

Data Structures

Asymptotic Complexity (2)

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How skilled programmers find the order
VERY fast?

Tips

- This code involves ~8 steps
- Ignore all constants.
- They don't affect us
- All these are a FIXED number
- So code is just $O(1)$

```
5 void ConstantOrder1() {  
6     // O(1)  
7     int start = 6;  
8     int end = 100;  
9  
10    int mid = (end - start) / 2;  
11  
12    if (mid % 2 == 0)  
13        --mid;  
14 }
```

Tips

- Too many steps?
- Yah BUT FIXED
 - Useless with very large N
- Ignore them
- Tip: Ignore anything that doesn't involve our factor N

```
void ConstantOrder2() {  
    int start = 7;  
    int end = 0;  
  
    for (int i = 0; i < 1000; ++i)  
        end += end * 2 + start;  
}
```

Tips

- Search for loops that is based on n
- A single loop is $O(n)$
- Nested loop is $O(n^2)$
- Triple nest loop is $O(n^3)$
- And so on
- On right: single loop
 - Inside it fixed operations = IGNORE

```
void linear1(int n) { // O(n)
    int sum = 0;
    for (int i = 0; i < n; i++) {
        // All below are O(1)
        int x = 2 + 3 * 4;
        sum += i;
        sum += 2;
        sum += x;
    }
}
```

Tips

- 2 parallel loops. Each is single loop
- Each depends on n
 - One $10n$ and one is $5n$
 - Ignore these constants
- Practically: $10n + 5n = 17n$
- Ignore constants $\Rightarrow O(n)$
- Tip: what is the deepest?
 - A single loop $\Rightarrow O(n)$

```
void linear2(int n) { // O(n)
    for (int i = 0; i < 10 * n; i++)
        ConstantOrder1();

    for (int i = 0; i < 5 * n; i++)
        ConstantOrder1();
}
```

Tips

- This is $5n \times 3n$ loop steps
 - Multiplied with some factor from all these FIXED steps
 - Overall $O(n^2)$
- Tip: nested loops $\Rightarrow O(n^2)$

```
void quadratic1(int n) { //  $O(n^2)$ 
    int cnt = 0;
    for (int i = 0; i < 5 * n; ++i) {
        for (int j = 0; j < 3 * n; ++j) {
            cnt++;
            ConstantOrder1();
        }
    }
}
```

Tips

- We have 2 parallel things
 - Nested loops: $O(n^2)$
 - Linear loop: $O(n)$
- Tip: focus on the biggest
 - As it dominates
 - $n^2 + n \Rightarrow n^2$

```
void quadratic2(int n) { //  $O(n^2)$ 
    int cnt = 0;
    for (int i = 0; i < 5 * n; ++i) {
        for (int j = 0; j < 3 * n; ++j) {
            cnt++;
            ConstantOrder1();
        }
    }
    for (int i = 0; i < 10 * n; i++)
        ConstantOrder1();
}
```


Tips

- Again 2 parallel things
 - 3 nested loop
 - 1 loop
- But in 3 nested loop
 - One loop is just fixed operation
 - Again ignore constant operations
 - This 3rd loop is useless
- Total: $15000 n^2 + 10n \Rightarrow n^2$

```
void quadratic3(int n) { // O(n^2)
    int cnt = 0;
    for (int i = 0; i < 5 * n; ++i) {
        for (int j = 0; j < 3 * n; ++j) {
            for (int k = 0; k < 1000; ++k) {
                cnt++;
                ConstantOrder1();
            }
        }
    }
    for (int i = 0; i < 10 * n; i++)
        ConstantOrder1();
}
```

Tips

- 2 parallel blocks
 - Singel loop
 - Single loop
- Ok then $O(n)$? No, there is a trick
- The 2nd loop is not linear in n
 - It moves $3n^2$ steps
- The order of the second loop is $O(n^2)$
- Tip: observe if the loop based on fixed, n , n^2 or what
 - Its value decides its order!

```
void quadratic4(int n) { //  $O(n^2)$ 
    for (int i = 0; i < 10 * n; i++)
        ConstantOrder1();

    for (int i = 0; i < 3 * n * n; i++)
        ConstantOrder1();
}
```

Tips

- As this code has 3 nested loops
 - Each depends on n
 - It is $O(n^3)$

```
void cubic1(int n = 1000) { // O(n^3)
    int cnt = 0;
    for (int i = 0; i < n; ++i) {
        for (int j = 0; j < n; ++j) {
            for (int k = 0; k < n; ++k) {
                cnt++;
            }
        }
    }
}
```

Tips

- 2 parallel loops
 - 3 nested loops $\Rightarrow n^3$
 - 2 nested loops $\Rightarrow n^2$
 - Don't be cheated by 1000 constant
 - IGNORE constants
- $n^3 + n^2 \Rightarrow O(n^3)$
 - Always focus on the biggest

```
void cubic2(int n) { // O(n^3)
    int cnt = 0;
    for (int i = 0; i < n; ++i)
        for (int j = 0; j < n; ++j)
            for (int k = 0; k < n; ++k)
                cnt++;

    for (int i = 0; i < 1000*n; ++i)
        for (int j = 0; j < 1000*n; ++j)
            cnt++;
}
```

Tips

- Why not $O(n^3)$?
 - First loop is n^2 , then n , then n^3
 - Total $O(n^6)$
- Again double check if loop is based on n or what?

```
void f(int n) { //  $O(n^6)$ 
    int cnt = 0;
    for (int i = 0; i < n * n; ++i)
        for (int j = 0; j < n; ++j)
            for (int k = 0; k < n * n * n; ++k)
                cnt++;
}
```

Tips

- We know f1 is $O(n^3)$
- Now f2 has a single loop: $O(n)$
 - But its body is NOT constant!
 - Its body is call that is $O(n^3)$
- Overall $O(n^4)$
- Tip
 - Imagine we copy-pasted f2 in f1
 - Now u see clearly n^4 total steps
- Tip
 - Double check if the body is FIXED or variable based on N

```
void f1(int n) { //  $O(n^3)$ 
    int cnt = 0;
    for (int i = 0; i < n * n; ++i)
        for (int j = 0; j < n; ++j)
            cnt++;
}

void f2(int n) { //  $O(n^4)$ 
    for (int i = 0; i < n; ++i)
        f1(i); //  $n^3$ 
}
```

Tips

- Sometimes our function depends on several variables
- Here we have total: $6nm$
- Drop constants $\Rightarrow O(nm)$

```
void f3(int n, int m) { // O(nm)
    int cnt = 0;
    for (int i = 0; i < 2 * n; ++i)
        for (int j = 0; j < 3 * m; ++j)
            cnt++;
}
```

Tips

- 2 parallel blocks
 - Block 1: $O(nm)$
 - Block 2: $O(n^2)$
- Which is bigger? We don't know
- Total: $O(nm + n^3)$

```
void f4(int n, int m) {    //  $O(nm + n^3)$ 
    int cnt = 0;
    for (int i = 0; i < 2 * n; ++i)
        for (int j = 0; j < 3 * m; ++j)
            cnt++;

    for (int i = 0; i < n * n * n; ++i)
        cnt++;
}
```


Polynomial Order

- Today we discussed polynomial order functions (format n^k)
 - $n^0 = 1$ (const), n^2 , n^3 and so on
- Intuition: code is doing some hundred million steps $\Rightarrow \sim 1$ second (not really)
- From the table,
The **bigger** your $O()$
The **slower** your code
- There are other worse Families (later)
 - E.g. $O(n^n)$ or $O(!n)$

	n=100	n=1000	n=1000000
$O(n)$	100	1000	1000000
$O(n^2)$	10000	1000000	Too much
$O(n^3)$	1000000	1000000000	Too much
$O(n^4)$	100000000	Too much	Too much

Overall

- Try to keep these tips in mind
- But be careful from tricky codes, so don't be so systematic
 - E.g. 3 nested while loops might actually just doing $10n$ steps NOT n^3
- Whenever you write a code from now one, always compute its order
 - This is how your skill will grow up
 - As it is your code, you know really what is happening

“Acquire knowledge and impart it to the people.”

“Seek knowledge from the Cradle to the Grave.”