When a query q is asked if a matching fact q' is known, return the unifier for each rule whose consequent q' matches q attempt to prove each premise of the rule by backward chaining

(Some added complications in keeping track of the unifiers)

(More complications help to avoid infinite loops)

Two versions: find <u>any</u> solution, find <u>all</u> solutions

Backward chaining is the basis for logic programming, e.g., Prolog

A simple example

- B^C=> G
- A^G=> I
- D^G=>J
- E=> C
- D^C=>K
- F=>C
- Q: I?

A simple example

- B^C=> G
- A^G=> I
- D^G=>J
- E=> C
- D^C=>K
- F=>C
- Q: I?

- 1. A^G
- 2. A?
 - 1. USER
- 3. G?
 - 1. B^C
 - 1. USER
 - 2. E v F

- Current knowledge:
 - hurts(x, head)
- What implications can lead to this fact?
 - kicked(x, head)
 - fell_on(x, head)
 - brain_tumor(x)
 - hangover(x)
- What facts do we need in order to prove these?

- The algorithm (available in detail in AIMA text):
 - a knowledge base KB
 - a desired conclusion c or question q
 - finds all sentences that are answers to q in KB *or* proves c
 - if q is directly provable by premises in KB, infer q and remember how q was inferred (building a list of answers).
 - find all implications that have q as a consequent.
 - for each of these implications, find out whether all of its premises are now in the KB, in which case infer the consequent and add it to the KB, remembering how it was inferred. If necessary, attempt to prove the implication also via backward chaining
 - premises that are conjuncts are processed one conjunct at a time

Backward chaining algorithm

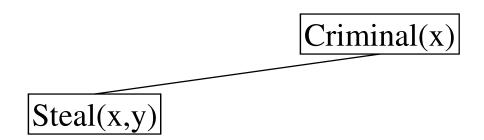
```
function FOL-BC-Ask(KB, goals, \theta) returns a set of substitutions inputs: KB, a knowledge base goals, a list of conjuncts forming a query \theta, the current substitution, initially the empty substitution \{\} local variables: ans, a set of substitutions, initially empty if goals is empty then return \{\theta\} q' \leftarrow \text{SUBST}(\theta, \text{FIRST}(goals)) for each r in KB where STANDARDIZE-APART(r) = (p_1 \land \ldots \land p_n \Rightarrow q) and \theta' \leftarrow \text{UNIFY}(q, q') succeeds ans \leftarrow \text{FOL-BC-Ask}(KB, [p_1, \ldots, p_n | \text{Rest}(goals)], \text{Compose}(\theta, \theta')) \cup ans return ans
```

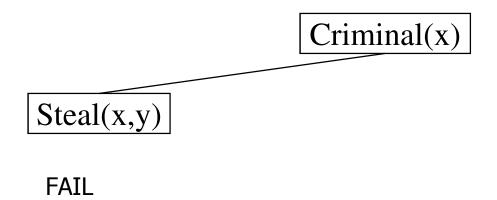
SUBST(COMPOSE(
$$\theta_1, \theta_2$$
), p) = SUBST(θ_2 , SUBST(θ_1, p))

- Question: Has Reality Man done anything criminal?
 - Criminal(Reality Man)
- Possible answers:
 - Steal(x, y) \Rightarrow Criminal(x)
 - $Kill(x, y) \Rightarrow Criminal(x)$
 - Grow(x, y) \land Illegal(y) \Rightarrow Criminal(x)
 - HaveSillyName(x) \Rightarrow Criminal(x)
 - Programmer(x) ∧ Emulator(y) ∧ People(z) ∧ Provide(x,z,y)
 ⇒Criminal(x)

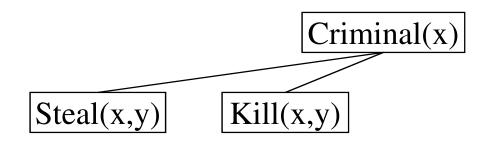
• Question: Has Reality Man done anything criminal?

Criminal(x)

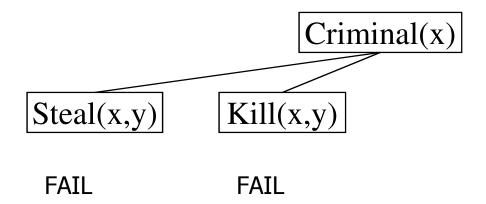


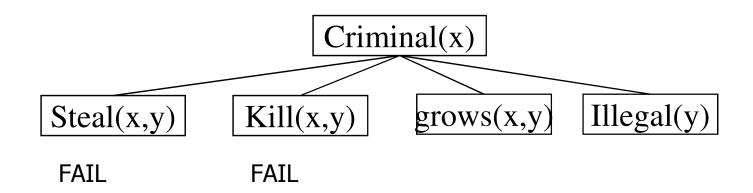


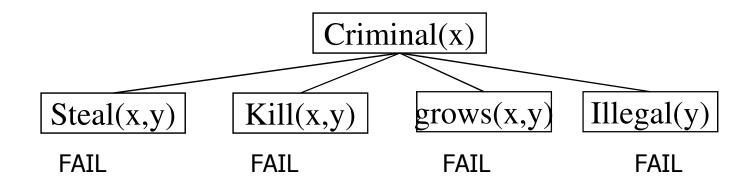
• Question: Has Reality Man done anything criminal?



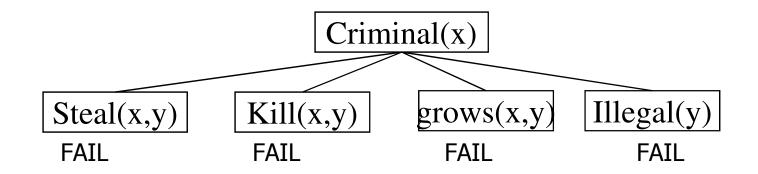
FAIL







Question: Has Reality Man done anything criminal?



 Backward Chaining is a depth-first search: in any knowledge base of realistic size, many search paths will result in failure.

- Question: Has Reality Man done anything criminal?
- We will use the same knowledge as in our forward-chaining version of this example:

```
Programmer(x) \land Emulator(y) \land People(z) \land Provide(x,z,y)\Rightarrow Criminal(x)

Use(friends, x) \land Runs(x, N64 games) \Rightarrow Provide(Reality Man, friends, x)

Software(x) \land Runs(x, N64 games) \Rightarrow Emulator(x)

Programmer(Reality Man)

People(friends)

Software(U64)

Use(friends, U64)

Runs(U64, N64 games)
```

• Question: Has Reality Man done anything criminal?

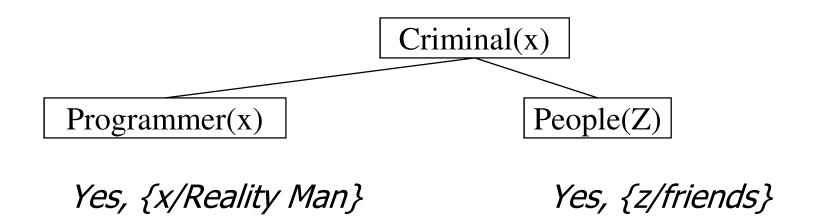
Criminal(x)

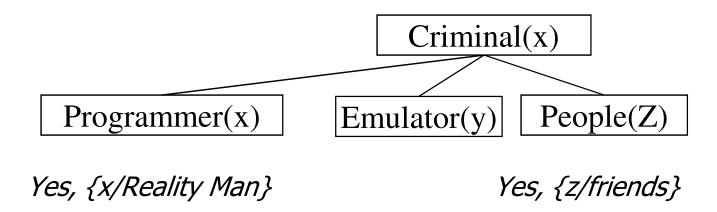
• Question: Has Reality Man done anything criminal?

Criminal(x)

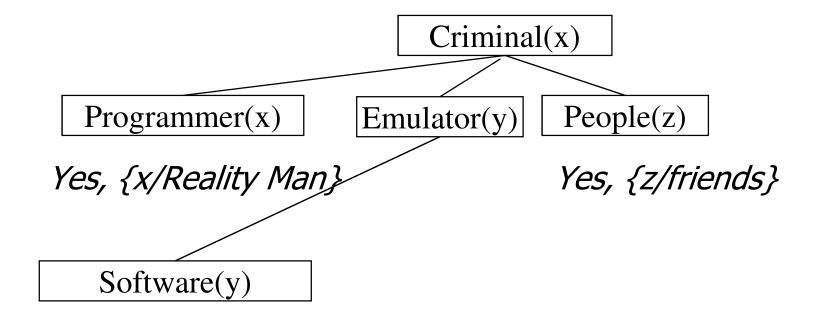
Programmer(x)

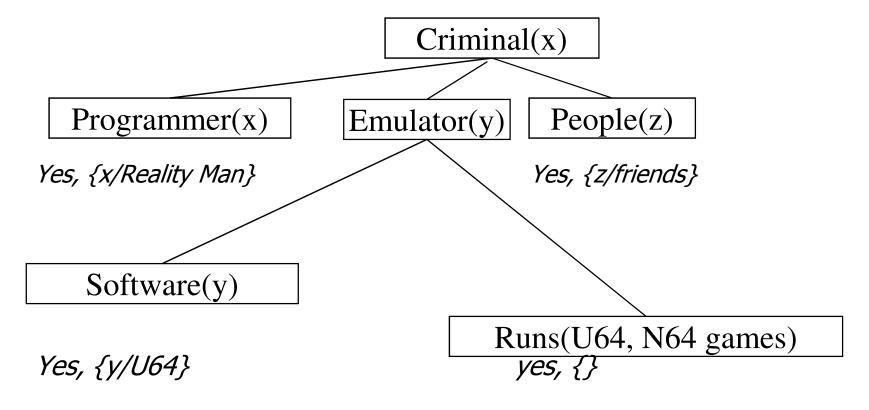
Yes, {x/Reality Man}

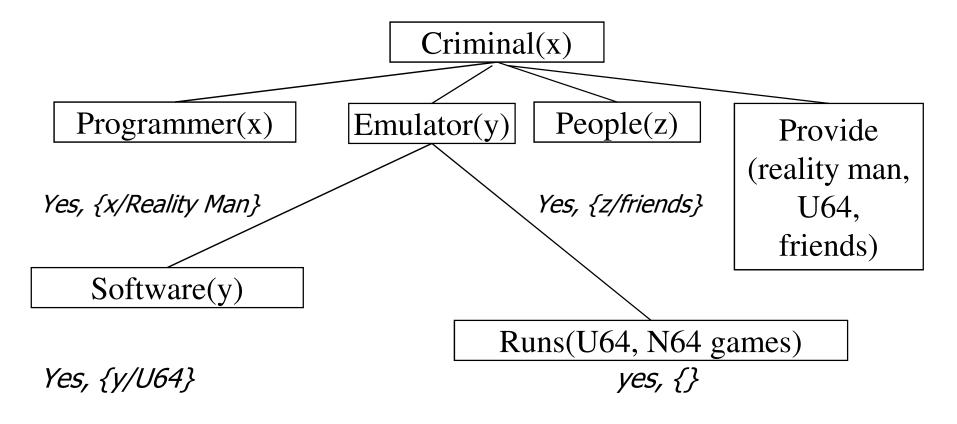


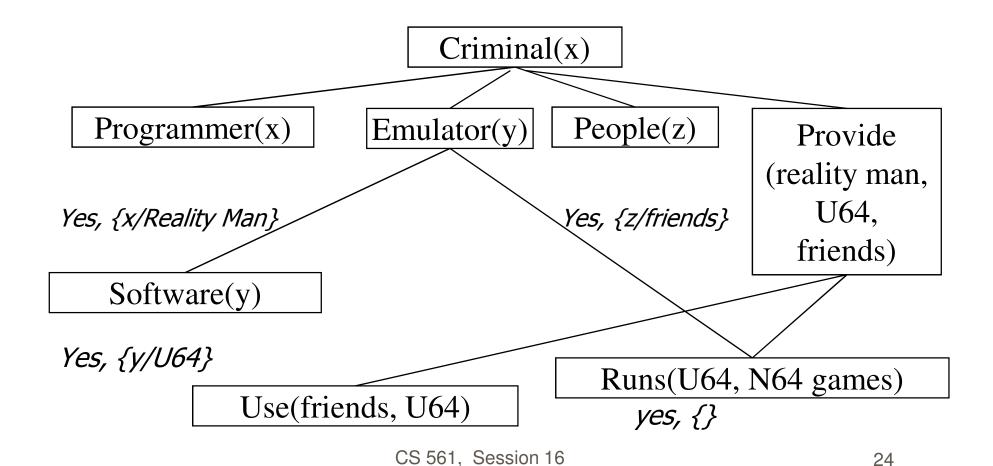


Yes, {*y/U64*}

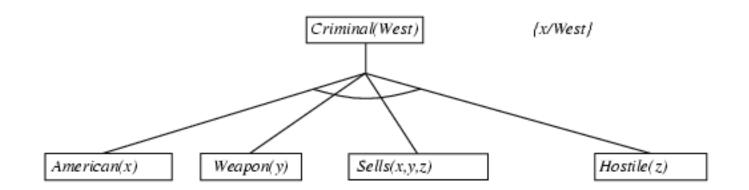


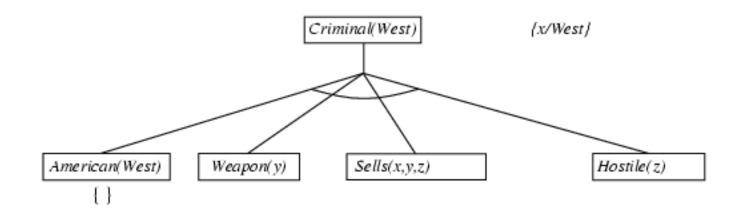


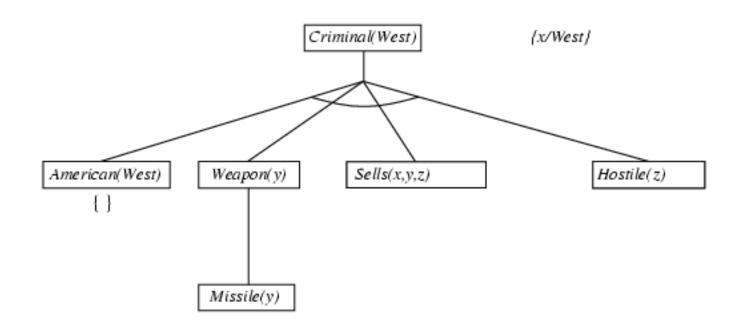


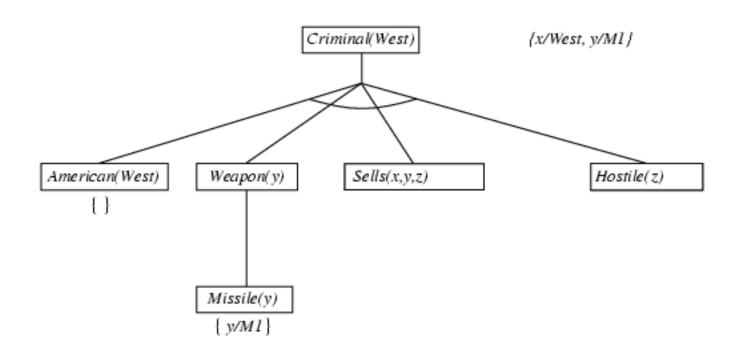


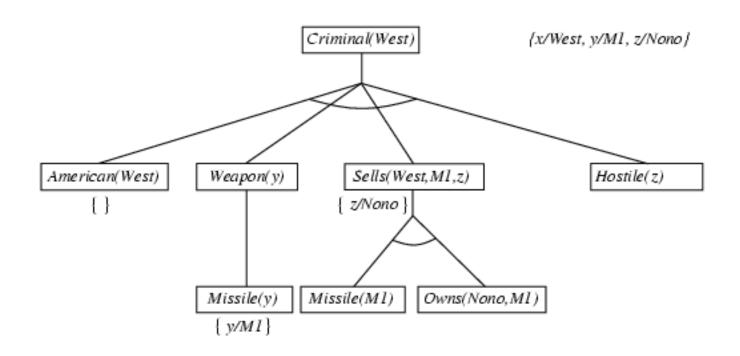
Criminal(West)

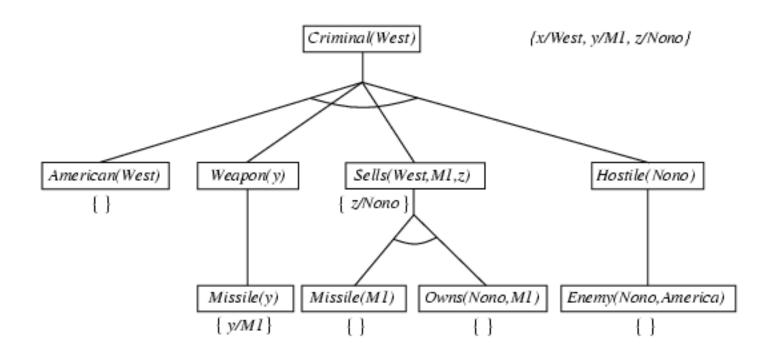


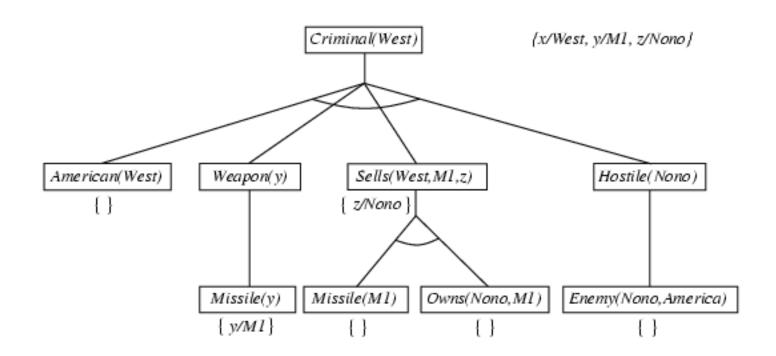












Properties of backward chaining

- Depth-first recursive proof search: space is linear in size of proof
- Incomplete due to infinite loops
 - → fix by checking current goal against every goal on stack
- Inefficient due to repeated subgoals (both success and failure)
 - → fix using caching of previous results (extra space)
- Widely used for logic programming