


EECS402 Lecture 20

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Savitch, Ch. 13
Recursion

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Printing a Singly-Linked List

```
class ListNodeClass
{
public:
    int val;
    ListNodeClass *next;
};
class LinkedListClass
{
public:
    ListNodeClass *head;    //data member

    LinkedListClass():head(0) {} //ctor
    void insertHead(int inVal);
    void printFwd()
    {
        ListNodeClass *tmp = head;
        cout << "Fwd List!" << endl;
        while (tmp != 0)
        {
            cout << "Val: " << tmp->val << endl;
            tmp = tmp->next;
        }
    }
};
```

From main():

```
LinkedListClass ll;
ll.insertHead(1);
ll.insertHead(5);
ll.insertHead(2);
ll.insertHead(9);
ll.printFwd();
```

Output:

Fwd List!

Val: 9


Val: 2

Val: 5

Val: 1

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- How would you write a printBwd() member function?

```
void LinkedListClass::printBwd()
{
    int i, j;
    int count = 0;
    ListNodeClass *temp = head;
    cout << "Bwd print!" << endl;
    while (temp != 0)
    {
        count++;
        temp = temp->next;
    }
    for (i = 0; i < count; i++)
    {
        temp = head;
        for (j = 0; j < (count - 1) - i; j++)
            temp = temp->next;
        cout << "Val: " << temp->val << endl;
    }
}
```

This works, but it is very inefficient, and very error-prone.

Leaving the "-1" out of the loop results in a seg fault, but it is not intuitive that it belongs there, etc...

- A recursive function is a function that calls itself either:
 - Directly, in the function body
 - Indirectly, by calling another function, which, in turn calls the recursive function
- Many iterative algorithms can be done using recursion
- Note: Just because something can be done recursively does not mean it is the best way to do it
- *Recursion is best used when a problem can be continually broken down into problems that are easier to solve than the previous*
- A recursive function **MUST** have a terminal condition, or else infinite recursion will occur
 - Always provide a condition that will end the recursive calls
 - Condition should be checked **BEFORE** making the recursive call
- Think of the call stack that is created when thinking of recursion

- First, a simple, incredibly common example - factorial
 - $n! = n * (n - 1) * (n - 2) * \dots * (n - (n-1))$
 - $4! = 4 * (4 - 1) * (4 - 2) * (4 - 3) = 4 * 3 * 2 * 1 = 24$
- Notice that $4! = 4 * 3! = 4 * 3 * 2! = 4 * 3 * 2 * 1!$
- This is a pattern that suggest recursion as a good solution

Iterative solution (NOT recursive):

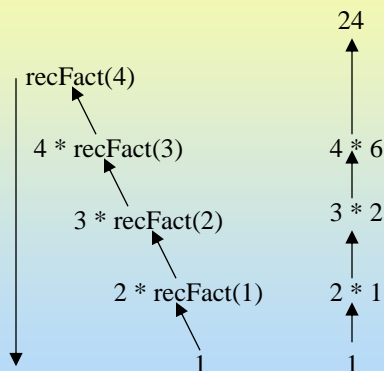
```
int factorial(int n)
{
    int i;
    int fact = 1;
    for (i = n; i >= 1; i--)
    {
        fact *= i;
    }
    return fact;
}
```

- Here is a recursive solution for the factorial problem:

```
int recFact(int n)
{
    if (n <= 1)
    {
        //this is the base case
        //(or terminal case)
        return 1;
    }
    else
    {
        //Otherwise, return the
        //current val times the
        //factorial of the next
        //lowest integer
        return n * recFact(n - 1);
    }
}
```

Function calls
(trace downwards)

Return values
(trace upwards)





Back To the Backwards List

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- As you may have guessed, we can use recursion to print out a list backwards
 - While you haven't reached the end of the list, call the function with the "next" node
 - When you are done with the if statement, print the value
- Need to pass in a ListNodeClass pointer
- Start with "head" but we don't want the user to worry about implementation
 - Implement two functions
 - 1. Public function available in the interface, with no params This function will just call function #2 with head as the param
 - 2. Private function to do recursive call, taking in a ListNodeClass *

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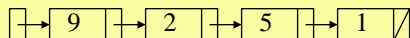
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Recursively Printing a List Backwards

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```

void recPrintBwd(ListNodeClass *n)
{
    bool endOfList = false;

    if (n != 0)
        recPrintBwd(n->next); //recursion
    else
        endOfList = true; //base case

    if (!endOfList)
        cout << "Val: " << n->val << endl;
}

void printBwd2()
{
    cout << "Bwd list!" << endl;
    recPrintBwd(head);
}
  
```

On the way down

```

For node 9:
endOfList=false
recPrintBwd(node 2)

For node 2:
endOfList=false
recPrintBwd(node 5)

For node 5:
endOfList=false
recPrintBwd(node 1)

For node 1:
endOfList=false
recPrintBwd(node NULL)

For node NULL:
endOfList=true
  
```

On the way up

```

Node NULL:
no print - returns
to node 1

Node 1:
prints 1 - returns
to node 5

Node 5:
prints 5 - returns
to node 2

Node 2:
prints 2 - returns
to node 9

Node 9:
prints 9 - returns
to calling
function
  
```

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- Recursion is never "necessary"
 - Anything that can be done recursively, can be done iteratively
 - Recursive solution may seem more logical
 - For example, printing the list - the iterative solution given is very awkward, and does not model the human way of doing the problem, if given a list
- The recursive solution did not use any nested loops, while the iterative solution did
- However, the recursive solution made many more function calls, which adds a lot of overhead
- Recursion is NOT an efficiency tool - use it only when it helps the logical flow of your program

- Avoid the following pitfalls demonstrated below:

```
int badRecursion1(int num)
{
    return num + badRecursion(num - 2);
}
```

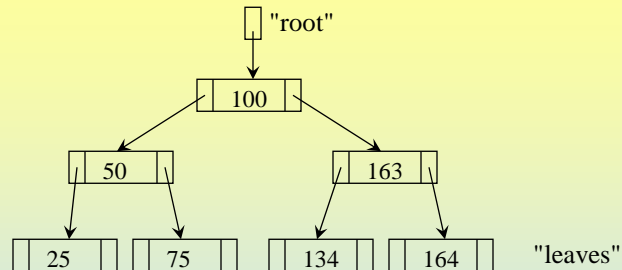
- Avoid the following pitfalls demonstrated below:

```
int badRecursion2(int num)
{
    if (num == 0)
    {
        return 1;    //base case
    }
    else
    {
        return num + badRecursion2(num - 2);
    }
}
```

- A binary tree is yet another data structure
- Nodes of a binary tree consist of a value and 2 pointers

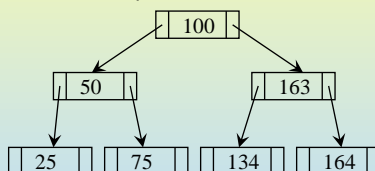
```
class BinaryTreeNodeClass
{
    int value;
    BinaryTreeNodeClass *left;
    BinaryTreeNodeClass *right;
};

BinaryTreeNodeClass *root;
```



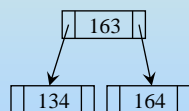
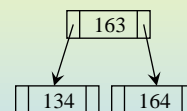
- This specific tree has a couple special properties
 - It is a complete binary tree
 - All levels of the tree are completely filled
 - It is a binary search tree
 - For every node:
 - Any node to the left has a lower value
 - Any node to the right has a higher value

- **Every** node of a binary search tree is the root of a subtree, which is also a binary search tree
- Unlike linked lists, binary search trees can be searched in an efficient way
 - Possibly equivalent to a binary search of a sorted array
 - Example: Search for 134 in the shown tree



Root node contains 100
Since $134 > 100$, it will be in the right subtree. No need to search left subtree, which disregards half of the tree elements!

Remaining Subtree:



Root node contains 163
Since $134 < 163$, it will be in the left subtree. Able to disregard half of the nodes that remain

Remaining Subtree:



Value found!

```
bool inBST(BSTNodeClass *root, int val)
{
    bool found = false;
    BSTNodeClass *tmp = root;

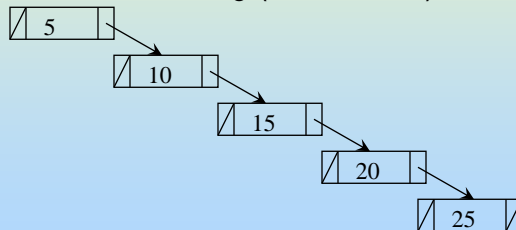
    while (tmp != 0 && !found)
    {
        if (tmp->value > val)
        {
            tmp = tmp->left;
        }
        else if (tmp->value < val)
        {
            tmp = tmp->right;
        }
        else
        {
            found = true;
        }
    }
    return found;
}
```

```
bool inBSTRec(BSTNodeClass *root, int val)
{
    bool found;

    if (root == 0)
    {
        found = false; //A base case
    }
    else if (root->value < val)
    {
        found = inBSTRec(root->right, val); //Recursive call
    }
    else if (root->value > val)
    {
        found = inBSTRec(root->left, val); //Recursive call
    }
    else if (root->value == val)
    {
        found = true; //Another base case
    }
    return found;
}
```

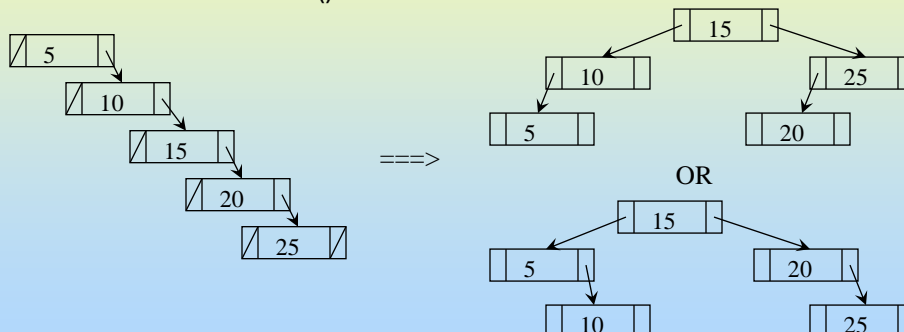

- Search and insert are efficient for a BST.
 - True with assumption that tree is "balanced"
- Consider the following operations


```
insertIntoBST(root, 5);
insertIntoBST(root, 10);
insertIntoBST(root, 15);
insertIntoBST(root, 20);
insertIntoBST(root, 25);
```
- Results in the following (unbalanced) tree



Note: Searching this BST results in a **linear search** instead of the binary search that was the primary advantage of a binary search tree!

- An unbalanced BST, as previously shown, has a search time of $O(n)$.
- A balanced BST has a search time of $O(\log n)$.
- How can you be sure you keep your BST balanced?
- Write a `balanceTree()` function!





BST Efficiency Discussions

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- Discuss the efficiency of the `balanceTree()` function
 - When should the function be called?
- Discuss the efficiency of deleting a node from a BST
 - Are there way to improve the efficiency of the delete operation?

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