

## CS 340 - Assignment 2

Total: 100pts

### 1 Written Part [18pts]

*Note: pages refer to the 2nd edition of the textbook.*

**Exercise 1:** *exercise 4.6 page 171* [4pts]

A *full node* is a node with two children. Prove that the number of full nodes plus one is equal to the number of leaves in a nonempty binary tree.

**Exercise 2:** *exercise 4.48 page 175* [6pts]

Suppose we want to add the operation `findKth` to our repertoire. The operation `findKth(k)` returns the  $k$ th smallest item in the tree. Assume all items are distinct. Explain how to modify the binary search tree to support this operation in  $O(\log N)$  average time, without sacrificing the time bounds of any other operation.

**Exercise 3:** *exercise 6.8 page 246* [8pts]

Show the following regarding the maximum item in the heap:

1. It must be at one of the leaves.
2. There are exactly  $\lceil N/2 \rceil$  leaves.
3. Every leaf must be examined to find it.

**Exercise 4:** *exercise 6.10 page 246, question a* [6pts]

Give an algorithm to find all nodes less than some value,  $X$ , in a binary heap. Your algorithm should run in  $O(K)$ , where  $K$  is the number of nodes output.

**Exercise 5:** *exercise 6.13 page 247, questions a and b* [8pts]

Each `deleteMin` operation uses  $2 \log N$  comparisons in the worst case.  $N$  is the number of nodes in the heap structure.

1. Propose a scheme so that the `deleteMin` operation uses only  $\log N + \log \log N + O(1)$  comparisons between elements. This need not imply less data movement.
2. Extend your scheme in the previous question so that only  $\log N + \log \log \log N + O(1)$  comparisons are performed.

## Programming Part: Min-Max Heap [68pts]

A **min-max heap** is a data structure that supports both `deleteMin` and `deleteMax` in  $O(\log N)$  per operation. The structure is identical to a binary heap, but the heap-order property is that for any node,  $X$ , at even depth, the element stored at  $X$  is smaller than the parent but larger than the grandparent (where this makes sense), and for any node  $X$  at odd depth, the element stored at  $X$  is larger than the parent but smaller than the grandparent (see figure 1).

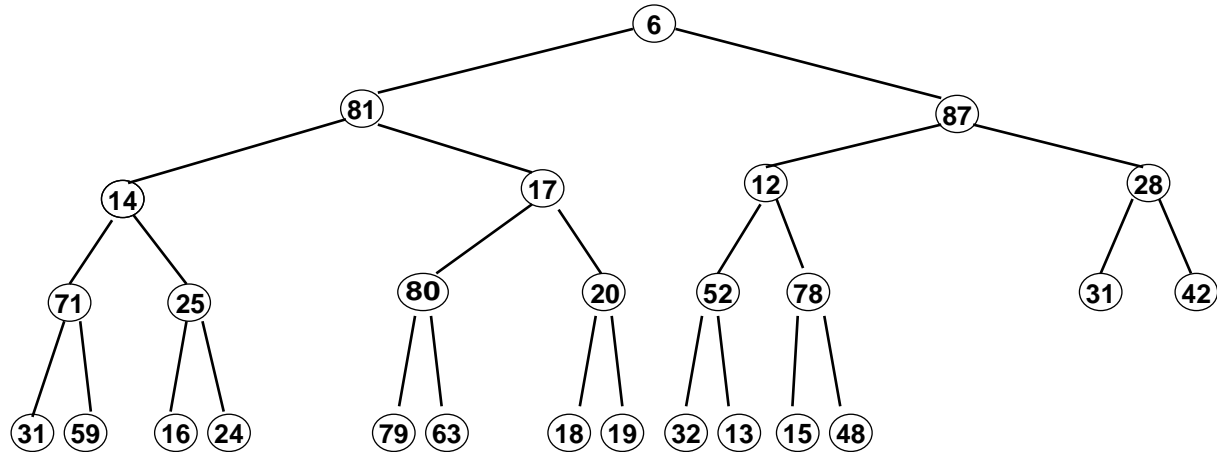


Figure 1: A min-max heap.

Using an array to represent the min-max heap structure (in the same way as for min heap or max heap), implement the following operations.

1. `buildHeap`: Builds a min-max heap from a list of naturals read from standard input.
2. `findMin` and `findMax`: Returns the minimum (resp the maximum) element.
3. `insertHeap`: Inserts a new element into the min-max heap.
4. `deleteMin` and `deleteMax`: Deletes the minimum (resp the maximum) element.

### Marking scheme

1. Readability : 10pts
2. Compiling and execution process : 10pts
3. Correctness : 48pts

## 2 Hand in

### 2.1 Written Part

- Submit electronic version: pdf(preferred) or word via URCourses. Your file should be named assign2.pdf(or assign2.doc).

### 2.2 Programming Part

#### 2.2.1 Single file submission for each programming part

Using webct, submit the file containing the C++ code of the programming part **assign2.cpp**. At the top of this file add the following comments :

1. the compiling command you have used
2. an example on how to execute the program
3. other comments describing the program

#### 2.2.2 Multiple files submission for each programming part

Using webct, submit all files required by each programming part :

1. README (file explaining the compilation and execution of your program including the format of input and other details)
2. headers (.h)
3. implementations (.cpp)
4. the Makefile :
  - should be named "**makefile**". In the makefile, the generated executable should be named : "**assign2**"

You can give any name to your source files. The marker will only have to run "**make**" to compile your program and "**assign2**" to execute it.