

# Compositional Semantics: Quantification and Underspecification

**Instructor:** Jackie CK Cheung

COMP-550

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# Outline

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Syntax-driven semantic composition

Quantifiers

Generalized quantifiers

Underspecification

# Power of Lambda Calculus

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They allow us to store partial computations of the MR, as we are composing the meaning of the sentence constituent by constituent.

*Whiskers disdained catnip.*

*disdained*  $\lambda x. \lambda y. disdained(y, x)$

*disdained catnip*  $(\lambda x. \lambda y. disdained(y, x)) catnip$   
 $= \lambda y. disdained(y, catnip)$

*Whiskers disdained catnip*

$(\lambda y. disdained(y, catnip)) Whiskers$   
 $= disdained(Whiskers, catnip)$

# Exercises

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What is the result of simplifying the following expressions in lambda calculus through beta reduction?

$(\lambda z. z)(\lambda y. y \ y)(\lambda x. x \ a)$

$((\lambda x. \lambda y. (x \ y))(\lambda y. y)) \ w)$

$(\lambda x. x \ x) (\lambda y. y \ x) \ z$

# Syntax-Driven Semantic Composition

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Augment CFG trees with lambda expressions

- Syntactic composition = function application

Semantic attachments:

$$A \rightarrow \alpha_1 \dots \alpha_n \qquad \{f(\alpha_j.sem, \dots, \alpha_k.sem)\}$$

syntactic composition

semantic attachment

# Proper Nouns

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Proper nouns are FOL constants

$$PN \rightarrow COMP550 \quad \{COMP550\}$$

Actually, we will **type-raise** proper nouns

$$PN \rightarrow COMP550 \quad \{\lambda x. x(COMP550)\}$$

- It is now a function rather than an argument.
- We will see why we do this.

NP rule:

$$NP \rightarrow PN \quad \{PN.sem\}$$

# Common Nouns

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Common nouns are predicates inside a lambda expression of type  $\langle e, t \rangle$

- Takes an entity, tells you whether the entity is a member of that class

$N \rightarrow student \quad \{\lambda x. Student(x)\}$

Let's talk more about common nouns next class when we also talk about quantifiers.

# Intransitive Verbs

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We introduce an *event variable*  $e$ , and assert that there exists a certain event associated with this verb, with arguments.

$$V \rightarrow rules \quad \{\lambda x. \exists e. Rules(e) \wedge Ruler(e, x)\}$$

Then, composition is

$$S \rightarrow NP VP \quad \{NP.sem(VP.sem)\}$$

Let's derive the representation of the sentence "*COMP-550 rules*"



# Neo-Davidsonian Event Semantics

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Notice that we have changed how we represent events

**Method 1:** multi-place predicate

*Rules(x)*

**Method 2:** Neo-Davidsonian version with event variable

$\exists e. Rules(e) \wedge Ruler(e, x)$

Reifying the event variable makes things more flexible

- Optional elements such as location and time, passives
- Add information to the event variable about tense, modality

# Transitive Verbs

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## Transitive verbs

$V \rightarrow enjoys$

$\{\lambda w. \lambda z. w(\lambda x. \exists e. Enjoys(e) \wedge Enjoyer(e, z) \wedge Enjoyee(e, x))\}$

$VP \rightarrow V NP \quad \{V.sem(NP.sem)\}$

$S \rightarrow NP VP \quad \{NP.sem(VP.sem)\}$

**Exercise:** verify that this works with the sentence “*Jackie enjoys COMP-550*”

# Quantifiers

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## Universal quantifiers

- *all, every*

*All students like COMP-550.*

$\forall x. \text{Student}(x) \rightarrow \text{Like}(x, \text{COMP-550})$

## Existential quantifiers

- *a, an, some*

*Some/A student likes COMP-550.*

$\exists x. \text{Student}(x) \wedge \text{Like}(x, \text{COMP-550})$

Why  $\rightarrow$  for the universal quantifier, but  $\wedge$  for the existential one?

# Russell (1905)'s Definite Descriptions

How to express “*the student*” in FOL?

e.g., *The student took COMP-550.*

Need to enforce three properties:

1. There is an entity who is the student.
2. There is at most one thing being referred to who is a student.
3. The student participates in some predicate, here, “*took COMP-550*”.

# The King of France is Bald

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Property 1 is important. Consider “*The King of France is bald.*”

## Solution 1:

- Define a new constant for KING-OF-FRANCE, much like for proper nouns.

FOL MR becomes *Bald*(KING—OF—FRANCE)

What is the problem with this solution?

# Definite Articles

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*The student took COMP-550:*

1. There is an entity who is the student.
2. There is at most one thing being referred to who is a student.
3. The student participates in some predicate.

What is the range of this existential quantifier?

$$\exists x. Student(x) \wedge \forall y. (Student(y) \rightarrow y = x) \\ \wedge took(x, COMP-550)$$

For simplicity, for now, assume took is a predicate, rather than use event variables.

# Incorporating into Syntax

Now, let's incorporate this to see how lambda calculus can deal with this compositionally.

Semantic attachment for lexical rule for *every*:

$Det \rightarrow every$                        $\{\lambda P. \lambda Q. \forall x. P(x) \rightarrow Q(x)\}$

What do  $P$  and  $Q$  represent?

# Every Student Likes COMP-550

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$Det \rightarrow every \quad \{\lambda P. \lambda Q. \forall x. P(x) \rightarrow Q(x)\}$

$NP \rightarrow Det N \quad \{Det.sem(N.sem)\}$

Let's do the derivation of *Every student likes COMP-550*.

Recall:

$VP \rightarrow V NP \quad \{V.sem(NP.sem)\}$

$S \rightarrow NP VP \quad \{NP.sem(VP.sem)\}$

$V \rightarrow likes$

$\{\lambda w. \lambda z. w(\lambda x. \exists e. Likes(e) \wedge Liker(e, z) \wedge Likee(e, x))\}$

Using explicit event variables again.



# Questions and Exercise

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What are the lexical rules with semantic attachments for *a*? For *the*?

Come up with the derivation of *COMP-550 likes every student*.

# Adjectives

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Can we figure out the pattern for adjectives?

*student*

$\lambda x. Student(x)$

*smart student*

$\lambda x. Smart(x) \wedge Student(x)$

*smart*

?

Also need an augmented rule for  $N \rightarrow A N$

# Scopal Ambiguity: Multiple Quantifiers

What are the possible readings for the following?

*Every student took a course.*

This is known as **scopal ambiguity**.

# Scopal Ambiguity

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*Every student took a course.*

*every* > *a*

$\forall x. \text{Student}(x)$

$\rightarrow (\exists y. \text{Course}(y) \wedge \exists e. \text{took}(e) \wedge \text{taker}(e, x) \wedge \text{takee}(e, y))$

*a* > *every*

$\exists y. \text{Course}(y)$

$\wedge (\forall x. \text{Student}(x) \rightarrow \exists e. \text{took}(e) \wedge \text{taker}(e, x) \wedge \text{takee}(e, y))$

Would like a way to derive **both** of these readings from the syntax. What would we get with our current method?

# Underspecification

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**Solution:** Derive a representation that allows for both readings

**Underspecified representation** – A meaning representation that can embody all *possible* readings without explicitly enumerating all of them.

Other cases where this is useful:

- We are genuinely missing some information (e.g., the tense information), so we choose not to include it in the meaning representation.

# Cooper Storage (1983)

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Associate a **store** with each FOL expression that allows both readings to be recovered.

*Every student took a course.*

$\exists e. took(e) \wedge taker(e, s_1) \wedge takee(e, s_2)$

$(\lambda Q. \forall x. Student(x) \rightarrow Q(x), 1),$

$(\lambda Q. \exists y. Course(y) \wedge Q(y), 2)$

# Recovering the Reading

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Once we know which reading we want (e.g., by looking at the context), recover the store:

1. Select order to incorporate quantifiers
2. For each quantifier:
  - Introduce lambda abstraction over the appropriate index variable
  - Do beta-reduction

# Example: 1, then 2

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*Every student took a course.*

$$\begin{aligned} & \exists e. \text{took}(e) \wedge \text{taker}(e, s_1) \wedge \text{takee}(e, s_2) \\ & (\lambda Q. \forall x. \text{Student}(x) \rightarrow Q(x), 1), \\ & (\lambda Q. \exists y. \text{Course}(y) \wedge Q(y), 2) \end{aligned}$$

1 first:

$$\begin{aligned} & (\lambda Q. \forall x. \text{Student}(x) \rightarrow Q(x)) \\ & (\lambda s_1. \exists e. \text{took}(e) \wedge \text{taker}(e, s_1) \wedge \text{takee}(e, s_2)) \\ & = \forall x. \text{Student}(x) \rightarrow \exists e. \text{took}(e) \wedge \text{taker}(e, x) \wedge \text{takee}(e, s_2) \end{aligned}$$

Then 2:

$$\begin{aligned} & (\lambda Q. \exists y. \text{Course}(y) \wedge Q(y)) \\ & (\lambda s_2. \forall x. \text{Student}(x) \rightarrow \exists e. \text{took}(e) \wedge \text{taker}(e, x) \wedge \text{takee}(e, s_2)) \\ & = \exists y. \text{Course}(y) \wedge \forall x. \text{Student}(x) \rightarrow \exists e. \text{took}(e) \wedge \text{taker}(e, x) \wedge \text{takee}(e, y) \end{aligned}$$



# Compositional Rules

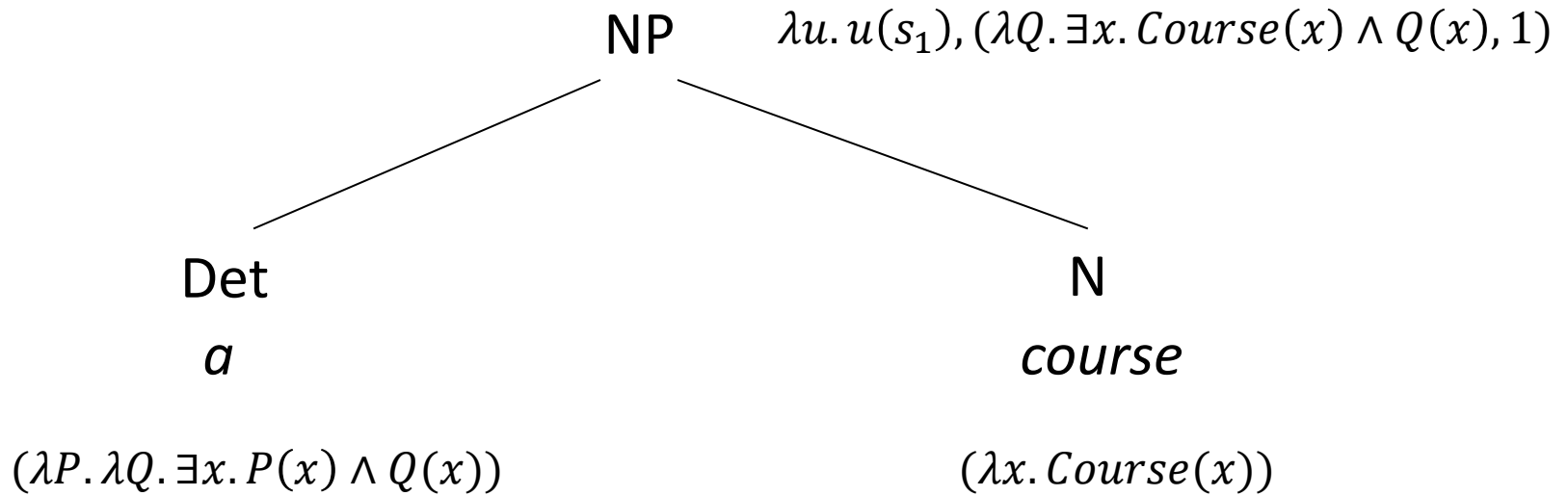
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We also need new rules with semantic attachments for our quantifiers:

- Composing quantifier with N is now modifying the *inside* part of a store
- An NP is now introduces a new index variable, which is wrapped in a lambda expression

# A Course

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What is the semantic attachment for NP  $\rightarrow$  Det N? Use .sem.store to access the store.

# Exercise

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Finish the derivation for the underspecified representation of *Every student took a course*.

Recall:

$VP \rightarrow V \ NP \quad \{V.sem(NP.sem)\}$

$S \rightarrow NP \ VP \quad \{NP.sem(VP.sem)\}$

$V \rightarrow took$

$\{\lambda w.\lambda z.w(\lambda x.\exists e.Took(e) \wedge Taker(e,z) \wedge Takee(e,x))\}$

$Det \rightarrow every \quad \{(\lambda P.\lambda Q.\forall x.P(x) \rightarrow Q(x))\}$

$N \rightarrow student \quad \{\lambda x.Student(x)\}$

# Question

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How would we disambiguate between the possible readings?