2 - the complexity of the move search we know that all the passible can ditates one at No= 27 now after making a guess we are getting K which is the agrade of the guess now after the guess each time we are pliping the bits of the candidates and making the Ust shorter the list is getting to NI= (k) = N! N SETTER (E) SETTERS (N-K) TO AND S will be 1== = Ni 1'smax $N_{e} = \frac{\sqrt{2\pi}n}{2\pi} \left(\frac{n}{2}\right)^{n} = \sqrt{\frac{2}{n}} \frac{2}{\sqrt{n}}$ naw If we exceptate the nation of the change between No and Ne we will get the the eamplexity -> No - VE In 2 In

3- as described in the solving mastermind. PSF in github repo you ean see a defaile salving steps and there I defined the states as I for example p 1=3 > The states will be 1000 > 100 > 1,000 >

not-xol

no the gues I as described in the solving-masterial. Political and how usage examples done there are there gates the Pistis the simple Hadamard gates that places as all the passible candidates the superposition Ahen I am using [] which takes Them then it sums the abtained the N.

1000 > 1000 > 1000 > 10000 > 10010> SUM [[] (010) (000) (000) then there 13 gat [P] gate Plis the gate that 1000 (aux 1000) changes the score of this candidate by Aliping the Pirst O bits in the first an gubit corres panding to the & zeros in the not-Kor (3rd gubit) and this was in case it the difference in seare of this candidate and the scare of the guess that is all those positive if the difference of the scare was pregative then we pap the first 1 bits ok now lets design this gate

Shis gate takes 2 inputs: First input is the difference (quess score - candilat score) second input is the candidat example 1000 1010/10to) now if the difference first inputis positive then need to to do the fallowing we need set up a parcillary gubits and make Them 10> then we need to creat another, gubit and Initialize it to a desect to the maximum (based on the difference number) like (00)=C now we have ancillary gates ancogancyance are do CNOT (Og anca) g (NoT (1 g anc1) g (NoT (Og anc2))

-> anc= @ 1 anc1= 0 anc2=1 now for each anci it- anci de== 1 and the c difference -> Plip bit in first gubit of input? now for case of negative difference.
we do the same but for every anci==0
then flip bit in the first gubit 5- you cand demonstration of the jates in solving mastermind. Pdf