Counting Binary Trees & Selection Trees

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Outline

1. Counting Binary Trees
2. Selection Trees

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1 Counting Binary Trees

2 Selection Trees

Counting Binary Trees

- Consider the following three disparate problems:
  1. The number of distinct binary trees having n nodes.
  2. The number of distinct permutations of the numbers from 1 to n obtainable by a stack.
  3. The number of distinct ways of multiplying $n+1$ matrices.

Counting Binary Trees

- Consider the following three disparate problems:
  1. The number of distinct binary trees having n nodes.
  2. The number of distinct permutations of the numbers from 1 to n obtainable by a stack.
  3. The number of distinct ways of multiplying $n+1$ matrices.
- Amazingly, these problems have the same solution!

Problem One

- The number of distinct binary trees having n nodes.

⋆ Example of $n = 3$.

Problem Two

- The number of distinct permutations of the numbers from 1 to n obtainable by a stack.

1. push 1 $\rightarrow$ pop $\rightarrow$ push 2 $\rightarrow$ pop $\rightarrow$ push 3 $\rightarrow$ pop $\Rightarrow$ 123.
2. push 1 $\rightarrow$ pop $\rightarrow$ push 2 $\rightarrow$ push 3 $\rightarrow$ pop $\rightarrow$ pop $\Rightarrow$ 132.
3. push 1 $\rightarrow$ push 2 $\rightarrow$ push 3 $\rightarrow$ pop $\rightarrow$ pop $\rightarrow$ pop $\Rightarrow$ 321.
4. push 1 $\rightarrow$ push 2 $\rightarrow$ pop $\rightarrow$ pop $\rightarrow$ push 3 $\rightarrow$ pop $\Rightarrow$ 213.
5. push 1 $\rightarrow$ push 2 $\rightarrow$ pop $\rightarrow$ push 3 $\rightarrow$ pop $\rightarrow$ pop $\Rightarrow$ 231.

$\star$ Example of $n = 3$.

Problem Three

- The number of distinct ways of multiplying $n+1$ matrices.

1. $(((M_1 \times M_2) \times M_3) \times M_4)$.
2. $((M_1 \times (M_2 \times M_3)) \times M_4)$.
3. $(M_1 \times ((M_2 \times M_3) \times M_4))$.
4. $(M_1 \times (M_2 \times (M_3 \times M_4)))$.
5. $((M_1 \times M_2) \times (M_3 \times M_4))$.

$\star$ Example of $n = 3$.

Stack Permutation (1/4)

- Recall: preorder, inorder and postorder traversal of a binary tree.
  - Each traversal requires a stack.

Every binary tree has a unique pair of preorder/inorder sequences.

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Every binary tree has a unique pair of preorder/inorder sequences.

- The number of distinct binary trees is equal to the number of inorder permutations obtainable from binary trees having the preorder permutation, $1, 2, \ldots, n$.

Stack Permutation (2/4)

- preorder: A B C E D G H F I
- inorder: B C A E D G H F I

- preorder: A B C (D E F G H I)
- inorder: B C A (E D F G H I)

Stack Permutation (3/4)

- We can show that

  the number of distinct permutations obtainable by passing the numbers $\{1, 2, \ldots, n\}$ through a stack is equal to the number of distinct binary trees with n nodes.

1. push 1 $\rightarrow$ pop $\rightarrow$ push 2 $\rightarrow$ pop $\rightarrow$ push 3 $\rightarrow$ pop $\Rightarrow$ 123.
2. push 1 $\rightarrow$ pop $\rightarrow$ push 2 $\rightarrow$ push 3 $\rightarrow$ pop $\rightarrow$ pop $\Rightarrow$ 132.
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Stack Permutation (4/4)

⋆ Output: inorder traversal:

- visit node → push to stack;
  - going left → keep visiting next node
  - going right → pop the stack
- the leaf → pop the stack until empty

(1, 2, 3) (1, 3, 2) (2, 1, 3) (2, 3, 1) (3, 2, 1)

Go Back to the Matrix Multiplication

- Computing the product of n matrices are related to the distinct binary tree problem.
- $n = 3$:
  1. $((M_1 \times M_2) \times M_3)$.
  2. $(M_1 \times (M_2 \times M_3))$.
- $n = 4$:
  1. $(((M_1 \times M_2) \times M_3) \times M_4)$
  2. $((M_1 \times (M_2 \times M_3)) \times M_4)$
  3. $(M_1 \times ((M_2 \times M_3) \times M_4))$
  4. $(M_1 \times (M_2 \times (M_3 \times M_4)))$
  5. $((M_1 \times M_2) \times (M_3 \times M_4))$

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  2. $(M_1 \times (M_2 \times M_3))$.
- $n = 4$:
  1. $(((M_1 \times M_2) \times M_3) \times M_4)$ (push, pop, push, pop, push, pop)
  2. $((M_1 \times (M_2 \times M_3)) \times M_4)$ (?)
  3. $(M_1 \times ((M_2 \times M_3) \times M_4))$ (?)
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  5. $((M_1 \times M_2) \times (M_3 \times M_4))$ (?)

Matrix Multiplication (2/2)

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- Trivially, $b_1 = 1$, $b_2 = 1$.
- We have also derived that $b_3 = 2$ and $b_4 = 5$.
- We can compute that

$$b_n = \sum_{i=1}^{n-1} b_i b_{n-i}, \text{ for } n > 1.$$

Distinct Binary Trees

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- But, how to compute b_n exactly?

The Generating Function Trick

- **Trick:** Let $B(x) = \sum_{i \geq 0} b_i x^i$ be the generating function for the number of binary trees.

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$$B(x) = \sum_{i \geq 0} b_i x^i = \sum_{i=0}^{\infty} \sum_{k=0}^{i-1} b_k b_{i-1-k} x^i = 1 + \sum_{i=1}^{\infty} \sum_{k=0}^{i-1} b_k b_{i-1-k} x^i$$

$$=$$

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&= 1 + x \sum_{i=1}^{\infty} \sum_{k=0}^{i-1} b_k b_{i-1-k} x^{i-1} \\
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&= 1 + x \sum_{j=0}^{\infty} \sum_{k=0}^{j} b_k b_{j-k} x^j \\
&= 1 + x \left(\sum_{j=0}^{\infty} b_j x^j \right)^2 = 1 + x B(x)^2.
\end{aligned}
$$

The Generating Function Trick

- By the recurrence relation we get:

$$xB(x)^2 = B(x) - 1.$$

- Solving the recurrence relation, we have

$$\begin{aligned} B(x) &= \frac{1-\sqrt{1-4x}}{2x} \\ &= \frac{1}{2x}\left(1 - \sum_{n\geq 0}\binom{1/2}{n}(-4x)^n\right) \\ &= \sum_{m\geq 0}\binom{1/2}{m+1}(-1)^m 2^{2m+1}x^m. \end{aligned}$$

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$$\therefore b_n = \frac{1}{n+1}\binom{2n}{n}.$$

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&= \left(\frac{1}{2}\right)^{n+1}\cdot\frac{1(2-1)(4-1)\cdots(2n-1)}{(n+1)n(n-1)(n-2)\cdots 2\cdot 1}2^{2n+1} \\
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&= \frac{1(1)(3)\cdots(2n-1)}{(n+1)n(n-1)(n-2)\cdots 2\cdot 1}\cdot 2^n\cdot\frac{2\cdot 4\cdot 6\cdots(2n)}{2^n(1\cdot 2\cdots n)} \\
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\end{aligned}
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⋆ supplementary: Stirling's approximation

Catalan number by monotonic lattice paths (1/5)

Catalan number by monotonic lattice paths (2/5)

First time breaking the diagonal

Catalan number by monotonic lattice paths (3/5)

More → than ↑

Catalan number by monotonic lattice paths (4/5)

Catalan number by monotonic lattice paths (5/5)

- The number of monotonic lattice paths not passing the diagonal is

$$\begin{aligned}
\binom{2n}{n} - \binom{2n}{n-1} &= \frac{(2n)!}{n!n!} - \frac{(2n)!}{(n-1)!(n+1)!} \\
&= (2n)!\left(\frac{n+1-n}{n!(n+1)!}\right) \\
&= \frac{1}{n+1}\frac{(2n)!}{n!n!} \\
&= \frac{1}{n+1}\binom{2n}{n}.
\end{aligned}$$

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1 Counting Binary Trees

2 Selection Trees

Scenarios of Using the Selection Trees

- External sorting.
- Data stored in each queue (run) is sorted.

Winner Selection Tree

- In the following figure, computing the first winner takes

Winner Selection Tree

- In the following figure, computing the first winner takes $k-1$ comparisons.

Winner Selection Tree

- In the following figure, computing the first winner takes $k-1$ comparisons. (not better?)
- But wait, how about the following iterations?

Winner Selection Tree

(Winner) Selection Tree:

- Setting up the selection tree: $O(k)$ time.
- Restructuring: $O(\lg k)$ time.
- merging all n items: $O(n \lg k)$ time.

Winner Selection Tree

(Winner) Selection Tree:

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Loser Selection Tree

Loser Tree (after the winner tree has been built)

- Sibling nodes represent the losers.
- The restructuring process can be simplified.

Loser Selection Tree

Loser Tree (after the winner tree has been built)

- Sibling nodes represent the losers.
- The restructuring process can be simplified.

Note for the Loser Selection Tree

- Comparison with the sibling is required for the first step construction.
- After the first construction, we only need to compare each node with its parent; "push" the smaller key value upward and left the "larger" key value as the loser.

Discussions