Linked List (III):

Equivalence Relations, Sparse Matrices & Doubly Linked Lists

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- Equivalence Relations
- Sparse Matrices Revisted
- Ooubly Linked Lists



- Equivalence Relations
- 2 Sparse Matrices Revisted
- Oubly Linked Lists



Equivalence Relation

A relation over a set S is said to be an equivalence relation over S iff it is symmetric, reflexive, and transitive over S.

- reflexive: $x \equiv x$ for each $x \in S$.
- symmetric: for $x, y \in S$, if $x \equiv y$, then $y \equiv x$.
- transitive: for x, y, z, if $x \equiv y$ and $y \equiv z$, then $x \equiv z$.



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Example

Given $0 \equiv 4$, $3 \equiv 1$, $6 \equiv 10$, $8 \equiv 9$, $7 \equiv 4$, $6 \equiv 8$, $3 \equiv 5$, $2 \equiv 11$, $11 \equiv 1$.



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Example

Given $0\equiv 4$, $3\equiv 1$, $6\equiv 10$, $8\equiv 9$, $7\equiv 4$, $6\equiv 8$, $3\equiv 5$, $2\equiv 11$, $11\equiv 1$. We have three equivalent classes:

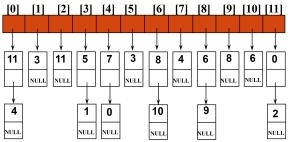
$${0, 2, 4, 7, 11}, {1, 3, 5}, {6, 8, 9, 10}.$$



Lists after Giving Pairs as the Input

```
0 \equiv 4, 3 \equiv 1, 6 \equiv 10, 8 \equiv 9, 7 \equiv 4,
6 \equiv 8, 3 \equiv 5, 2 \equiv 11, 11 \equiv 0.
```

```
typedef struct node *nodePointer;
typedef struct node {
    int data;
    nodePointer link;
};
```







- Equivalence Relations
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Issues for Previous Representation

- When we performed matrix operations such as +, -, or *, the number of nonzero terms varied.
- The sequential representation of sparse matrices suffered from the same inadequacies as the similar representation of polynomials.

Solution:

• Linked list representation for sparse matrices.

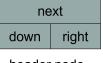


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Solution:

- Linked list representation for sparse matrices.
- Two types of nodes in the representation: header nodes and element nodes.

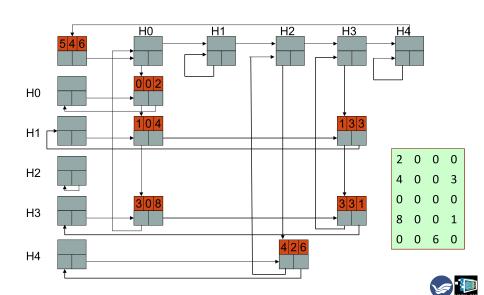


header node



element node





Sparse Matrix Representation

- We represent each column (row) of a sparse matrix as a circularly linked list with a header node.
- The header node for row i is also the header node for column i. The number of header nodes is $\max\{\text{numRows}, \text{numCols}\}$.
- Each element node is simultaneously linked into two lists: a row list, and a column list.
- Each head node is belonged to three lists: a row list, a column list, and a header node list.



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Issues for Singly Linked Lists

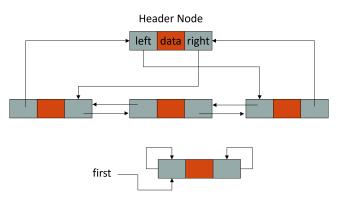
- The only way to find the node that precedes some node *p* is to start at the beginning of the list.
- Sometimes it is necessary to move in either direction.

Doubly linked lists:

```
typedef struct node *nodePointer;
typedef struct node {
   nodePointer llink;
   element data;
   nodePointer rlink;
};
```



ptr = ptr->llink->rlink = ptr->rlink->llink

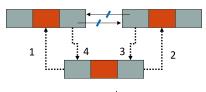


Empty doubly linked circular list with header node



Insertion into a doubly linked circular List

node

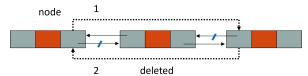


newnode



Insertion into a doubly linked circular List

```
void d_LCL_delete(nodePointer node, nodePointer deleted) {
/* delete from the doubly linked list */
   if (node == deleted)
      printf("Deletion of header node not permitted.\n");
   else {
      deleted->llink->rlink = deleted->rlink;^^I^^I// 1
       deleted->rlink->llink = deleted->llink;^^I^^I// 2
      free(deleted);
   }
}
```





Discussions

