

Fundamentals/ICY: Databases

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WEEK 7 –Friday

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Reminder of Monday

(necessarily non-symmetric) **1:M** recursive:
“EMPLOYEE Manages EMPLOYEE”

FIGURE 4.23 IMPLEMENTATION OF THE 1:M “EMPLOYEE MANAGES EMPLOYEE” RECURSIVE RELATIONSHIP

Table name: EMPLOYEE_V2

Database name: Ch04_PartCo

	EMP_CODE	EMP_LNAME	EMP_MANAGER
►	101	Waddell	102
	102	Orincona	
	103	Jones	102
	104	Reballoh	102
	105	Robertson	102
	106	Deltona	102

Just a standard 1:M implementation except linking a table to itself.
No redundancy problem.

non-symmetric M:N recursive: “PART Contains PART”

FIGURE 4.21 IMPLEMENTATION OF THE M:N RECURSIVE “PART CONTAINS PART” RELATIONSHIP

Table name: COMPONENT

Database name: Ch04_PartCo

	COMP_CODE	PART_CODE	COMP_PARTS_NEEDED
▶	C-130	AA21-6	4
	C-130	AB-121	2
	C-130	E129	1
	C-131A2	E129	1
	C-130	X10	4
	C-131A2	X10	1
	C-130	X34AW	2
	C-131A2	X34AW	2

Table name: PART

	PART_CODE	PART_DESCRIPTION	PART_IN_STOCK
▶	AA21-6	2.5 cm. washer, 1.0 mm. rim	432
	AB-121	Cotter pin, copper	1,034
	C-130	Rotor assembly	36
	E129	2.5 cm. steel shank	128
	X10	10.25 cm. rotor blade	345
	X34AW	2.5 cm. hex nut	879

The COMPONENT entity type is just a bridging type, linking PART to itself. NB: its first two columns both refer to PART's PK but must be differently named. No redundancy problem.

symmetric (1:1) recursive relationship:
“EMPLOYEE Married to EMPLOYEE”

Suppose you tried the following:

FIGURE 4.19 THE 1:1 RECURSIVE RELATIONSHIP “EMPLOYEE IS MARRIED TO EMPLOYEE”

Table name: EMPLOYEE_V1

Database name: Ch04_PartCo

	EMP_NUM	EMP_LNAME	EMP_FNAME	EMP_SPOUSE
►	345	Ramirez	James	347
	346	Jones	Anne	349
	347	Ramirez	Louise	345
	348	Delaney	Robert	
	349	Shapiro	Anton	346

Redundancy problem!!

New

Symmetry is the Problem

- ◆ *A non-symmetric 1-1 relationship would not have the problem shown on previous slide.*
- ◆ *A symmetric M:N relationship would have a redundancy problem, whether implemented as in the 1-1 case or by a bridging table.*
 - *E.g.: being-sibling-of.*

symmetric (1:1) recursive relationship: redundant & non-redundant implementations

FIGURE 4.48 VARIOUS IMPLEMENTATIONS OF A 1:1 RECURSIVE RELATIONSHIP

Table name: EMPLOYEE_V1

	EMP_NUM	EMP_LNAME	EMP_FNAME	EMP_SPOUSE
▶	345	Ramirez	James	347
	346	Jones	Anne	349
	347	Ramirez	Louise	345
	348	Delaney	Robert	
	349	Shapiro	Anton	346

Database name: Ch04_PartCo

First implementation

Table name: EMPLOYEE

	EMP_NUM	EMP_LNAME	EMP_FNAME
▶	345	Ramirez	James
	346	Jones	Anne
	347	Ramirez	Louise
	348	Delaney	Robert
	349	Shapiro	Anton

Table name: MARRIED_V1

	EMP_NUM	EMP_SPOUSE
▶	345	347
	346	349
	347	345
	349	346

Second implementation

Table name: MARRIAGE

		MAR_NUM	MAR_DATE
▶	+	1	04-Mar-03
	+	2	02-Feb-99

Table name: MARPART

	MAR_NUM	EMP_NUM
▶	1	345
	1	347
	2	346
	2	349

Table name: EMPLOYEE

	EMP_NUM	EMP_LNAME	EMP_FNAME
▶	345	Ramirez	James
	346	Jones	Anne
	347	Ramirez	Louise
	348	Delaney	Robert
	349	Shapiro	Anton

- 1) As previously—redundant.
- 2) MARRIED_V1 is just a bridging entity type: still redundant.
- 3) MARRIAGE together with MARPART act as a sort of bridge.
Non-redundant.

Symmetric M:N, etc.

- ◆ Method 3 on previous slide can straightforwardly be generalized to:
 - ¶ symmetric recursive $M:N$ relationships
 - ¶ (((partially-)symmetric&recursive, etc. relationships that link more entities together, whatever the connectivity --- $M:N:P$, $1:1:1$, $M:1:P$, ...))

Summary: Creating ERM/ERDs

- ◆ Designing an ER model for a database is an *iterative process*, because, e.g.:
 - As you proceed, you think of **new ways of conceiving** what's going on (much as in ordinary programming)
 - **Multivalued attributes** need to be re-represented eventually
 - **Supertype/subtype** notation needs eventually to be converted into more standard diagram notation
 - **M:N relationships** can be included as such at an early stage, but usually need to be replaced by means of bridging entity types later
 - **1:1 relationships** raise a **red flag**: may indicate poor design, especially when mandatory both ways.
 - *But standard in supertype/subtype representation, and natural in recursive relationships like “married-to”.*
 - **Symmetric** recursive relationships usually need **special handling**.
 - Conversion of ERM portions to a **“Normal Form” – LATER.**

**Next major databases topic
is
Normalization
but for now ...**

MATHEMATICAL BACKGROUND TO TABLES

Mathematical “sets”: Basics

- ◆ A “**set**” is an **unordered** collection of items of any sorts (people, numbers, numerals, shoes, atoms, strings of characters, databases, sets, blades of grass, ...) **without any duplication of items**.

The items are called “elements” or “members”.

- ◆ $S = \{34, \text{JAB}, 59, \text{UoB}\}$, where “JAB” is a name for me and “UoB” is a name for this university,

means that

S is the set consisting of (exactly) the following four items:

the abstract number 34, me, the abstract number 59, this university.

Basics, contd

- ◆ $\{34, \text{JAB}, 59, \text{UoB}\} = \{\text{UoB}, 59, 34, \text{JAB}, \text{JAB}, 34\}$
 - **Order of writing the members doesn't matter; duplication in the writing doesn't duplicate the member.**
- ◆ A set can be **infinite** (e.g., the set of all whole numbers).
- ◆ A set can contain just **one member** (e.g. the set whose only element is your favourite pencil). **Singleton set.**
 - It's **different** from the member itself..
- ◆ There's a set **with no members at all**: the “**empty set**”, usually notated as \emptyset , but can also be written $\{ \}$.
 - Somewhat analogous to zero, or a new committee which has no members yet.
 - There is only one empty set (rather than an empty set of numbers, an empty set of pencils, etc.)

Another Notation

◆ $\{n \mid n \text{ is an integer, } n > 301\} =$
“The set of n such that n is an integer and $n > 301$.”

(Actually, this notation is a slight simplification.)

The set is the *same* as that denoted by, for instance,
 $\{n \mid n \text{ is an integer, } n \geq 302\}.$

Some More Examples

- ◆ $\{\text{JAB}, \text{"JAB"}\}$ has 2 members: me, and a 3-char string.
- ◆ $\{3, \{4,5\}, 4, 6\}$ has 4 members, one of which is a set.
- ◆ $\{3, \{5,4\}, 4, 6\}$ is that same set.
- ◆ $\{\{4,5\}\}$ has 1 member, which is a set.
 $\{4,5\}$ has 2 members, both numbers.
- ◆ $\{\emptyset\}$ is a singleton set. Its only member is the empty set.
- ◆ $\{\{\emptyset\}\}$ is a different singleton set.