

3: Runtime Analysis

a) void f1(int n) {
 create int
 while (i < n) {
 something O(1)
 i = i + 1
 }
 }

$\sum_i O(1)$
 $= O(\log n)$

b) void f2(int n)
 for (i: 1 → n)
 if (i % √n == 0)
 for (k: 0 → i³)
 O(1) something
] O(i³)

$\sum_i [O(1) + O(\sum_{k=0}^{i^3} O(1))]$

$n = 9$

i	% 3 =
1	1
2	2
3	0
4	1
5	2
6	0
7	1
8	2
9	0

 $\frac{n}{\sqrt{n}} = \sqrt{n}$

$n = 16$

i	% 4 =
1	1
2	2
3	3
4	0
5	1
6	2
7	3
8	0
12	0
16	0

$n = 2$

i	% 2 =
1	1
2	0

$$\begin{aligned}
 & \sum_i [O(1) + O(\sum_{k=0}^{i^3} O(1))] \\
 &= \sum_i [O(1)] + \sum_i \sum_{k=0}^{i^3} O(1) \\
 &= O(n) + \sum_i \sum_{k=0}^{i^3} O(1) \\
 &= O(n) + \sum_i O(i^3 \cdot n^{3/2}) \\
 &= O(n) + O(n^{3/2} (\sqrt{n})^4) \\
 &= O(n^{7/2})
 \end{aligned}$$

$i = \sqrt{n} \quad 2\sqrt{n} \quad 3\sqrt{n}$
 $i = \sqrt{n}j$

c) for (i: 1 → n)
 for (k: 1 → n)
 if (A[k] == i)
 for (m: 1 → n, m % 2 = 2)
 O(1)

can happen at most n times

$A = \{1, 1, \dots, 1\}$ $A = \{1, 2, 3, \dots, n\}$

$\sum_{i=1}^n \sum_{k=1}^n [O(1) + O(\sum_{m=1}^n O(1))]$

$\sum_{i=1}^n [O(n)] + \sum_{k=1}^n O(\log n) = O(n^2) + O(n \log n) = O(n^2)$

d) $f(n)$

10 1
15 2

```

new int a[10]
int size = 10
for (i := 0 → n)
  if (i == size)
    newSize =  $\frac{3}{2}$  size
    b = new int[newSize]
    for (j = 0 → size)
      b[j] = a[j]  $\theta(1)$ 
    delete [] a
    a = b
    size = newSize
    a[i] = i2

```

triggers
 $10(\frac{3}{2})^k = n$
 $k = \log_{\frac{3}{2}}(\frac{n}{10})$
 times

$\sum_j^{\text{size}} \theta(1) = \theta(\text{size})$

$$\begin{aligned}
 & \sum_i^n [\theta(1) + \sum_i \theta(\text{size})] \\
 &= \sum_i [\theta(1)] + \sum_k^{\log_{\frac{3}{2}}(\frac{n}{10})} \theta(10(\frac{3}{2})^k) \\
 &= \theta(n) + \theta(\frac{3}{2}^{\log_{\frac{3}{2}}(\frac{n}{10})}) \\
 &= \theta(n) + \theta(n) \\
 &= \theta(n)
 \end{aligned}$$