

# Final Project

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## 1 HOSPITALS AND RESIDENTS STABLE MATCHING PROBLEM

### 1.1 THE ALGORITHM

In the hospitals and residents stable matching problem, the goal is to assign residents to hospitals given the preferences of both sides so that all assignments are stable. In this context, the term "stability" means that for each resident, there is no hospital that is available that is higher on a resident's list compared to that resident's current assignment. The reason stability is in the terms of the residents is because the residents propose to the hospitals on their preference lists and the hospitals have the ability to either provisionally accept or reject the residents based on their resident preferences and current capacity. In other words, hospitals may have a preferred resident that is available, but that resident may not want to go to that hospital and, therefore, will end up elsewhere.

### 1.2 ASYMPTOTIC ANALYSIS

The pseudocode for the hospitals and residents stable matching algorithm is provided in Algorithm 1. The algorithm starts off by assigning all residents and hospitals to be completely free, which are  $O(r)$  and  $O(h)$  operations, respectively. Next, line 8 is the condition for a while loop that runs until the residents list is empty. Since residents are being picked off one-by-one as done in line 9, the loop will run for each resident, which makes it run on average  $r$  times. Line 10, similar to line 8, also defines a while loop. This time, however, it is iterating over the resident's preferences, which means the while loop runs  $h$  times for each iteration of the outer loop. The if-statement block on lines 12-16 is inside of the inner loop, and contains a statement to get the least preferred assigned resident, which is an  $O(r)$  operation as it may have to loop through all of the residents at worst case. The condition and other 2 assignments in the if-statement block are constant time operations. The assignment on line 17 sets the hospital assignment for the resident, which runs in constant time. Next, just like line 12, the check if the hospital is full a constant time check. Line 19 is the same as line 13 and is an  $O(r)$  operation. Next, there is a loop defined on line 20 that iterates through the remaining residents in the list of hospital preferences, which at worst case is about  $r$  iterations. It also has an  $O(r)$  operation that is executed once to get the starting index. Lastly, there is a constant time assignment on line 21, and the removal function calls are  $O(h)$  and  $O(r)$  operations, respectively, as the hospitals and residents are being iterated through in each call. However, as shown in Listing 1 on lines 57-65, a space for time tradeoff was made as the index of each resident and hospital was stored to prevent having to iterate through the lists

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**Algorithm 1** Hospitals and Residents Stable Matching Algorithm

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```
1: procedure STABLEMATCHORIGINAL(residents, hospitals)
2:   for r of residents do
3:     r.assignment  $\leftarrow$  null    // Residents start off unassigned
4:   end for
5:   for h of hospitals do
6:     h.assignments  $\leftarrow$  [ ]    // Hospitals initially have no assignments
7:   end for
8:   while !residents.isEmpty() do
9:     r  $\leftarrow$  residents.dequeue()    // Get the next resident in line to be assigned
10:    while r.assignment == null && !r.preferences.isEmpty() do
11:      h  $\leftarrow$  r.preferences.dequeue()    // Try the resident's next top preference
12:      if h.isFull() then
13:        r'  $\leftarrow$  h.getLeastPreferredAssignedResident()
14:        r'.assignment  $\leftarrow$  null    // Set the least preferred assigned resident to be free
15:        residents.enqueue(r)    // Add the resident back to the list to be reassigned
16:      end if
17:      r.assignment  $\leftarrow$  h    // Provisionally assign r to h
18:      if h.isFull() then
19:        s  $\leftarrow$  h.getLeastPreferredAssignedResident()
20:        for i  $\leftarrow$  h.preferences.indexOf(s) + 1, len(h.preferences) - 1 do
21:          s'  $\leftarrow$  h.preferences[i]
22:          s'.preferences.remove(h)    // Remove h from preferences of s'
23:          h.preferences.remove(s')    // Remove s' from preferences of h
24:        end for
25:      end if
26:    end while
27:  end while
28: end procedure
```

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to get the appropriate position. This changes the runtime of the entire loop to be  $O(r)$  with the condition and body running in constant time. Overall, when putting it all together, the runtime of the original stable matching problem for residents and hospitals is  $r + h + r * h * (r + r + r) = r + h + rh * (3r) = r + h + 3r^2h$ . This can be simplified to  $O(r^2)$  because there are typically a lot more residents than hospitals, so  $r^2$  becomes the dominant term in the expression.

### 1.3 EXAMPLE AND WALKTHROUGH

Listing 2 contains a sample test case for the algorithm using 4 residents and 3 hospitals. The initial state of the algorithm is shown below:

| Entity Name | Preferences |
|-------------|-------------|
| r1          | h1 h2 h3    |
| r2          | h1 h3 h2    |
| r3          | h1 h2       |
| r4          | h2 h3       |
| h1 (1)      | r2 r3       |
| h2 (1)      | r4 r3       |
| h3 (1)      | r1 r2 r4    |

The residents go in order and try to propose to their preferred hospitals. The first resident to propose is r1, which begins by proposing to h1. Although r1 is not on h1's preference list, it will be provisionally assigned there for now because h1 has no other residents to take in. The table will be updated as shown below.

| Entity Name | Preferences     |
|-------------|-----------------|
| r1          | <b>h1</b> h2 h3 |
| r2          | h1 h3 h2        |
| r3          | h1 h2           |
| r4          | h2 h3           |
| h1 (1)      | r2 r3           |
| h2 (1)      | r4 r3           |
| h3 (1)      | r1 r2 r4        |

Next, r2 will propose to its top choice of h1. Since h1 is full and r2 is still on the hospital's preference list, h1 will find its lowest preferred resident that is assigned to it (r1) and unassign it to make room to r2. Once r2 gets assigned, all residents that are below r2 on h1's preference list are removed and also remove h1 from their list.

| Entity Name | Preferences         |
|-------------|---------------------|
| r1          | <del>h1</del> h2 h3 |
| r2          | <b>h1</b> h3 h2     |
| r3          | <del>h1</del> h2    |
| r4          | h2 h3               |
| h1 (1)      | r2 <del>r3</del>    |
| h2 (1)      | r4 r3               |
| h3 (1)      | r1 r2 r4            |

It is now r3's turn to propose and, since h1 was removed from its list in the last step, it will turn to h2 for an assignment. Since r3 is on h2's preference list and the hospital is not at capacity, r3 will be assigned to h2. The hospital is now full, so all residents below r3 on h2's list get removed and also remove h2 from their respective lists. The updated table is shown below.

| Entity Name | Preferences                    |
|-------------|--------------------------------|
| r1          | <del>h1</del> <del>h2</del> h3 |
| r2          | <b>h1</b> h3 <del>h2</del>     |
| r3          | <del>h1</del> <b>h2</b>        |
| r4          | h2 h3                          |
| h1 (1)      | r2 <del>r3</del>               |
| h2 (1)      | r4 r3                          |
| h3 (1)      | r1 r2 r4                       |

r4 will now propose to its top choice of h2, which still has r4 on its list. Since h2 is already full, r3 will have to be removed prior to handling r4's assignment. Once r4 gets assigned to h2, all residents below it on h2's list will be removed and will also remove h2 from their lists.

| Entity Name | Preferences                    |
|-------------|--------------------------------|
| r1          | <del>h1</del> <del>h2</del> h3 |
| r2          | <b>h1</b> h3 <del>h2</del>     |
| r3          | <del>h1</del> <del>h2</del>    |
| r4          | <b>h2</b> h3                   |
| h1 (1)      | r2 <del>r3</del>               |
| h2 (1)      | r4 <del>r3</del>               |
| h3 (1)      | r1 r2 r4                       |

The resident queue goes back to r1, which only has h3 left in its preference list. Thus, r1 will propose to h3 and all resident and hospital lists will be adjusted accordingly as h3 as it will now be full with r1 being assigned to it.

| Entity Name | Preferences                           |
|-------------|---------------------------------------|
| r1          | <del>h1</del> <del>h2</del> <b>h3</b> |
| r2          | <b>h1</b> <del>h3</del> <del>h2</del> |
| r3          | <del>h1</del> <del>h2</del>           |
| r4          | <b>h2</b> <del>h3</del>               |
| h1 (1)      | r2 <del>r3</del>                      |
| h2 (1)      | r4 <del>r3</del>                      |
| h3 (1)      | r1 <del>r2</del> <del>r4</del>        |

Lastly, r3 is reached again, but none of its preferences are available, so it will remain unassigned. The algorithm will then terminate in the state presented above.

## 2 VARIATION OF HOSPITALS AND RESIDENTS STABLE MATCHING PROBLEM

### 2.1 THE ALGORITHM

One variation of the hospitals and residents stable matching problem is when only the residents rank the hospitals and the hospitals do not rank the residents. When making the assignments, stability is defined by not having a resident who is assigned to one hospital that would rather be in another hospital that is still available. Similar to the original algorithm, hospital capacity would be nice to have, but is not a necessity if not enough residents prefer the hospital to meet its capacity. In other words, overall resident happiness is the most important variable to consider when making the assignments. As displayed in Algorithm 2, the method needed to make the assignments requires a greedy approach and taking the first possible outcome that works for the data. It is for this reason that each resident is initially assigned to their top choice hospital and then

gets trimmed down to meet the respective capacity of each hospital. The asymptotic analysis of the runtime of this algorithm is explored in detail in Section 2.2.

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**Algorithm 2** Hospitals and Residents Stable Matching Algorithm

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```

1: procedure STABLEMATCHVARIATION(residents, hospitals)
2:   for r of residents do
3:     topHosp  $\leftarrow$  r.preferences.head
4:     r.assignment  $\leftarrow$  topHosp // Residents start with their top choice regardless of capacity
5:   end for
6:   hasToCheck  $\leftarrow$  true
7:   while hasToCheck do
8:     residentsToReassign  $\leftarrow$  newQueue
9:     for h of hospitals do
10:      h.sortResidents() // Sort the residents from worst to best for being able to leave
11:      while h.isOverCapacity() do
12:        residentToRemove  $\leftarrow$  h.assignments[h.numAssigned - 1] // Get best one to reassign
13:        h.numAssigned  $\leftarrow$  h.numAssigned - 1
14:        residentToRemove.curPreferenceIndex  $\leftarrow$  residentToRemove.curPreferenceIndex + 1
15:        residentsToReassign.enqueue(residentToRemove) // Add it to the queue
16:      end while
17:    end for
18:    if !residentsToReassign.isEmpty() then
19:      while !residentsToReassign.isEmpty() do
20:        res  $\leftarrow$  residentsToReassign.dequeue()
21:        if res.curPreferenceIndex < NUM_LEVELS then // Add to next preference
22:          res.getHospitalPreferences[res.curPreferenceIndex].addResident(res)
23:        end if
24:      end while
25:    else
26:      hasToCheck  $\leftarrow$  false // We have a stable situation
27:    end if
28:  end while
29:  computeSwaps(residents) // Perform swaps to get best possible outcome
30: end procedure

```

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## 2.2 ASYMPTOTIC ANALYSIS

Listing 3 provides the C++ implementation of the variation algorithm to the stable matching problem in addition to the swapping algorithm in Listing 5 and how the residents compare themselves to each other in Listing 4. Before going through the algorithm, in addition to  $r$  being the number of residents and  $h$  being the number of hospitals,  $p$  will be used as the notation for the number of preferences ( $p \leq h$ ) as all residents get a fixed number of hospital preferences. First, lines 5-8 of the variation algorithm contains a for-loop that provisionally assigns all residents to their top choice hospital, which is an  $O(r)$  operation as the actual adding is a constant time operation that runs  $r$  times. Next, line 14 begins a while-loop that is dependent on the *hasToCheck* variable. The number of times the loop will run is unknown for now and will be temporarily noted with an  $x$ . Inside the loop, after created a linked list on line 17, there is a loop that iterates through all of the hospitals. First, the provisionally assigned residents get sorted by the availability of their lower preference hospitals using a quicksort algorithm. The sort itself runs in  $O(r \log_2 r)$  time, which means the residents are compared about  $r \log_2 r$  times.

The comparison algorithm in Listing 4 presents how each resident is compared with one another. After

doing 2  $O(p)$  operations by iterating through the hospitals in the preference lists for both residents on lines 7-21 and some assignments on lines 24 and 25, both the if-statement branch and else-if branch contain the same bodies, but for different residents. The main section is the nested for-loop on lines 37-53. The outer loop iterates through the resident's preferences and the middle loop goes through each resident provisionally assigned to the hospital from the outer loop. Overall, this will run at worst case  $r - 1$  times as all other residents will be compared to each other. The inner-most loop considers the remaining preferences for the middle man being looked at, which can be  $p - 1$  iterations at worst case. The body of the inner-most loop contains assignments and calculations and runs in  $O(1)$  time. Thus, the nested loops as a whole run in  $O(r * p)$  time because it would have to go through all preferences for all residents. Overall, the comparison algorithm runs in  $O(2p + rp)$  time, which is  $O(rp)$  as  $rp$  is the dominant term. Therefore the sorting algorithm runs in  $O(rp * r * \log_2 r)$  is  $O(pr^2 \log_2 r)$ .

Once the residents are sorted on line 22 of Listing 2, the while-loop on lines 25-36 iterates until the hospital is no longer over capacity, which will be around  $r$  times. The body of the while-loop has several assignments, which run in constant time. Thus, the overall loop on lines 20-37 runs in  $O(h * (pr^2 \log_2 r + r))$  is  $O(hpr^2 \log_2 r + hr)$ , which simplifies to  $O(hpr^2 \log_2 r)$ . Next, the body of the if-block defined on lines 40-52 contains a while-loop that goes through all of the residents that were removed from their hospital and reassigns them to their next preference, which is an  $O(r)$  operation. However, the residents are only able to be reassigned if there is still a hospital in their preference list, which means that this part of the code will only be called about  $p$  times, which is the number of times the big while-loop runs and the value of  $x$ . Therefore, the while-loop on lines 15-56 runs in  $O(p * (hpr^2 \log_2 r + r))$ , which is  $O(hp^2 r^2 \log_2 r)$  as  $pr$  is not dominant and can be removed for simplification purposes. Next, the loops on lines 59-63 formally assign each resident to a hospital, which is another  $O(r)$  operation.

The final part of the algorithm is on line 85, which calls the swapping function. This function can be found in Listing 5. Line 4 of the swapping function defines a loop that runs until no swaps are made in an iteration. The body of the while-loop contains nested for-loops that each go through all residents, which means the inner body runs about  $r^2$  times. The entire body, however, is only comprised of assignments, math operations, and comparisons, which are all constant time operations. Lastly, the if-else block on lines 74-92 also has assignments in the body, which runs in constant time. Overall, the swapping algorithm runs in  $O(r^2)$  time.

The runtime for the variation algorithm is  $O(hp^2 r^2 \log_2 r + r + r^2)$ , which is  $O(hp^2 r^2 \log_2 r)$ .

## 2.3 EXAMPLE AND WALKTHROUGH

Listing 6 contains a sample data set for the variation algorithm that contains 8 residents, 7 hospitals, and 6 preferences per resident. The first step of the algorithm is to assign each resident to their top choice hospital. The initial assignments will be as follows:

| Hospital | Capacity | Residents |
|----------|----------|-----------|
| h1       | 1        | r1        |
| h2       | 1        | r2, r6    |
| h3       | 1        | r3        |
| h4       | 2        | r4        |
| h5       | 1        | r5        |
| h6       | 1        | r7        |
| h7       | 1        | r8        |

As shown in the table of initial results, h2 is over capacity as it has 2 residents, r2 and r6, provisionally assigned to the hospital even though h2 has a capacity of only 1. Therefore, when comparing r2 and r6 for the sort, r6 is initially favored to leave h2 because it has a preference of h4, which has availability, and r2 does not. However, through the loop structure on lines 37-53 in Listing 4, it is determined that r3 can act as a middle man and would be able to fill the open space in h4 with r2 taking the r3's spot in h3. This would

create a better overall resident happiness as 2 residents being in their second choice is better than 1 resident being in their top choice and the other being in their sixth choice. Therefore, r2 will be added to the list of residents to be reassigned and r6 stays put in h2. None of the other hospitals are over capacity, so there are no other changes. r2 will then be assigned to h3, which is its second choice hospital, which will cause the assignments to be as shown in the table below.

| Hospital | Capacity | Residents |
|----------|----------|-----------|
| h1       | 1        | r1        |
| h2       | 1        | r6        |
| h3       | 1        | r3, r2    |
| h4       | 2        | r4        |
| h5       | 1        | r5        |
| h6       | 1        | r7        |
| h7       | 1        | r8        |

In this next iteration, h3 is now over capacity. When comparing r3 and r2, r3 has a preference for h4, which has an availability, and none of r2's preferences are open, so r3 is prioritized to move out of h3 and will be reassigned to h4. The assignment table will look as follows.

| Hospital | Capacity | Residents |
|----------|----------|-----------|
| h1       | 1        | r1        |
| h2       | 1        | r6        |
| h3       | 1        | r2        |
| h4       | 2        | r4, r3    |
| h5       | 1        | r5        |
| h6       | 1        | r7        |
| h7       | 1        | r8        |

In the current state, all hospitals are at capacity and all residents have matched the definition of stability by being assigned only to hospitals on their preference list. Additionally, there is no swap between 2 residents that will make the overall resident happiness increase and, therefore, the algorithm terminates in this state.

### 3 APPENDIX

#### 3.1 ORIGINAL ALGORITHM

Listing 1: Original Stable Matching Algorithm (C++)

```

1 void generateStableMatches(ResidentArr* residents, HospitalArr* hospitals) {
2     Queue<Resident*> residentQueue;
3
4     // Add the residents to a queue for determining the next resident to propose
5     for (int i = 0; i < residents->length; i++) {
6         Node<Resident*>* n = new Node(&residents->arr[i]);
7         residentQueue.enqueue(n);
8     }
9
10    std::cout << std::endl;
11
12    while (!residentQueue.isEmpty()) {
13        residentQueue.printQueue();
14
15        Node<Resident*>* resident = residentQueue.dequeue();
16
17        while (resident->data->getAssignment() == nullptr && !resident->data->
18            getHospitalPreferences()->isEmpty()) {
19            for (int j = 0; j < hospitals->length; j++) {

```

```

19         std::cout << resident->data->getPreferencesArr()[j] << "\n";
20     }
21     std::cout << std::endl;
22
23     // Get the hospital at the front of the preference list
24     Node<Hospital*>* preference = resident->data->getHospitalPreferences()->dequeue
        ();
25
26     // The hospital is full with more preferable residents, so the current resident
        is not getting in
27     if (resident->data->getPreferencesArr()[preference->data->getIndex()] == 0) {
28         // Clean up memory and try the next hospital
29         delete preference;
30         continue;
31     }
32
33     if (preference->data->isFull()) {
34         // Get the lowest preferred assigned resident index
35         preference->data->setLowestPreferredAssignedResidentIndex();
36
37         // If the hospital is full, then replace the lowest preferred resident with
        the current resident
38         Resident* lowest = &residents->arr[preference->data->getAssignedResidents()[
        preference->data->getLowestPreferredAssignedResidentIndex()].getIndex()
        ];
39         lowest->setAssignment(nullptr);
40
41         preference->data->replaceLowest(resident->data);
42         // Since the lowest preferred resident was replaced, we need to come back to
        it later
43         residentQueue.enqueue(new Node<Resident*>(lowest));
44     } else {
45         // Add the new resident to the
46         preference->data->addResident(resident->data);
47     }
48
49     // Print the assignment as it happens
50     std::cout << "Assigned_" << resident->data->getName() << "_to_" << resident->
        data->getAssignment()->getName() << std::endl;
51
52     if (preference->data->isFull()) {
53         // Get the lowest preferred assigned resident index
54         preference->data->setLowestPreferredAssignedResidentIndex();
55
56         // We need to iterate through all of the residents
57         for (int i = 0; i < residents->length; i++) {
58             // Check if the current resident is less preferred compared to the
        lowest preferred resident
59             if (preference->data->getResidentPreferences()[residents->arr[i].
        getIndex()] > preference->data->getResidentPreferences()[preference
        ->data->getAssignedResidents()[preference->data->
        getLowestPreferredAssignedResidentIndex()].getIndex()]) {
60
61                 // Take the resident out of the running if that is the case
62                 residents->arr[i].getPreferencesArr()[preference->data->getIndex()]
        = 0;
63                 preference->data->getResidentPreferences()[residents->arr[i].
        getIndex()] = INT_MAX;
64             }
65         }
66     }
67
68     // Clean up memory because the node is no longer needed
69     delete preference;
70 }

```



```

71
72     // The resident is done for now and a new node has already been created if needed
73     delete resident;
74 }
75
76 std::cout << "Finished_algo" << std::endl << std::endl;
77
78 // Print the final results
79 std::cout << "Final_results:" << std::endl;
80 for (int i = 0; i < residents->length; i++) {
81     if (residents->arr[i].getAssignment() != nullptr) {
82         std::cout << "(" << residents->arr[i].getName() << ",_" << residents->arr[i].
            getAssignment()->getName() << ")" << std::endl;
83     } else {
84         std::cout << "(" << residents->arr[i].getName() << ",_nullptr)" << std::endl;
85     }
86 }
87 }

```

Listing 2: Original Algorithm Sample Test Data

```

1 r1: h1 h2 h3
2 r2: h1 h3 h2
3 r3: h1 h2
4 r4: h2 h3
5 h1: 1 - r2 r3
6 h2: 1 - r4 r3
7 h3: 1 - r1 r2 r4

```

### 3.2 VARIATION ALGORITHM

Listing 3: Variation Stable Matching Algorithm (C++)

```

1 void generateStableMatches(ResidentArr* residents, HospitalArr* hospitals) {
2     std::cout << std::endl;
3
4     // Iterate through all residents
5     for (int i = 0; i < residents->length; i++) {
6         Hospital* cur = residents->arr[i].getHospitalPreferences()[residents->arr[i].
            getCurPreferenceIndex()];
7
8         // Add each resident to their top choice hospital
9         cur->addResident(&residents->arr[i]);
10    }
11
12    // We have to make sure that the loop iterates at least once, so initialize the variable
        to true
13    bool hasToCheck = true;
14
15    while (hasToCheck) {
16        // Create a linked list of the residents that get booted from their current choice
            hospital
17        List<Resident*> residentsToReassign;
18
19        // Iterate through all of the hospitals
20        for (int i = 0; i < hospitals->length; i++) {
21            // Sort the residents by keeping the best one to leave at the end of the list
22            hospitals->arr[i].sortResidentAssignments();
23
24            // Continue until the hospital is no longer over capacity
25            while (hospitals->arr[i].getNumAssigned() - hospitals->arr[i].getCapacity() > 0)
                {
26                hospitals->arr[i].printAssignments();
27
28                // The last one in the assignment array will get booted

```

```

29         Resident* residentToRemove = hospitals->arr[i].getAssignments()[hospitals->
30             arr[i].getNumAssigned() - 1];
31         hospitals->arr[i].setNumAssigned(hospitals->arr[i].getNumAssigned() - 1);
32         // Update the resident preference to try the next hospital on its list and
33             add it to the linked list
34         residentToRemove->setCurPreferenceIndex(residentToRemove->
35             getCurPreferenceIndex() + 1);
36         Node<Resident*>* resNode = new Node<Resident*>(residentToRemove);
37         residentsToReassign.enqueue(resNode);
38     }
39 }
40 // Check to make sure there are residents that have to be reassigned
41 if (residentsToReassign.getSize() > 0) {
42     while (!residentsToReassign.isEmpty()) {
43         // Remove the residents one by one
44         Node<Resident*>* res = residentsToReassign.dequeue();
45         // Add the resident to its next choice if possible
46         if (res->data->getCurPreferenceIndex() < NUM_PREFERENCES) {
47             res->data->getHospitalPreferences()[res->data->getCurPreferenceIndex()]
48                 ->addResident(res->data);
49         }
50         delete res;
51     }
52 } else {
53     // If the linked list was empty, then we know all residents are in a stable
54     position
55     hasToCheck = false;
56 }
57 }
58 // Formally assign each resident to the hospitals
59 for (int i = 0; i < hospitals->length; i++) {
60     for (int j = 0; j < hospitals->arr[i].getNumAssigned(); j++) {
61         hospitals->arr[i].getAssignments()[j]->setAssignment(&hospitals->arr[i]);
62     }
63 }
64
65 std::cout << "Finished_main_algo" << std::endl << std::endl;
66
67 // Print the final results
68 std::cout << "Initial_results:" << std::endl;
69 for (int i = 0; i < residents->length; i++) {
70     if (residents->arr[i].getAssignment() != nullptr) {
71         std::cout << "(" << residents->arr[i].getName() << ",_" << residents->arr[i].
72             getAssignment()->getName() << ")" << std::endl;
73     } else {
74         std::cout << "(" << residents->arr[i].getName() << ",_nullptr)" << std::endl;
75     }
76 }
77 std::cout << std::endl;
78
79 // Compute the happiness indices for both residents and hospitals
80 std::cout << "Resident_Happiness:_" << computeResidentHappiness(residents) << std::endl;
81 std::cout << "Hospital_Happiness:_" << computeHospitalHappiness(hospitals) << std::endl;
82
83 std::cout << std::endl;
84
85 // Try to do swaps to boost the performance of the algorithm
86 performSwaps(residents);
87
88 // Print the final results

```

```

88     std::cout << std::endl;
89     std::cout << "Final_results:" << std::endl;
90     for (int i = 0; i < residents->length; i++) {
91         if (residents->arr[i].getAssignment() != nullptr) {
92             std::cout << "(" << residents->arr[i].getName() << ",_" << residents->arr[i].
                getAssignment()->getName() << ")" << std::endl;
93         } else {
94             std::cout << "(" << residents->arr[i].getName() << ",_nullptr)" << std::endl;
95         }
96     }
97     std::cout << std::endl;
98
99     // Compute the happiness indices for both residents and hospitals
100    std::cout << "Resident_Happiness:_" << computeResidentHappiness(residents) << std::endl;
101    std::cout << "Hospital_Happiness:_" << computeHospitalHappiness(hospitals) << std::endl;
102 }

```

Listing 4: Resident Comparison Function (C++)

```

1 int Resident::compare(Resident* compResident) {
2     // Get the current hospitals being compared
3     int thisCur = curPreferenceIndex + 1;
4     int otherCur = compResident->getCurPreferenceIndex() + 1;
5
6     // Get the first hospital in the preference list that is still available
7     while (thisCur < NUM_PREFERENCES) {
8         if (hospitalPreferences[thisCur]->isFull()) {
9             thisCur++;
10        } else {
11            break;
12        }
13    }
14
15    while (otherCur < NUM_PREFERENCES) {
16        if (compResident->getHospitalPreferences()[otherCur]->isFull()) {
17            otherCur++;
18        } else {
19            break;
20        }
21    }
22
23    // Compute the difference if each resident went to their next available choice
24    int thisDiff = thisCur - curPreferenceIndex;
25    int otherDiff = otherCur - compResident->getCurPreferenceIndex();
26
27    // Assume they are the same
28    int out = 0;
29
30    if (thisDiff < otherDiff) {
31        // Lower difference should be put at the end of the list
32        out = 1;
33
34        Hospital* target = hospitalPreferences[thisCur];
35        bool found = false;
36        // Iterate back through all possible moves to go down for the other resident
37        for (int i = compResident->getCurPreferenceIndex() + 1; i < NUM_PREFERENCES && !
            found; i++) {
38            // Iterate through all the residents assigned to the hospital
39            for (int j = 0; j < compResident->getHospitalPreferences()[i]->getNumAssigned()
                && !found; j++) {
40                Resident* middleMan = compResident->getHospitalPreferences()[i]->
                    getAssignments()[j];
41                for (int k = middleMan->getCurPreferenceIndex() + 1; k < NUM_PREFERENCES &&
                    !found; k++) {
42                    if (middleMan->getHospitalPreferences()[k] == target) {

```

```

43         int middleDiff = k - middleMan->getCurPreferenceIndex();
44         int newOtherDiff = i - compResident->getCurPreferenceIndex();
45         double avg = (double) (middleDiff + newOtherDiff) / 2;
46         if (avg < thisDiff) {
47             out = -1;
48             found = true;
49         }
50     }
51 }
52 }
53 }
54 } else if (thisDiff > otherDiff) {
55     // Higher difference should be kept near the front
56     out = -1;
57
58     Hospital* target = compResident->getHospitalPreferences()[otherCur];
59     bool found = false;
60     // Iterate back through all possible moves to go down for the other resident
61     for (int i = curPreferenceIndex + 1; i < NUM_PREFERENCES && !found; i++) {
62         // Iterate through all the residents assigned to the hospital
63         for (int j = 0; j < hospitalPreferences[i]->getNumAssigned() && !found; j++) {
64             Resident* middleMan = hospitalPreferences[i]->getAssignments()[j];
65             for (int k = middleMan->getCurPreferenceIndex() + 1; k < NUM_PREFERENCES &&
66                 !found; k++) {
67                 if (middleMan->getHospitalPreferences()[k] == target) {
68                     int middleDiff = k - middleMan->getCurPreferenceIndex();
69                     int newThisDiff = i - curPreferenceIndex;
70                     double avg = (double) (middleDiff + newThisDiff) / 2;
71                     if (avg < thisDiff) {
72                         out = -1;
73                         found = true;
74                     }
75                 }
76             }
77         }
78     } else {
79         // If the differences are the same, prioritize keeping the one that is happier as it
80         // is
81         if (curPreferenceIndex < compResident->getCurPreferenceIndex()) {
82             out = -1;
83         } else if (curPreferenceIndex > compResident->getCurPreferenceIndex()) {
84             out = 1;
85         }
86     }
87     return out;
88 }

```

Listing 5: Swapping Function (C++)

```

1 void performSwaps(ResidentArr* residents) {
2     // Make any needed adjustments to increase resident happiness
3     bool swapped = true;
4     while (swapped) {
5         // Initialize the residents to swap to be nothing
6         Resident* res1 = nullptr;
7         Resident* res2 = nullptr;
8         double bestDiff = 0;
9
10        for (int i = 0; i < residents->length; i++) {
11            for (int j = i + 1; j < residents->length; j++) {
12                // Get the current assignments
13                Hospital* iHosp = residents->arr[i].getAssignment();
14                Hospital* jHosp = residents->arr[j].getAssignment();

```

```

15
16     bool canSwap = true;
17     // Make sure jHosp is a preference for resident i
18     if (jHosp != nullptr && residents->arr[i].getPreferencesArr()[jHosp->
19         getIndex()] == 0) {
20         // Cannot swap if resident i does not want to go to jHosp
21         canSwap = false;
22     } else if (iHosp != nullptr && residents->arr[j].getPreferencesArr()[iHosp->
23         getIndex()] == 0) {
24         // Same but for resident j and iHosp
25         canSwap = false;
26     }
27
28     if (canSwap) {
29         // Compute the current average happiness among the 2 residents
30         int iCurHappiness = 0;
31         if (iHosp != nullptr) {
32             iCurHappiness = residents->arr[i].getPreferencesArr()[iHosp->
33                 getIndex()];
34         }
35
36         int jCurHappiness = 0;
37         if (jHosp != nullptr) {
38             jCurHappiness = residents->arr[j].getPreferencesArr()[jHosp->
39                 getIndex()];
40         }
41
42         double curHappiness = (double) (iCurHappiness + jCurHappiness) / 2;
43
44         // Compute the average happiness if the 2 residents swapped
45         int iSwapHappiness = 0;
46         if (jHosp != nullptr) {
47             iSwapHappiness = residents->arr[i].getPreferencesArr()[jHosp->
48                 getIndex()];
49         }
50
51         int jSwapHappiness = 0;
52         if (iHosp != nullptr) {
53             jSwapHappiness = residents->arr[j].getPreferencesArr()[iHosp->
54                 getIndex()];
55         }
56
57         double swapHappiness = (double) (iSwapHappiness + jSwapHappiness) / 2;
58
59         // Make sure the new happiness is better than it was before
60         if (swapHappiness > curHappiness) {
61             // Get the gained happiness
62             double diff = swapHappiness - curHappiness;
63
64             // Check to see if it is better than the current best
65             if (diff > bestDiff) {
66                 // Update the variables
67                 bestDiff = diff;
68                 res1 = &residents->arr[i];
69                 res2 = &residents->arr[j];
70             } else if (diff == bestDiff && rand() % 2 == 0) {
71                 // If equal, randomly pick one or the other to use as the best
72                 // difference
73                 bestDiff = diff;
74                 res1 = &residents->arr[i];
75                 res2 = &residents->arr[j];
76             }
77         }
78     }
79 }

```

```

73     // Make the swap if a possible swap was found
74     if (res1 != nullptr && res2 != nullptr) {
75         // Get the hospitals being used
76         Hospital* h1 = res1->getAssignment();
77         Hospital* h2 = res2->getAssignment();
78
79         // Replace the residents
80         h1->replace(res1, res2);
81         h2->replace(res2, res1);
82
83         // Update the assignments
84         res2->setAssignment(h1);
85         res1->setAssignment(h2);
86
87         std::cout << "Swapped_" << res1->getName() << "_and_" << res2->getName() << std
            ::endl;
88     } else {
89         // No swaps were made, so can set the flag to false and end the loop
90         std::cout << "No_swaps_made" << std::endl;
91         swapped = false;
92     }
93 }
94 }

```

Listing 6: Variation Algorithm Sample Test Data

```

1  r1: h1 h3 h2 h4 h5 h7
2  r2: h2 h3 h6 h5 h1 h7
3  r3: h3 h4 h1 h2 h5 h7
4  r4: h4 h5 h1 h2 h3 h7
5  r5: h5 h4 h1 h2 h3 h7
6  r6: h2 h1 h3 h5 h7 h4
7  r7: h6 h1 h2 h3 h4 h7
8  r8: h7 h1 h2 h3 h4 h5
9  h1: 1
10 h2: 1
11 h3: 1
12 h4: 2
13 h5: 1
14 h6: 1
15 h7: 1

```