

Problem 1:

Pseudo code:

Greedy Algorithm (A):

denominations = [0.25, 0.1, 0.05, 0.01]

while $i < \text{length of denominations}$ while $A \geq \text{denominations}[i]$ $A -= \text{denominations}[i]$ $i += 1$

We can be greedy in assuming we can use the highest denomination and continuously subtract it from the value 'A'.

If the largest denomination doesn't fit in the value, switch to the next largest, until all denominations have been used. Time Complexity is $O(A)$.

Problem 2:

Pseudo code:

Huffman (Symbol Frequency Object as SFO for short):

heap = Heapify(SFO)

while heap length > 1

left = pop heap (smallest frequency)

right = pop heap (smallest frequency)

left.code = 0; right.code = 1

create node z = (left.frequency + right.frequency, left = left, right = right, symbol = None)

heap.push(z)

return heap

Prefix Encode (message, heap):

codes = Store Codes (heap)

encoded Message = None

for letter in message:

encodedMessage += codes[letter]

return encodedMessage

Store Codes (node, value)

newValue = value + node.code

if (node.left)

Store Codes (node.left, newValue)

if (node.right)

Store Codes (node.right, newValue)

if not node.left and not node.right

codes[~~new~~Symbol] = newValue

Prefix Decode (message, heap)

decoded Message = None

for ~~number~~^{number} in message

if number == 0

heap = heap.left

else

heap = heap.right

if heap.left and heap.right is None

decodedMessage += heap.symbol

return decodedMessage

Huffman Algorithm runs in $O(n \log n)$ time. The storing codes algorithm ~~also~~ runs in ~~$O(n)$~~ $O(n)$ time, in which Prefix Encode will run in $O(|n| + |v|)$ time, where v is every letter in the message. Prefix Decode will run in $O(|e|)$ time, where it will back track each node to ~~decode~~ decode, where e is the length of encoded message. ~~$|e|$~~

Each letter in the message will get its own unique code depending on its frequency. This helps compress the message, as 8 bits per letter isn't necessary.