

# Fundamentals/ICY: Databases

## 2013/14

### ***WEEK 7 – Monday***

*John Barnden*

*Professor of Artificial Intelligence*

*School of Computer Science*

*University of Birmingham, UK*



**Reminder of Friday**



# Entity Supertypes and Subtypes

- ◆ Generalization (or: specialization) hierarchy
  - A group of relationships each of which is between a higher-level “supertype” entity (e.g. EMPLOYEE) and a lower-level “subtype” entity (e.g., PROFESSOR)
- ◆ Supertype
  - Contains attributes shared by all its subtypes
- ◆ Subtype
  - Contains special attributes: ones that not all sister subtypes have.
- ◆ Primary key of a subtype = that of the supertype (normally)



# Disjoint (or: Non-Overlapping) Subtypes

- ◆ Each entity in the supertype can appear in **at most** one of the subtypes
- ◆ Overlapping = a given entity can be in more than one subtype.

## Exhaustive Subtypes

- ◆ Each entity in the supertype must appear in **at least** one of the subtypes
- ◆ Other terminology:
  - exhaustiveness = total completeness (!!)  
= mandatoriness [of being in some subtype]
  - non-exhaustiveness = partial completeness (!!)  
= optionality [of being in som subtype]



**New**

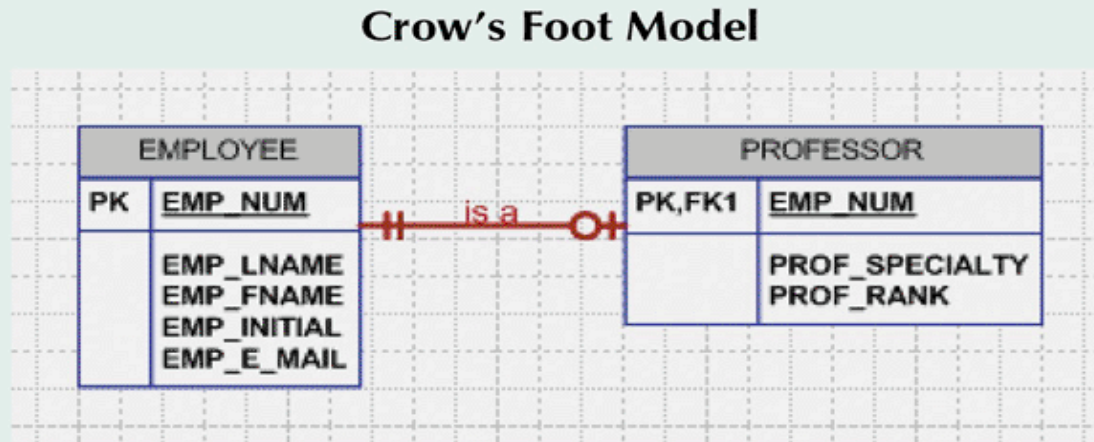
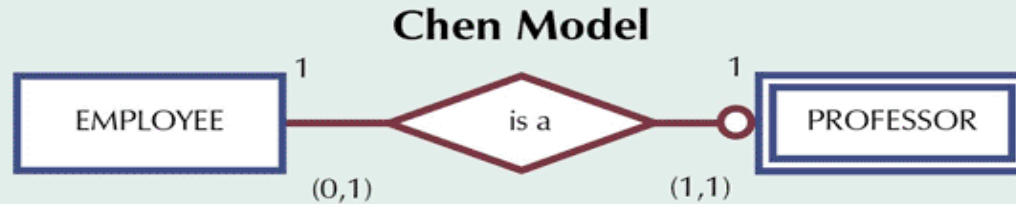


# Why Consider Supertypes and Subtypes?

- ◆ We would not just want to have a table for the supertype and none for the subtypes, because of the resulting poor table structure.
- ◆ And we would not *just* want to use separate entity types for the subtypes, because of
  - Attributes (etc.) in common and relationships in common
  - Need an extra entity type anyway when the subtypes are not exhaustive.

# Actual Realization of Subtyping

FIGURE 4.37 A SUPERTYPE/SUBTYPE RELATIONSHIP IN AN ERD



A supertype has a *1:1 relationship* with each subtype.

*Mandatory* in the *sub-to-super* direction.

*Optional* in the *super-to-sub* direction, *even in* exhaustive case.



# The EMPLOYEE/PILOT Supertype/Subtype Relationship

**FIGURE 4.29 THE EMPLOYEE/PILOT SUPERTYPE/SUBTYPE RELATIONSHIP**

**Table name: EMPLOYEE (the supertype)**

**Database name: Ch04\_AirCo**

		EMP_NUM	EMP_LNAME	EMP_HIRE_DATE
►	+	100	Kolmycz	15-Mar-88
	+	101	Lewis	25-Apr-89
	+	102	Vandam	20-Dec-93
	+	103	Jones	28-Aug-03
	+	104	Lange	20-Oct-97
	+	105	Williams	08-Nov-97
	+	106	Duzak	05-Jan-04
	+	107	Diante	02-Jul-97
	+	108	Wiesenbach	18-Nov-95
	+	109	Travis	14-Apr-01
	+	110	Genkazi	01-Dec-03

**Table name: PILOT (the subtype)**

		EMP_NUM	PIL_LICENSE	PIL_RATINGS	PIL_MED_TYPE
►	+	101	ATP	SEL/MEL/Instr/CFII	1
	+	104	ATP	SEL/MEL/Instr	1
	+	105	COM	SEL/MEL/Instr/CFI	2
	+	106	COM	SEL/MEL/Instr	2
	+	109	COM	SEL/MEL/SES/Instr/CFII	1



# But Special ERD Notation is Desirable

- ◆ To clearly bring out the structure in a conceptual ERD.
- ◆ Can refine later to get rid of the special notation and instead use the basic ERD facilities as in the “Actual Realization” slide above.
- ◆ Those basic facilities don’t clearly show that there is a specialization/generalization at all, and also don’t show (non-)disjointness or (non-)exhaustiveness.



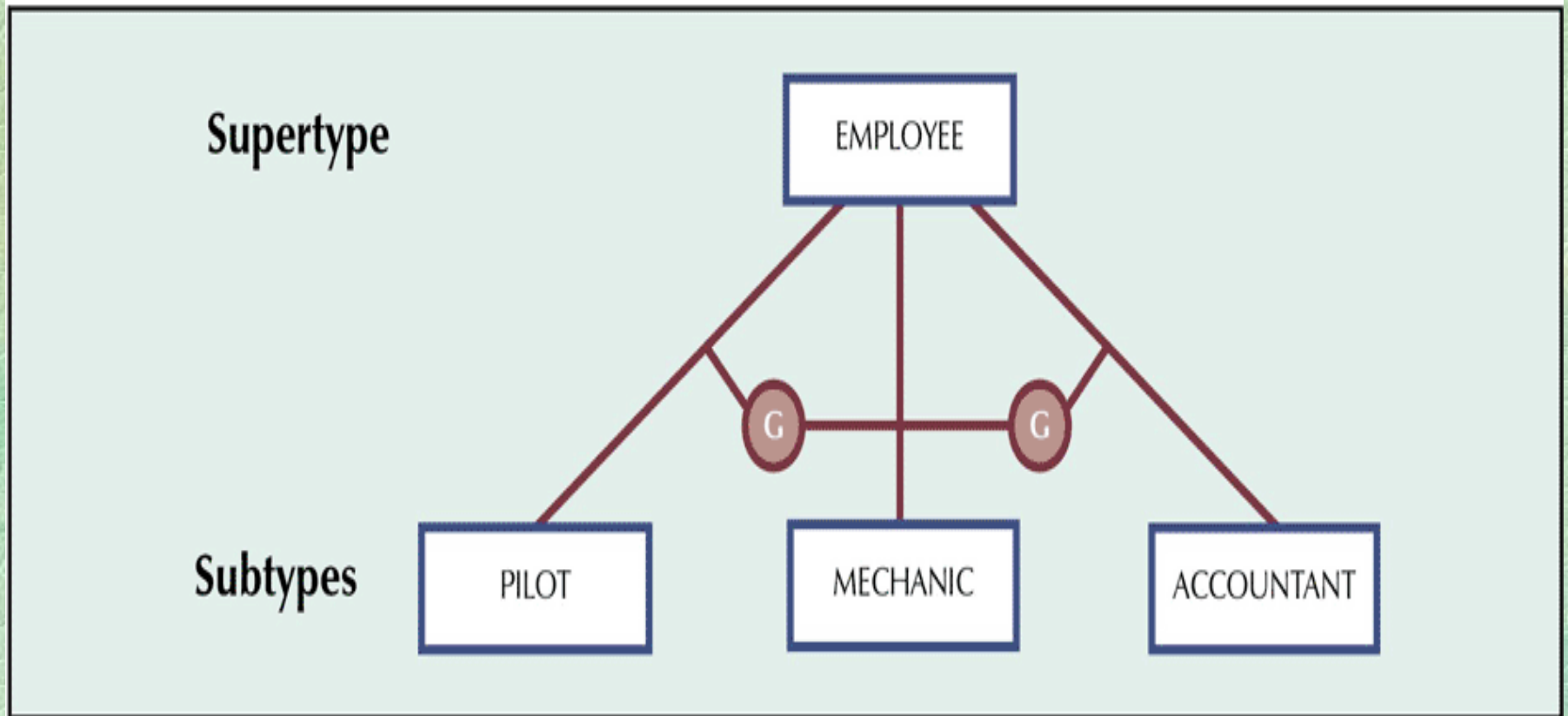
**Special Notation in following slides  
is from a previous textbook.**

**See other, *better*, notations in  
current textbook and in  
Additional Notes**



# Disjoint Subtypes

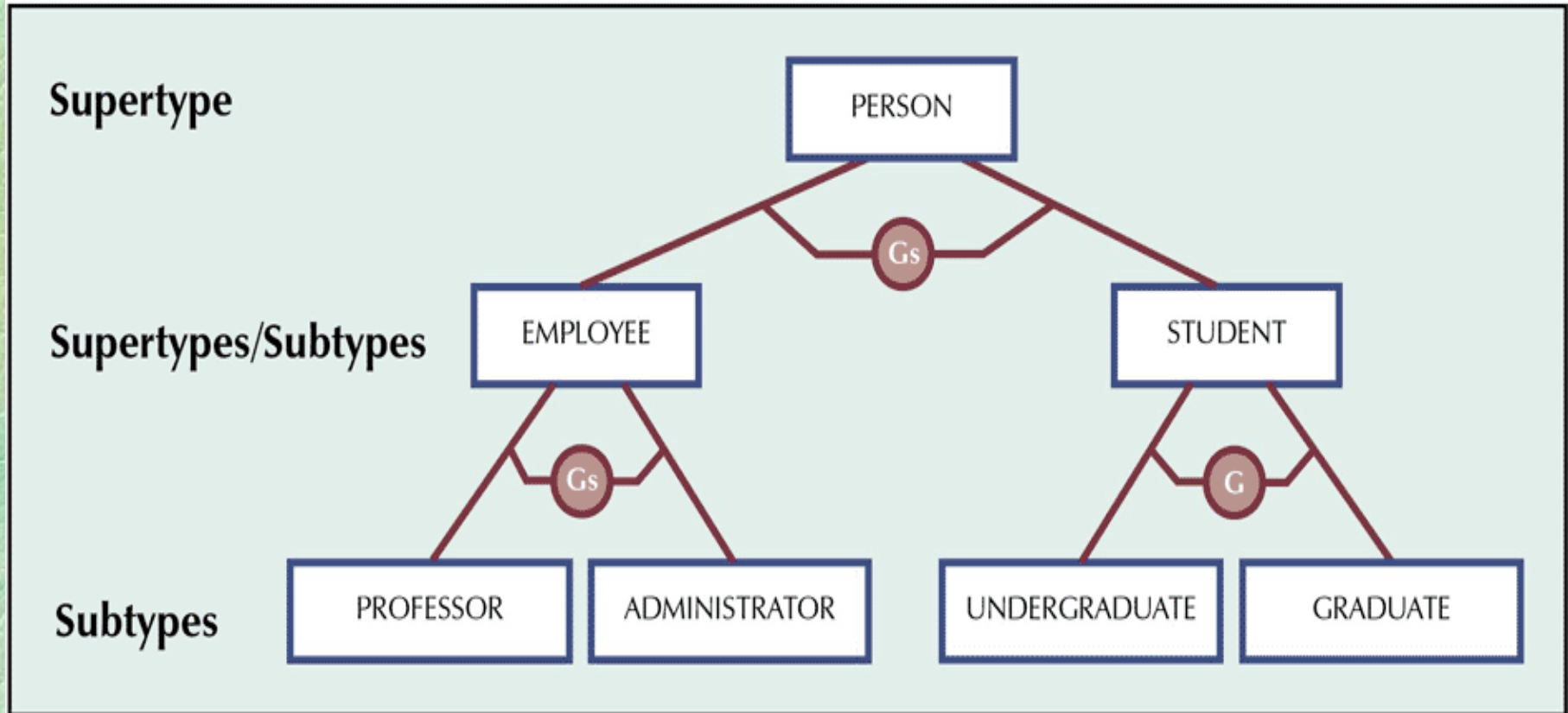
FIGURE 4.28 A GENERALIZATION HIERARCHY





# A Generalization Hierarchy with Overlapping Subtypes

FIGURE 4.30 A GENERALIZATION HIERARCHY WITH OVERLAPPING SUBTYPES



*G = disjoint, GS = overlapping.*

*NB: No representation of (non-)exhaustiveness.*



# **More-Than-Binary Relationships in ERM<sub>s</sub> and ERD<sub>s</sub>**



# Relationship Degree

## ◆ **Binary** relationship [*my* definition]

- **Two** entities (“entity occurrences”) are associated by each instance of the relationship, as in all previous examples in lecture slides.

## ◆ **Ternary** relationship [*my* definition]

- **Three** entities are associated by each relationship instance.

## ◆ Etc.

## ◆ NB: the entities associated with each other need not be of different *types*.

- Indeed, could have an entity (occurrence) associated with itself: e.g. an “employs” relationship where someone can employ him/herself. Still **binary**.



# Relationship Degree: Terminology Problems

- ◆ The standard terminology & definitions relating to relationship degree (see the textbook) are *mathematically anomalous*.
- ◆ I believe the degree of a relationship should be the number of **entities** (“entity **occurrences**”) that are associated by each instance of the relationship, no matter how many entity types are involved.
- ◆ But the definitions used in the books count the number of entity **types** (“entities” in the abbreviated language used), departing from normal mathematical practice.



# Terminology Problems contd.

- ◆ A “**unary**” relationship is standardly defined as being a relationship where the entity occurrences are all within the same entity type (“entity”).
  - E.g., a “manages” relationship between employees.
- ◆ A better name is “recursive” (also used in the books).
  - Above example is **binary** recursive, under **my** definition.
- ◆ The books define a **binary** relationship as one relating *two different types*, even when there are more than two entity occurrences per relationship instance.
- ◆ Similarly **ternary** and **three different** types.



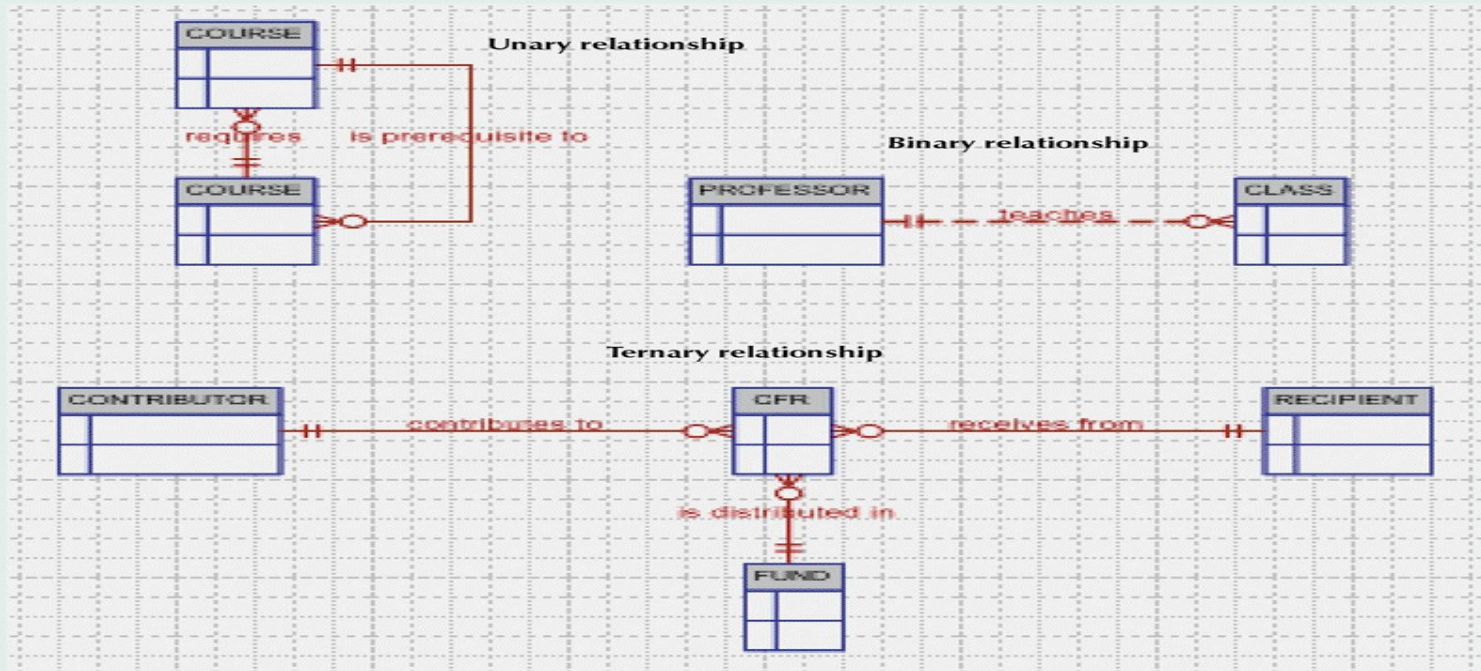
# Different Degrees (using Book Terminology)

FIGURE 4.16 THREE TYPES OF RELATIONSHIPS

## Chen Models



## Crow's Foot Models



*The "unary" case is badly named. It's really a type of binary.  
The word "recursive" [later] is better.*



# Caution about Next Slide

- ◆ There is an error in the following slide.
- ◆ The FUND table should just be descriptions of different funds, such as Heart fund, a Cancer fund, and so on. FUND\_ID should be the PK.
- ◆ It shouldn't have the CONTRIB\_ID and FUND\_AMOUNT columns.



# Tables for a Ternary Relationship

FIGURE 4.17 THE IMPLEMENTATION OF A TERNARY RELATIONSHIP

Database name: Ch04\_MedCo

Table name: CONTRIBUTOR

	CONTRIB_ID	CONTRIB_LNAME
►	C1	Brown
	C2	Iglesas
	C3	Smith

Table name: FUND

	FUND_ID	FUND_NAME	CONTRIB_ID	FUND_AMOUNT
►	F1	Heart	C1	\$50,000.00
	F1	Heart	C2	\$10,000.00
	F2	Cancer	C1	\$10,000.00
	F2	Cancer	C2	\$5,000.00
	F2	Cancer	C3	\$10,000.00

Table name: RECIPIENT

	REC_ID	REC_TYPE
►	R1	Rogers
	R2	Chen
	R3	Oshanski

Table name: CFR

	FUND_ID	CON_ID	REC_ID	CFR_AMOUNT
►	F1	C1	R2	\$30,000.00
	F1	C1	R3	\$20,000.00
	F1	C2	R2	\$10,000.00
	F2	C1	R1	\$10,000.00
	F2	C2	R1	\$5,000.00

*CFR is just like the bridging entity types you've seen before, but has 3 links to other types instead of 2*



# **Recursive Relationships in ERM and ERDs**



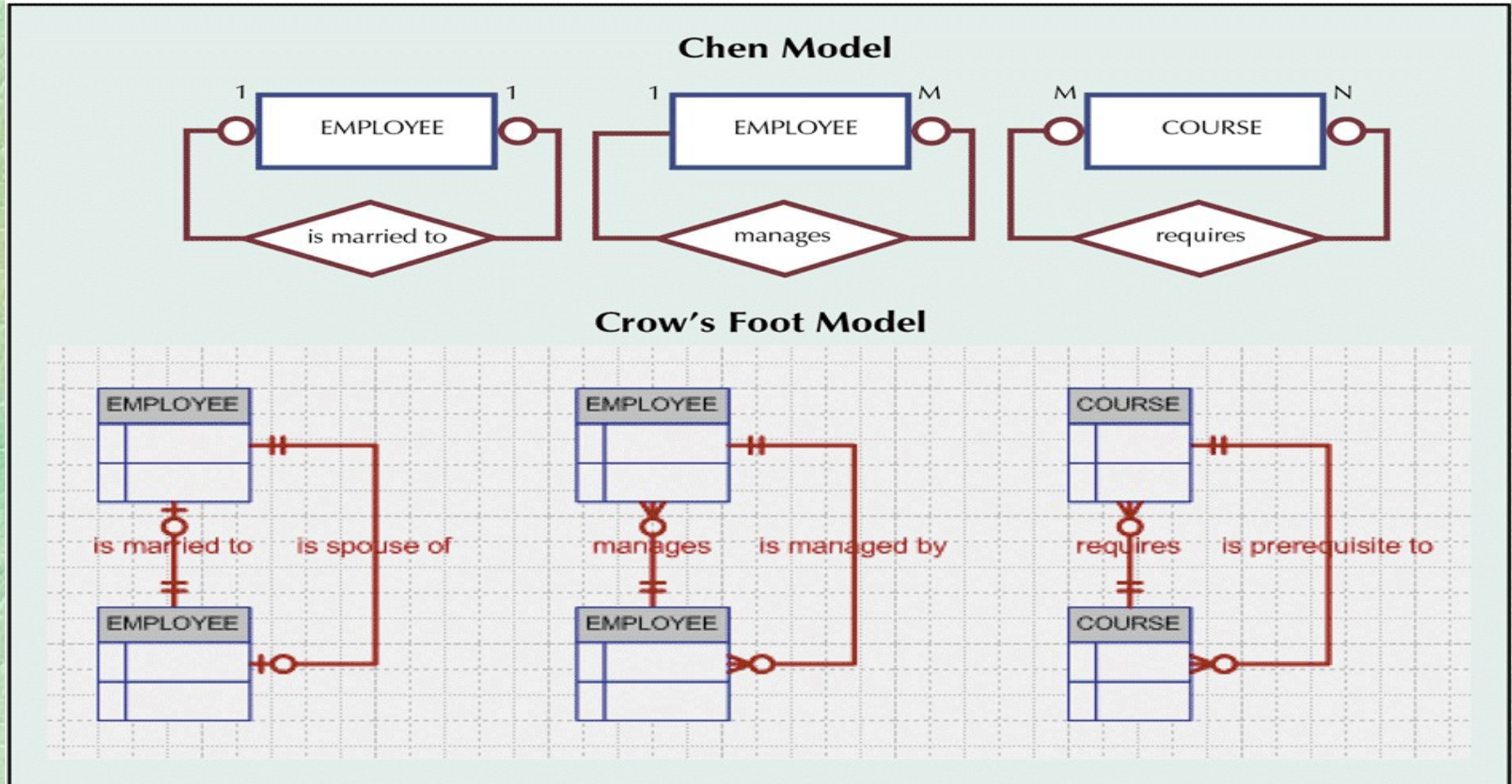
# Recursive Relationships

- ◆ A *recursive* relationship links entities of the same type.
  - *E.g.:* marriage, management, parthood, ...
- ◆ Can have partial recursion: just *some* of the entity types involved in a relationship (that is ternary or of higher degree in my sense) could be the same.



# An ER Representation of Recursive Relationships

FIGURE 4.18 AN ER REPRESENTATION OF RECURSIVE RELATIONSHIPS



I don't know why the Crow's Ft model uses two links in each case



# Recursive Relationships: Symmetry

- ◆ A relationship  $R$  between entity types  $E, F$  (possibly the same) is *symmetric* iff:

if  $eRf$  then  $fRe$  (i.e., IF  $R$  relates entity  $e$  of type  $E$  to entity  $f$  of type  $F$ , then it must ALSO relate  $f$  to  $e$ .)

*E.g.: marriage, being-sibling-of.*

- ◆ *Recursive relationships cause major redundancy problems when ALSO symmetric.*
- ◆ *Symmetry only makes sense in the 1:1 and M:N cases.*
- ◆ ((Can generalize the points to partly-recursive cases.))



(necessarily non-symmetric) **1:M** recursive:  
“EMPLOYEE Manages EMPLOYEE”

**FIGURE 4.23** IMPLEMENTATION OF THE 1:M “EMPLOYEE MANAGES EMPLOYEE” RECURSIVE RELATIONSHIP

Table name: EMPLOYEE\_V2

Database name: Ch04\_PartCo

	EMP_CODE	EMP_LNAME	EMP_MANAGER
►	101	Waddell	102
	102	Orincona	
	103	Jones	102
	104	Reballoh	102
	105	Robertson	102
	106	Deltona	102

*Just a standard 1:M implementation except linking a table to itself.*  
*No redundancy problem.*



# non-symmetric M:N recursive: “PART Contains PART”

FIGURE 4.21 IMPLEMENTATION OF THE M:N RECURSIVE “PART CONTAINS PART” RELATIONSHIP

Table name: COMPONENT

Database name: Ch04\_PartCo

	COMP_CODE	PART_CODE	COMP_PARTS_NEEDED
▶	C-130	AA21-6	4
	C-130	AB-121	2
	C-130	E129	1
	C-131A2	E129	1
	C-130	X10	4
	C-131A2	X10	1
	C-130	X34AW	2
	C-131A2	X34AW	2

Table name: PART

	PART_CODE	PART_DESCRIPTION	PART_IN_STOCK
▶	AA21-6	2.5 cm. washer, 1.0 mm. rim	432
	AB-121	Cotter pin, copper	1,034
	C-130	Rotor assembly	36
	E129	2.5 cm. steel shank	128
	X10	10.25 cm. rotor blade	345
	X34AW	2.5 cm. hex nut	879

*The COMPONENT entity type is just a bridging type, linking PART to itself. NB: its first two columns both refer to PART's PK but must be differently named. No redundancy problem.*



symmetric (1:1) recursive relationship:  
“EMPLOYEE Married to EMPLOYEE”

*Suppose you tried the following:*

FIGURE 4.19 THE 1:1 RECURSIVE RELATIONSHIP “EMPLOYEE IS MARRIED TO EMPLOYEE”

Table name: EMPLOYEE\_V1

Database name: Ch04\_PartCo

	EMP_NUM	EMP_LNAME	EMP_FNAME	EMP_SPOUSE
►	345	Ramirez	James	347
	346	Jones	Anne	349
	347	Ramirez	Louise	345
	348	Delaney	Robert	
	349	Shapiro	Anton	346

*Redundancy problem!!*