# University of Cambridge

## MATHEMATICS TRIPOS

Part III

## Combinatorics

Example Sheet I

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#### Introduction

These are written solutions to Combinatorics Example Sheet I. Solutions are written based on those seen in examples classes and may contain errors, likely due to the author. Solutions may be incomplete and do not usually include starred questions. These are to be used as a reference for revision **after** examples classes and should never be used beforehand. Doing so will severely impair your ability to learn and study mathematics.

### Questions

**Question** (Question 1). Let P = (V, <) be a finite poset. Recall that a subset  $U \subset V$  is a chain if any two elements of U are comparable, and it is an antichain if no two elements of U are comparable. Show that the maximal size of an antichain in P is equal to the minimal number of chains in P that cover V.

Solution. Let  $N_1=$  maximum size of antichain,  $N_2=$  minimum number of chains that cover V.

 $\frac{N_2 \geq N_1 \text{ Given } A_1, A_2, ..., A_{N_2} \text{ minimal number of chains covering } V. \text{ Any antichain } B \text{ can contain at most one element from each } A_i \text{ so } N_1 \geq |B| \geq N_2.$ 

 $N_1 \geq N_2$  We prove this by induction on n, the size of the partially ordered set. If P is empty the theorem is vacuously true. Thus, assume P has at least one element and let a be a maximal element in P which exists since P is finite. By induction, assume  $\exists k: P' := P \ a$  can be covered by k disjoint chains  $C_1,...,C_k$  and there is an antichain  $A_0$  of size at least k. Have  $A_0 \cap C_i \neq \emptyset$ . Let  $x_i$  be the maximal element of  $C_i$  belonging to an antichain of length at least k.

**Remark** (Claim). Let  $A_0 = \{x_1, x_2, ..., x_k\}$ , then A is an antichain

Proof of Claim. Let  $A_i$  be an antichain of size k that contains  $x_i$ , fix  $i \neq j$  arbitrarily. Then  $A_i \cup C_j \neq \emptyset$ . Suppose  $y \in A_i \cup C_j$ . Then  $y \leq x_j$  since  $x_j$  is maximal in  $C_j$ . Thus  $x_i \ngeq x_j$  since  $x_i \ngeq y$ . Exchanging i,j gives  $x_i \not \geqslant x_i$ .

Now suppose  $a \geq x_i$  for some  $1 \leq i \leq k$ . Then set

$$K = \{a\} \cup \{z \in C_i : z \le x_i\}$$

Then by choice of  $x_i$ , P K does not have an antichain of size k and so by induction  $P \setminus K$  can be covered by k-1 disjoint chains as A  $x_i$  is an antichain of size k-1 in  $P \setminus K$ . Thus P can be covered by k disjoint chains.

Else, suppose instead that  $a \ngeq x_i$  for all  $1 \le i \le k$ . The  $A \cup \{a\}$  is an antichain of size k+1 in P and P can be covered by k+1 chains  $\{a\}, C_1, C_2, ..., C_k$ .

**Remark.** This proof is tedious and a very difficult Question 1. The ideas are, however, important and should be understood.

Question (Question 2). Let (V, <) be a finite ranked poset with non-empty level sets  $V_0V_1, ...., V_n$ . Suppose for  $0 < i \le n$  every  $v \in V_i$  dominates exactly  $d_i \ge 1$  elements of  $V_{i-1}$ , for  $0 \le i < n$  every  $v \in V_i$  is dominated by exactly  $e_i \ge 1$  elements of  $V_{i+1}$ , and the partial order on  $V = \bigcup_{i=0}^{n} V_i$  is induced by these relations.

Show that if  $U \subset V$  is an antichain then

$$\sum_{0}^{n} \frac{|U \cap V_i|}{|V_i|} \le 1$$

Idea. Count number of chains of maximal length in two ways

Solution. Must have  $|V_i|e_i = |V_{i+1}|d_{i+1}$  for all  $0 \le i < n$ . Thus there are  $|V_0|e_0e_1...e_{n-1} = d_1...d_k|V_k|e_k...e_{n-1}$  chains of maximal length in V. For each maximal chain C we have  $|C \cap U| \le 1$  as U is an antichain. Every element in  $V_k$  is contained in exactly  $(d_kd_{k-1}...d_1)(e_k...e_{n-1})$  maximal chains. Putting both of these together gives:

$$\sum_{0}^{n}|U\cap V_{k}|(d_{k}...d_{1})(e_{k}...e_{n-1})=\#\text{maximal chains}=|V_{0}|e_{0}...e_{n-1}$$

which upon dividing the LHS by the RHS yields the required result.

**Remark.** Counting arguments like these are popular. The counting itself is not difficult, but knowing what to count often is.