

# Randomized Algorithms

## Chernoff and Hoeffding Bounds

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# Moment Generating Functions

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$$\mathbf{E}[X^n] = M_X^{(n)}(0) \quad \text{The } n\text{th derivative of } M_X(t) \text{ at } t = 0.$$

# Example

- Consider a geometric random variable  $X$  with parameter  $p$ .
- For  $t < -\ln(1-p)$ ,

$$\begin{aligned} M_X(t) &= \mathbf{E}[e^{tX}] \\ &= \sum_{k=1}^{\infty} (1-p)^{k-1} p e^{tk} \\ &= \frac{p}{1-p} \sum_{k=1}^{\infty} (1-p)^k e^{tk} \\ &= \frac{p}{1-p} ((1-(1-p)e^t)^{-1} - 1). \end{aligned}$$

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$$\begin{aligned} M_X(t) &= \mathbf{E}[e^{tX}] & \therefore M_X^{(1)}(t) &= p(1 - (1-p)e^t)^{-2}e^t, \\ &= \sum_{k=1}^{\infty} (1-p)^{k-1} p e^{tk} & M_X^{(2)}(t) &= 2p(1-p)(1 - (1-p)e^t)^{-3}e^{2t} \\ &= \frac{p}{1-p} \sum_{k=1}^{\infty} (1-p)^k e^{tk} & &+ p(1 - (1-p)e^t)^{-2}e^t. \\ &= \frac{p}{1-p} ((1 - (1-p)e^t)^{-1} - 1). \end{aligned}$$

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# MGF for sum of independent r.v.'s

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- *Proof.*

$$M_{X+Y}(t) = \mathbf{E}[e^{t(X+Y)}] = \mathbf{E}[e^{tX} e^{tY}] = \mathbf{E}[e^{tX}] \cdot \mathbf{E}[e^{tY}] = M_X(t) \cdot M_Y(t).$$

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- *Generalization:*

$$M_{X_1+X_2+\dots+X_k}(t) = M_{X_1}(t) \cdot M_{X_2}(t) \cdots M_{X_k}(t).$$

# Chernoff bounds: Applying Markov's inequality to $e^{tX}$

- From Markov's inequality,

For any  $t > 0$ ,

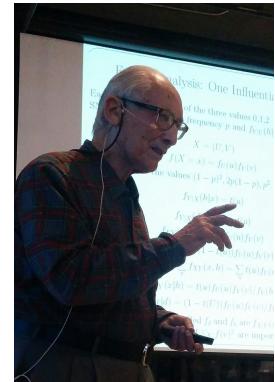
$$\Pr[X \geq a] = \Pr[e^{tX} \geq e^{ta}] \leq \frac{\mathbf{E}[e^{tX}]}{e^{ta}}.$$

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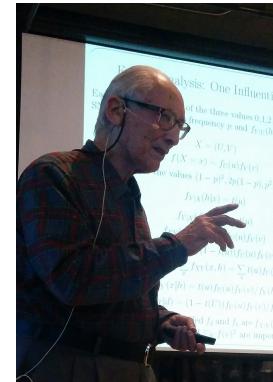
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- Choose appropriate values for  $t$  for specific distributions.



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# Chernoff bounds for sum of Poisson trials

- Poisson trials:
  - ≈ Bernoulli trials
  - while the trials are **not necessarily identical.**
- $X_1, \dots, X_n$ : independent Poisson trials with  $\Pr[X_i = 1] = p_i$ .
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$$\mu = \mathbf{E}[X] = \mathbf{E}\left[\sum_{i=1}^n X_i\right] = \sum_{i=1}^n \mathbf{E}[X_i] = \sum_{i=1}^n p_i.$$

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$$\begin{aligned} M_{X_i}(t) &= \mathbf{E}[e^{tX_i}] & \therefore M_X(t) &= \prod_{i=1}^n M_{X_i}(t) \\ &= p_i e^t + (1 - p_i)e^0 & &\leq \prod_{i=1}^n e^{p_i(e^t - 1)} \\ &= 1 + p_i(e^t - 1) & &= \exp \left\{ \sum_{i=1}^n p_i(e^t - 1) \right\} \\ &\leq e^{p_i(e^t - 1)}. & &= e^{(e^t - 1)\mu}. \end{aligned}$$

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Let  $X = \sum_{i=1}^n X_i$  and  $\mu = \mathbf{E}[X]$ . Then the following Chernoff bounds hold:

1. For  $\delta > 0$ ,

$$\Pr[X \geq (1 + \delta)\mu] \leq \left( \frac{e^\delta}{(1 + \delta)^{(1+\delta)}} \right)^\mu;$$

2. For  $0 < \delta \leq 1$ ,

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# Proof sketch

- Applying Markov's inequality for  $t > 0$ ,
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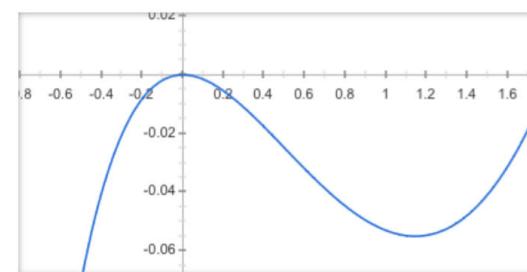
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Graph for  $x-(1+x)\ln(1+x)+x^{2/3}$



[More info](#)

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- Theorem. Let  $X_1, \dots, X_n$  be independent Poisson trials with  $\Pr[X_i = 1] = p_i$ .

Let  $X = \sum_{i=1}^n X_i$  and  $\mu = \mathbf{E}[X]$ . Then the following Chernoff bounds hold:

For  $0 < \delta < 1$ ,

$$\Pr[X \leq (1 - \delta)\mu] \leq \left( \frac{e^\delta}{(1 - \delta)^{(1-\delta)}} \right)^\mu.$$

$$\Pr[X \leq (1 - \delta)\mu] \leq e^{-\mu\delta^2/2}.$$

- ✓ Therefore we have:

$$\Pr[|X - \mu| \geq \delta\mu] \leq 2e^{-\mu\delta^2/3}.$$

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$$\text{Var}[X] = \text{Var} \left[ \sum_{i=1}^n X_i \right] = \sum_{i=1}^n \text{Var}[X_i] = \frac{n}{4}.$$

$$\Pr[X \geq 3n/4] \leq \Pr[|X - \mathbf{E}[X]| \geq n/4] \leq \frac{\text{Var}[X]}{(n/4)^2} = \frac{n/4}{(n/4)^2} = \frac{4}{n}.$$

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	$n=50$	$n=100$	$n=200$
$2/n$	0.2	0.02	0.01
$e^{-n/24}$	0.125	0.016	0.00025

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$$\Pr[X \leq n/2] = \Pr\left[X - \frac{2n}{3} \leq -\frac{n}{6}\right] \leq e^{-(2n/3) \cdot (1/4)^2 \cdot (1/2)} = e^{-n/48} < 0.016.$$

$X_i$  : 1 if  $i$ th test is correct and 0 otherwise

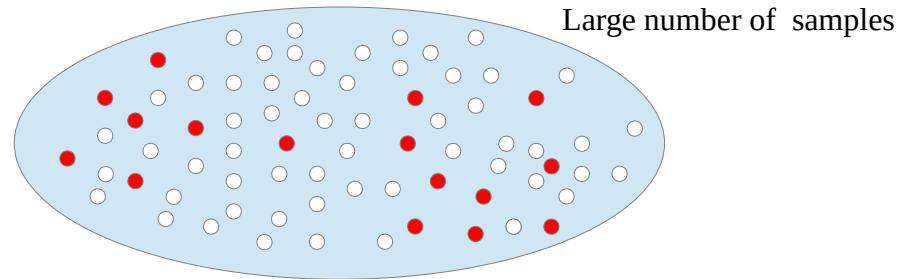
# An application: Parameter Estimation

- **Goal:** evaluating the probability that a particular gene mutation occurs in the population.
- A lab test can determine if a DNA sample carries the mutation.
- However, the test is very **expensive**, so we want to obtain a relatively reliable estimate from a **small** number of samples.



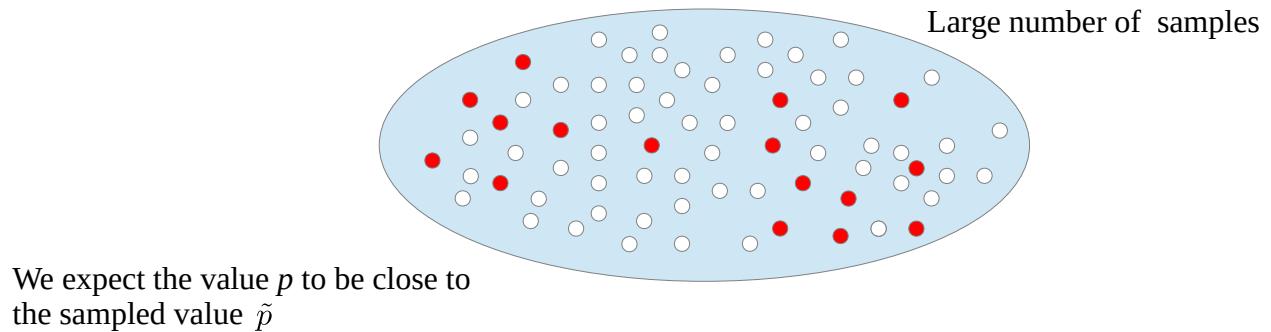
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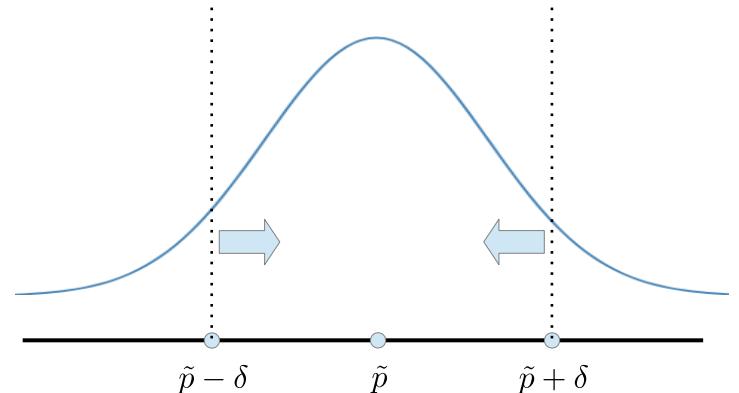
# An application: Parameter Estimation

- Definition. A  $1-\gamma$  **confidence interval** for a parameter  $p$  is an interval

$$[\tilde{p} - \delta, \tilde{p} + \delta]$$

such that

$$\Pr[p \in [\tilde{p} - \delta, \tilde{p} + \delta]] \geq 1 - \gamma.$$



We need to find values of  $\delta$  and  $\gamma$  such that

$$\Pr[p \in [\tilde{p} - \delta, \tilde{p} + \delta]] = \Pr[np \in [n(\tilde{p} - \delta), n(\tilde{p} + \delta)]] \geq 1 - \gamma.$$

# An application: Parameter Estimation

- Apply the Chernoff bound:

$$\begin{aligned}\Pr[p \notin [\tilde{p} - \delta, \tilde{p} + \delta]] &= \Pr\left[X < np\left(1 - \frac{\delta}{p}\right)\right] + \Pr\left[X > np\left(1 + \frac{\delta}{p}\right)\right] \\ &< e^{-np(\delta/p)^2/2} + e^{-np(\delta/p)^2/3} \\ &= e^{-n\delta^2/2p} + e^{-n\delta^2/3p}.\end{aligned}$$

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- But we do not know the value of  $p$ , so it's not useful...
- Take  $p \leq 1$ ,

$$\Pr[p \notin [\tilde{p} - \delta, \tilde{p} + \delta]] < e^{-n\delta^2/2} + e^{-n\delta^2/3}.$$

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Setting  $\gamma = e^{-n\delta^2/2} + e^{-n\delta^2/3}$ , we obtain a trade-off between  $\delta$  and  $n$ .

# Example

- Set  $\gamma = 0.05$ ,  $\delta = 0.03$ .

$$\begin{aligned} e^{-n(0.03)^2/2} + e^{-n(0.03)^2/3} &< 2e^{-n(0.03)^2/3} < 0.05 \\ \Rightarrow e^{-n(0.03)^2/3} &< 0.025 \\ \Rightarrow -n(0.03)^2/3 &< \ln(0.025) \approx -3.6889 \\ \Rightarrow n > 3.6889 \cdot 3/(0.03)^2 &\approx 12296.33. \end{aligned}$$

# The Hoeffding Bound

Wassily Hoeffding (1914–1991)

refer to <https://tinyurl.com/mzz7x8pb>



- Extending the Chernoff bound to general random variables with a **bounded range**.

**Hoeffding's Lemma:** Let  $X$  be a random variable such that  $\Pr[X \in [a, b]] = 1$  and  $\mathbf{E}[X] = 0$ . Then for every  $\lambda > 0$ ,

$$\mathbf{E}[e^{\lambda X}] \leq e^{\lambda^2(b-a)^2/8}.$$

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- Extending the Chernoff bound to general random variables with a **bounded range**.

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# The Hoeffding Bound

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- Extending the Chernoff bound to general random variables with a **bounded range**.

**Theorem:** Let  $X_1, \dots, X_n$  be independent random variables such that  $\mathbf{E}[X_i] = \mu_i$  and  $\Pr[a_i \leq X_i \leq b_i] = 1$  for constant  $a_i$  and  $b_i$ . Then,

$$\Pr \left[ \left| \sum_{i=1}^n X_i - \sum_{i=1}^n \mu_i \right| \geq \epsilon \right] \leq 2e^{-2\epsilon^2 / \sum_{i=1}^n (b_i - a_i)^2}.$$

# Proofs

# Proof of Hoeffding's Lemma

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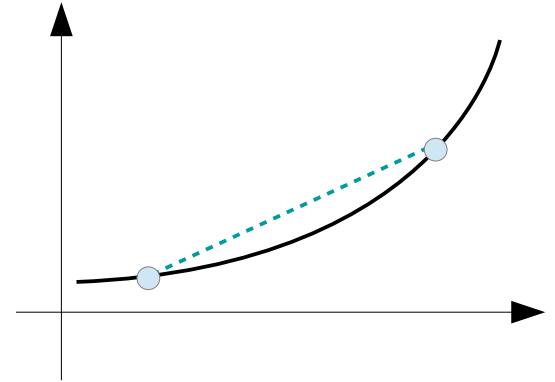
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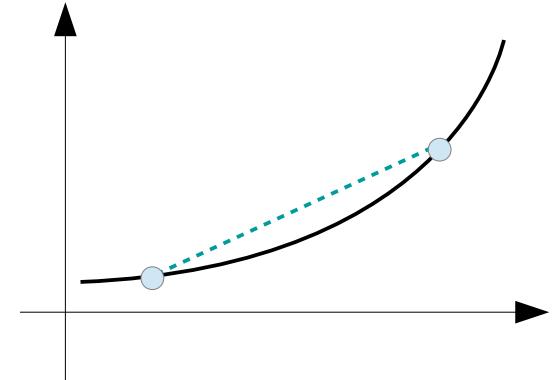
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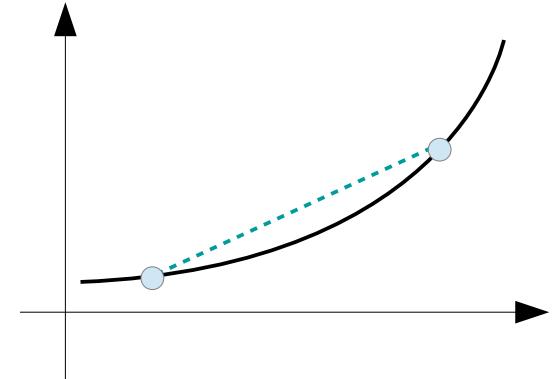
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Let  $\phi(t) = -\theta t + \ln(1 - \theta + \theta e^t)$ .

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For  $t = \lambda(b - a)$ ,

$$\phi(\lambda(b - a)) \leq \frac{\lambda^2(b - a)^2}{8}.$$

Thus,  $\mathbf{E}[e^{\lambda X}] \leq e^{\phi(\lambda(b-a))} \leq e^{\lambda^2(b-a)^2/8}$ .

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$$\begin{aligned}\Pr[Z \geq \epsilon] &= \Pr[e^{\lambda Z} \geq e^{\lambda \epsilon}] \leq \frac{\mathbf{E}[e^{\lambda Z}]}{e^{\lambda \epsilon}} &= \frac{\prod_{i=1}^n \mathbf{E}[e^{\lambda Z_i/n}]}{e^{\lambda \epsilon}} \\ &\leq \frac{\prod_{i=1}^n e^{\lambda^2(b-a)^2/(8n^2)}}{e^{\lambda \epsilon}} \\ &\leq \frac{e^{\lambda^2(b-a)^2/8n}}{e^{\lambda \epsilon}}.\end{aligned}$$

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**Note:**  $Z_i/n \in [(a - \mu)/n, (b - \mu)/n]$ .

Setting  $\lambda = \frac{4n\epsilon}{(b-a)^2}$ ,

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