Counting Binary Trees & Selection Trees

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Outline

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1. Counting Binary Trees
2. Selection Trees

Counting Binary Trees

- Consider the following three disparate problems:
  1. The number of distinct binary trees having n nodes.
  2. The number of distinct permutations of the numbers from 1 to n obtainable by a stack.
  3. The number of distinct ways of multiplying $n+1$ matrices.

Counting Binary Trees

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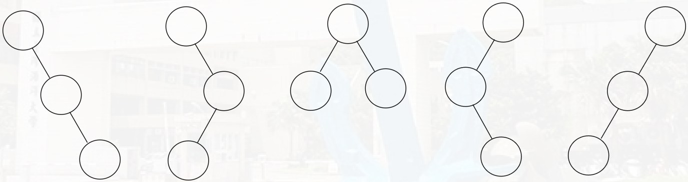
  1. The number of distinct binary trees having n nodes.
  2. The number of distinct permutations of the numbers from 1 to n obtainable by a stack.
  3. The number of distinct ways of multiplying $n+1$ matrices.
- Amazingly, these problems have the same solution!

Problem One

- The number of distinct binary trees having n nodes.



$\star$ Example of $n = 3$.

Problem Two

- The number of distinct permutations of the numbers from 1 to n obtainable by a stack.

1. push 1 $\rightarrow$ pop $\rightarrow$ push 2 $\rightarrow$ pop $\rightarrow$ push 3 $\rightarrow$ pop $\Rightarrow$ 123.
2. push 1 $\rightarrow$ pop $\rightarrow$ push 2 $\rightarrow$ push 3 $\rightarrow$ pop $\rightarrow$ pop $\Rightarrow$ 132.
3. push 1 $\rightarrow$ push 2 $\rightarrow$ push 3 $\rightarrow$ pop $\rightarrow$ pop $\rightarrow$ pop $\Rightarrow$ 321.
4. push 1 $\rightarrow$ push 2 $\rightarrow$ pop $\rightarrow$ pop $\rightarrow$ push 3 $\rightarrow$ pop $\Rightarrow$ 213.
5. push 1 $\rightarrow$ push 2 $\rightarrow$ pop $\rightarrow$ push 3 $\rightarrow$ pop $\rightarrow$ pop $\Rightarrow$ 231.

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Problem Three

- The number of distinct ways of multiplying $n+1$ matrices.

1. $((M_1 \times M_2) \times M_3) \times M_4$.
2. $(M_1 \times (M_2 \times M_3)) \times M_4$.
3. $M_1 \times ((M_2 \times M_3) \times M_4)$.
4. $M_1 \times (M_2 \times (M_3 \times M_4))$.
5. $(M_1 \times M_2) \times (M_3 \times M_4)$.

$\star$ Example of $n = 3$.

Stack Permutation (1/4)

- Recall: preorder, inorder and postorder traversal of a binary tree.
  - Each traversal requires a stack.

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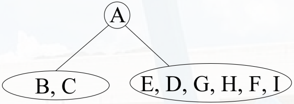
Every binary tree has a unique pair of preorder/inorder sequences.

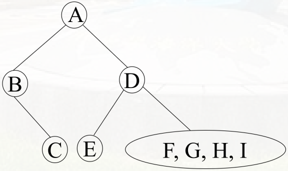
- The number of distinct binary trees is equal to the number of inorder permutations obtainable from binary trees having the preorder permutation, $1, 2, \ldots, n$.

Stack Permutation (2/4)

- preorder: A B C E D G H F I
- inorder: B C A E D G H F I



- preorder: A B C (D E F G H I)
- inorder: B C A (E D F G H I)



Stack Permutation (3/4)

- We can show that

  the number of distinct permutations obtainable by passing the numbers $\{1, 2, \ldots, n\}$ through a stack is equal to the number of distinct binary trees with n nodes.

1. push 1 $\rightarrow$ pop $\rightarrow$ push 2 $\rightarrow$ pop $\rightarrow$ push 3 $\rightarrow$ pop $\Rightarrow$ 123.
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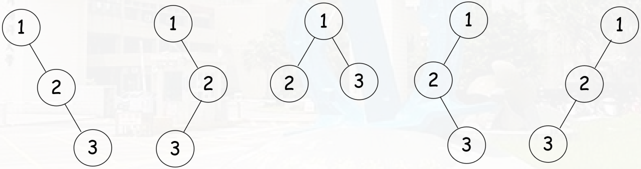
Stack Permutation (4/4)

- visit node → push to stack;
  - going left → keep visiting next node
  - going right → pop the stack
- the leaf → pop the stack until empty



(1, 2, 3) (1, 3, 2) (2, 1, 3) (2, 3, 1) (3, 2, 1)

Go Back to the Matrix Multiplication

- Computing the product of n matrices are related to the distinct binary tree problem.
- $n = 3$:
  1. $(M_1 \times M_2) \times M_3$.
  2. $M_1 \times (M_2 \times M_3)$.
- $n = 4$:
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- $n = 4$:
  1. $((M_1 \times M_2) \times M_3) \times M_4$ (push, pop, push, pop, push, pop)
  2. $(M_1 \times (M_2 \times M_3)) \times M_4$ (?)
  3. $M_1 \times ((M_2 \times M_3) \times M_4)$ (?)
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Matrix Multiplication (2/2)

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- b_n: the number of different ways to compute the product of n matrices.
- Trivially, $b_1 = 1$, $b_2 = 1$.
- We have also derived that $b_3 = 2$ and $b_4 = 5$.
- We can compute that

$$b_n = \sum_{i=1}^{n-1} b_i b_{n-i}, \text{ for } n > 1.$$

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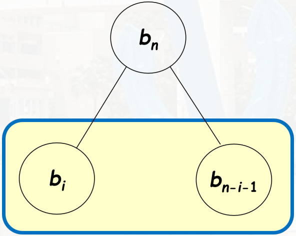

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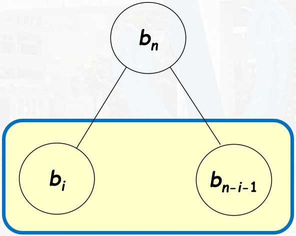

- But, how to compute b_n exactly?

The Generating Function Trick

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$$=$$

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&= \quad 1 + x \sum_{i=1}^{\infty} \sum_{k=0}^{i-1} b_k b_{i-1-k} x^{i-1} \\
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&= 1 + x \sum_{j=0}^{\infty} \sum_{k=0}^{j} b_k b_{j-k} x^j \\
&= 1 + x \left(\sum_{j=0}^{\infty} b_j x^j \right)^2 = 1 + xB(x)^2.
\end{aligned}
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The Generating Function Trick

- By the recurrence relation we get:

$$xB(x)^2 = B(x) - 1.$$

- Solving the recurrence relation, we have

$$\begin{aligned} B(x) &= \frac{1-\sqrt{1-4x}}{2x} \\ &= \frac{1}{2x}\left(1 - \sum_{n\geq 0}\binom{1/2}{n}(-4x)^n\right) \\ &= \sum_{m\geq 0}\binom{1/2}{m+1}(-1)^m 2^{2m+1} x^m. \end{aligned}$$

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$$\therefore b_n = \frac{1}{n+1}\binom{2n}{n}.$$

By the Binomial Theorem...

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$\star$ supplementary: Stirling's approximation

Outline

Scenarios of Using the Selection Trees

- External sorting.
- Data stored in each queue (run) is sorted.

Winner Selection Tree

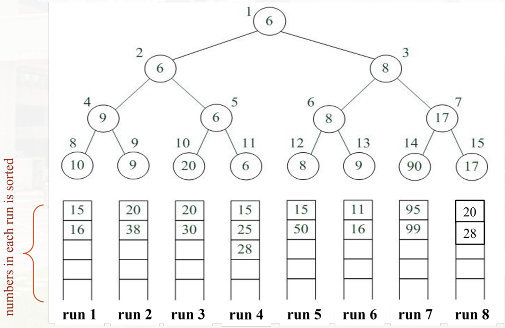
- In the following figure, computing the first winner takes

Winner Selection Tree

- In the following figure, computing the first winner takes $k-1$ comparisons.

Winner Selection Tree

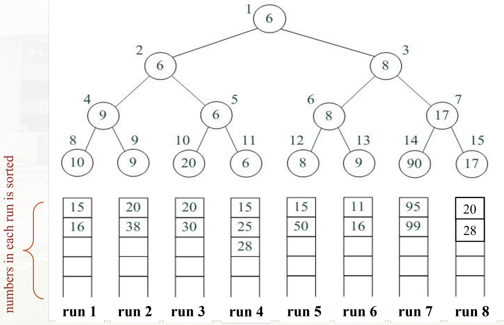
- In the following figure, computing the first winner takes $k - 1$ comparisons. (not better?)
- But wait, how about the following iterations?



Winner Selection Tree

(Winner) Selection Tree:

- Setting up the selection tree: $O(k)$ time.
- Restructuring: $O(\lg k)$ time.
- merging all n items: $O(n \lg k)$ time.



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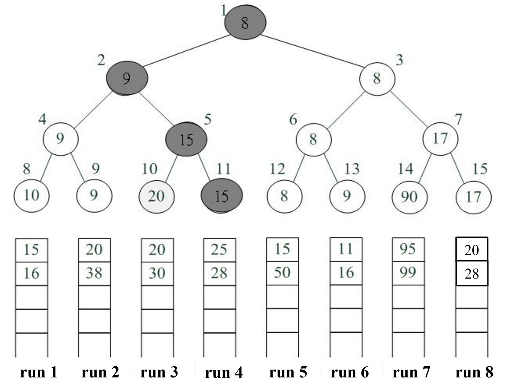
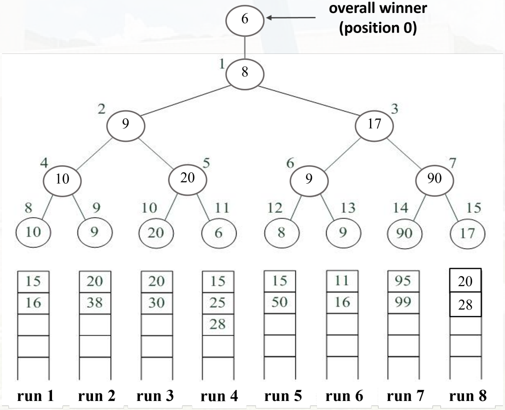


Loser Selection Tree

Loser Tree (after the winner tree has been built)

- Sibling nodes represent the losers.
- The restructuring process can be simplified.



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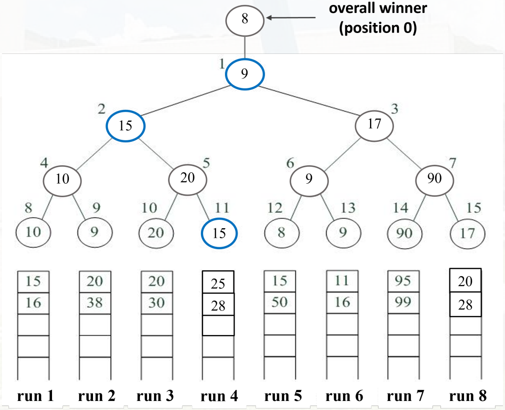

Note for the Loser Selection Tree

- Comparison with the sibling is required for the first step construction.
- After the first construction, we only need to compare each node with its parent; “push” the smaller key value upward and left the “larger” key value as the loser.

Discussions