1.

STATEMENT S(i): If  $1 \le i < n$ , then

$$T(n) = T(n-i) + \sum_{j=0}^{i-1} g(n-j).$$

BASIS. The basis is i = 1, as given by the lower bound of the interval. Then we have the inductive definition of the recurrence hence the basis holds.

INDUCTION. If  $i \ge n$  then there is nothing to prove. Suppose that i + 1 < n. From the inductive hypothesis we expand T(n-i) to get

$$T(n) = T(n-i-1) \sum_{j=0}^{i-1} g(n-j) + g(n-i)$$
$$= T(n-i-1) + \sum_{j=0}^{i} g(n-j).$$

This is the statement S(i+1) and we have proved the inductive step. We conclude S(i) is true for  $1 \le i < n$ .

2.

a) We have

$$T(n) = T(n-1) + g(n)$$

$$= T(n-2) + g(n-1) + g(n)$$

$$= T(n-3) + g(n-2) + g(n-1) + g(n).$$

Thus we have

$$T(n) = T(n-i) + \sum_{j=0}^{i-1} g(n-j).$$

But we know this already by Exercise 3.11.1.

Let i = n - 1. Then we have the recurrence in terms of the basis, which is

$$T(n) = a + \sum_{j=0}^{n-2} g(n-j).$$
 (1)

But  $g(n) = n^2$ . This is a summation of squares. By Exercise 2.3.1(b) we know that

$$\sum_{j=1}^{n} j^2 = n(n+1)(2n+1)/6.$$

We adjust the bounds to get the equivalent summation

$$\sum_{j=1}^{n} j^2 = \sum_{j=0}^{n-1} (n-j)^2.$$

But this has the extra term for when j = n - 1, that is g(1). We subtract 1 and get

$$T(n) = a + \sum_{j=0}^{n-1} (n-j)^2 - 1,$$

the same as Equation (1). We could have just taken the summation and decomposed it with the (n-2)th index then wrote the remaining summation as the summation of squares. But now that we have the last term separate, we rewrite the summation in terms of the identity

$$T(n) = a + n(n+1)(2n+1)/6 - 1.$$

Clearly T(n) is cubic, hence T(n) is  $O(n^3)$  if  $g(n) = n^2$ .

b) We have

$$T(n) = a + \sum_{j=0}^{n-2} g(n-j)$$

where  $g(n) = n^2 + 3n$ . Then we have two summations, one of the  $n^2$  terms and the other of 3n terms. Write

$$T(n) = a + \sum_{j=0}^{n-2} (n-j)^2 + \sum_{j=0}^{n-2} 3(n-j).$$

The left summation is n(n+1)(2n+1)/6-1 as given by (a), and the right summation is

$$\sum_{j=0}^{n-2} 3(n-j) = 3\sum_{j=0}^{n-2} (n-j) = 3(n-1)(n+2)/2.$$

Thus

$$T(n) = a + n(n+1)(2n+1)/6 - 1 + 3(n-1)(n+2)/2.$$

We see that T(n) is cubic. We conclude that T(n) is  $O(n^3)$ .

- c) Again the solution to T(n) is  $T(n) = a + \sum_{j=0}^{n-2} g(n-j)$ . Here  $g(n) = n^{3/2}$ . In (a) we found that there are n-1 terms of  $n^2$  and the same follows in (b). We can treat the summation as n-1 terms of  $n\sqrt{n}$  with lower terms. Clearly  $(n-1)n\sqrt{n}$  is  $O(n^{5/2})$ . We conclude that T(n) is  $O(n^{5/2})$ .
- d) Following (c), there are n-1 terms of  $n \log n$  plus lower terms. Hence  $(n-1)(n \log n)$  is  $O(n^2 \log n)$ . We conclude that T(n) is  $O(n^2 \log n)$ .
- e) Following (d), there are n-1 terms of  $2^n$  plus lower terms. Therefore T(n) is  $O(2^n n)$ .

For this recurrence we can assume that for any g(n) that T(n) is O(ng(n)). But that is just a guess.

3. Let us first solve this recurrence. We write some terms in the sequence and get