

FINAL EXAM: CS3342b Tuesday, 25 April 2017, 2pm, Room FEB GYM

NAME AS APPEARS ON STUDENT ID:

STUDENT ID NUMBER:

GAUL/CONFLUENCE USER NAME:

REMINDERS:

1. (from course outline) The final exam will be closed book, closed notes, with no electronic devices allowed, with particular reference to any electronic devices that are capable of communication and/or storing information.
2. Write neatly. If the marker can't read it, it is wrong.
3. This exam shouldn't take long to write. On the other hand, time will pass. It is a 3 hour exam with 50 questions. If you complete a question every 3 minutes (or 10 questions every half hour), you will still have a half hour at the end to double check that everything is in order.
4. While you are not allowed to open the exam booklet until the proctor says you can, you can fill out the information on the cover page. You should also get out your student id and make sure your pencils and pens are in order. If you need to get something out of your jacket or knapsack once the exam has started, raise your hand and wait til a proctor comes to you to oversee the matter.

1. In Ruby, the evaluation of arguments to a message are handled by the object sending the message. In Haskell, the runtime environment decides when and how much to evaluate an argument to a function. In Io, the evaluation of the arguments to a message is made by ANSWER.
  - the receiver of the message
2. In Clojure, the value of (repeat 1) is ANSWER.
  - an infinite sequence of 1s
  - a lazy infinite sequence of 1s
3. Using the straightforward statement translation scheme in the ICD textbook, if I were to TransStat('if true then z := 1 else z := 2', vtable, ftable), newlabel() will be invoked ANSWER times.
  - 3
4. In the ICD textbook's example interpreter for evaluating expressions, in the row labelled let id = Exp1 in Exp2, we have the code: v1 = EvalExp(Exp1, vtable, ftable); vtableP = bind(vtable, getname(id), v1), EvalExp(Exp2, vtableP, ftable). The bind function changes vtable into vtableP by ANSWER.
  - inserting the association of getname(id) with the value v1 into the table
  - inserting the binding of getname(id) with the value v1 into the table
5. In Haskell, instead of defining second by writing something like second x = head( tail( x ) ), you can write this without introducing the parameter x by using function composition. Doing that, you would define second by ANSWER.
  - second = head . tail
6. In Ruby, the mixin is used to solve the object-oriented programming problem of ANSWER.
  - multiple inheritance
7. In Prolog, the expression hi(X, 4) = hi(3, X) causes X to have the value ANSWER.
  - X will not be bound and the expression will fail
  - X will not be bound
8. When a function is invoked, if the language passes a copy of the value of each parameter to the code that performs the function, this is called ANSWER.
  - call-by-value
  - pass-by-value
9. In Prolog, the most natural way to express the rule that 'I am an ancestor of you if I am a parent of an ancestor of you' is ANSWER.
  - ancestor(I, You) :- parent(I, Ancestor), ancestor(Ancestor, You).
10. Matz, the creator of Ruby, thinks that it is less important to optimize the execution (efficiency) of a programming language and more important to optimize the efficiency of ANSWER.
  - the programmers
11. In the ICD textbook's example interpreter for evaluating expressions, in the row labelled id, we have the code: v = lookup(vtable, getname(id)) ; if v = unbound then error() else v. It says getname(id) instead of id, because ANSWER.
  - id indicates a token with a type and value field

12. In Ruby, by convention, the `?` in the method `me?` is used to indicate that `me` is ANSWER.
- `boolean`
13. In Ruby, the `@` is used to indicate that the variable `@me` is ANSWER.
- an instance variable
14. In Haskell's `do` notation for working with monads, assignment uses the ANSWER operator.
- `←`
15. When the structure of the syntax tree is used to determine which object corresponds to a name, this is called ANSWER.
- static scoping
  - lexical scoping
16. In the context-free grammar  $A \rightarrow BA$ ,  $B \rightarrow AB$ ,  $A \rightarrow B$ ,  $A \rightarrow a$ ,  $B \rightarrow b$ , and  $B \rightarrow$  the value of `Nullable(A)` is ANSWER.
- `true`
17. The technical term for the compiler design methodology where the translation closely follows the syntax of the language is ANSWER.
- syntax-directed translation
18. When type checking done during program execution, the type system is called ANSWER.
- dynamic typing
19. Another method of parameter passing, whose technical name is ANSWER, is implemented by passing the address of the variable (or whatever the given parameter is). Assigning to such a parameter would then change the value stored at the address.
- call-by-reference
  - pass-by-reference
20. In the context-free grammar  $A \rightarrow BA$ ,  $B \rightarrow AB$ ,  $A \rightarrow B$ ,  $A \rightarrow a$ ,  $B \rightarrow b$ , and  $B \rightarrow$  the value of `FIRST(A)` is ANSWER.
- $\{a, b\}$
21. ANSWER is the data structure used in language translation to track the binding of variables and functions to their type.
- A symbol table
22. The `loop` and `recur` constructs are in Clojure to guide ANSWER.
- tail recursion optimization
  - tail recursion elimination
23. In Haskell, if we want to define a local named function inside a function definition, we use the keyword ANSWER.
- `where`
24. In Scala, the type that every type is a subtype of is called ANSWER.

- Any
25. The central idea of context-free grammars is to define a language by productions. These productions say that a nonterminal symbol can be replaced by ANSWER.
    - a sequence of terminals and nonterminals
    - a sequence of symbols
  26. In Ruby, normally, when you try to add a String to a Fixnum, you get an error message saying that a String can't be coerced to a Fixnum. This is because Ruby is ANSWER typed.
    - strongly
  27. In the Erlang community, ANSWER code refers to replacing pieces of your application without stopping your application.
    - hot-swapping
  28. In Erlang, you can link two processes together. Then when one dies, it sends ANSWER to its twin.
    - an exit signal
  29. In Haskell, instead of writing something like `if x == 0 then 1 else fact ( x - 1 ) * x`, you can write a series of lines starting with `| x > 1 = x * factorial ( x - a)`. This second style is called ANSWER.
    - using guards
  30. In most languages, a function definition like `f a b = a : (f (a + b) b)` would result in an infinite recursion. However, in Haskell we can partially evaluate functions like this because Haskell is based on ANSWER.
    - lazy evaluation
  31. One of the three most significant parts of a monad is called ANSWER, which wraps up a function and puts it in the monad's container.
    - return
  32. The way Haskell handles functions with more than one parameter is called ANSWER.
    - currying
  33. Three concepts related to concurrency were discussed with regards to the language Io. ANSWER was presented as a general mechanism for sending a message to an object that would cause that object to respond to the message as a separate process running asynchronously.
    - Actors
  34. Io is known for taking ANSWER -based approach to object-oriented programming.
    - a prototype
  35. In the Ruby community, the acronym DSL is an abbreviation for ANSWER.
    - domain specific language
  36. ANSWER typing is when the language implementation ensures that the arguments of an operation are of the type the operation is defined for.
    - Strong

37. One approach to speeding up an interpreter is to translate pieces of the code being interpreted directly into machine code during program execution, this is called ANSWER.
  - just-in-time compilation
38. In the chapter on Scala, we get the following interesting quote: ANSWER is the most important thing you can do to improve code design for concurrency.
  - Immutability
39. Since a compiler may have to look up what object is associated with a name many times, it is typical to use ANSWER to avoid linear search times.
  - hash tables
40. The context-free grammar  $A \rightarrow BA$ ,  $B \rightarrow AB$ ,  $A \rightarrow a$ ,  $B \rightarrow b$ ,  $B \rightarrow$  is not LL(1) specifically because ANSWER.
  - FIRST(BA) and FIRST(a) both include a, so we do not know which A rule to use
41. The specifications of how to group characters into meaningful basic units of a programming language are generally implemented in code that has the abstract form of ANSWER.
  - a finite automata
  - a finite state machine
42. Since Haskell doesn't have traditional error handling mechanisms, by convention, people use the ANSWER monad to distinguish a valid return from an error return.
  - Maybe
43. In Io, the basic method for creating a new object is ANSWER.
  - clone
44. Each named object will have ANSWER, where the name is defined as a synonym for the object.
  - a declaration
45. Another design goal for Scala was to have its programs easily interoperate with those written in ANSWER.
  - Java
46. In automatically generating the code that reads characters and outputs the part of a programming language that is analogous to its words, we start with a specification and then traditionally convert it into code in two stages. The main problem that can arise in moving from the first stage to the second stage is ANSWER.
  - an exponential explosion in the number of states needed
47. ANSWER data structures have the property that no operation on the structure will destroy or modify it.
  - persistent
  - functional
  - immutable
48. Using the straightforward expression translation scheme in the ICD textbook, if I were to TransExp('3 \* x + 1', vtable, ftable), newvar() will be invoked ANSWER times.

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49. Unlike most Lisp systems, Clojure doesn't use its own custom virtual machine. It was originally designed to compile to code that would run on the ANSWER.

- JVM
- Java Virtual Machine

50. Scala uses few type declarations because its compiler does ANSWER.

- type inferencing