional case. Let's first formally state the hypothesis class

$$\mathcal{H} = \{h_{(a_i, b_i)} | a_i \leq b_i, h_{(a_i, b_i)}(x_1, \dots, x_d) = \prod_{i=1}^d \mathbf{1}_{a_i \leq x_i \leq b_i}\}$$

Consider set  $\{\mathbf{x}_1, \dots, \mathbf{x}_{2d}\}$ , where  $\mathbf{x}_i = \mathbf{e}_i$  for  $1 \leq i \leq d$  and  $\mathbf{x}_i = -\mathbf{e}_{i-d}$  for  $d+1 \leq i \leq 2d$ . For any labeling  $(y_1, \dots, y_{2d})$ , pick  $a_i = -2$  if  $y_{d+i} = 1$  and  $a_i = -0.5$  otherwise. Similarly, pick  $b_i = 2$  if  $y_i = 1$  and  $b_i = 0.5$  otherwise. Then  $h_{(a_i, b_i)}(\mathbf{x}_i) = y_i$  and hence  $\text{VCdim}(\mathcal{H}) \geq 2d$ .

For a set  $C$  of size  $2d+1$ , by the pigeonhole principle there exists an element  $\mathbf{x}$  s.t.  $\forall j \in [d]$  there exist  $\mathbf{x}', \mathbf{x}'' \in C : x'_j \leq x_j \leq x''_j$ . This means that labeling with only  $\mathbf{x}$  negative and all other elements positive cannot be obtained and therefore  $\text{VCdim}(\mathcal{H}) \leq 2d$ .

**6.8** Let's prove the lemma first.

$$\begin{aligned} \sin(2^m \pi x) &= \sin(2^m \pi \cdot (0.x_1 x_2 \dots)) = \sin(2\pi \cdot (x_1 x_2 \dots x_{m-1} x_m x_{m+1} \dots)) \\ &= \sin(2\pi \cdot (0.x_m x_{m+1} \dots)) \end{aligned}$$

For  $x_m = 0$ , we know that  $\exists k \geq m$  s.t.  $x_k = 1$ , i.e. the number  $0.0x_{m+1} \dots$  is nonzero. This means that  $2\pi \cdot (0.0x_{m+1} \dots) \in (0, \pi)$ , where  $\sin(x)$  is positive, which gives the label 1. For  $x_m = 1$ , we get  $2\pi \cdot (0.1x_{m+1} \dots) \in (\pi, 2\pi)$ , where  $\sin(x)$  is negative, which gives the label 0. Proof completed.

To prove that  $\mathcal{H}$  has infinite VC-dimension, we need to show that for any  $n$  there is a set  $x$  of  $n$  points in  $\mathbb{R}$  on which we can obtain all  $2^n$  possible labelings. Consider  $x_1, \dots, x_n \in [0, 1]$  so that first  $2^n$  bits of their binary expansions give all possible labelings.

Example for  $n = 3$ :

$$\begin{array}{rcccccccccc} x_1 & 0. & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & \dots \\ x_2 & 0. & 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & \dots \\ x_3 & 0. & 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 & \dots \end{array}$$

Using the lemma, invoking the function  $\lceil \sin(2^i \pi x) \rceil$  on the set  $\{x_1, \dots, x_n\}$  for  $1 \leq i \leq 2^n$  allows to obtain all possible labelings. Hence,  $\mathcal{H}$  shatters the set  $\{x_1, \dots, x_n\}$

**6.9**  $\text{VCdim}(\mathcal{H}) = 3$ . In order to prove it, let's recall the unsigned intervals class  $\mathcal{H}_+$ , which was studied during the class. It can be seen that if labeling  $(y_0, y_1, \dots)$  is obtained by  $h_{a,b} \in \mathcal{H}_+$ , then  $h_{a,b,+} \in \mathcal{H}$  gives the same labeling and  $h_{a,b,-} \in \mathcal{H}$  gives its inverse  $(1-y_0, 1-y_1, \dots)$ . Labeling  $(0, 1, 0)$  can be obtained by an interval, so signed intervals can label  $(1, 0, 1)$  and therefore  $\text{VCdim}(\mathcal{H}) \geq 3$ .

Consider the set of 4 points. Labels  $(0, 1, 0, 1)$  and  $(1, 0, 1, 0)$  cannot be obtained with any signed interval, so  $\text{VCdim}(\mathcal{H}) \leq 3$ , which concludes the proof.

**7.3** (a) For any  $h \in \mathcal{H}$  and given  $n(h), |\mathcal{H}_{n(h)}|$ , we can set  $w(h) = \frac{2^{-n(h)}}{|\mathcal{H}_{n(h)}|}$ . This gives

$$\sum_{h \in \mathcal{H}} w(h) = \sum_{h \in \mathcal{H}} \frac{2^{-n(h)}}{|\mathcal{H}_{n(h)}|} = \sum_{n \in \mathbb{N}} \frac{2^{-n}}{|\mathcal{H}_n|} \sum_{\substack{h \in \mathcal{H}_n \\ h \notin \mathcal{H}_{n'}, n' < n}} 1 \leq \sum_{n \in \mathbb{N}} \frac{2^{-n}}{|\mathcal{H}_n|} \sum_{h \in \mathcal{H}_n} 1 = \sum_{n \in \mathbb{N}} 2^{-n} = 1.$$

The equality is achieved when all  $\mathcal{H}_n$  are disjoint

- (b) Since  $\mathcal{H}_n$  is countable, we can enumerate all  $h \in \mathcal{H}_n$  as  $h_{n,1}, h_{n,2}, \dots$ . Consider  $w(h_{n,k}) = 2^{-n}2^{-k}$ . Similarly to the previous exercise, we get

$$\sum_{h \in \mathcal{H}} w(h) \leq \sum_{n \in \mathbb{N}} 2^{-n} \sum_{k \in \mathbb{N}} 2^{-k} = 1.$$

It should be noted that for some  $\mathcal{H}_n$  hypotheses  $h_{n,k}$  may not exist for sufficiently big  $k$  (e.g.  $\mathcal{H}_n$  is finite), but we are only interested in upper bound, so it does not change anything.

### Exercise 1

1.  $f(x) = \max_{1 \leq i \leq m} f_i(\mathbf{x})$  where  $f_i(\mathbf{x}) = \mathbf{a}_i^T \mathbf{x} + b_i$  is convex differentiable with gradient  $\nabla f_i(\mathbf{x}) = \mathbf{a}_i$ . By Claim 14.6, it follows that  $\forall \mathbf{x} : \mathbf{a}_j \in \partial f(\mathbf{x})$  where  $j \in \arg \max_i f_i(\mathbf{x})$ .
2.  $f(x) = \max_{1 \leq i \leq m} f_i(\mathbf{x})$  where  $f_i(\mathbf{x}) = |\mathbf{a}_i^T \mathbf{x} + b_i|$  is convex subdifferentiable. Fix  $\mathbf{x}$ , let  $j \in \arg \max_i f_i(\mathbf{x})$  and choose  $\mathbf{v} \in \partial f_j(\mathbf{x})$  as follows:

$$\mathbf{v} = \begin{cases} -\mathbf{a}_j & \text{if } \mathbf{a}_j^T \mathbf{x} + b_j < 0, \\ 0 & \text{if } \mathbf{a}_j^T \mathbf{x} + b_j = 0, \\ +\mathbf{a}_j & \text{if } \mathbf{a}_j^T \mathbf{x} + b_j > 0. \end{cases}$$

A straightforward generalization of Claim 14.6 shows that  $\mathbf{v}$  is a subgradient of  $f$  at  $\mathbf{x}$ .

3. Note that the sup is really a maximum as  $t \mapsto p(t, \mathbf{x})$  is a continuous function on a compact. Hence  $f(\mathbf{x}) = \max_{t \in [0,1]} p(t, \mathbf{x})$  and  $\forall t \in [0, 1] : \nabla_{\mathbf{x}} p(t, \mathbf{x}) = [1, t, \dots, t^{n-1}]^T \in \mathbb{R}^n$ . A straightforward generalization of Claim 14.6 shows that  $[1, t(\mathbf{x}), \dots, t(\mathbf{x})^{n-1}]^T \in \partial f(\mathbf{x})$ , where  $t(\mathbf{x}) \in \arg \max_{t \in [0,1]} p(t, \mathbf{x})$ .

### Exercise 2

1.  $v$  is a subgradient of  $f$  at 0 if  $\forall u > 0 : f(u) \geq f(0) + (u - 0)v$ , i.e.,

$$\forall u > 0 : 0 \geq 1 + uv. \tag{1}$$

Clearly  $v$  must be negative for the later to hold, and if  $v$  is negative then  $0 \geq 1 + uv \Leftrightarrow u \geq 1/|v|$ . Whatever  $v$ , (1) cannot hold on the whole interval  $[0, +\infty)$ . Hence  $f$  is not subdifferentiable at 0.

2.  $v$  is a subgradient of  $f$  at 0 if  $\forall u > 0 : f(u) \geq f(0) + (u - 0)v$ , i.e.,

$$\forall u > 0 : -1 \geq \sqrt{uv}. \tag{2}$$

Clearly  $v$  must be negative for the later to hold, and if  $v$  is negative then  $-1 \geq \sqrt{uv} \Leftrightarrow u \geq 1/v^2$ . Whatever  $v$ , (2) cannot hold on the whole interval  $[0, +\infty)$ . Hence  $f$  is not subdifferentiable at 0.



### Exercise 3

Fix  $\mathbf{w}, \mathbf{u}$ . The function  $f$  is  $\lambda$ -strongly convex, so for all  $\alpha \in [0, 1]$  we have:

$$\begin{aligned} f((1-\alpha)\mathbf{w} + \alpha\mathbf{u}) &\leq (1-\alpha)f(\mathbf{w}) + \alpha f(\mathbf{u}) - \frac{\lambda}{2}\alpha(1-\alpha)\|\mathbf{w} - \mathbf{u}\|^2 \\ \Leftrightarrow f(\mathbf{w} + \alpha(\mathbf{u} - \mathbf{w})) - f(\mathbf{w}) &\leq \alpha\left(f(\mathbf{u}) - f(\mathbf{w}) - \frac{\lambda}{2}(1-\alpha)\|\mathbf{w} - \mathbf{u}\|^2\right) \end{aligned} \quad (3)$$

Let  $\mathbf{v} \in \partial f(\mathbf{w})$ . Then,  $\forall \alpha \in [0, 1] : f(\mathbf{w} + \alpha(\mathbf{u} - \mathbf{w})) \geq f(\mathbf{w}) + \langle \alpha(\mathbf{u} - \mathbf{w}), \mathbf{v} \rangle$ . Combining this inequality and (3) gives:

$$\begin{aligned} \langle \alpha(\mathbf{u} - \mathbf{w}), \mathbf{v} \rangle &\leq \alpha\left(f(\mathbf{u}) - f(\mathbf{w}) - \frac{\lambda}{2}(1-\alpha)\|\mathbf{w} - \mathbf{u}\|^2\right) \\ \Leftrightarrow \langle \mathbf{u} - \mathbf{w}, \mathbf{v} \rangle &\leq f(\mathbf{u}) - f(\mathbf{w}) - \frac{\lambda}{2}(1-\alpha)\|\mathbf{w} - \mathbf{u}\|^2 \\ \Leftrightarrow \langle \mathbf{w} - \mathbf{u}, \mathbf{v} \rangle &\geq f(\mathbf{w}) - f(\mathbf{u}) + \frac{\lambda}{2}(1-\alpha)\|\mathbf{w} - \mathbf{u}\|^2 \end{aligned}$$

Taking the limit  $\alpha \rightarrow 0+$  ends the proof:  $\langle \mathbf{w} - \mathbf{u}, \mathbf{v} \rangle \geq f(\mathbf{w}) - f(\mathbf{u}) + \frac{\lambda}{2}\|\mathbf{w} - \mathbf{u}\|^2$ .

### Exercise 4

To prove that  $\pi_C(\cdot)$  is Lipschitzian, we first show an important property of projection onto a closed convex set:

**Lemma 1.** *If  $C$  is a non-empty closed convex subset of a Hilbert space  $H$  then  $\forall(\mathbf{x}, \mathbf{y}) \in H \times C : \langle \mathbf{x} - \pi_C(\mathbf{x}), \mathbf{y} - \pi_C(\mathbf{x}) \rangle \leq 0$ .*

*Proof.* Let  $\alpha \in (0, 1)$ . By definition of  $\pi_C(\cdot)$ , we have:

$$\begin{aligned} 0 &\leq \|\mathbf{x} - (1-\alpha)\pi_C(\mathbf{x}) - \alpha\mathbf{y}\|^2 - \|\mathbf{x} - \pi_C(\mathbf{x})\|^2 \\ &= \|\mathbf{x} - \pi_C(\mathbf{x}) - \alpha(\mathbf{y} - \pi_C(\mathbf{x}))\|^2 - \|\mathbf{x} - \pi_C(\mathbf{x})\|^2 \\ &= \alpha^2\|\mathbf{y} - \pi_C(\mathbf{x})\|^2 - 2\alpha\langle \mathbf{x} - \pi_C(\mathbf{x}), \mathbf{y} - \pi_C(\mathbf{x}) \rangle. \end{aligned}$$

Dividing the final inequality by  $\alpha$  and taking the limit  $\alpha \rightarrow 0$  ends the proof.  $\square$

We can now prove that  $\pi_C(\cdot)$  is 1-Lipschitz.  $\forall \mathbf{x}_0, \mathbf{x}_1 :$

$$\begin{aligned} \|\pi_C(\mathbf{x}_0) - \pi_C(\mathbf{x}_1)\|^2 &= \langle \pi_C(\mathbf{x}_0) - \pi_C(\mathbf{x}_1), \pi_C(\mathbf{x}_0) - \pi_C(\mathbf{x}_1) \rangle \\ &= \underbrace{\langle \pi_C(\mathbf{x}_0) - \mathbf{x}_0, \pi_C(\mathbf{x}_0) - \pi_C(\mathbf{x}_1) \rangle}_{\leq 0} + \langle \mathbf{x}_0 - \pi_C(\mathbf{x}_1), \pi_C(\mathbf{x}_0) - \pi_C(\mathbf{x}_1) \rangle \\ &\leq \langle \mathbf{x}_0 - \pi_C(\mathbf{x}_1), \pi_C(\mathbf{x}_0) - \pi_C(\mathbf{x}_1) \rangle \\ &\leq \underbrace{\langle \mathbf{x}_1 - \pi_C(\mathbf{x}_1), \pi_C(\mathbf{x}_0) - \pi_C(\mathbf{x}_1) \rangle}_{\leq 0} + \langle \mathbf{x}_0 - \mathbf{x}_1, \pi_C(\mathbf{x}_0) - \pi_C(\mathbf{x}_1) \rangle \\ &\leq \langle \mathbf{x}_0 - \mathbf{x}_1, \pi_C(\mathbf{x}_0) - \pi_C(\mathbf{x}_1) \rangle \\ &\leq \|\mathbf{x}_0 - \mathbf{x}_1\| \|\pi_C(\mathbf{x}_0) - \pi_C(\mathbf{x}_1)\| \quad (\text{Cauchy-Schwarz inequality}) \end{aligned}$$

It directly implies  $\|\pi_C(\mathbf{x}_0) - \pi_C(\mathbf{x}_1)\| \leq \|\mathbf{x}_0 - \mathbf{x}_1\|$ . Note that for  $\mathbf{x}_0, \mathbf{x}_1 \in C$  this inequality is an equality, hence it cannot be improved.

### Exercise 1

a) Fix  $A, B \in \mathcal{S}_n^+$  and  $\alpha \in [0, 1]$ . Let  $\mathbf{e} \in \mathbb{R}^n$  a unit-norm eigenvector of  $\alpha A + (1 - \alpha)B$  associated to the maximum eigenvalue, i.e.,  $(\alpha A + (1 - \alpha)B)\mathbf{e} = \lambda_{\max}(\alpha A + (1 - \alpha)B)\mathbf{e}$  and  $\|\mathbf{e}\| = 1$ . We have:

$$\begin{aligned} f(\alpha A + (1 - \alpha)B) &= \mathbf{e}^T(\alpha A + (1 - \alpha)B)\mathbf{e} = \alpha \mathbf{e}^T A \mathbf{e} + (1 - \alpha) \mathbf{e}^T B \mathbf{e} \\ &\leq \alpha \lambda_{\max}(A) + (1 - \alpha) \lambda_{\max}(B) \\ &= \alpha f(A) + (1 - \alpha) f(B). \end{aligned}$$

This shows that  $f$  is convex.

b) Let  $A \in \mathcal{S}_n^+$ . A subgradient of  $f$  at  $A$  is a matrix  $V \in \mathbb{R}^{n \times n}$  that satisfies:

$$\forall B \in \mathcal{S}_n^+ : f(B) \geq f(A) + \text{Tr}((B - A)^T V).$$

Consider any  $\mathbf{e} \in \mathbb{R}^n$  which is a unit-norm eigenvector of  $A$  associated to the maximum eigenvalue, i.e.,  $A\mathbf{e} = \lambda_{\max}(A)\mathbf{e}$  and  $\|\mathbf{e}\| = 1$ . Then for all  $B \in \mathcal{S}_n^+$ :

$$\begin{aligned} f(A) &= \lambda_{\max}(A) = \mathbf{e}^T A \mathbf{e} = \mathbf{e}^T B \mathbf{e} + \mathbf{e}^T (A - B) \mathbf{e} \leq \lambda_{\max}(B) + \mathbf{e}^T (A - B) \mathbf{e} \\ &= f(B) + \text{Tr}(\mathbf{e}^T (A - B) \mathbf{e}) \\ &= f(B) + \text{Tr}((A - B)^T \mathbf{e} \mathbf{e}^T). \end{aligned}$$

In the last equality we used that  $(A - B)^T = A - B$  and that the trace is preserved by cyclic permutations. We see that  $\mathbf{e} \mathbf{e}^T$  satisfies the definition of a subgradient:  $\mathbf{e} \mathbf{e}^T \in \partial f(A)$ .

### Exercise 2

a)  $\min_{\|\mathbf{w}\| \leq \|\mathbf{w}^*\|} f(\mathbf{w}) \leq f(\mathbf{w}^*) \leq 0$  because  $\forall i \in [m] : y_i \langle \mathbf{w}^*, \mathbf{x}_i \rangle \geq 1$ . Suppose there exists  $\mathbf{w}$  satisfying both  $\|\mathbf{w}\| \leq \|\mathbf{w}^*\|$  and  $f(\mathbf{w}) < 0$ . Then  $\mathbf{w}$  can be slightly modify to obtain a vector  $\tilde{\mathbf{w}}$  such that  $\|\tilde{\mathbf{w}}\| < \|\mathbf{w}^*\|$ , while still having  $f(\tilde{\mathbf{w}}) \leq 0$ . It contradicts  $\mathbf{w}^*$ 's definition, hence  $\min_{\|\mathbf{w}\| \leq \|\mathbf{w}^*\|} f(\mathbf{w}) \geq 0$ . It proves  $\min_{\|\mathbf{w}\| \leq \|\mathbf{w}^*\|} f(\mathbf{w}) = 0$ .

b) If  $f(\mathbf{w}) < 1$  then  $\forall i \in [m] : y_i \langle \mathbf{w}^*, \mathbf{x}_i \rangle > 0$ , i.e.,  $\mathbf{w}$  separates the examples.

c) For all  $i \in [m]$  the gradient of  $f_i : \mathbf{w} \mapsto 1 - y_i \langle \mathbf{w}, \mathbf{x}_i \rangle$  is  $-y_i \mathbf{x}_i$ . Applying Claim 14.6, we get that a subgradient of  $f$  at  $\mathbf{w}$  is given by  $-y_{i^*} \mathbf{x}_{i^*}$  where  $i^* \in \arg \max_{i \in [m]} \{1 - y_i \langle \mathbf{w}, \mathbf{x}_i \rangle\}$ .

d) The algorithm is inialized with  $\mathbf{w}^{(1)} = 0$ . At each iteration, if  $f(\mathbf{w}^{(t)}) \geq 1$  then it chooses  $i^* \in \arg \min_{i \in [m]} \{y_i \langle \mathbf{w}^{(t)}, \mathbf{x}_i \rangle\}$  and updates  $\mathbf{w}^{(t+1)} = \mathbf{w}^{(t)} + \eta y_{i^*} \mathbf{x}_{i^*}$ . Otherwise, if

$f(\mathbf{w}^{(t)}) < 1$ ,  $\mathbf{w}^{(t)}$  separates all the examples and we stop. To analyze the speed of convergence of the subgradient algorithm, first notice that  $\langle \mathbf{w}^*, \mathbf{w}^{(t+1)} \rangle - \langle \mathbf{w}^*, \mathbf{w}^{(t)} \rangle = \eta y_{i^*} \langle \mathbf{w}^*, \mathbf{x}_{i^*} \rangle \geq \eta$ . Therefore, after performing  $T$  iterations, we have

$$\langle \mathbf{w}^*, \mathbf{w}^{(T+1)} \rangle = \langle \mathbf{w}^*, \mathbf{w}^{(T+1)} \rangle - \langle \mathbf{w}^*, \mathbf{w}^{(1)} \rangle = \sum_{t=1}^T \langle \mathbf{w}^*, \mathbf{w}^{(t+1)} \rangle - \langle \mathbf{w}^*, \mathbf{w}^{(t)} \rangle \geq \eta T. \quad (1)$$

Besides,  $\|\mathbf{w}^{(t+1)}\|^2 = \|\mathbf{w}^{(t)}\|^2 + \eta^2 y_{i^*}^2 \|\mathbf{x}_{i^*}\|^2 + 2\eta y_{i^*} \langle \mathbf{w}^{(t)}, \mathbf{x}_{i^*} \rangle \leq \|\mathbf{w}^{(t)}\|^2 + \eta^2 R^2$ . The last inequality follows from  $\|\mathbf{x}_i\| \leq R$  and  $y_i \langle \mathbf{w}^{(t)}, \mathbf{x}_{i^*} \rangle \leq 0$  (we update only if  $f(\mathbf{w}^{(t)}) \geq 1$ ). Then

$$\|\mathbf{w}^{(T+1)}\| \leq \eta R \sqrt{T}. \quad (2)$$

Combining Cauchy-Schwarz inequality, (1) and (2), we obtain

$$1 \geq \frac{\langle \mathbf{w}^*, \mathbf{w}^{(T+1)} \rangle}{\|\mathbf{w}^{(T+1)}\| \|\mathbf{w}^*\|} \geq \frac{\sqrt{T}}{R \|\mathbf{w}^*\|}. \quad (3)$$

The subgradient algorithm must stop in less than  $R^2 \|\mathbf{w}^*\|^2$  iterations. We see that  $\eta$  does not affect the speed of convergence. The algorithm is almost identical to the Batch Perceptron algorithm with two modifications. First, the Batch Perceptron updates with any example for which  $y_i \langle \mathbf{w}^{(t)}, \mathbf{x}_i \rangle \leq 0$ , while the current algorithm chooses the example for which  $y_i \langle \mathbf{w}^{(t)}, \mathbf{x}_i \rangle$  is minimal. Second, the current algorithm employs the parameter  $\eta$ . However, the only difference with the case  $\eta = 1$  is that it scales  $\mathbf{w}^{(t)}$  by  $\eta$ .

### Exercise 3

We prove the following Theorem:

**Theorem 1.** *Let  $B, \rho > 0$ . Let  $f$  be a convex function and let  $\mathbf{w}^* \in \arg \min_{\mathbf{w}: \|\mathbf{w}\| \leq B} f(\mathbf{w})$ . Assume that SGD is run for  $T$  iterations with  $\eta_t = \frac{B}{\rho \sqrt{t}}$ . Assume also that for all  $t$ ,  $\mathbb{E} \|\mathbf{v}_t\|^2 \leq \rho^2$ . Then*

$$\mathbb{E}_{\mathbf{v}_{1:T}}[f(\bar{\mathbf{w}})] - f(\mathbf{w}^*) \leq \frac{3\rho B}{\sqrt{T}}$$

*Proof.* By Jensen's inequality, we have:

$$\mathbb{E}_{\mathbf{v}_{1:T}}[f(\bar{\mathbf{w}})] - f(\mathbf{w}^*) \leq \mathbb{E}_{\mathbf{v}_{1:T}} \left[ \frac{1}{T} \sum_{t=1}^T f(\mathbf{w}^{(t)}) - f(\mathbf{w}^*) \right]. \quad (4)$$

As  $\forall t : \mathbb{E}[\mathbf{v}_t | \mathbf{w}^{(t)}] \in \partial f(\mathbf{w}^{(t)})$ , we can reproduce what is done in Theorem 14.8 to get the inequality:

$$\mathbb{E}_{\mathbf{v}_{1:T}} \left[ \frac{1}{T} \sum_{t=1}^T f(\mathbf{w}^{(t)}) - f(\mathbf{w}^*) \right] \leq \mathbb{E}_{\mathbf{v}_{1:T}} \left[ \frac{1}{T} \sum_{t=1}^T \langle \mathbf{w}^{(t)} - \mathbf{w}^*, \mathbf{v}_t \rangle \right]. \quad (5)$$

We now have to prove an upper bound on the right-hand side of (5). This is similar to what is done in Lemma 14.10, except that we have to take into account the time-dependence of

the steps  $\eta_t$ . For all  $t \in \{1, \dots, T\}$ :

$$\begin{aligned}
\langle \mathbf{w}^{(t)} - \mathbf{w}^*, \mathbf{v}_t \rangle &= \frac{1}{\eta_t} \langle \mathbf{w}^{(t)} - \mathbf{w}^*, \eta_t \mathbf{v}_t \rangle = \frac{1}{2\eta_t} (\|\mathbf{w}^{(t)} - \mathbf{w}^*\|^2 - \|\mathbf{w}^{(t)} - \mathbf{w}^* - \eta_t \mathbf{v}_t\|^2 + \eta_t^2 \|\mathbf{v}_t\|^2) \\
&= \frac{1}{2\eta_t} (\|\mathbf{w}^{(t)} - \mathbf{w}^*\|^2 - \|\mathbf{w}^{(t+1/2)} - \mathbf{w}^*\|^2 + \eta_t^2 \|\mathbf{v}_t\|^2) \\
&\leq \frac{1}{2\eta_t} (\|\mathbf{w}^{(t)} - \mathbf{w}^*\|^2 - \|\mathbf{w}^{(t+1)} - \mathbf{w}^*\|^2) + \frac{\eta_t}{2} \|\mathbf{v}_t\|^2. \quad (6)
\end{aligned}$$

Let  $\mathcal{H} = \{\mathbf{w} : \|\mathbf{w}\| \leq B\}$ . The last inequality follows from  $\mathbf{w}^{(t+1)} = \pi_{\mathcal{H}}(\mathbf{w}^{(t+1/2)})$  and the 1-Lipschitzianity of  $\pi_{\mathcal{H}}$  (see Homework 4, Exercise 4):

$$\|\pi_{\mathcal{H}}(\mathbf{w}^{(t+1/2)}) - \mathbf{w}^*\| = \|\pi_{\mathcal{H}}(\mathbf{w}^{(t+1/2)}) - \pi_{\mathcal{H}}(\mathbf{w}^*)\| \leq \|\mathbf{w}^{(t+1/2)} - \mathbf{w}^*\|.$$

Summing the inequality (6) over  $t$ , we have:

$$\begin{aligned}
\sum_{t=1}^T \langle \mathbf{w}^{(t)} - \mathbf{w}^*, \mathbf{v}_t \rangle &\leq \sum_{t=1}^T \frac{1}{2\eta_t} (\|\mathbf{w}^{(t)} - \mathbf{w}^*\|^2 - \|\mathbf{w}^{(t+1)} - \mathbf{w}^*\|^2) + \frac{\eta_t}{2} \|\mathbf{v}_t\|^2 \\
&= \frac{1}{2\eta_1} \|\mathbf{w}^{(1)} - \mathbf{w}^*\|^2 + \sum_{t=1}^{T-1} \frac{\|\mathbf{w}^{(t+1)} - \mathbf{w}^*\|^2}{2} \left( \frac{1}{\eta_{t+1}} - \frac{1}{\eta_t} \right) \\
&\quad - \frac{1}{2\eta_T} \|\mathbf{w}^{(T+1)} - \mathbf{w}^*\|^2 + \sum_{t=1}^T \frac{\eta_t}{2} \|\mathbf{v}_t\|^2 \\
&\leq \frac{1}{2\eta_1} \|\mathbf{w}^{(1)} - \mathbf{w}^*\|^2 + \sum_{t=1}^{T-1} \frac{\|\mathbf{w}^{(t+1)} - \mathbf{w}^*\|^2}{2} \left( \frac{1}{\eta_{t+1}} - \frac{1}{\eta_t} \right) + \sum_{t=1}^T \frac{\eta_t}{2} \|\mathbf{v}_t\|^2 \\
&\leq 2B^2 \left( \frac{1}{\eta_1} + \sum_{t=1}^{T-1} \frac{1}{\eta_{t+1}} - \frac{1}{\eta_T} \right) + \sum_{t=1}^T \frac{\eta_t}{2} \|\mathbf{v}_t\|^2 \\
&= \frac{2B^2}{\eta_T} + \sum_{t=1}^T \frac{\eta_t}{2} \|\mathbf{v}_t\|^2. \quad (7)
\end{aligned}$$

Taking the expectation of inequality (7) and dividing by  $T$ , we obtain:

$$\mathbb{E}_{\mathbf{v}_{1:T}} \left[ \frac{1}{T} \sum_{t=1}^T \langle \mathbf{w}^{(t)} - \mathbf{w}^*, \mathbf{v}_t \rangle \right] \leq \frac{2B^2}{T\eta_T} + \sum_{t=1}^T \frac{\eta_t}{2T} \mathbb{E} \|\mathbf{v}_t\|^2 \leq \frac{2\rho B}{\sqrt{T}} + \frac{\rho^2}{2T} \sum_{t=1}^T \eta_t. \quad (8)$$

The last inequality follows from the assumption  $\mathbb{E} \|\mathbf{v}_t\|^2 \leq \rho^2$  and  $\eta_T$ 's definition. Besides

$$\sum_{t=1}^T \eta_t = \frac{B}{\rho} \sum_{t=1}^T \frac{1}{\sqrt{t}} \leq \frac{B}{\rho} \left( 1 + \sum_{t=2}^T \int_{t-1}^t \frac{dx}{\sqrt{x}} \right) = \frac{B}{\rho} \left( 1 + \int_1^T \frac{dx}{\sqrt{x}} \right) = \frac{B}{\rho} (2\sqrt{T} - 1).$$

Combining this last inequality with (4), (5) and (8), we finally obtain:

$$\mathbb{E}_{\mathbf{v}_{1:T}} [f(\bar{\mathbf{w}})] - f(\mathbf{w}^*) \leq \frac{2\rho B}{\sqrt{T}} + \frac{\rho B}{2T} (2\sqrt{T} - 1) \leq \frac{3\rho B}{\sqrt{T}}.$$

It concludes the proof.  $\square$

#### Exercise 4

$\mathcal{H}_{n\text{-parity}}$  is a finite class, therefore (see paragraph 6.3.4):

$$\text{VCdim}(\mathcal{H}_{n\text{-parity}}) \leq \log_2 |\mathcal{H}_{n\text{-parity}}| = \log_2 2^n = n.$$

We now show that this upperbound on  $\text{VCdim}(\mathcal{H}_{n\text{-parity}})$  is tight, i.e., there exists  $n$  points in  $\{0, 1\}^n$  that are shattered by  $\mathcal{H}_{n\text{-parity}}$ . Let  $\mathbf{e}^{(j)} \in \{0, 1\}^n$  be such that  $\mathbf{e}_j^{(j)} = 1$  and  $\forall i \neq j : \mathbf{e}_i^{(j)} = 0$ . The subset  $C = \{\mathbf{e}^{(j)}\}_{j=1}^n$  of  $n$  points is shattered by  $\mathcal{H}_{n\text{-parity}}$ . Indeed, given  $(y_1, \dots, y_n) \in \{0, 1\}^n$ , we can define  $J = \{j \in \{1, \dots, n\} : y_j = 1\}$  and see that:

$$\forall j \in \{1, \dots, n\} : h_J(\mathbf{e}^{(j)}) = \sum_{i \in J} \mathbf{e}_i^{(j)} \mod 2 = \sum_{i=1}^n \mathbf{e}_i^{(j)} y_i \mod 2 = y_j.$$

Hence  $\text{VCdim}(\mathcal{H}_{n\text{-parity}}) = n$ .

**Problem 1**

- 1) The joint distribution is (up to normalisation factors of Gaussians)

$$p(y_1, \dots, y_m, x_1, \dots, x_m, w_1, \dots, w_p) \propto \prod_{i=1}^m e^{-\frac{1}{2\sigma^2}(y_i - \sum_{a=1}^p w_a x_i^a)^2} \prod_{i=1}^m P_0(x_i) \prod_{a=1}^p e^{-\alpha w_a^2}$$

- 2) Here the  $x_i$  is a parent of  $y_i$  (for all  $i = 1, \dots, m$ ) and  $w_1, \dots, w_p$  are parents of each  $y_i, i = 1, \dots, m$ .
- 3) The ML principle says that you maximize the log-likelihood  $\log P(\text{data} \mid w_1, \dots, w_p)$ . Since

$$P(\text{data} \mid w_1, \dots, w_p) \propto \prod_{i=1}^m e^{-\frac{1}{2\sigma^2}(y_i - \sum_{a=1}^p w_a x_i^a)^2} \prod_{i=1}^m P_0(x_i)$$

this is equivalent to minimising

$$\mathcal{E}_{\text{data}}(f) = \frac{1}{m} \sum_{i=1}^m (y_i - f(x_i))^2$$

over functions in the class  $\mathcal{H} \ni f(x) = \sum_{a=1}^p w_a x^a$ .

- 4) The posterior distribution is

$$\begin{aligned} P(w_1, \dots, w_p \mid \text{data}) &= \frac{p(y_1, \dots, y_m, x_1, \dots, x_m, w_1, \dots, w_p)}{\int \prod_{a=1}^p dw_a p(y_1, \dots, y_m, x_1, \dots, x_m, w_1, \dots, w_p)} \\ &= \frac{\prod_{i=1}^m e^{-\frac{1}{2\sigma^2}(y_i - \sum_{a=1}^p w_a x_i^a)^2} \prod_{a=1}^p e^{-\alpha w_a^2}}{\int \prod_{a=1}^p dw_a \prod_{i=1}^m e^{-\frac{1}{2\sigma^2}(y_i - \sum_{a=1}^p w_a x_i^a)^2} \prod_{a=1}^p e^{-\alpha w_a^2}} \end{aligned}$$

The MAP principle says you maximise the posterior which is equivalent to minimizing

$$\frac{1}{m} \sum_{i=1}^m (y_i - f(x_i))^2 + 2\alpha\sigma^2 \sum_{a=1}^p w_a^2$$

over the functions in the class  $\mathcal{H} \ni f(x) = \sum_{a=1}^p w_a x^a$ .

- 5) The optimal regression function is  $f_{\text{regr}}(x) = \mathbb{E}_{w|data} \mathbb{E}_{y|x,w}[y]$ . From the model it is clear that

$$\mathbb{E}_{y|x,w}[y] = \sum_{a=1}^p w_a x^a$$

Further average over the posterior gives

$$f_{\text{regr}}(x) = \sum_{a=1}^p \mathbb{E}_{w|data}[w_a] x^a$$

## Problem 2

- 1)  $a \perp\!\!\!\perp b|c$  because  $p(a, b|c) = \frac{p(a,b,c)}{p(c)} = \frac{p(a|c)p(b|c)p(c)}{p(c)} = p(a|c)p(b|c)$ . But  $a, b$  are not independent because  $p(a, b) = \sum_c p(a|c)p(b|c) \neq p(a)p(b)$ .
- 2)  $a \perp\!\!\!\perp b|c$  because  $p(a, b|c) = \frac{p(a,b,c)}{p(c)} = \frac{p(b|c)p(c|a)p(a)}{p(c)} = p(b|c) \frac{p(c|a)p(a)}{p(c)} = p(b|c)p(a|c)$ . But  $a, b$  are not independent because  $p(a, b) = \sum_c p(a)p(c|a)p(b|c) = p(a)p(b|a) \neq p(a)p(b)$ .
- 3)  $a \perp\!\!\!\perp b$  because

$$p(a, b) = \sum_{c,d} p(a, b, c, d) = \sum_{c,d} p(a)p(b)p(c|a, b)p(d|c) = p(a)p(b) \sum_{c,d} p(c|a, b)p(d|c) = p(a)p(b).$$

However, we don't have  $a \perp\!\!\!\perp b|c$  because  $p(a, b|c) = \frac{p(a)p(b)p(c|a,b)}{p(c)}$  cannot be decomposed.

## Problem 3

The left hand side is

$$p(x_i | \mathbf{x}_{\sim i}) = \frac{p(\mathbf{x})}{\int dx_i p(\mathbf{x})} \quad (1)$$

where

$$p(\mathbf{x}) = p(x_i | \{x_v\}_{v \in \text{pa}(i)}) \prod_{k \in \text{child}(j)} p(x_j | \{x_v\}_{v \in \text{pa}(k)}) \prod_{\substack{l \neq i \\ l \neq \text{child}(i)}} p(x_l | \{x_v\}_{v \in \text{pa}(l)}).$$

The product  $\prod_{\substack{l \neq i \\ l \neq \text{child}(i)}} p(x_l | \{x_v\}_{v \in \text{pa}(l)})$  is independent of  $x_i$ . It cancels with the same factor in the denominator of (1). So we have

$$p(x_i | \mathbf{x}_{\sim i}) = \frac{p(x_i | \{x_v\}_{v \in \text{pa}(i)}) \prod_{k \in \text{child}(j)} p(x_j | \{x_v\}_{v \in \text{pa}(k)})}{\int dx_i p(x_i | \{x_v\}_{v \in \text{pa}(i)}) \prod_{k \in \text{child}(j)} p(x_j | \{x_v\}_{v \in \text{pa}(k)})} \quad (2)$$

On the other hand, the right hand side is

$$p(x_i | \{x_v\}_{v \in \text{MB}(i)}) = \frac{p(x_i, \{x_v\}_{v \in \text{MB}(i)})}{\int dx_i p(x_i, \{x_v\}_{v \in \text{MB}(i)})} \quad (3)$$

where

$$\begin{aligned}
& p(x_i, \{x_v\}_{v \in \text{MB}(i)}) \\
&= \int d\mathbf{x}_{\sim i, \text{MB}(i)} p(\mathbf{x}) \\
&= \int d\mathbf{x}_{\sim i, \text{MB}(i)} p(x_i | \{x_v\}_{v \in \text{pa}(i)}) \prod_{k \in \text{child}(j)} p(x_j | \{x_v\}_{v \in \text{pa}(k)}) \prod_{\substack{l \neq i \\ l \neq \text{child}(i)}} p(x_l | \{x_v\}_{v \in \text{pa}(l)}) \\
&= p(x_i | \{x_v\}_{v \in \text{pa}(i)}) \prod_{k \in \text{child}(j)} p(x_j | \{x_v\}_{v \in \text{pa}(k)}) \left[ \int d\mathbf{x}_{\sim i, \text{MB}(i)} \prod_{\substack{l \neq i \\ l \neq \text{child}(i)}} p(x_l | \{x_v\}_{v \in \text{pa}(l)}) \right]
\end{aligned}$$

We identify  $\int d\mathbf{x}_{\sim i, \text{MB}(i)} \prod_{\substack{l \neq i \\ l \neq \text{child}(i)}} p(x_l | \{x_v\}_{v \in \text{pa}(l)})$  independent of  $x_i$ . It cancels with the same factor in the denominator of (3). So (3) is reduced to the same expression as (2).

#### Problem 4 (Bishop, p.371 & 419, Exercise 8.7)

Using  $\mathbb{E}[x_i] = \sum_{j \in \text{pa}(i)} w_{ij} \mathbb{E}[x_j] + b_i$  gives

$$\begin{aligned}
\mu_1 &= \sum_{j \in \emptyset} w_{1j} \mathbb{E}[x_j] + b_1 = b_1 \\
\mu_2 &= \sum_{j \in \{1\}} w_{2j} \mathbb{E}[x_j] = w_{21} b_1 + b_2 \\
\mu_3 &= \sum_{j \in \{2\}} w_{3j} \mathbb{E}[x_j] + b_3 = w_{32} (w_{21} b_1 + b_2) + b_3
\end{aligned}$$

Using  $\text{cov}[x_i, x_j] = \sum_{k \in \text{pa}(j)} w_{jk} \text{cov}[x_i, x_k] + I_{ij} v_j$  for  $i \leq j$  and  $\text{cov}[x_i, x_j] = \text{cov}[x_j, x_i]$  gives

$$\begin{aligned}
\text{cov}[x_1, x_1] &= \sum_{k \in \emptyset} w_{1k} \text{cov}[x_1, x_k] + v_1 = v_1 \\
\text{cov}[x_1, x_2] &= \sum_{k \in \{1\}} w_{2k} \text{cov}[x_1, x_k] = w_{21} v_1 \\
\text{cov}[x_1, x_3] &= \sum_{k \in \{2\}} w_{3k} \text{cov}[x_1, x_k] = w_{32} (w_{21} v_1) \\
\text{cov}[x_2, x_2] &= \sum_{k \in \{1\}} w_{2k} \text{cov}[x_2, x_k] + v_2 = w_{21} (w_{21} v_1) + v_2 \\
\text{cov}[x_2, x_3] &= \sum_{k \in \{2\}} w_{3k} \text{cov}[x_2, x_k] = w_{32} (w_{21}^2 v_1 + v_2) \\
\text{cov}[x_3, x_3] &= \sum_{k \in \{2\}} w_{3k} \text{cov}[x_3, x_k] + v_3 = w_{32}^2 (w_{21}^2 v_1 + v_2) + v_3
\end{aligned}$$



**Problem 5 (Barber, p.75, Exercise 4.4)**

1) First note that

$$p(\mathbf{h}|\mathbf{v}) \propto e^{(\mathbf{v}^\top \mathbf{W} + \mathbf{b}^\top) \mathbf{h}} = \prod_i e^{h_i(b_i + \sum_j W_{ji} v_j)}$$

So  $p(\mathbf{h}|\mathbf{v}) = \prod_i p(h_i|\mathbf{v})$ . Recall  $h_i \in \{0, 1\}$ . Thus we have

$$p(h_i = 1|\mathbf{v}) = \frac{e^{b_i + \sum_j W_{ji} v_j}}{\sum_{h_i \in \{0,1\}} e^{h_i(b_i + \sum_j W_{ji} v_j)}} = \sigma\left(b_i + \sum_j W_{ji} v_j\right).$$

2)

$$p(\mathbf{v}|\mathbf{h}) = \prod_i p(v_i|\mathbf{h}), \quad \text{with } p(v_i = 1|\mathbf{h}) = \sigma\left(a_i + \sum_j W_{ij} h_j\right)$$

3) No. Because the term  $\mathbf{v}^\top \mathbf{W} \mathbf{h}$  in  $p(\mathbf{v}, \mathbf{h})$  introduces dependence between  $\mathbf{v}$  and  $\mathbf{h}$ .

4) For a general  $\mathbf{W}$  there is no known efficient way to compute  $Z$  efficiently. The dependence between  $\mathbf{v}$  and  $\mathbf{h}$  does not allow always decomposition of  $p(\mathbf{v}, \mathbf{h})$ .

**Problem 6 (Barber, p.77, Exercise 4.14)**

We write

$$\begin{aligned} \phi_{ij}(x_i, x_j) &= e^{\ln \phi_{ij}(x_i, x_j)} \\ &= e^{\mathbb{I}(x_i=0, x_j=0) \ln \phi_{ij}(0,0) + \mathbb{I}(x_i=0, x_j=1) \ln \phi_{ij}(0,1) + \mathbb{I}(x_i=1, x_j=0) \ln \phi_{ij}(1,0) + \mathbb{I}(x_i=1, x_j=1) \ln \phi_{ij}(1,1)} \end{aligned}$$

With  $x_i \in \{0, 1\}$  we can replace  $\mathbb{I}[\cdot]$  by

$$\begin{aligned} \mathbb{I}(x_i = 0, x_j = 0) &= (1 - x_i)(1 - x_j), & \mathbb{I}(x_i = 0, x_j = 1) &= (1 - x_i)x_j, \\ \mathbb{I}(x_i = 1, x_j = 0) &= x_i(1 - x_j), & \mathbb{I}(x_i = 1, x_j = 1) &= x_i x_j. \end{aligned}$$

So  $\phi_{ij}(x_i, x_j)$  is in the form  $e^{W_{ij}x_i x_j + b_i x_i + b_j x_j + \text{constant}}$  and  $p(\mathbf{x}) = \frac{1}{Z} e^{\sum_{ij \in \mathcal{E}} W_{ij} x_i x_j + \sum_i \deg(i) b_i x_i}$  is the Boltzmann machine.

**Problem 7**

Fix a subset  $S \subseteq V$ . We have:

$$\begin{aligned} p(\mathbf{x}_S, \mathbf{x}_{V \setminus S}) &= p(\mathbf{x}) = \frac{1}{Z} \prod_{\substack{C' \in \mathcal{C}: \\ S \cap C' = \emptyset}} \psi_{C'}(\mathbf{x}_{C'}) \cdot \prod_{\substack{C \in \mathcal{C}: \\ S \cap C \neq \emptyset}} \psi_C(\mathbf{x}_C); \\ p(\mathbf{x}_{V \setminus S}) &= \sum_{\mathbf{x}_S} p(\mathbf{x}) = \frac{1}{Z} \prod_{\substack{C' \in \mathcal{C}: \\ S \cap C' = \emptyset}} \psi_{C'}(\mathbf{x}_{C'}) \cdot \left( \sum_{\mathbf{x}_S} \prod_{\substack{C \in \mathcal{C}: \\ S \cap C \neq \emptyset}} \psi_C(\mathbf{x}_C) \right). \end{aligned}$$

Therefore, the conditional distribution of  $\mathbf{x}_S$  given  $\mathbf{x}_{V \setminus S}$  reads:

$$p(\mathbf{x}_S | \mathbf{x}_{V \setminus S}) = \frac{p(\mathbf{x}_S, \mathbf{x}_{V \setminus S})}{p(\mathbf{x}_{V \setminus S})} = \frac{\prod_{C: S \cap C \neq \emptyset} \psi_C(\mathbf{x}_C)}{\sum_{\tilde{\mathbf{x}}_S} \prod_{C: S \cap C \neq \emptyset} \psi_C(\tilde{\mathbf{x}}_C)}. \quad (4)$$

To write the denominator in the last equality, we implicitly introduced  $\tilde{\mathbf{x}} = (\tilde{\mathbf{x}}_S, \mathbf{x}_{V \setminus S})$ , while  $\mathbf{x} = (\mathbf{x}_S, \mathbf{x}_{V \setminus S})$ .

Consider any maximal clique  $C$  such that  $S \cap C \neq \emptyset$  and let  $i \in S \cap C$ . If  $j \in C \setminus S$  then  $j \in \partial S$  because  $\{i, j\} \in E$  ( $i \in C$  and  $C$  is a clique). Therefore  $C \subseteq S \cup \partial S$ . It follows:

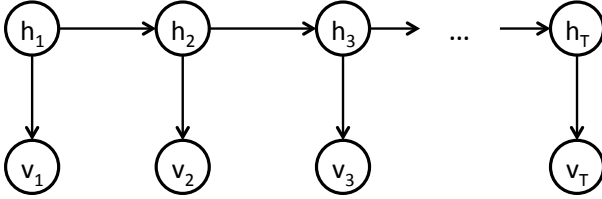
$$\begin{aligned} p(\mathbf{x}_S, \mathbf{x}_{\partial S}) &= \sum_{\mathbf{x}_{V \setminus (S \cup \partial S)}} p(\mathbf{x}) = \frac{1}{Z} \prod_{\substack{C \in \mathcal{C}: \\ S \cap C \neq \emptyset}} \psi_C(\mathbf{x}_C) \cdot \left( \sum_{\mathbf{x}_{V \setminus (S \cup \partial S)}} \prod_{\substack{C' \in \mathcal{C}: \\ S \cap C' = \emptyset}} \psi_{C'}(\mathbf{x}_{C'}) \right); \\ p(\mathbf{x}_{\partial S}) &= \sum_{\mathbf{x}_{V \setminus (S \cup \partial S)}} p(\mathbf{x}_{V \setminus S}) = \sum_{\mathbf{x}_{V \setminus (S \cup \partial S)}} \frac{1}{Z} \prod_{\substack{C' \in \mathcal{C}: \\ S \cap C' = \emptyset}} \psi_{C'}(\mathbf{x}_{C'}) \cdot \left( \sum_{\mathbf{x}_S} \prod_{\substack{C \in \mathcal{C}: \\ S \cap C \neq \emptyset}} \psi_C(\mathbf{x}_C) \right) \\ &= \frac{1}{Z} \left( \sum_{\mathbf{x}_{V \setminus (S \cup \partial S)}} \prod_{\substack{C' \in \mathcal{C}: \\ S \cap C' = \emptyset}} \psi_{C'}(\mathbf{x}_{C'}) \right) \left( \sum_{\mathbf{x}_S} \prod_{\substack{C \in \mathcal{C}: \\ S \cap C \neq \emptyset}} \psi_C(\mathbf{x}_C) \right). \end{aligned}$$

It comes

$$p(\mathbf{x}_S | \mathbf{x}_{\partial S}) = \frac{p(\mathbf{x}_S, \mathbf{x}_{\partial S})}{p(\mathbf{x}_{\partial S})} = \frac{\prod_{C: S \cap C \neq \emptyset} \psi_C(\mathbf{x}_C)}{\sum_{\tilde{\mathbf{x}}_S} \prod_{C: S \cap C \neq \emptyset} \psi_C(\tilde{\mathbf{x}}_C)}. \quad (5)$$

The final equalities in (4) and (5) are the same, thus proving that  $p(\mathbf{x}_S | \mathbf{x}_{V \setminus S})$  and  $p(\mathbf{x}_S | \mathbf{x}_{\partial S})$  are equal.

#### Problem 8 (Barber, p.99, Exercise 5.4)



1)

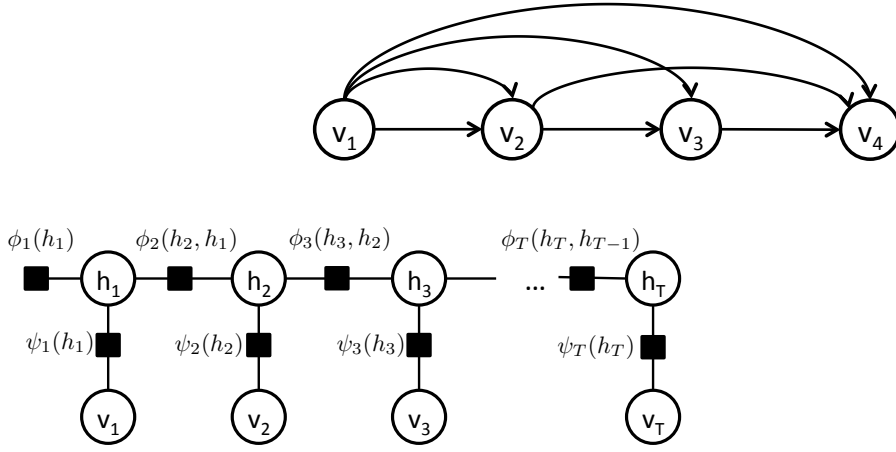
2) A simple linear chain for  $p(\mathbf{h})$  can be easily seen from

$$p(\mathbf{h}) = \sum_{\mathbf{v}} p(\mathbf{v}, \mathbf{h}) = p(h_1) \prod_{t=2}^T p(h_t | h_{t-1})$$

On the other hand,  $p(\mathbf{v})$  is a fully connected cascade belief network because the marginal probability does not admit any decomposition. For example  $T = 4$ ,

$$\begin{aligned} p(v_1, v_2, v_3, v_4) &= \sum_{h_1, h_2, h_3, h_4} p(v_1, v_2, v_3, v_4, h_1, h_2, h_3, h_4) \\ &= \sum_{h_4} p(v_4 | h_4) \sum_{h_3} \left( p(v_3, h_4 | h_3) \sum_{h_2} (p(v_2, h_3 | h_2) p(v_1, h_2)) \right) \end{aligned}$$

We see that  $v_1, v_2, v_3, h_4$  are all coupled.



3)

The factors are  $\psi_t(h_t) = p(v_t|h_t)$ ,  $\phi_1(h_1) = p(h_1)$  and  $\phi_t(h_t, h_{t-1}) = p(h_t|h_{t-1})$  for  $t \geq 2$ .

4) Suppose our observation is  $\hat{\mathbf{v}} = (\hat{v}_1, \dots, \hat{v}_T)$ . Since

$$p(\mathbf{h}|\mathbf{v} = \hat{\mathbf{v}}) \propto p(\mathbf{h}, \mathbf{v} = \hat{\mathbf{v}}),$$

we can use a sum-product algorithm to compute the marginal  $p(h_t, \hat{\mathbf{v}})$  and then it is easy to obtain  $p(h_t|\hat{\mathbf{v}}) = \frac{p(h_t, \hat{\mathbf{v}})}{\sum_{h_t} p(h_t, \hat{\mathbf{v}})}$ . Recall that

$$\begin{aligned} p(\hat{\mathbf{v}}, h_t) &= \sum_{\mathbf{h} \sim_t} p(\hat{\mathbf{v}}, \mathbf{h}) = \sum_{\mathbf{h} \sim_t} p(h_1) p(\hat{v}_1|h_1) \prod_{i=2}^T p(\hat{v}_i|h_i) p(h_i|h_{i-1}) \\ &= \sum_{\mathbf{h} \sim_t} \phi_1(h_1) \psi_1(h_1) \prod_{i=2}^T \psi_i(h_i) \phi_i(h_i, h_{i-1}) \end{aligned}$$

To compute the sum efficiently we define messages propagating from the two ends of the factor graph. For the forward propagation we define the factor-to-variable message

$$\mu_{\psi_i \rightarrow h_i}(h_i) = \psi(h_i), \quad \mu_{\phi_i \rightarrow h_i}(h_i) = \sum_{h_{i-1}} \phi_i(h_i, h_{i-1}) \mu_{h_{i-1} \rightarrow \phi_i}(h_{i-1}) \text{ with } \phi_1(h_1, h_0) \triangleq \phi_1(h_1)$$

and variable-to-factor message

$$\mu_{h_i \rightarrow \phi_{i+1}}(h_i) = \mu_{\psi_i \rightarrow h_i}(h_i) \mu_{\phi_i \rightarrow h_i}(h_i)$$

We compute the messages in the order  $(\mu_{\psi_1 \rightarrow h_1}, \mu_{\phi_1 \rightarrow h_1}) \rightarrow \mu_{h_1 \rightarrow \phi_2} \rightarrow (\mu_{\psi_2 \rightarrow h_2}, \mu_{\phi_2 \rightarrow h_2}) \rightarrow \mu_{h_2 \rightarrow \phi_3} \rightarrow \dots \rightarrow (\mu_{\psi_t \rightarrow h_t}, \mu_{\phi_t \rightarrow h_t})$ . So we have

$$\mu_{\phi_t \rightarrow h_t} = \sum_{h_1, \dots, h_{t-1}} \psi_1(h_1) \psi_1(h_1) \prod_{i=2}^t \psi_i(h_i) \phi_i(h_i, h_{i-1})$$

It does not harm to continue the forward propagation up to  $(\mu_{\psi_T \rightarrow h_T}, \mu_{\phi_T \rightarrow h_T})$  but here it is unnecessary. Next, we start the backward propagation with factor-to-variable message

$$\mu_{\phi_i \rightarrow h_{i-1}}(h_{i-1}) = \sum_{h_i} \phi_i(h_i, h_{i-1}) \mu_{h_i \rightarrow \phi_i}(h_i)$$

and variable-to-factor message

$$\mu_{h_i \rightarrow \phi_i}(h_i) = \mu_{\psi_i \rightarrow h_i}(h_i) \mu_{\phi_{i+1} \rightarrow h_i}(h_i) \text{ with } \mu_{\phi_{T+1} \rightarrow h_T}(h_T) \triangleq 1$$

We proceed with  $\mu_{\psi_T \rightarrow h_T} \rightarrow \mu_{h_T \rightarrow \phi_T} \rightarrow (\mu_{\psi_{T-1} \rightarrow h_{T-1}}, \mu_{\phi_T \rightarrow h_{T-1}}) \rightarrow \mu_{h_{T-1} \rightarrow \phi_{T-1}} \rightarrow \dots \rightarrow \mu_{\phi_{t+1} \rightarrow h_t}$ . So we have

$$\mu_{\phi_t \rightarrow h_t}(h_t) = \sum_{h_{t+1}, \dots, h_T} \prod_{i=t+1}^T \psi_i(h_i) \phi_i(h_i, h_{i-1})$$

and therefore

$$p(h_t, \hat{\mathbf{v}}) = \mu_{\phi_t \rightarrow h_t}(h_t) \mu_{\psi_t \rightarrow h_t}(h_t) \mu_{\phi_{t+1} \rightarrow h_t}(h_t),$$

$$p(h_t | \hat{\mathbf{v}}) = \frac{\mu_{\phi_t \rightarrow h_t}(h_t) \mu_{\psi_t \rightarrow h_t}(h_t) \mu_{\phi_{t+1} \rightarrow h_t}(h_t)}{\sum_{h_t} \mu_{\phi_t \rightarrow h_t}(h_t) \mu_{\psi_t \rightarrow h_t}(h_t) \mu_{\phi_{t+1} \rightarrow h_t}(h_t)}.$$

- 5) Like the starting argument in the last question, we need to compute  $\sum_{\mathbf{h}_{\sim t, t+1}} p(h_t, h_{t+1}, \hat{\mathbf{v}})$  where  $\mathbf{h}_{\sim t, t+1}$  means  $h_t$  and  $h_{t+1}$  are excluded. So with the same message passing rules we obtain

$$p(h_t, h_{t+1} | \hat{\mathbf{v}}) \propto \mu_{\phi_t \rightarrow h_t}(h_t) \mu_{\psi_t \rightarrow h_t}(h_t) \phi_{t+1}(h_t, h_{t+1}) \mu_{\phi_{t+2} \rightarrow h_{t+1}}(h_{t+1}) \mu_{\psi_{t+1} \rightarrow h_{t+1}}(h_{t+1})$$

### Problem 9 (Barber, p.98, Exercise 5.1)

The underlying undirected graph of a singly connected network with  $N$  nodes is a tree. We denote the tree with  $N$  nodes by  $\mathcal{T}_N$ . By definition it contains a leaf  $i$  which is connected to node  $j$ . The tree structure ensures the decomposition

$$Z = \sum_{\mathbf{x}_{\sim i}} \prod_{\substack{k \sim l \\ k \neq i \\ l \neq i}} \phi_{k,l}(x_k, x_l) \sum_{x_i} \phi_{i,j}(x_i, x_j).$$

where  $\mathbf{x}_{\sim i}$  means  $x_i$  is excluded. So we can start the following recursion with  $\mathcal{T}_N$ .

1. Find a leaf  $i$  which is connected to node  $j$ .
2. Compute  $\psi_{i,j}(x_j) = \sum_{x_i} \phi_{i,j}(x_i, x_j)$ .
3. If node  $j$  has another neighbor node  $k$ ,
  - 3a. obtain  $\mathcal{T}_{n-1}$  by removing node  $i$  and updating  $\phi_{j,k}(x_j, x_k) \rightarrow \psi_{i,j}(x_j) \phi_{j,k}(x_j, x_k)$ , and go to step 1 with  $\mathcal{T}_{n-1}$ ;
  - 3b. otherwise, there remain only node  $i$  and  $j$ , so we output  $Z = \sum_{x_j} \psi_{i,j}(x_j)$ .

The above algorithm ends with  $N$  iterations and therefore the time complexity is  $O(N)$ .

### Problem 10 (Bishop, p.397 & 421, Exercise 8.16 & 8.17)

- 1) Given the observation  $x_N = \hat{x}_N$ , the initial message for  $\beta$ -recursion becomes

$$\mu_{\beta}(x_{N-1}) = \phi_{N-1,N}(x_{N-1}, \hat{x}_N).$$

Note that this initial message does not sum over  $x_N$ . The other message passing equations are unchanged. This message passing allows us to compute  $p(x_n | x_N = \hat{x}_N)$ .

2) Given the observation  $x_3 = \hat{x}_3$ , the algorithm suggests

$$p(x_2) = \frac{1}{Z} \mu_\alpha(x_2) \mu_\beta(x_2)$$

where

$$\begin{aligned} \mu_\beta(x_2) &= \phi_{2,3}(x_2, \hat{x}_3) \mu_\beta(\hat{x}_3), \\ Z &= \sum_{x_2} \mu_\alpha(x_2) \mu_\beta(x_2) = \sum_{x_2} \mu_\alpha(x_2) \phi_{2,3}(x_2, \hat{x}_3) \mu_\beta(\hat{x}_3). \end{aligned}$$

We can simplify the expression to

$$p(x_2) = \frac{\mu_\alpha(x_2) \phi_{2,3}(x_2, \hat{x}_3)}{\sum_{x_2} \mu_\alpha(x_2) \phi_{2,3}(x_2, \hat{x}_3)}.$$

Different  $x_5$  will rescale  $\mu_\beta(\hat{x}_3)$  but it changes nothing on  $p(x_2)$ . This aligns with the fact that  $x_2 \perp\!\!\!\perp x_5 | x_3$ .

**Problem 1**

1) For every  $i \in [K]$ ,  $\underline{d}_i$  is the  $i^{\text{th}}$  canonical basis vector of  $\mathbb{R}^K$  and we define the latent random vector  $\underline{h} \in \{\underline{d}_i : i \in [K]\}$  whose distribution is  $\forall i \in [K] : \mathbb{P}(\underline{h} = \underline{d}_i) = w_i$ . Finally, let  $\underline{x} = \sum_{i=1}^K h_i \underline{a}_i + \underline{z}$  where  $\underline{z} \sim \mathcal{N}(0, \sigma^2 I_{D \times D})$  is independent of  $\underline{h}$ . The random vector  $\underline{x}$  has a probability density function  $p(\cdot)$ . We have:

$$\begin{aligned} \mathbb{E}[\underline{x}] &= \sum_{i=1}^K \mathbb{E}[h_i] \underline{a}_i + \mathbb{E}[\underline{z}] = \sum_{i=1}^K w_i \underline{a}_i \quad ; \\ \mathbb{E}[\underline{x} \underline{x}^T] &= \mathbb{E}[\underline{z} \underline{z}^T] + \sum_{i=1}^K \mathbb{E}[h_i] \underbrace{\mathbb{E}[\underline{z}]^T}_{=0} \underline{a}_i^T + \mathbb{E}[h_i] \underline{a}_i \mathbb{E}[\underline{z}]^T + \sum_{i,j=1}^K \underbrace{\mathbb{E}[h_i h_j]}_{=w_i \delta_{ij}} \underline{a}_i \underline{a}_j^T \\ &= \sigma^2 I_{D \times D} + \sum_{i=1}^K w_i \underline{a}_i \underline{a}_i^T . \end{aligned}$$

Finally, to compute the third moment tensor, note that  $\mathbb{E}[\underline{z} \otimes \underline{z} \otimes \underline{z}] = 0$  and that for every  $(i, j) \in [K]^2$ :  $\mathbb{E}[\underline{a}_i \otimes \underline{a}_j \otimes \underline{z}] = \mathbb{E}[\underline{a}_i \otimes \underline{z} \otimes \underline{a}_j] = \mathbb{E}[\underline{z} \otimes \underline{a}_i \otimes \underline{a}_j] = 0$ . Hence:

$$\begin{aligned} \mathbb{E}[\underline{x} \otimes \underline{x} \otimes \underline{x}] &= \sum_{i,j,k=1}^K \underbrace{\mathbb{E}[h_i h_j h_k]}_{=w_i \delta_{ij} \delta_{ik}} \underline{a}_i \otimes \underline{a}_j \otimes \underline{a}_k \\ &\quad + \sum_{i=1}^K \mathbb{E}[h_i] \mathbb{E}[\underline{a}_i \otimes \underline{z} \otimes \underline{z}] + \mathbb{E}[h_i] \mathbb{E}[\underline{z} \otimes \underline{a}_i \otimes \underline{z}] + \mathbb{E}[h_i] \mathbb{E}[\underline{z} \otimes \underline{z} \otimes \underline{a}_i] \\ &= \sum_{i=1}^K w_i \underline{a}_i \otimes \underline{a}_i \otimes \underline{a}_i + \sigma^2 \sum_{j=1}^D \sum_{i=1}^K w_i (\underline{a}_i \otimes \underline{e}_j \otimes \underline{e}_j + \underline{e}_j \otimes \underline{e}_j \otimes \underline{a}_i + \underline{e}_j \otimes \underline{a}_i \otimes \underline{e}_j) . \end{aligned}$$

2) Let  $A = [\underline{a}_1, \underline{a}_2, \dots, \underline{a}_K] \in \mathbb{R}^{D \times K}$  and  $A' = [\underline{a}'_1, \underline{a}'_2, \dots, \underline{a}'_K] \in \mathbb{R}^{D \times K}$ . By definition,  $\tilde{R} = \Sigma^{-1} R \Sigma$  where  $\Sigma$  is the diagonal matrix such that  $\Sigma_{ii} = \sqrt{w_i}$  and  $A' = A \tilde{R}^T$ . We can directly apply the formula of question 1) to compute the second moment matrix of the new mixture of Gaussians:

$$\begin{aligned} \mathbb{E}[\underline{x} \underline{x}^T] &= \sigma^2 I_{D \times D} + A' \Sigma^2 A'^T = \sigma^2 I_{D \times D} + A \tilde{R}^T \Sigma^2 \tilde{R} A^T \\ &= \sigma^2 I_{D \times D} + A \Sigma R^T R \Sigma A^T = \sigma^2 I_{D \times D} + A \Sigma^2 A^T . \end{aligned}$$

**Problem 2: Examples of tensors and their rank**

1) The matrices corresponding to  $B$ ,  $P$ ,  $E$  are:

$$B = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} ; P = \begin{bmatrix} 1 & 1 \\ 1 & 1 \end{bmatrix} ; E = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix}$$

The frontal slices of  $G$  and  $W$  are:

$$G_1 = \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}, G_2 = \begin{bmatrix} 0 & 0 \\ 0 & 1 \end{bmatrix}; W_1 = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}, W_2 = \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}.$$

The matricizations of  $G$  and  $W$  are:

$$G_{(1)} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}; G_{(2)} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}; G_{(3)} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix};$$

$$W_{(1)} = \begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix}; W_{(2)} = \begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix}; W_{(3)} = \begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix}.$$

**2)**  $B$  and  $E$  are clearly rank-2 matrices, while  $P = (e_0 + e_1) \otimes (e_0 + e_1)$  is a rank-1 matrix.

By its definition,  $G$  is at most rank 2. Assume it is rank 1:  $G = a \otimes b \otimes c$  with  $a, b, c \in \mathbb{R}^2$ . We have  $a_1 b_1 c_1 = G_{111} = 1$  and  $a_2 b_1 c_1 = G_{211} = 0$  so we must have  $a_2 = 0$ . Besides,  $a_2 b_2 c_2 = G_{222} = 1$  and  $a_1 b_2 c_2 = G_{122} = 0$  so  $a_1 = 0$ . Hence  $a^T = (0, 0)$  and  $G$  is the all-zero tensor. This is a contradiction and we conclude that  $G$  is rank 2.

By its definition,  $W$  is at most rank 3. To prove the rank cannot be smaller than 3, we will proceed by contradiction:

- Assume  $W$  is rank 1:  $W = a \otimes b \otimes c$  with  $a, b, c \in \mathbb{R}^2$ . We have  $a_1 b_1 c_1 = W_{111} = 0$  and  $a_2 b_1 c_1 = W_{211} = 1$  so  $a_1 = 0$ . Besides,  $a_1 b_1 c_2 = W_{112} = 1$  and  $a_2 b_1 c_2 = W_{212} = 0$  so  $a_2 = 0$ . Then  $a = (0, 0)^T$  and  $W$  is the all-zero tensor, which is a contradiction.
- Assume  $W$  is rank 2:  $W = a \otimes b \otimes c + d \otimes e \otimes f$ . We claim that  $a$  and  $d$  must be linearly independent. Indeed, suppose they are parallel and take a vector  $x$  perpendicular to both  $a$  and  $d$ . Then

$$W(x, I, I) = (x^T a) b \otimes c + (x^T d) e \otimes f = 0$$

but also

$$W(x, I, I) = (x^T e_0) e_0 \otimes e_1 + (x^T e_0) e_1 \otimes e_0 + (x^T e_1) e_0 \otimes e_0 = \begin{bmatrix} x^T e_1 & x^T e_0 \\ x^T e_0 & 0 \end{bmatrix}$$

which cannot be zero since  $x$  cannot be perpendicular to both  $e_0$  and  $e_1$ . Now, we take  $x$  perpendicular to  $d$ . We have

$$W(x, I, I) = (x^T a) b \otimes c$$

which is rank one. Therefore, we must have  $x^T e_0 = 0$  which implies that  $x$  is parallel to  $e_1$  and thus  $d$  parallel to  $e_0$ . Now, if we take  $x$  perpendicular to  $a$ , the matrix

$$W(x, I, I) = (x^T d) e \otimes f$$

is rank one and, once again, we must have  $x^T e_0 = 0$ , which implies  $x$  parallel to  $e_1$  and thus  $a$  parallel to  $e_0$ . Hence, we have shown that  $a$  and  $d$  are linearly independent but also that both are parallel to  $e_0$ . This is a contradiction.

**3)** Writing everything in terms of matrix product, it comes:

$$(Oe_0) \otimes (Oe_0) + (Oe_1) \otimes (Oe_1) = Oe_0 e_0^T O^T + Oe_1 e_1^T O^T = OO^T = B.$$

so  $B$  does not have a unique decomposition.

For  $G$  we have  $G = \underline{a}_1 \otimes \underline{b}_1 \otimes \underline{c}_1 + \underline{a}_2 \otimes \underline{b}_2 \otimes \underline{c}_2$  with

$$A = [\underline{a}_1, \underline{a}_2] = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}; B = [\underline{b}_1, \underline{b}_2] = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}; C = [\underline{c}_1, \underline{c}_2] = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}.$$

$A, B, C$  are full column rank and  $G$  has rank 2: by Jennrich's algorithm, the decomposition is unique (up to trivial rank permutation and feature scaling).

For  $W$  we have  $W = \underline{a}_1 \otimes \underline{b}_1 \otimes \underline{c}_1 + \underline{a}_2 \otimes \underline{b}_2 \otimes \underline{c}_2 + \underline{a}_3 \otimes \underline{b}_3 \otimes \underline{c}_3$  with

$$A = [\underline{a}_1, \underline{a}_2, \underline{a}_3] = \begin{bmatrix} 1 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}; \quad B = [\underline{b}_1, \underline{b}_2, \underline{b}_3] = \begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix}; \quad C = [\underline{c}_1, \underline{c}_2, \underline{c}_3] = \begin{bmatrix} 0 & 1 & 1 \\ 1 & 0 & 0 \end{bmatrix}.$$

$A, B, C$  are not full column rank: Jennrich's theorem does not allow to conclude that the decomposition of  $W$  is unique.

4) We expand the tensor products in the definition of  $D_\epsilon$ :

$$\begin{aligned} D_\epsilon &= \frac{1}{\epsilon} \left[ (e_0 + \epsilon e_1) \otimes (e_0 + \epsilon e_1) \otimes (e_0 + \epsilon e_1) - e_0 \otimes e_0 \otimes e_0 \right] \\ &= \frac{1}{\epsilon} \left[ e_0 \otimes e_0 \otimes e_0 + \epsilon e_0 \otimes e_0 \otimes e_1 + \epsilon e_0 \otimes e_1 \otimes e_0 + \epsilon e_1 \otimes e_0 \otimes e_0 \right. \\ &\quad \left. + \epsilon^2 e_1 \otimes e_1 \otimes e_0 + \epsilon^2 e_1 \otimes e_0 \otimes e_1 + \epsilon^2 e_0 \otimes e_1 \otimes e_1 + \epsilon^3 e_1 \otimes e_1 \otimes e_1 - e_0 \otimes e_0 \otimes e_0 \right] \\ &= e_0 \otimes e_0 \otimes e_1 + e_0 \otimes e_1 \otimes e_0 + e_1 \otimes e_0 \otimes e_0 \\ &\quad + \epsilon(e_1 \otimes e_1 \otimes e_0 + e_1 \otimes e_0 \otimes e_1 + e_0 \otimes e_1 \otimes e_1) + \epsilon^2 e_1 \otimes e_1 \otimes e_1 \\ &= W + \epsilon(e_1 \otimes e_1 \otimes e_0 + e_1 \otimes e_0 \otimes e_1 + e_0 \otimes e_1 \otimes e_1) + \epsilon^2 e_1 \otimes e_1 \otimes e_1. \end{aligned}$$

Hence  $\lim_{\epsilon \rightarrow 0} D_\epsilon = 0$ .

### Problem 3

1) There cannot be an analogous general result for tensors. Indeed, the order-3 tensor  $W$  of Problem 2 is rank 3 and we show in 4) that  $\lim_{\epsilon \rightarrow 0} \|W - D_\epsilon\|_F = 0$ . So there is no minimum attained in the space of rank 2 tensors. In this sense, there is simply no *best* rank-two approximation of  $W$ .

2) Let  $M$  a matrix of rank  $R + 1$  with singular values  $\sigma_1 \geq \sigma_2 \cdots \geq \sigma_R \geq \sigma_{R+1} > 0$ . By the Eckart-Young-Mirsky theorem, the minimum of  $\|M - \widehat{M}\|_F$  over rank  $R$  matrices  $\widehat{M}$  is equal to  $\sigma_{R+1} > 0$ . Therefore, there cannot be a sequence of matrices  $M_n$  given by a sum of  $R$  rank-one matrices such that  $\lim_{n \rightarrow +\infty} \|M - M_n\|_F = 0$ .

4) In the real-valued case, we have:

$$|T(R_1, R_2, R_3)^{\alpha\beta\gamma}|^2 = \sum_{\delta, \epsilon, \zeta, \delta', \epsilon', \delta'} R_1^{\delta\alpha} R_1^{\delta'\alpha} R_2^{\epsilon\beta} R_2^{\epsilon'\beta} R_3^{\zeta\gamma} R_3^{\zeta'\gamma} T^{\delta\epsilon\zeta} T^{\delta'\epsilon'\zeta'}.$$

Summing over  $\alpha, \beta, \gamma$  and using the orthogonality of rotation matrices, we find:

$$\sum_{\alpha} R_1^{\delta\alpha} R_1^{\delta'\alpha} = \delta_{\delta\delta'}, \quad \sum_{\beta} R_2^{\epsilon\beta} R_2^{\epsilon'\beta} = \delta_{\beta\beta'}, \quad \sum_{\gamma} R_3^{\zeta\gamma} R_3^{\zeta'\gamma} = \delta_{\zeta\zeta'}.$$

The result directly follows:

$$\|T(R_1, R_2, R_3)\|_F^2 = \sum_{\delta\epsilon\zeta} |T(R_1, R_2, R_3)^{\alpha\beta\gamma}|^2 = \sum_{\delta\epsilon\zeta} |T^{\delta\epsilon\zeta}|^2 = \|T\|_F^2.$$



#### Problem 4

1) To show that  $A \odot_{KhR} B$  is full column rank, we have to prove that the kernel of the linear application  $\underline{x} \mapsto (A \odot_{KhR} B)\underline{x}$  is  $\{0\}$ . Let  $\underline{x} \in \mathbb{R}^R$  with components  $(x^1, x^2, \dots, x^R)$  be such that  $(A \odot_{KhR} B)\underline{x} = 0$ . Then,  $\forall \alpha \in [I_1]$ :

$$\sum_{r=1}^R a_r^\alpha x^r \underline{b}_r = 0.$$

Because  $B$  is full column rank,  $\sum_{r=1}^R a_r^\alpha x^r \underline{b}_r = 0$  implies that  $\forall r \in [R] : a_r^\alpha x^r = 0$ . Note that:

$$\forall \alpha \in [I_1], \forall r \in [R] : a_r^\alpha x^r = 0 \Leftrightarrow A\underline{x} = 0.$$

$A$  is full column rank and  $A\underline{x} = 0$ , hence  $\underline{x} = 0$ .  $A \odot_{KhR} B$  is full column rank.

2) Suppose we are given a tensor (the weights  $\lambda_r$  that usually appear in the sum are absorbed in the vectors  $\underline{a}_r$ )

$$\mathcal{X} = \sum_{r=1}^R \underline{a}_r \otimes \underline{b}_r \otimes \underline{c}_r, \quad (1)$$

where  $A = [\underline{a}_1, \underline{a}_2, \dots, \underline{a}_R] \in \mathbb{R}^{I_1 \times R}$ ,  $B = [\underline{b}_1, \underline{b}_2, \dots, \underline{b}_R] \in \mathbb{R}^{I_2 \times R}$  and  $C = [\underline{c}_1, \underline{c}_2, \dots, \underline{c}_R] \in \mathbb{R}^{I_3 \times R}$  are full column rank. By Jennrich's algorithm, the decomposition (1) is unique up to trivial rank permutation and feature scaling and Jennrich's algorithm is a way to recover this decomposition. At the end of the step (5) of the algorithm, we have computed  $A, B$  and it remains to recover  $C$ . We now show how the result in question 1) allows to recover  $C$  uniquely. For each  $\gamma \in [I_3]$ , define the slice  $\mathcal{X}_\gamma$  as the  $I_1 \times I_2$  matrix with entries  $(\mathcal{X}_\gamma)^{\alpha\beta} = \mathcal{X}^{\alpha\beta\gamma}$  and denote  $F(\mathcal{X}_\gamma)$  the  $I_1 I_2$  column vector with entries  $F(\mathcal{X}_\gamma)^{\beta+I_2(\alpha-1)} = \mathcal{X}^{\alpha\beta\gamma}$ . We have:

$$\forall (\alpha, \beta) \in [I_1] \times [I_2] : F(\mathcal{X}_\gamma)^{\beta+I_2(\alpha-1)} = \sum_{r=1}^R a_r^\alpha b_r^\beta c_r^\gamma = \sum_{r=1}^R (A \odot_{KhR} B)^{\beta+I_2(\alpha-1), r} c_r^\gamma.$$

Therefore, the  $I_1 I_2 \times I_3$  matrix  $F(\mathcal{X}) = [F(\mathcal{X}_1), F(\mathcal{X}_2), \dots, F(\mathcal{X}_{I_3})]$  satisfies:

$$F(\mathcal{X}) = (A \odot_{KhR} B) C^T.$$

Because  $A \odot_{KhR} B$  is full column rank, we can invert the system with the Moore-Penrose pseudoinverse:  $C^T = (A \odot_{KhR} B)^\dagger F(\mathcal{X})$ .

#### Problem 5

1) To apply Jennrich's algorithm we need to prove that the matrix  $E = [\underline{c}_1 \otimes_{Kro} \underline{d}_1, \dots, \underline{c}_R \otimes_{Kro} \underline{d}_R]$  is full column rank ( $A, B$  are full column rank by assumption). Note that the same proof as the one in Problem 4 question 1 applies. Nevertheless we repeat the argument here.

Let  $\underline{v} \in \mathbb{R}^R$  a column vector in the kernel of  $E$ , i.e.,  $E\underline{v} = 0$ . Then:

$$\forall \gamma \in [I_3] : \sum_{r=1}^R (c_r^\gamma v^r) \underline{d}_r = 0 \Rightarrow \forall \gamma \in [I_3], \forall r \in [R] : c_r^\gamma v^r = 0 \Rightarrow C\underline{v} = 0 \Rightarrow \underline{v} = 0.$$

The first implication follows from  $D$  being full column rank and the third one from  $C$  being full column rank. We conclude that the kernel of  $E$  is  $\{0\}$ :  $E$  is full column rank.

We can therefore apply Jennrich's algorithm.

2) We recover the rank  $R$  as well as  $A$ ,  $B$  and  $E$  by applying Jennsen's algorithm to  $\tilde{T}$ . From  $E$  we can then determine  $C$  and  $D$ . Fix  $r \in [R]$ . Since  $C$  is full column rank, there exists  $\alpha \in [I_3]$  such that  $c_r^\alpha \neq 0$ . As  $c_r^\alpha \neq 0$ , we can use the  $I_4$ -dimensional column vector  $c_r^\alpha \underline{d}_r$  contained in the  $r^{\text{th}}$  column of  $E$  to recover  $\underline{d}_r$ . Doing this for every  $r \in [R]$  we recover the matrix  $D$ . Finally, for every  $r \in R$ , pick  $\beta \in I_4$  such that  $d_r^\beta \neq 0$  (such  $\beta$  exists because  $D$  is full column rank) and use the entries  $c_r^\alpha d_r^\beta$ ,  $\alpha \in [I_3]$ , to recover  $\underline{c}_r$ .  $C$  has then been recovered.

### Problem 6

1) Define  $\Sigma^\dagger$  as the  $N \times M$  diagonal matrix with diagonal entries:

$$\forall i \in \{1, 2, \dots, \min\{M, N\}\} : (\Sigma^\dagger)_{ii} \begin{cases} 1/\Sigma_{ii} & \text{if } \Sigma_{ii} \neq 0 ; \\ 0 & \text{otherwise.} \end{cases}$$

Then both  $\Sigma^\dagger \Sigma \in \mathbb{C}^{N \times N}$  and  $\Sigma \Sigma^\dagger \in \mathbb{C}^{M \times M}$  are diagonal square matrices with diagonal entries:

$$\begin{aligned} \forall i \in [N] : (\Sigma^\dagger \Sigma)_{ii} &= \begin{cases} 1 & \text{if } i \leq \min\{M, N\} \text{ and } \Sigma_{ii} \neq 0 ; \\ 0 & \text{otherwise.} \end{cases} \\ \forall i \in [M] : (\Sigma \Sigma^\dagger)_{ii} &= \begin{cases} 1 & \text{if } i \leq \min\{M, N\} \text{ and } \Sigma_{ii} \neq 0 ; \\ 0 & \text{otherwise.} \end{cases} \end{aligned}$$

It is then easy to check that  $\Sigma^\dagger$  satisfies the first two conditions of the Moore-Penrose pseudoinverse:  $\Sigma \Sigma^\dagger \Sigma = \Sigma$  and  $\Sigma^\dagger \Sigma \Sigma^\dagger = \Sigma^\dagger$ . Besides,  $\Sigma^\dagger \Sigma$  and  $\Sigma \Sigma^\dagger$  being real diagonal matrices, the last two conditions are clearly satisfied too.

2) We can check that the matrix  $V \Sigma^\dagger U^*$  satisfies the four conditions of the Moore-Penrose pseudoinverse, i.e.,  $A^\dagger = V \Sigma^\dagger U^*$ :

$$\begin{aligned} A[V \Sigma^\dagger U^*]A &= U \Sigma (V^* V) \Sigma^\dagger (U^* U) \Sigma V^* = U \Sigma \Sigma^\dagger \Sigma V^* = U \Sigma V^* = A ; \\ [V \Sigma^\dagger U^*]A[V \Sigma^\dagger U^*] &= V \Sigma^\dagger (U^* U) \Sigma (V^* V) \Sigma^\dagger U^* = V \Sigma^\dagger \Sigma \Sigma^\dagger U^* = V \Sigma^\dagger U^* ; \\ (A V \Sigma^\dagger U^*)^* &= (U \Sigma \Sigma^\dagger U^*)^* = U (\Sigma \Sigma^\dagger)^* U^* = U \Sigma \Sigma^\dagger U^* = A V \Sigma^\dagger U^* ; \\ (V \Sigma^\dagger U^* A)^* &= (V \Sigma^\dagger \Sigma V^*)^* = V (\Sigma^\dagger \Sigma)^* V^* = V \Sigma^\dagger \Sigma V^* = V \Sigma^\dagger U^* A . \end{aligned}$$

3)  $A$  is full column rank, therefore  $A^* A$  is a full rank  $N \times N$  matrix and has a unique inverse  $(A^* A)^{-1}$ . The matrix  $(A^* A)^{-1} A^*$  satisfies the four conditions:

$$\begin{aligned} A[(A^* A)^{-1} A^*]A &= A ; [(A^* A)^{-1} A^*]A[(A^* A)^{-1} A^*] = (A^* A)^{-1} A^* ; \\ (A[(A^* A)^{-1} A^*])^* &= A[(A^* A)^{-1} A^*] ; ([A[(A^* A)^{-1} A^*]A]^* = A^* A (A^* A)^{-1} = I_{N \times N} = ([A[(A^* A)^{-1} A^*]A)^* . \end{aligned}$$

Hence  $A^\dagger = (A^* A)^{-1} A^*$ .

4)  $A$  is full row rank, therefore  $AA^*$  is a full rank  $M \times M$  matrix and has a unique inverse  $(AA^*)^{-1}$ . The matrix  $A^*(AA^*)^{-1}$  satisfies the four conditions:

$$\begin{aligned} A[A^*(AA^*)^{-1}]A &= A ; [A^*(AA^*)^{-1}]A[A^*(AA^*)^{-1}] = A^*(AA^*)^{-1} ; \\ (A[A^*(AA^*)^{-1}])^* &= (AA^*)^{-1} AA^* = I_{M \times M} = AA^\dagger ; ([A^*(AA^*)^{-1}]A)^* = A^*(AA^*)^{-1} A . \end{aligned}$$

Hence  $A^\dagger = A^*(AA^*)^{-1}$ .

5) We have  $AA^{-1}A = A$ ,  $A^{-1}AA^{-1} = A^{-1}$ ,  $(AA^{-1})^* = I_{M \times M} = AA^{-1}$ ,  $(A^{-1}A)^* = I_{N \times N} = A^{-1}A$ . Hence  $A^\dagger = A^{-1}$ .

6)  $A$  is full column rank so  $A^\dagger A = I_{M \times M}$  and  $B$  is full column rank so  $BB^\dagger = I_{N \times N}$ . Therefore:

$$\begin{aligned} (AB)(B^\dagger A^\dagger)(AB) &= A(BB^\dagger)(A^\dagger A)B = AI_{M \times M}I_{N \times N}B = AB; \\ (B^\dagger A^\dagger)(AB)(B^\dagger A^\dagger) &= B^\dagger(A^\dagger A)(BB^\dagger)A^\dagger = B^\dagger I_{N \times N}I_{M \times M}A^\dagger = B^\dagger A^\dagger; \\ (ABB^\dagger A^\dagger)^* &= (AI_{N \times N}A^\dagger)^* = (AA^\dagger)^* = AA^\dagger = (AB)(B^\dagger A^\dagger); \\ (B^\dagger A^\dagger AB)^* &= (B^\dagger I_{M \times M}B)^* = (B^\dagger B)^* = B^\dagger B = (B^\dagger A^\dagger)(AB). \end{aligned}$$

Hence  $(AB)^\dagger = B^\dagger A^\dagger$ .