```
1.
void inorder(TREE T)
{
    if (T != NULL) {
        inorder(T->leftchild);
        printf("%c\n", T->label);
        inorder(T->rightchild);
    }
}
2.
BOOLEAN isoperator(char sym) {
    if (sym == '+' || sym == '-' || sym == '*' || sym == '/')
        return TRUE;
    else return FALSE;
}
void inorder(TREE T)
    if (T != NULL) {
        if (isoperator(T->label)) {
            putchar('(');
            inorder(T->leftchild);
            printf("%c", T->label);
            inorder(T->rightchild);
            putchar(')');
        }
        else {
            inorder(T->leftchild);
            printf("%c", T->label);
            inorder(T->rightchild);
        }
    }
}
3.
BOOLEAN isoperator(char sym) {
    if (sym == '+' || sym == '-' || sym == '*' || sym == '/')
        return TRUE;
    else return FALSE;
}
BOOLEAN need_parentheses(char sym) {
    if (sym == '*' || sym == '/')
        return TRUE;
    else return FALSE;
```

```
}
void inorder(TREE T)
    if (T != NULL) {
        if (need_parentheses(T->label)) {
            if (T->leftchild != NULL && isoperator(T->leftchild->label)) {
                putchar('(');
                inorder(T->leftchild);
                putchar(')');
            }
            else inorder(T->leftchild);
            printf("%c", T->label);
            if (T->rightchild != NULL && isoperator(T->rightchild->label)) {
                putchar('(');
                inorder(T->rightchild);
                putchar(')');
            else inorder(T->rightchild);
        }
        else {
            inorder(T->leftchild);
            printf("%c", T->label);
            inorder(T->rightchild);
        }
    }
}
4.
int max(int a, int b) {
    if (a >= b) return a;
    else return b;
}
BOOLEAN isleaf(TREE T) {
    if (T->leftchild == NULL && T->rightchild == NULL)
        return TRUE;
    else return FALSE;
}
int height(TREE T) {
    if (T == NULL) return 0;
    else if (isleaf(T)) return 0;
    else return 1 + max(height(T->leftchild), height(T->rightchild));
}
```

5. We prove the following statement by structural induction.

STATEMENT S(T): The number of full nodes in T is 1 fewer than the number of leaves

BASIS. Where T is a single node n, n is a leaf, and there are zero full nodes.

INDUCTION. Let n be the root of T. Assume the inductive hypothesis holds for the children of n. Let fn(c) be the number of full nodes in the binary tree rooted at c. Let leaves(c) be the number of leaves in the binary tree rooted at c.

Suppose n is a full node. Then by the inductive hypothesis we must have

$$1 + fn(c_1) + fn(c_2) = 1 + (leaves(c_1) - 1) + (leaves(c_2) - 1)$$
$$= leaves(c_1) + leaves(c_2) - 1$$

full nodes in T, which is 1 fewer than the number of leaves.

Suppose n is not a full node. Then n must have one and only one child. Therefore by the inductive hypothesis we must have

$$fn(c_1) = leaves(c_1) - 1$$

full nodes in T, which gives the same result. This proves the inductive step. \blacklozenge

6. We shall prove the following statement by structural induction.

STATEMENT S(T): The number of NULL pointers in T is 1 greater than the number of nodes.

BASIS. Where T is a single node n, then n has 2 NULL pointers.

INDUCTION. Let n be the root of T having at least one child. Let np(c) be the number of NULL pointers in the binary tree rooted at c. Let nodes(c) be the number of nodes in the binary tree rooted at c.

Consider the case where n has exactly one child c_1 . Since c_1 is the root of a binary tree, then the inductive hypothesis holds for c_1 . We know that there are $1 + nodes(c_1)$ nodes in T. Thus by the inductive hypothesis we must have

$$1 + nodes(c_1) =$$