

Computer systems

course “Essentials of computing systems”

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- 1 Levels of a computer
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computer

- A computer is a device that can do work for its users by executing instructions given to it.
- A computing device is considered as a **computer** if it has the following characteristics:
 - **Data-dependent instruction sequence.** A computer must include conditional branch constructors, whose execution depends on values that can only be calculated during execution time.
 - **Data-dependent data selection.** A computer must be able to use data to determine the actual data to be used in calculations.
- It is a **calculator** if it lacks either or both characteristics.



program

- The sequence of instructions that describe the task to be performed is the **program**.
- The electronic circuits of each computer are able to directly execute a limited set of simple instructions, into which all its programs must be converted before they can be executed.
- These simple instructions are operations like:
 - add two numbers;
 - test if a number is equal to zero;
 - copy a piece of data from one part to another;
 - decide what instruction to execute next based on the evaluation of some logical condition.

machine language

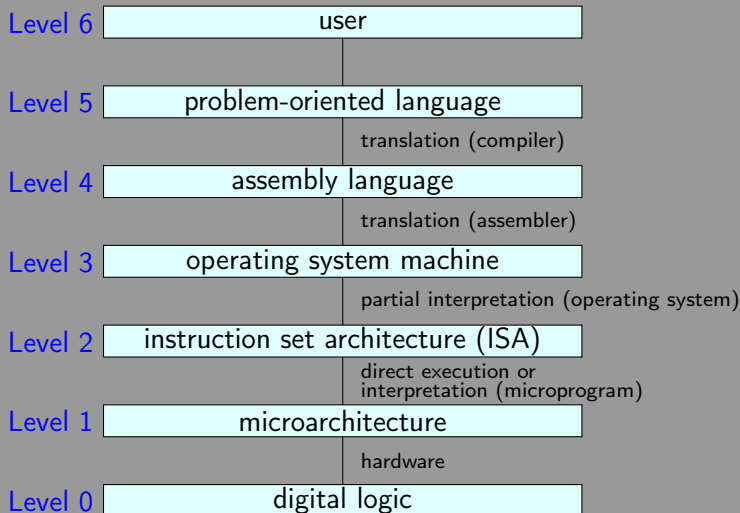
```

HERBYMON (V1.1)
SP 21 1F 1F HL DE BC A
.. DB21 DE92 F38B BFA3 0007 0300 00
.. D000 F3 DI
.. D001 2C 40 LD H,#40
.. D002 2C 92 LD A,#92
.. D003 CD 24 00 CALL #0024
.. D008
HERBYMON (V1.1)
SP 21 1F 1F HL DE BC A
.. DB21 DE92 F38B BFA3 0007 0300 00
.. M 4000 4008
.. 4000 E5 CD 39 54 "T"
.. 4004 44 4D ED 78 "DM.x"
$ AAAA
%1010101010101010 $AAAA #43690
.. D 4000 4004
.. 4000 E5
.. 4001 CD 39 54 PUSH HL
CALL #5439

```

- The primitive instructions constitute a language, called a **machine language**.
- Programmers use these instructions to operate the computer.
- Programming with machine languages is usually difficult and tedious for people.
- The machine language exists mainly to be executed by the computer.
- New languages can be constructed on top of the machine language to rise the abstraction level.

a seven-level computer



Levels 4 and 3

- A fundamental frontier exists between levels 4 and 3.
- The lowest three levels are intended for running the interpreters and translators needed to support the highest levels.
- These interpreters/translators are developed by systems programmers (specialised in constructing virtual machines).
- Levels 4 and above are intended for the applications programmer with a given problem to solve in a specific domain.
- The machine languages of levels 1, 2, and 3 are numeric.
- Programs in these levels consist of long series of numbers (bits), which are appropriate for machines.
- At level 4 and above, the languages contain words and mnemonics that humans are able to easily understand.

programmer's perspective

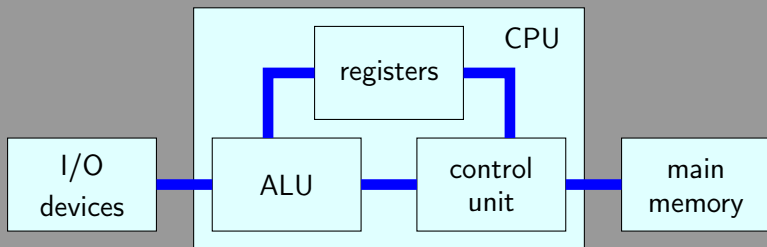
- Each level represents a distinct abstraction of the computer, with different concepts, objects and operations.
- This division is convenient as it permits those that work at a given level to be unaware of what happens in the other levels.
- From a programmer's perspective, it is relevant to understand how computers work at the lower levels.
- Whenever a program written in a high-level language is not behaving as expected, it is seldom necessary to analyse what happens at the lower levels.

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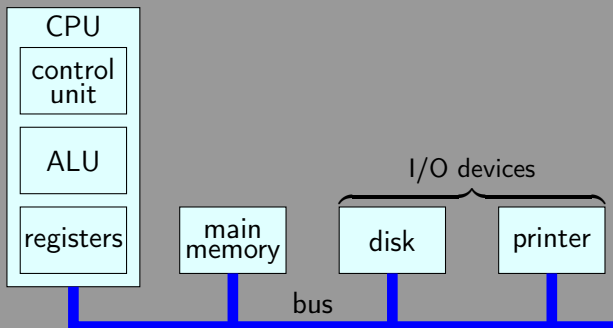
von Neumann architecture

- A computer is a complex system that contains millions of electronic components.
- To describe a computer, it is essential to follow an hierarchical approach.
- Almost all uniprocessor (or scalar) computers follow the so-called **von Neumann architecture**.



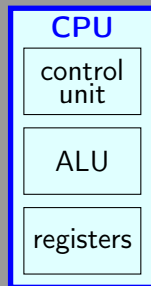
von Neumann architecture

- The von Neumann architecture has been extended in several ways to address some of its limitations.
- A well-known proposal is the inclusion of a **system bus** that connects the components of the computer.
- The system bus has three main elements: data bus, address bus, and control bus.



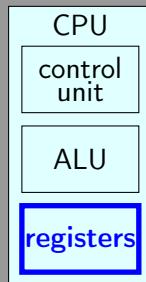
central processing unit (CPU)

- The **CPU** is the “brain” of the computer.
- The CPU is responsible for the execution of the program, stored as a sequence of machine language instructions in the main memory.
- Each instruction directs the CPU to perform some basic task.
- The CPU does all this mechanically, without understanding what is the purpose for executing the instruction.
- The CPU is composed of several distinct components: control unit, ALU, and registers.



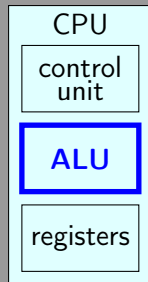
registers

- The CPU contains a small number of high-speed memory registers.
- A **register** is used to store temporary results and status information.
- All the registers have usually the same size.
- Registers can be read and written at a very high speed, since they are internal to the CPU.
- The **instruction pointer (IP)** indicates the address of the next instruction to be executed.
- The **instruction register (IR)** stores the instruction currently being executed.



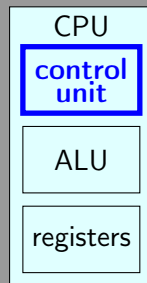
arithmetic logic unit (ALU)

- The **ALU** performs basic operations such as addition, subtractions, and comparisons.
- Usually, an ALU has two data inputs and one data output.
- The operations carried on by the ALU affect the status registers, which indicate attributes related to the last arithmetic or logic operation.
- The operation that the ALU performs is signalled by the control unit, based on the instruction that is executed in each instant.



control unit

- The **control unit** is responsible for orchestrating the components of the computer.
- It ensures that the instruction that is being executed produces its expected effects.
- This is accomplished by activating in the right sequence the control signals that are sent to the computer components.
- This process is repeated indefinitely, while the computer is powered on.
- The CPU receives an instruction from the memory, decodes it, and executes it using data obtained from the memory or stored in the registers.
- When the execution of an instruction is finished, the cycle starts again for the next instruction.



fetch-decode-execute cycle

- The control unit can be seen as a cyclic (micro)program that goes through the set of steps to execute the instructions of another program.
- The **fetch-decode-execute cycle** is central to the operation of computers, since they all just execute instructions:
 - ① Fetch the instruction from memory and store it into the IR.
 - ② Determine the type of instruction.
 - ③ If a word in memory is used by the instruction, determine its location.
 - ④ Store the word into a CPU register, if needed.
 - ⑤ Produce the effects of the instruction, i.e., perform some simple operation dictated by the instruction.
 - ⑥ Change the IP to refer to the next instruction.

clock

- Most computer systems have a **clock** signal that acts like a heartbeat.
- The clock regulates the passage of time within the computer.
- A clock emits periodically a pulse.
- A **clock cycle** is the time between two consecutive ticks (represents one discrete time unit).
- The different operations performed by a processor are synchronised by the clock.
- All these operations begin with the pulse of the clock.
- The speed at which the instructions are executed depends on the frequency of the clock (cycles per second or Hertz).

main memory

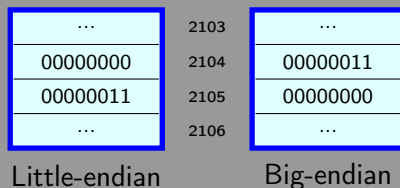
- The **main memory** is the component of the computers, where programs and data are both stored.
- A memory consists of a number of cells or locations, each of which can store a piece of information.
- Each cell has a number, called its **address**, by which programs can refer to it.
- A memory with n cells uses addresses from 0 to $n - 1$.
- All cells in a memory contain the same number of bits.
- The main memory is a **random-access memory (RAM)**.
- RAM memories are volatile, which means that the information disappears if the supply voltage is turned off.

bytes and words

- The cell is the smallest addressable unit.
- Almost all modern computers use 8-bit (1 byte) cells.
- Bytes are grouped into **words** (a computer with a 32-bit word has 4 bytes/word).
- The word is also important since most instructions operate on entire words, for example, adding two words together.
- A 32-bit machine has 32-bit registers and instructions for manipulating 32-bit words.

little- and big-endian

- When storing a multibyte value in a byte-addressable memory, the order of the individual bytes is relevant.
- The number $11\ 0000\ 0000_2$ requires two bytes of memory: $0000\ 0000_2$ and $0000\ 0011_2$.

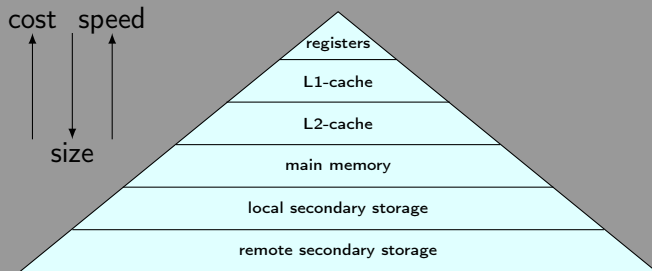


stored-program computer

- Digital computers possess a great versatility, as they can execute whatever programs their users wish.
- This flexibility is related to the ability of a computer to load the programs in its internal memory.
- By changing the program in the memory, one changes the problem the computer is solving.
- This justifies the notion of a **stored-program computer**, one that can be used for any problem.
- During a computation, a given instruction may also be modified, i.e., changed to a different instruction
- This gives rise to **self-modifying code**, which is straightforward to perform at the machine-level.

memory hierarchy

- The main memory of a computer is never enough.
- The traditional solution to store a great quantity of data is to use a **memory hierarchy**.



- The aim is to establish a system with cost almost as low as the cheapest level and speed almost as fast as the fastest level.
- All data in a given level can be found in the next lower level.

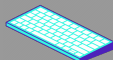
memory hierarchy

level	1	2	3	4
name	registers	cache	main memory	disk storage
typical size	< 1 kB	< 16 MB	< 512 GB	> 1 TB
implementation	CMOS	SRAM	DRAM	magnetic disk
access time (ns)	0.25–0.5	0.5–25	50–250	5,000,000
bandwidth (MB/s)	50,000–500,000	5000–20,000	2500–10,000	50–500

Adapted from [Hennessy and Patterson; 2007]

I/O devices

- Computers are built so that they can be expanded with new I/O (input/output) devices: hard disk, keyboard, mouse, monitor, printer, network interface, scanner.
- An **I/O device** is a hardware system used by humans or systems to communicate with a computer.
- The CPU must communicate with and control these devices.
- For each device, there is a **device driver**, which is software that the CPU executes when it has to deal with that device.
- A new device can be installed by connecting the device into the computer and installing its device driver software.
- The device driver allows the CPU to communicate with the respective device.



polling

- When a given device produces data that needs to be processed, somehow the CPU needs to take some action.
- With **polling**, the CPU keeps checking for incoming data over and over.
- Whenever it detects that data was made available by the device, it processes it.
- Polling is very simple, but is also very inefficient.

interrupts

- Interrupts are generally used instead of polling.
- An **interrupt** is a signal sent by a device to the CPU to request its attention.
- The CPU reacts by suspending what it is doing to answer to the interrupt.
- The CPU saves information about what it is currently doing so that it can return to it afterwards.
- Then the CPU jumps to some fixed memory location and begins executing the instructions stored there.
- Those instructions correspond to the interrupt handler that does the processing necessary to respond to the interrupt.
- Once the interrupt is processed, the CPU returns to the process that was suspended.

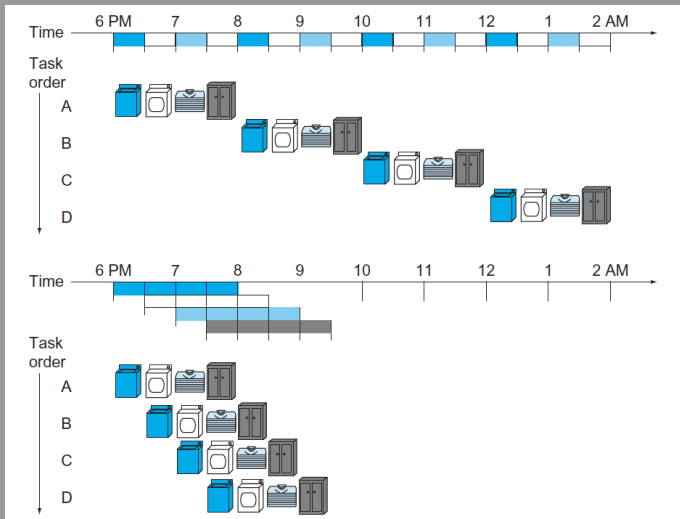
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improving computers

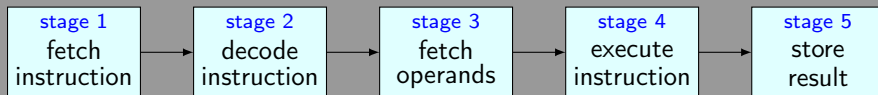
- Improving the characteristics of computers, like performance, cost, reliability, energy consumption, is a concern for computer engineers and architects.
- Performance is an inescapable issue, since making computers run faster is always relevant.
- One possible approach is to exploit parallelism, which can have different forms: instruction-, data-, and processor-level.
- **Instruction-level parallelism with pipelining** explores individual instructions to obtain higher throughput.

pipeline analogy



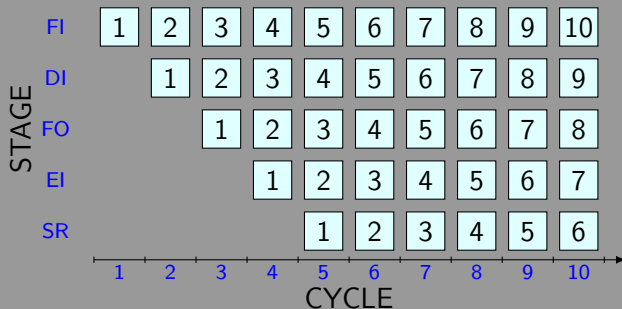
pipeline

- With **pipelining**, the process of executing instructions is divided in stages.
- Multiple instructions can be overlapped in execution.
- Each stage is responsible for a part of the process and all can run in parallel, thus speeding up the process.



ideal operation of the pipeline

- Stages operate in parallel, which accelerates the process.
- Pipelining increases the number of instructions that are simultaneously under execution.
- **Instruction throughput** (number of instructions completed per time unit) is improved.
- Latency (time that is required to complete each individual instruction) is maintained.

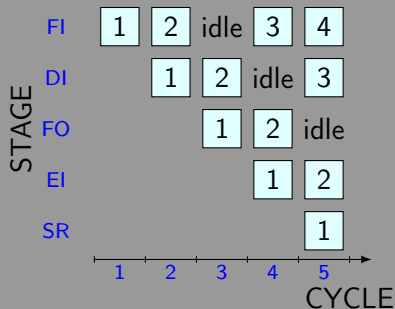
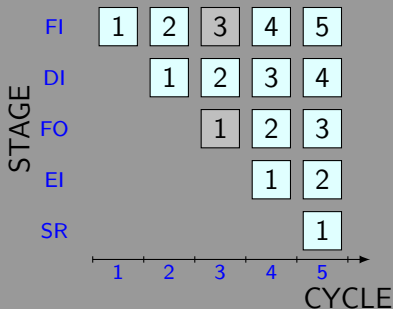


aspects affecting pipeline

- The maximum theoretical speed up is very unlikely to occur in practice, because it is difficult to balance the time it takes to complete each stage.
- All stages must be equally balanced, otherwise the faster stages must wait for the slower ones.
- Three major aspects negatively affect the theoretical speed up:
 - ① resource conflicts,
 - ② data dependencies,
 - ③ conditional branch statements.

Resource conflicts

- A **resource conflict** occurs whenever two pipeline stages need to simultaneously access the same resource (e.g., memory).



data dependency

- A **data dependency** occurs whenever the result of one instruction, not yet available, is also the operand of a subsequent instruction.
- Some solutions exist to handle these conflicts.
- By inserting a delay (e.g., a `nop` instruction) into the pipeline, sufficient time passes and the conflict disappears.
- Some architectures assume that this problem is solved by the compiler, which must reorder instructions.

branch instruction

- Problems in the pipeline may also be caused by a **branch instruction** that modifies the normal flow of a program.
- If the conditional branch is taken, the subsequent instructions must not be executed and the pipeline must be emptied.
- Some architectures predict the outcome of a conditional branch to try to identify the instructions that will be executed next.
- In some cases, compilers try to solve branching issues by rearranging the machine code to cause a delayed branch.
- Another alternative is to initiate fetches on both paths of the branch. When the branch is actually executed, the correct path is identified and the pipeline can be safely continued.