

Not having a good time with the `mathpartir` package here.

Exercise 4.1

Calls to `g` declare a reference `counter`, set its contents, then return its contents. The information between calls is lost because `counter` is different on each call. That is, the reference data structure represents a different location. The first program was made such that the environment of the closure had access to `counter`, and thus was able to reference that variable.

Exercise 4.2

$$\frac{(\text{value-of } exp_1 \ \rho \ \sigma_0) = (val_1, \sigma_1)}{(\text{value-of } (\text{zero?-exp } exp_1) \ \rho \ \sigma_0) = \begin{array}{ll} ((\text{bool-val } \#t), \sigma_1) & \text{if } [val_1] = 0 \\ ((\text{bool-val } \#f), \sigma_1) & \text{if } [val_1] \neq 0 \end{array}}$$

Exercise 4.3

$$\frac{\begin{array}{l} (\text{value-of } exp_1 \ \rho \ \sigma_0) = (val_1, \sigma_1) \\ (\text{value-of } exp_2 \ \rho \ \sigma_1) = (val_2, \sigma_2) \end{array}}{(\text{value-of } (\text{call-exp } exp_1 \ exp_2) \ \rho \ \sigma_0) = (\text{apply-procedure } (\text{expval} \rightarrow \text{proc } val_1) \ val_2 \ \sigma_2)}$$

$$\frac{val_1 = (\text{procedure } var \ body \ \rho)}{(\text{apply-procedure } val_1 \ val_2 \ \rho_0) = (\text{value-of } body \ [var = val_2] \rho \ \sigma_0)}$$

Exercise 4.4

$$\begin{array}{c}
(\text{value-of } exp_1 \ \rho \ \sigma_0) = (val_1, \sigma_1) \\
\vdots \\
(\text{value-of } exp_n \ \rho \ \sigma_{n-1}) = (val_n, \sigma_n) \\
\hline
(\text{value-of } (\text{begin } exp_1 \dots exp_n) \ \rho \ \sigma_0) = (val_n, \sigma_n)
\end{array}$$

Exercise 4.5

$$\begin{array}{c}
(\text{value-of } exp_1 \ \rho \ \sigma_0) = (val_1, \sigma_1) \\
\vdots \\
(\text{value-of } exp_n \ \rho \ \sigma_{n-1}) = (val_n, \sigma_n) \\
\hline
(\text{value-of } (\text{list-exp } exp_1 \dots exp_n) \ \rho \ \sigma_0) \\
= ((\text{pair-val } val_1 \ (\dots (\text{pair-val } val_n \ (\text{emptylist-val})) \dots)), \sigma_n)
\end{array}$$

Exercise 4.6

$$\begin{array}{c}
(\text{value-of } exp_1 \ \rho \ \sigma_0) = (l, \sigma_1) \\
(\text{value-of } exp_2 \ \rho \ \sigma_1) = (val, \sigma_2) \\
\hline
(\text{value-of } (\text{setref-exp } exp_1 \ exp_2) \ \rho \ \sigma_0) = (val, [l = val]\sigma_2)
\end{array}$$

Exercise 4.7

$$\begin{array}{c}
(\text{value-of } exp_1 \ \rho \ \sigma_0) = (l, \sigma_1) \\
(\text{value-of } exp_2 \ \rho \ \sigma_1) = (val, \sigma_2) \\
\hline
(\text{value-of } (\text{setref-exp } exp_1 \ exp_2) \ \rho \ \sigma_0) = (\sigma_0(l), [l = val]\sigma_2)
\end{array}$$

Exercise 4.8

`newref`, `deref`, and `setref!` take linear time.

Exercise 4.9

`newref` uses `new-store-longer-by-one` which takes linear time because I could not find a built-in procedure for copying vectors. `deref` takes as much time as `vector-ref`. `setref!` takes as much time as `vector-length`.

```
(define empty-store
  (lambda ()
    (make-vector 0)))

;; initialize-store! : () -> Sto
;; usage: (initialize-store!) sets the-store to the empty-store
(define initialize-store!
  (lambda ()
    (set! the-store (empty-store))))

;; reference? : SchemeVal -> Bool
(define reference?
  (lambda (v)
    (integer? v)))

;; new-store-longer-by-one : Sto -> Sto
(define new-store-longer-by-one
  (lambda (store)
    (let ((new-store (make-vector (+ 1 (vector-length store)))))
      (letrec ((inner
                  (lambda (current-index stop)
                     (if (= current-index stop)
                         new-store
                         (begin (vector-set!
                                new-store
                                current-index
```

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                                (vector-ref store current-index))
                                (inner (+ current-index 1) stop))))))
    (inner 0 (vector-length store))))))

;; newref : ExpVal -> Ref
(define newref
  (lambda (val)
    (let* ((next-ref (vector-length the-store))
           (new-store (new-store-longer-by-one the-store)))
      (vector-set! new-store next-ref val)
      (set! the-store new-store)
      (when (instrument-newref)
        (eopl:printf
         "newref: allocating location ~s with initial contents ~s~
         next-ref val))
      next-ref)))

;; deref : Ref -> ExpVal
(define deref
  (lambda (ref)
    (vector-ref the-store ref)))

;; setref! : Ref * ExpVal -> Unspecified
(define setref!
  (lambda (ref val)
    (if (> (vector-length the-store) ref)
        (vector-set! the-store ref val)
        (report-invalid-reference ref the-store))))

```

Exercise 4.10

```

;; the-grammar
(expression
  ("begin" expression (arbno ";" expression) "end")

```

```

begin-exp)

;; value-of
(begin-exp (exp1 rest-exps)
  (letrec ((begin-inner
    (lambda (exps final-val)
      (if (null? exps)
          final-val
          (begin-inner (cdr exps)
                        (value-of (car exps)
                                  env))))))
    (begin-inner rest-exps (value-of exp1 env))))

```

Exercise 4.11

```

;; the-grammar
(expression
  ("list" "(" (separated-list expression ",") ")")
  list-exp)

;; value-of
(list-exp (exps)
  (letrec ((list-inner
    (lambda (exps)
      (if (null? exps)
          (emptylist-val)
          (pair-val (value-of (car exps) env)
                     (list-inner (cdr exps))))))
    (list-inner exps)))

(define-datatype expval expval?
  (num-val
    (value number?))
  (bool-val

```

```

    (boolean boolean?))
(proc-val
  (proc proc?))
(ref-val
  (ref reference?))
(pair-val
  (val1 expval?)
  (val2 expval?))
(emptylist-val))

```

Exercise 4.12

The fragment of the interpreter given in the book does not have enough context. The meaning of `apply-store` is unclear and `deref` takes only one argument which makes the program look like the store is still a global variable.

```

;; value-of-program : Program -> ExpVal
(define value-of-program
  (lambda (pgm)
    (cases program pgm
      (a-program (exp1)
        (cases answer (value-of exp1
                                (init-env)
                                (init-store))
          (an-answer (val store)
                     val)))))))

;; value-of : Exp * Env -> Answer
(define value-of
  (lambda (exp env store)
    (cases expression exp
      (const-exp (num) (an-answer (num-val num) store))
      (var-exp (var)
        (an-answer (apply-env env var)

```

```

                                store))
(diff-exp (exp1 exp2)
  (cases answer (value-of exp1 env store)
    (an-answer (v1 store1)
      (cases answer (value-of exp2
                            env
                            store1)
        (an-answer (v2 store2)
          (an-answer
            (num-val
              (- (expval->num v1)
                (expval->num v2))))
            store2))))))
(zero?-exp (exp1)
  (cases answer (value-of exp1 env store)
    (an-answer (val1 store1)
      (an-answer
        (if (zero? (expval->num val1))
          (bool-val #t)
          (bool-val #f))
        store1))))
(if-exp (exp1 exp2 exp3)
  (cases answer (value-of exp1 env store)
    (an-answer (val1 store1)
      (if (expval->bool val1)
        (value-of exp2 env store1)
        (value-of exp3 env store1))))))
(let-exp (var exp1 body)
  (cases answer (value-of exp1 env store)
    (an-answer (val1 store1)
      (value-of body
        (extend-env var val1 env)

```

```

                                store1))))
(proc-exp (var body)
  (an-answer (proc-val (procedure var body env))
    store))
(call-exp (rator rand)
  (cases answer (value-of rator env store)
    (an-answer (proc-val store1)
      (cases answer (value-of rand
                                env
                                store1)
        (an-answer (arg store2)
          (apply-procedure
            (expval->proc proc-val)
            arg
            store2))))))
(letrec-exp (p-names b-vars p-bodies letrec-body)
  (value-of letrec-body
    (extend-env-rec* p-names
      b-vars
      p-bodies
      env)
    store))
(begin-exp (exp1 exps)
  (letrec
    ((begin-inner
      (lambda (e1 es store)
        (cases answer (value-of e1 env store)
          (an-answer (v1 store1)
            (if (null? es)
              (an-answer v1 store1)
              (begin-inner
                (car es)

```



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(cdr es)
store1))))))
(begin-inner exp1 exps store))
(newref-exp (exp1)
  (cases answer (value-of exp1 env store)
    (an-answer (val1 store1)
      (let ((newref-pair
              (newref val1 store1)))
        (let ((ref (car newref-pair))
              (store2 (cdr newref-pair)))
          (an-answer (ref-val ref)
                     store2))))))
(deref-exp (exp1)
  (cases answer (value-of exp1 env store)
    (an-answer (val1 store1)
      (let ((ref1 (expval->ref val1)))
        (an-answer (deref store1 ref1)
                   store1))))))
(setref-exp (exp1 exp2)
  (cases answer (value-of exp1 env store)
    (an-answer (val1 store1)
      (cases answer (value-of exp2
                              env
                              store1)
        (an-answer (val2 store2)
          (an-answer
            (num-val 23)
            (setref
              (expval->ref val1)
              val2
              store2))))))))))

```

```

;; apply-procedure : Proc * ExpVal * Sto -> ExpVal

;; uninstrumented version
(define apply-procedure
  (lambda (proc1 arg store)
    (cases proc proc1
      (procedure (bvar body saved-env)
        (value-of body
                    (extend-env bvar arg saved-env)
                    store))))))

;; instrumented version
(define apply-procedure
  (lambda (proc1 arg store)
    (cases proc proc1
      (procedure (var body saved-env)
        (let ((r arg))
          (let ((new-env (extend-env var r saved-env)))
            (when (instrument-let)
              (begin
                (eopl:printf
                  "entering body of proc ~s with env =~%"
                  var)
                (pretty-print (env->list new-env))
                (eopl:printf "store =~%")
                (pretty-print
                  (store->readable
                    (get-store-as-list store))))
                (eopl:printf "~%"))
              (value-of body new-env store)))))))

;; empty-store : () -> Sto

```

```
(define empty-store
  (lambda () '()))

;; init-store : () -> Sto
(define init-store
  (lambda ()
    (empty-store)))

;; newref : ExpVal * Sto -> Cons(Ref,Sto)
(define newref
  (lambda (val store)
    (let ((next-ref (length store)))
      (when (instrument-newref)
        (eopl:printf
          "newref: allocating location ~s with initial contents ~s~n"
          next-ref val)))
      (cons next-ref (append store (list val))))))

;; deref : Ref * Sto -> ExpVal
(define deref
  (lambda (store ref)
    (list-ref store ref)))

;; setref : Ref * ExpVal * Sto -> Sto
(define setref
  (lambda (ref val store)
    (letrec ((setref-inner
      (lambda (r s)
        (cond ((null? s)
              (report-invalid-reference r store))
              ((zero? r)
               (cons val (cdr s)))
              (else
               (setref-inner (add1 r) (cdr s)))))))
      (setref-inner ref store))))
```

```

                                (else
                                (cons
                                (car s)
                                (setref-inner (- r 1) (cdr s)))))))))
    (setref-inner ref store)))

(define report-invalid-reference
  (lambda (ref the-store)
    (eopl:error 'setref
      "illegal reference ~s in store ~s"
      ref the-store)))

;; get-store-as-list : Sto -> Listof(List(Ref,Expval))
(define get-store-as-list
  (lambda (store)
    (letrec
      ((inner-loop
        (lambda (sto n)
          (if (null? sto)
              '()
              (cons
               (list n (car sto))
               (inner-loop (cdr sto) (+ n 1)))))))
      (inner-loop store 0))))

(define-datatype answer answer?
  (an-answer
   (val expval?)
   (store store?)))

;; store? : Schemeval -> Bool
(define store?

```

```
(lambda (scmval)
  ((list-of expval?) scmval)))
```