



# Lecture 2

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## Relational Database Model



# Data Modeling

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- Data models are representations of complex real-world data structures
- Iterative and progressive process of creating a specific data model for a determined problem
- Abstraction for the creation of a good database
- Facilitate communication between the designer, applications programmer, and end user



# Data Models

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- Entity: Unique and distinct object used to collect and store data
- Attribute: Characteristic of an entity
- Relationship: Describes the association between entities
  - One-to-many (1:M)
  - Many-to-many (M:N)
  - One-to-one (1:1)
- Constraint: Set of rules placed on the data to ensure data integrity



# Business Rules

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- Description of policies, procedures, or principles within a specific organization
- Enables the defining of basic building blocks
- Describes the main and distinguishing characteristics of the data
- Ensures accurate data models
- Must be in writing and kept up to date



# Discovering Business Rules

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- Sources of business rules:
  - Company managers
  - Policy makers
  - Department managers
  - Written documentation
    - Procedures
    - Standards
    - Operations manuals
  - Direct interviews with end users



# Translating Business Rules

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- Nouns translate into entities
- Verbs translate into relationships among entities
- Relationships are bidirectional
- Questions to identify the relationship type:
  - How many instances of B are related to one instance of A?
  - How many instances of A are related to one instance of B?



# Entities

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- Person, place, or thing for which data can be collected and stored
- Relation is an entity (table) in a relational model
- Consists of a matrix of intersecting tuples (rows) and attributes (columns)



# Entities

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- An instance or occurrence is a tuple (row) of real data in an entity (table)
- Each entity name must be distinct within a database
- A tuple should be distinct with no duplicates





# Attributes

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- Piece of data to be stored about each entity instance
- Descriptive property or characteristic of the entity
- An attribute name must be distinct within an entity



# Attributes

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- Compound attribute
  - Attribute that can be subdivided to yield additional attributes
  - Example:  
    Address  
    (Street, City, Province)



# Attributes

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- Single-valued attribute
  - Attribute that has only a single value
  - Cannot be subdivided
  - Example:  
Birth\_Date



# Attributes

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- Multi-valued attribute
  - Attribute that has many values
  - Example:  
Student\_Contacts



# Attributes

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- Derived attribute

- Attribute whose value is calculated from other attributes using an algorithm
- Often not stored in the database

- Example:

$$\text{Annual} = \text{Monthly} * 12$$



# Constraints

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- Data Type

- Type of data that can be stored in the attribute
- Example: Character, Numeric, Date, Binary

- Attribute Domain

- Allowable attribute values
- Example: (M/F, Y/N, 1200-1499)

- Default Value

- Value that will be recorded if not specified
- Example: Invoice date defaults to today's date



# Constraints

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- NULL / NOT NULL

- Specifies if attribute is required (NOT NULL) or not required (NULL)
- Example: Birth\_Date NOT NULL

- UNIQUE

- No duplicate value exist for the attribute in the entity
- Example: SIN

# Characteristics of a Relational Table

TABLE 3.1

## CHARACTERISTICS OF A RELATIONAL TABLE

1	A table is perceived as a two-dimensional structure composed of rows and columns.
2	Each table row ( <b>tuple</b> ) represents a single entity occurrence within the entity set.
3	Each table column represents an attribute, and each column has a distinct name.
4	Each intersection of a row and column represents a single data value.
5	All values in a column must conform to the same data format.
6	Each column has a specific range of values known as the attribute <b>domain</b> .
7	The order of the rows and columns is immaterial to the DBMS.
8	Each table must have an attribute or combination of attributes that uniquely identifies each row.





# Schema and Subschema

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- Schema

- Conceptual organization of the entire database as viewed by the database administrator

- Subschema

- Portion of the database seen by application programs which produce the desired information using the data contained within the database



# Level Architecture

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- External model represents the end user's view of the data environment
- Conceptual model describes what data is stored and its relationships for the entire database
- Internal model represents the database as seen by the DBMS
- Physical model describes how the data is stored physically



# External Model

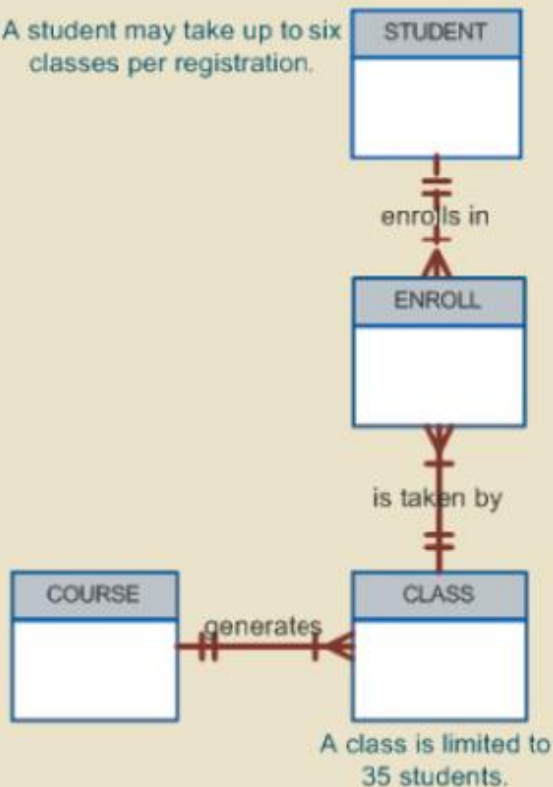
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- End user's view of the data environment
- Identify specific data required to support each business unit's operations
- External schema
  - Specific representation of an external view

# External Models for Tiny College

## Student Registration

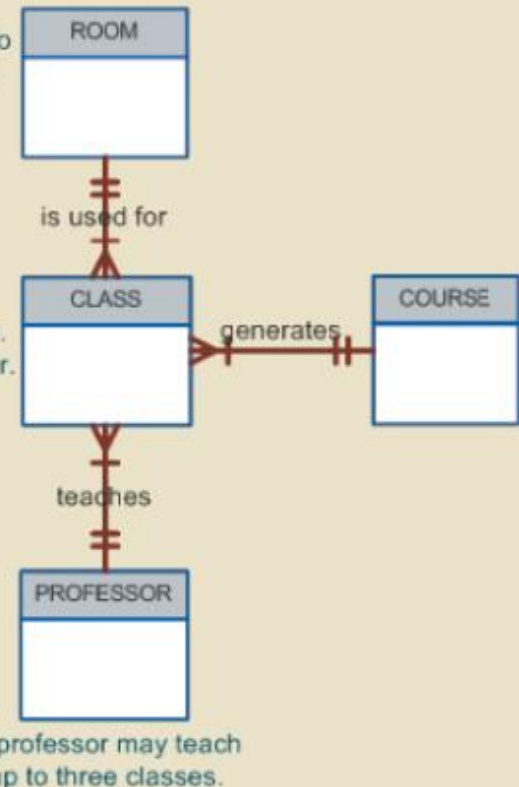
A student may take up to six classes per registration.



## Class Scheduling

A room may be used to teach many classes.

Each class is taught in only one room.  
Each class is taught by one professor.



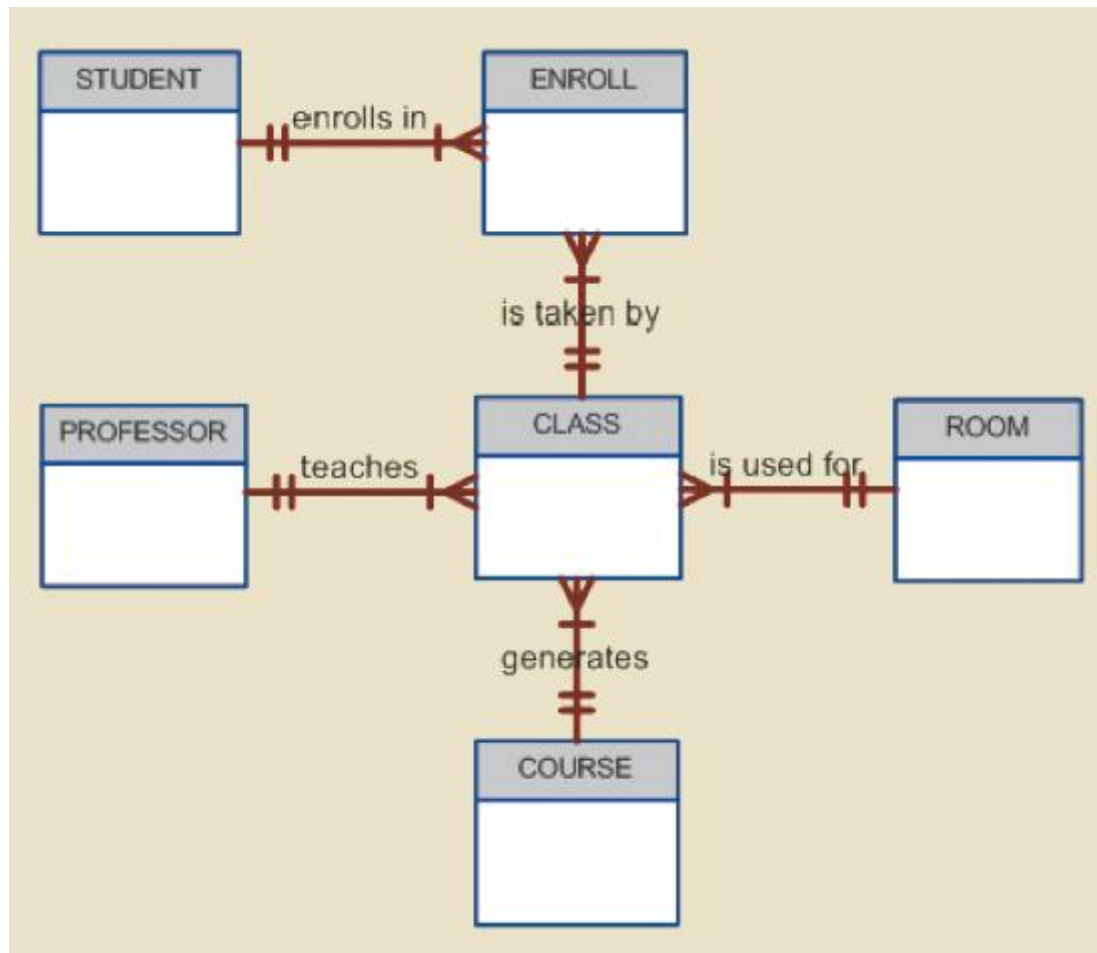


# Conceptual Model

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- Represents a global view of the entire database by the entire organization
- Conceptual schema
  - All external views integrated into a single global view of the data in the enterprise
- Software and hardware independence
- Logical design
  - Task of creating a conceptual data model

# Conceptual Model For Tiny College



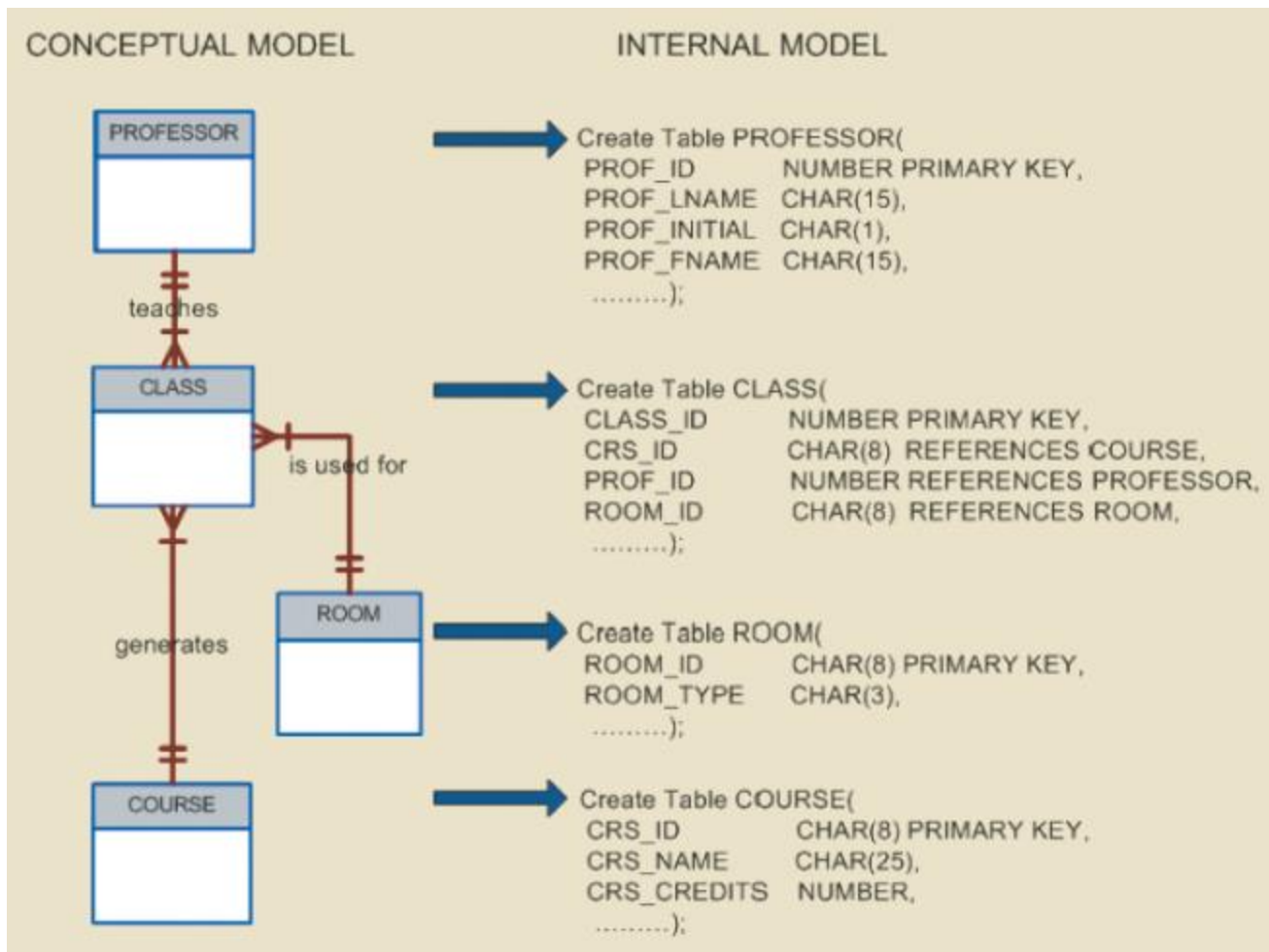


# Internal Model

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- Representation of the database as seen by the DBMS mapping conceptual model to the DBMS
- Internal schema
  - Specific representation of an internal model
- Software dependent and hardware independent
- Logical independence
  - Change internal model without affecting conceptual model

# Internal Model for Tiny College





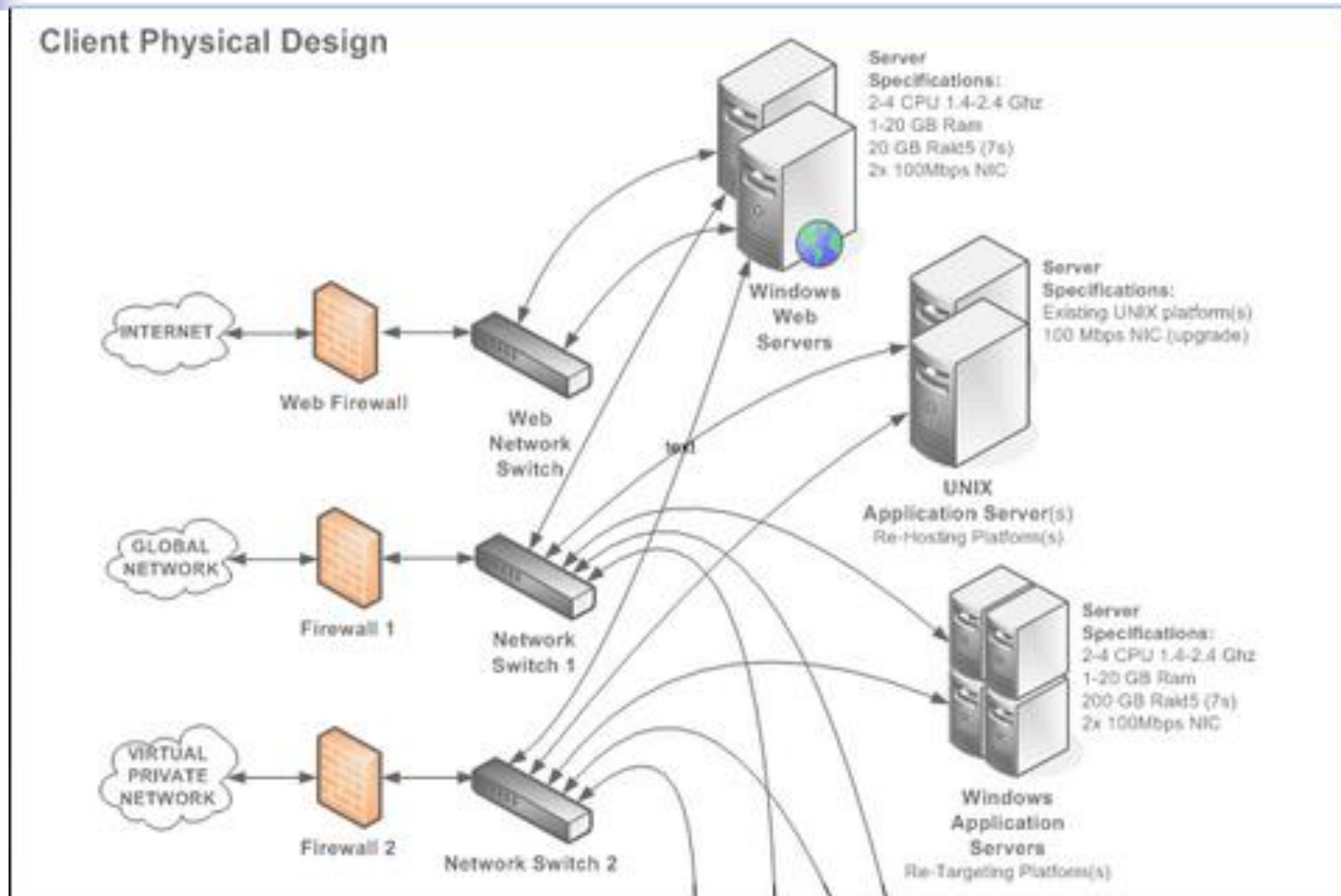


# Physical Model

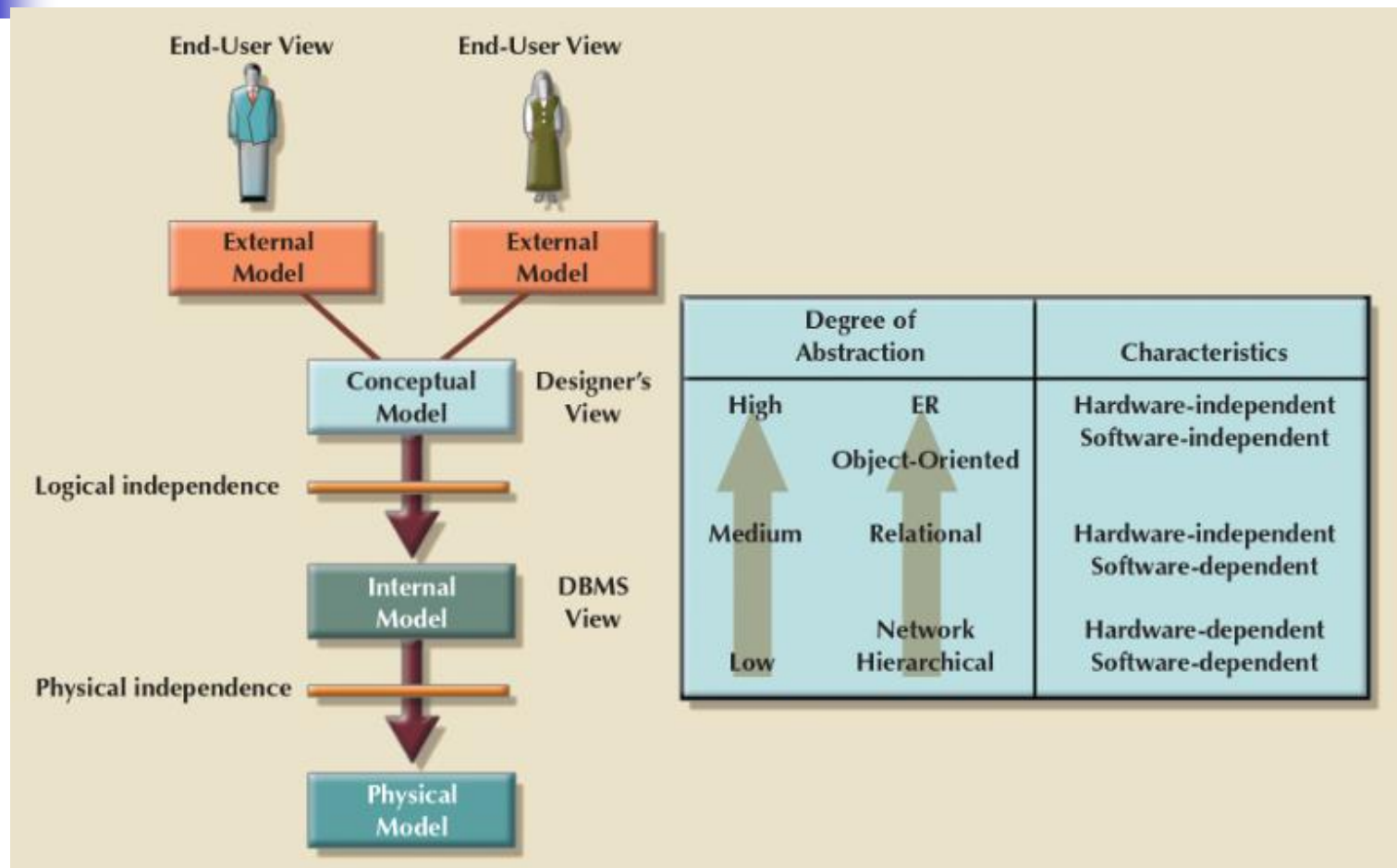
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- Operates at lowest level of abstraction
- Describes the way data is saved on storage media
- Requires the definition of physical storage and data access methods
- Physical independence
  - Changes in physical model do not affect internal model

# Physical Model



# Data Abstraction Levels





# Keys

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- Consists of one or more attributes that uniquely identify a tuple in an entity
- Ensures each tuple in an entity is uniquely identifiable
- Establishes relationships among the entities
- Ensures integrity of the data



# Keys

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- Candidate key
- Primary key
- Composite key
- Secondary key
- Foreign key



# Keys

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- Candidate key
  - Attributes that may be selected to become the primary identifier of instances of an entity
  - Example: Emp\_Id and SIN
- Primary key
  - Attribute selected to uniquely identify each entity instance
  - Cannot contain null entries (absence of any data value)
  - Example: Emp\_Id



# Keys

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- Composite key
  - Composed of more than one attribute to uniquely identify an instance of an entity
- Secondary key
  - Any attribute or combination of attributes used strictly for data retrieval purposes
- Foreign key
  - Attribute in one entity whose value either match the primary key in another entity or is null

# Example - Simple Relational Database

Table name: **PRODUCT**

Primary key: **PROD\_CODE**

Foreign key: **VEND\_CODE**

Database name: **Ch03\_SaleCo**

PROD_CODE	PROD_DESCRIPTOR	PROD_PRICE	PROD_ON_HAND	VEND_CODE
001278-AB	Claw hammer	12.95	23	232
123-21UUY	Houselite chain saw, 16-in. bar	189.99	4	235
QER-34256	Sledge hammer, 16-lb. head	18.63	6	231
SRE-657UG	Rat-tail file	2.99	15	232
ZZX/3245Q	Steel tape, 12-ft. length	6.79	8	235

link

Table name: **VENDOR**

Primary key: **VEND\_CODE**

Foreign key: none

VEND_CODE	VEND_CONTACT	VEND_AREACODE	VEND_PHONE
230	Shelly K. Smithson	608	555-1234
231	James Johnson	615	123-4536
232	Annelise Crystall	608	224-2134
233	Candice Wallace	904	342-6567
234	Arthur Jones	615	123-3324
235	Henry Ortozo	615	899-3425





# Integrity Rules

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- Domain Constraints

- Values are restricted by the domain of the attribute
  - Range of permissible values for an attribute

- Entity Integrity

- All values in the primary keys must be unique
  - No part of the primary key may be null

- Referential Integrity

- Foreign key matches a primary key value in an entity to which it is related or is a null entry



# Database Management System

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- Structured Query Language (SQL) is the most commonly used query language
- Uses Data Definition Language (DDL) to define the database and database objects
- Uses Data Manipulation Language (DML) to retrieve, insert, update, and delete data stored in the database
- Uses Data Control Language (DCL) to control access to data stored in a database



# Relational Model

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- Developed by E. F. Codd (IBM) in 1970
- Based on mathematical theory
- Introduces normalization concept
- Set-oriented using relations / tables
- Matrix consisting of the intersection of row and column
- In 1985, Codd published a list of 12 rules to define a relational database system



# Relational Algebra

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- Defines a theoretical way of manipulating a table contents using relational operators
- Helps in the understanding of the Data Manipulation Language (DML)
- Structured Query Language (SQL) is based on relational algebra



# Property of Closure

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- Operations that work on one or more relations produces a new relation without changing the original relation
- Result of an algebraic operation on a SQL command produces another temporary relation



# Relational Operators

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- Select
- Project
- Union
- Intersect
- Difference
- Product
- Join

# Select

- Yields a horizontal subset of a table

**Original table**

P_CODE	P_DESCRIPT	PRICE
123456	Flashlight	5.26
123457	Lamp	25.15
123458	Box Fan	10.99
213345	9v battery	1.92
254467	100W bulb	1.47
311452	Powerdrill	34.99

**SELECT ALL yields**

**New table**

P_CODE	P_DESCRIPT	PRICE
123456	Flashlight	5.26
123457	Lamp	25.15
123458	Box Fan	10.99
213345	9v battery	1.92
254467	100W bulb	1.47
311452	Powerdrill	34.99

**SELECT only PRICE less than \$2.00 yields**

P_CODE	P_DESCRIPT	PRICE
213345	9v battery	1.92
254467	100W bulb	1.47

**SELECT only P\_CODE = 311452 yields**

P_CODE	P_DESCRIPT	PRICE
311452	Powerdrill	34.99

# Project

- Yields a vertical subset of a table

**Original table**

P_CODE	P_DESCRIPT	PRICE
123456	Flashlight	5.26
123457	Lamp	25.15
123458	Box Fan	10.99
213345	9v battery	1.92
254467	100W bulb	1.47
311452	Powerdrill	34.99

**PROJECT PRICE yields**

**New table**

PRICE
5.26
25.15
10.99
1.92
1.47
34.99

**PROJECT P\_DESCRIPT and PRICE yields**

P_DESCRIPT	PRICE
Flashlight	5.26
Lamp	25.15
Box Fan	10.99
9v battery	1.92
100W bulb	1.47
Powerdrill	34.99

**PROJECT P\_CODE and PRICE yields**

P_CODE	PRICE
123456	5.26
123457	25.15
123458	10.99
213345	1.92
254467	1.47
311452	34.99



# Union

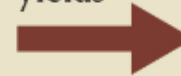
- Combines all the rows from two tables excluding duplicate rows
- Tables must be union-compatible

P_CODE	P_DESCRIPTION	PRICE
123456	Flashlight	5.26
123457	Lamp	25.15
123458	Box Fan	10.99
213345	9v battery	1.92
254467	100W bulb	1.47
311452	Powerdrill	34.99

UNION

P_CODE	P_DESCRIPTION	PRICE
345678	Microwave	160.00
345679	Dishwasher	500.00
123458	Box Fan	10.99

yields



P_CODE	P_DESCRIPTION	PRICE
123456	Flashlight	5.26
123457	Lamp	25.15
123458	Box Fan	10.99
213345	9v battery	1.92
254467	100W bulb	1.47
311452	Powerdrill	34.99
345678	Microwave	160
345679	Dishwasher	500

# Intersect

- Yields only the rows that appear in both tables
- Tables must be union-compatible

STU_FNAME	STU_LNAME
George	Jones
Jane	Smith
Peter	Robinson
Franklin	Johnson
Martin	Lopez

**INTERSECT**

EMP_FNAME	EMP_LNAME
Franklin	Lopez
William	Turner
Franklin	Johnson
Susan	Rogers

**yields**



STU_FNAME	STU_LNAME
Franklin	Johnson

# Difference

- Yields all the rows in one relation not found in the other relation
- Relations must be union-compatible

STU_FNAME	STU_LNAME
George	Jones
Jane	Smith
Peter	Robinson
Franklin	Johnson
Martin	Lopez

DIFFERENCE

EMP_FNAME	EMP_LNAME
Franklin	Lopez
William	Turner
Franklin	Johnson
Susan	Rogers

yields



STU_FNAME	STU_LNAME
George	Jones
Jane	Smith
Peter	Robinson
Martin	Lopez

# Product

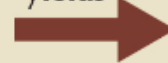
- Yields all possible pairs of rows from two relations
- Also known as Cartesian Product

P_CODE	P_DESCRIPTION	PRICE
123456	Flashlight	5.26
123457	Lamp	25.15
123458	Box Fan	10.99
213345	9v battery	1.92
254467	100W bulb	1.47
311452	Powerdrill	34.99

**PRODUCT**

STORE	aisle	shelf
23	W	5
24	K	9
25	Z	6

yields



P_CODE	P_DESCRIPTION	PRICE	STORE	aisle	shelf
123456	Flashlight	5.26	23	W	5
123456	Flashlight	5.26	24	K	9
123456	Flashlight	5.26	25	Z	6
123457	Lamp	25.15	23	W	5
123457	Lamp	25.15	24	K	9
123457	Lamp	25.15	25	Z	6
123458	Box Fan	10.99	23	W	5
123458	Box Fan	10.99	24	K	9
123458	Box Fan	10.99	25	Z	6
213345	9v battery	1.92	23	W	5
213345	9v battery	1.92	24	K	9
213345	9v battery	1.92	25	Z	6
311452	Powerdrill	34.99	23	W	5
311452	Powerdrill	34.99	24	K	9
311452	Powerdrill	34.99	25	Z	6
254467	100W bulb	1.47	23	W	5
254467	100W bulb	1.47	24	K	9
254467	100W bulb	1.47	25	Z	6



# Joins

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- Essential operation in relational algebra
- Builds a relation from two specified relations on the basis of matching a common attribute over the same domain
- Natural Join
  - Joins tables by selecting only rows with common values in their common attributes
- Outer Join
  - Matched pairs are retained and unmatched values in the other table are left null

# Inner Join

- Returns matched rows only from the tables that are being joined

Table name: CUSTOMER

CUS_CODE	CUS_LNAME	CUS_ZIP	AGENT_CODE
1132445	Walker	32145	231
1217782	Adares	32145	125
1312243	Rakowski	34129	167
1321242	Rodriguez	37134	125
1542311	Smithson	37134	421
1657399	Vanloo	32145	231

Table name: AGENT

AGENT_CODE	AGENT_PHONE
125	6152439887
167	6153426778
231	6152431124
333	9041234445

CUS_CODE	CUS_LNAME	CUS_ZIP	CUSTOMER.AGENT_CODE	AGENT.AGENT_CODE	AGENT_PHONE
1217782	Adares	32145	125	125	6152439887
1321242	Rodriguez	37134	125	125	6152439887
1312243	Rakowski	34129	167	167	6153426778
1132445	Walker	32145	231	231	6152431124
1657399	Vanloo	32145	231	231	6152431124

# Left Outer Join

- Yields all rows in the first table including those that do not have a matching value in the second table

Table name: CUSTOMER

CUS_CODE	CUS_LNAME	CUS_ZIP	AGENT_CODE
1132445	Walker	32145	231
1217782	Adares	32145	125
1312243	Rakowski	34129	167
1321242	Rodriguez	37134	125
1542311	Smithson	37134	421
1657399	Vanloo	32145	231

Table name: AGENT

AGENT_CODE	AGENT_PHONE
125	6152439887
167	6153426778
231	6152431124
333	9041234445

CUS_CODE	CUS_LNAME	CUS_ZIP	CUSTOMER.AGENT_CODE	AGENT.AGENT_CODE	AGENT_PHONE
1217782	Adares	32145	125	125	6152439887
1321242	Rodriguez	37134	125	125	6152439887
1312243	Rakowski	34129	167	167	6153426778
1132445	Walker	32145	231	231	6152431124
1657399	Vanloo	32145	231	231	6152431124
1542311	Smithson	37134	421		

# Right Outer Join

- Yields all rows in the second relation including those that do not have a matching value in the first relation

Table name: CUSTOMER

CUS_CODE	CUS_LNAME	CUS_ZIP	AGENT_CODE
1132445	Walker	32145	231
1217782	Adares	32145	125
1312243	Rakowski	34129	167
1321242	Rodriguez	37134	125
1542311	Smithson	37134	421
1657399	Vanloo	32145	231

Table name: AGENT

AGENT_CODE	AGENT_PHONE
125	6152439887
167	6153426778
231	6152431124
333	9041234445

CUS_CODE	CUS_LNAME	CUS_ZIP	CUSTOMER.AGENT_CODE	AGENT.AGENT_CODE	AGENT_PHONE
1217782	Adares	32145	125	125	6152439887
1321242	Rodriguez	37134	125	125	6152439887
1312243	Rakowski	34129	167	167	6153426778
1132445	Walker	32145	231	231	6152431124
1657399	Vanloo	32145	231	231	6152431124
				333	9041234445





# To Do

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- Assignments:
  - Exercise for Week 2
  - Review Questions for Week 2
  - Discussion Topic for Week 2
- Reading:
  - Review Chapters 2 and 3
  - Read Chapter 4