My approximation results seem to be better then optimal. I believe this is caused because, when I traverse through my nodes in the Preorder Traversal, it isn't adding the weight between one of the last children into a branch to the next branch's parent. The pathony however is correct

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The brute force algorithm runs in O(n!) as it takes all possible permutations form the root of TSP.

TSP (G, r):

min weight = inf
all permutations = permutations of all vertices in G except r

for permutation in all permutations:

weight = 0

for in permutation!

weight = euclidean distance of i coords to root

if weight < minuseight:

minuseight = weight

The approximation algorithm is much faster with $O(v^2)$ time. First we find Prim's algorithm's MST, then traverse adding weights of all children in a DFS ctyle.

Prim (G,r):

PQ = priority queue

MST = Graph with an vertex

colorcil = white for all u & G.v

Visit (r, color, G, PQ)

while PQ is not empty

weight, = PQ. extract min

if color (w) = white:

Visit (u, color, a, PQ)

MSt add edge

return MS+

Visit (u, color, a, PQ).

for each v, weight in Gradyacentlu)

if color[u] = white:

PQ. insert (weight, u=>u)

color(u) = black

Approx TSP Tour (G, x):

root = Prime (G, F)

H = Traversal (next)

return H

Traversal (node):

if node is entry return

for each child in node children

R=R+ Traversal (child)

return R.