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- Module SelfRefPuzzle
A self-referential logic puzzle, based on
https://davecturner.github.io/2018/10/22/kitty-grundman-self-referential-puzzle.html
EXTENDS Naturals, FiniteSets
VARIABLES X, S We define X and S as variables so that we could use TLC
Puzzle \triangleq
      \land~X \in 1 \dots 10
      \land S \in [1..10 \rightarrow \text{BOOLEAN}]
      \wedge LET Statements \stackrel{\triangle}{=} DOMAIN S The statement indices
               IsTrue(x) \triangleq S[x]
               IsFalse(x) \stackrel{\triangle}{=} \neg S[x]
                                 \triangleq \{i \in Statements : IsTrue(i)\}
                True
                                 \triangleq \{i \in Statements : IsFalse(i)\}
                False
                IsEven(x) \stackrel{\triangle}{=} x\%2 = 0
                                 \triangleq x\%2 \neq 0
                IsOdd(x)
                                 \stackrel{\Delta}{=} \neg (a \equiv b) Exclusive or
               a \otimes b
             \land \neg \forall i \in Statements : IsTrue(i)
             \land \neg \forall i \in Statements : IsFalse(i)
             \land S[1] = \text{Let } sum[s \in \text{Subset } Nat] \stackrel{\triangle}{=} \text{ if } s = \{\} \text{ then } 0
                                                                                       ELSE LET x \stackrel{\triangle}{=} \text{CHOOSE } x \in s : \text{TRUE}
                                                                                               IN x + sum[s \setminus \{x\}]
                            IN X = sum[False]
             \wedge S[2] = (X < Cardinality(False) \wedge IsTrue(10))
             \land S[3] = ((Cardinality(True) = 3) \otimes IsFalse(1))
             \wedge S[4] = (1 \dots 3 \subseteq False \vee IsTrue(9))
             \wedge S[5] = (IsOdd(X) \otimes IsTrue(7))
             \land S[6] = \text{LET } Odds \stackrel{\triangle}{=} \{i \in Statements : IsOdd(i)\}
                                   Cardinality(Odds \cap False) = 2
             \wedge S[7] = IsTrue(X)
             \land S[8] = \text{LET } Evens \stackrel{\triangle}{=} \{i \in Statements : IsEven(i)\}
                            \text{IN} \quad \textit{Evens} \subseteq \textit{True} \vee \textit{Evens} \subseteq \textit{False}
             \land S[9] = \text{LET } First(s) \stackrel{\triangle}{=} \text{CHOOSE } x \in s : \forall y \in s : y \geq x
                            IN (X = 3 * First(True)) \lor IsFalse(4)
             \wedge S[10] = (IsEven(X) \vee IsTrue(6))
To get TLC to solve the puzzle for us, we have to put it in the form of a temporal specification and
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an invariant. A counterexample of the invariant would be the solution. To verify its uniqueness, we then let TLC check that the solution is an invariant (it is). $A = \frac{\Delta}{2} \left[\frac{1}{2} \left(\frac{1}{2} \frac{1}{2$

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Spec \qquad \stackrel{\triangle}{=} Puzzle \land \Box [\text{UNCHANGED } \langle X, S \rangle]_{\langle X, S \rangle} \\ Invariant \stackrel{\triangle}{=} X \notin 1 ... 10 \text{ Taunt TLC into proving us wrong} \\ \land S = \langle \text{False, false, true, true, false, true, false, true, false, true, false} \rangle \\ \land X = 9
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