Master's Thesis

Incremental Cata Computation for Generic Data Types

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1 Abstract

2 Introduction

2.1 Problem statement

3 Background

3.1 Regular

4 Implementation

Implementing the idea using a generic programming library, would be the ultimate goal. But first a **proof-of-concept** was made to show that the implementation is a viable product. A prototype-language is created, which is based on the notion of *pattern functors*.

4.1 Prototype language

The definition of the pattern functor only leads to shallow recursion. Meaning that pattern functor can only be used to observe a single layer of recursion. To apply a function over the complete data structure, deep recursion is used. To implement deep recursion, the fix point is introduced.

```
data Fix f = In { unFix :: f (Fix f) }
```

The fix point is then used to describe the recursion of the datatype on a type-level basis. Using pattern functors and fix point most of the Haskell datatypes can be represented. For example:

Because the generic representation of the Haskell datatypes can be represented using pattern functors, we can use Functors. Using the Functor class a cata function can be defined, which is a generic fold function.

To store the intermediate results of cata, we want the structure of the data to be hashed. This way we can easily compare if the data structure has changed over time, without completely

recomputing the resulting digests. To do this, first a fix point is introduced which additionally stores the digest.

```
type Merkle f = Fix (f :*: K Digest)
```

Then to convert the fix point to a fix point containing the structural digest, the Merkelize class is introduced.

```
class Merkelize f where
    merkleG :: Merkelize g => f (Fix g) -> (f :*: K Digest) (Fix (g :*: K Digest))
merkle :: Merkelize f => Fix f -> Merkle f
merkle = In . merkleG . unFix
```

Using the new fix point with its structural digest, a new cata function can be defined which can store its intermediate values in a Map Digest a.

5 Results

5.1 Memory

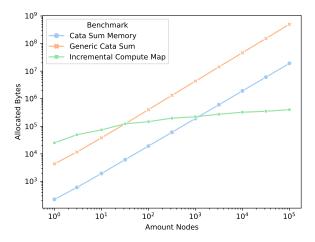


Figure 1: Overview Bytes Allocation

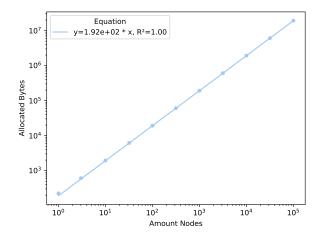


Figure 2: Bytes allocated for Cata Sum

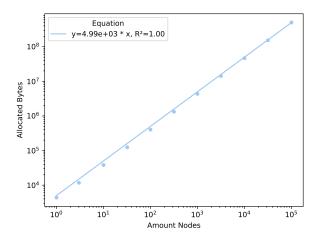


Figure 3: Bytes allocated for Generic Cata Sum

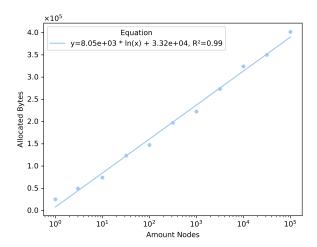


Figure 4: Bytes allocated for Incremental Cata Sum

6 Conclusion

7 Future Work

8 Appendix

A Implementation Regular

A.1 Definition Generic Datatypes

A.2 Implementation Hashable

```
class Hashable f where
 hash :: f (Fix (g :*: K Digest)) -> Digest
instance Hashable U where
 hash _ = digest "U"
instance (Show a) => Hashable (K a) where
 hash (K x) = digestConcat [digest "K", digest x]
instance Hashable I where
  hash (I x) = digestConcat [digest "I", getDigest x]
   where
      getDigest :: Fix (f :*: K Digest) -> Digest
     getDigest (In (_ :*: K h)) = h
instance (Hashable f, Hashable g) => Hashable (f :+: g) where
  hash (L x) = digestConcat [digest "L", hash x]
  hash (R x) = digestConcat [digest "R", hash x]
instance (Hashable f, Hashable g) => Hashable (f :*: g) where
  hash (x : *: y) = digestConcat [digest "P", hash x, hash y]
instance (Hashable f) => Hashable (C c f) where
  hash (C x) = digestConcat [digest "C", hash x]
```

A.3 Implementation Merkle

A.4 Implementation Cata Merkle

A.5 Implementation Zipper Merkle