

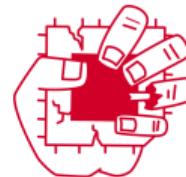
# Deepen the Defenses: A Case for Microarchitectural Isolation

Jo Van Bulck

Cybersec Europe, FutureLab Stage, Brussels, May 11, 2022

🏡 imec-DistriNet, KU Leuven, Belgium 📩 jo.vanbulck@cs.kuleuven.be 🐦 jovanbulck

- Postdoctoral **researcher** @imec-DistriNet, KU Leuven, Belgium  
→ PhD “Microarchitectural Side-Channel Attacks for Privileged Software Adversaries”
- **Trust across the system stack:** App > compiler > OS > CPU >  $\mu$ -arch



Side-channel analysis

Transient-execution attacks  
(Intel x86 SGX)

Embedded trust  
(TI MSP430)

**Hardware (noun.)** — *The part of a computer that you can kick.*

**Software (noun.)** — *The reason you want to kick the hardware.*

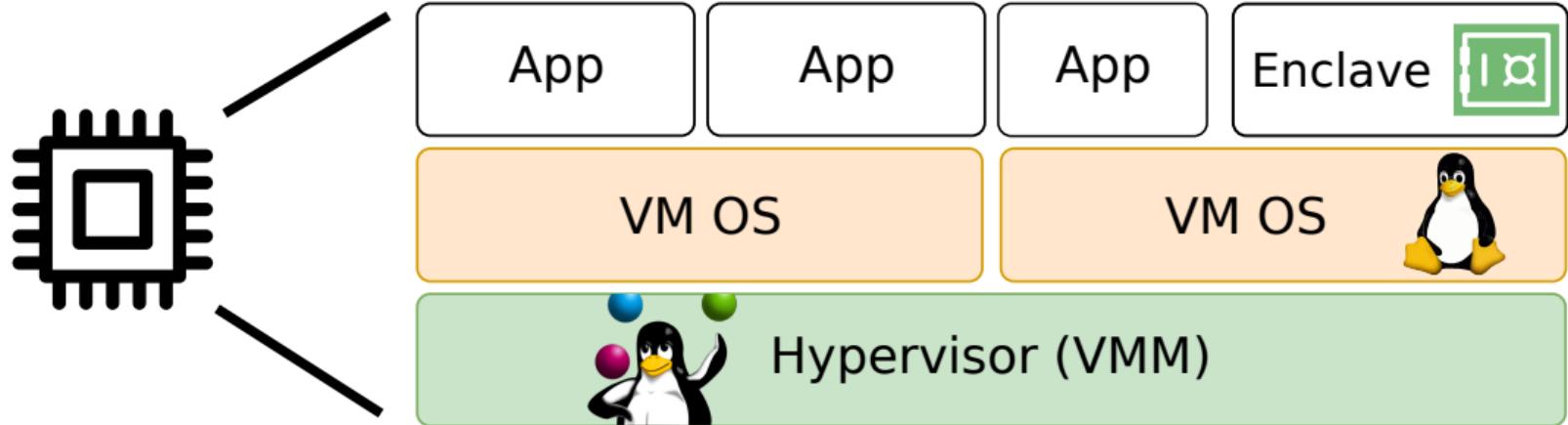
# Software Engineer vs Hardware Engineer



Job Title	
Software engineer	Hardware engineer
Job Description	
Develop, design and test software or construct, maintain computer networks and programs	Research, develop and test hardware or computer equipment
Education	
Software Engineering or Computer Science Degree	Electrical & Computer Engineering Degree
Skill Sets	
Technology Design, Complex Problem Solving, Critical Thinking, etc.	Troubleshooting, Problem Solving, Systems Evaluation, etc.
Salary	
\$107,840	\$112,760
Number of Jobs	
>1,128,000	>87,000

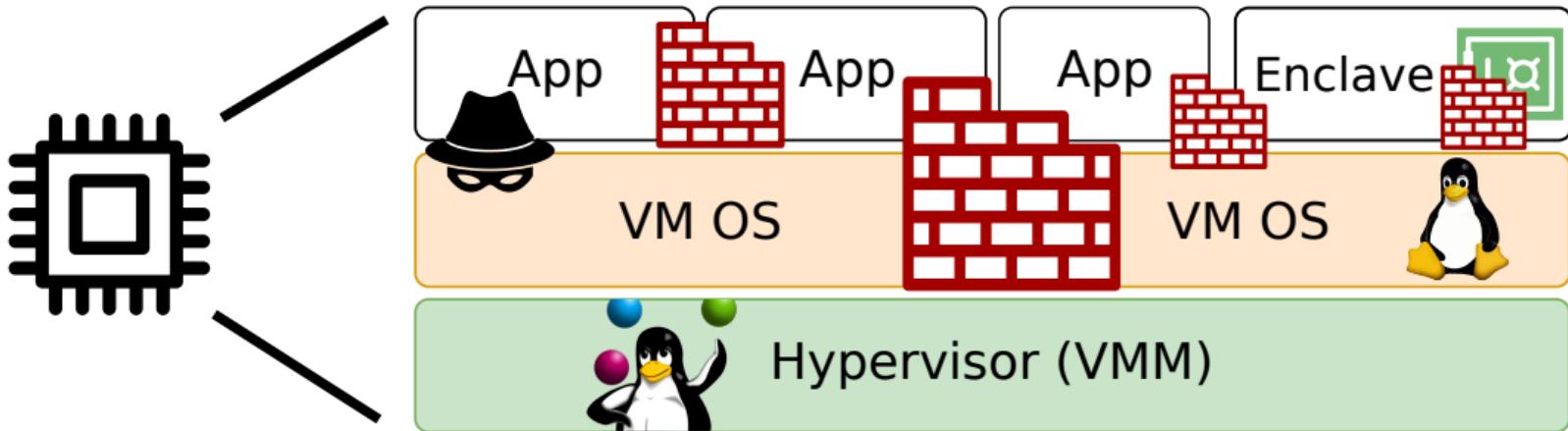


# Processor security: Hardware isolation mechanisms



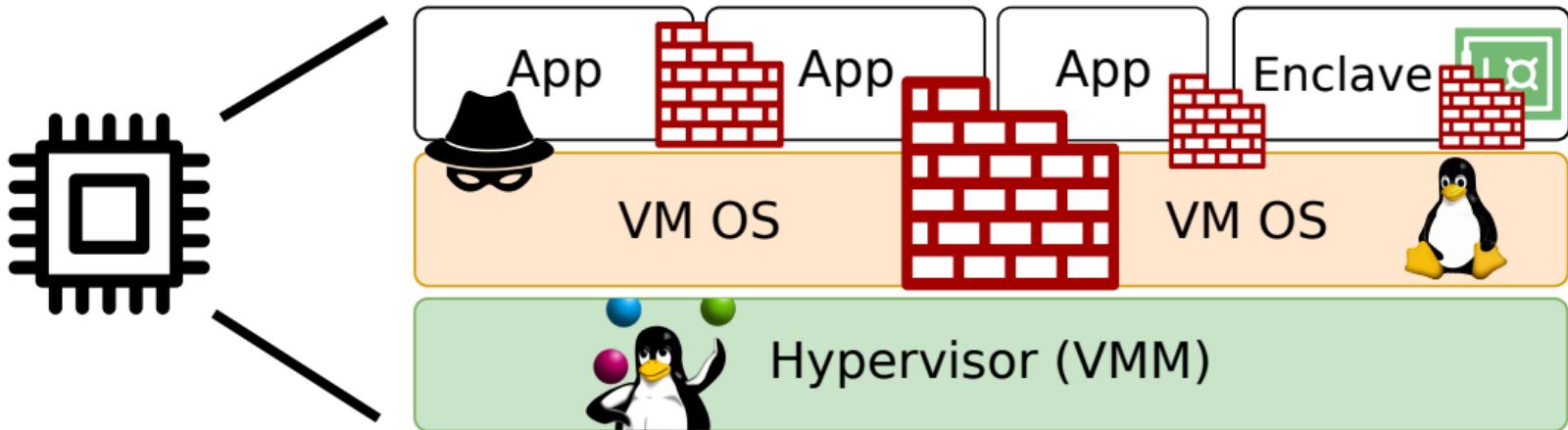
- Different software **protection domains**: Processes, VMs, enclaves

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- CPU builds “walls” for **memory isolation** between apps and privilege levels

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- Different software **protection domains**: Processes, VMs, enclaves
- CPU builds “walls” for **memory isolation** between apps and privilege levels  
↔ Architectural protection walls permeate **microarchitectural side channels!**



**VAULT DOOR**

WEIGHT: 22 1/2 Tons

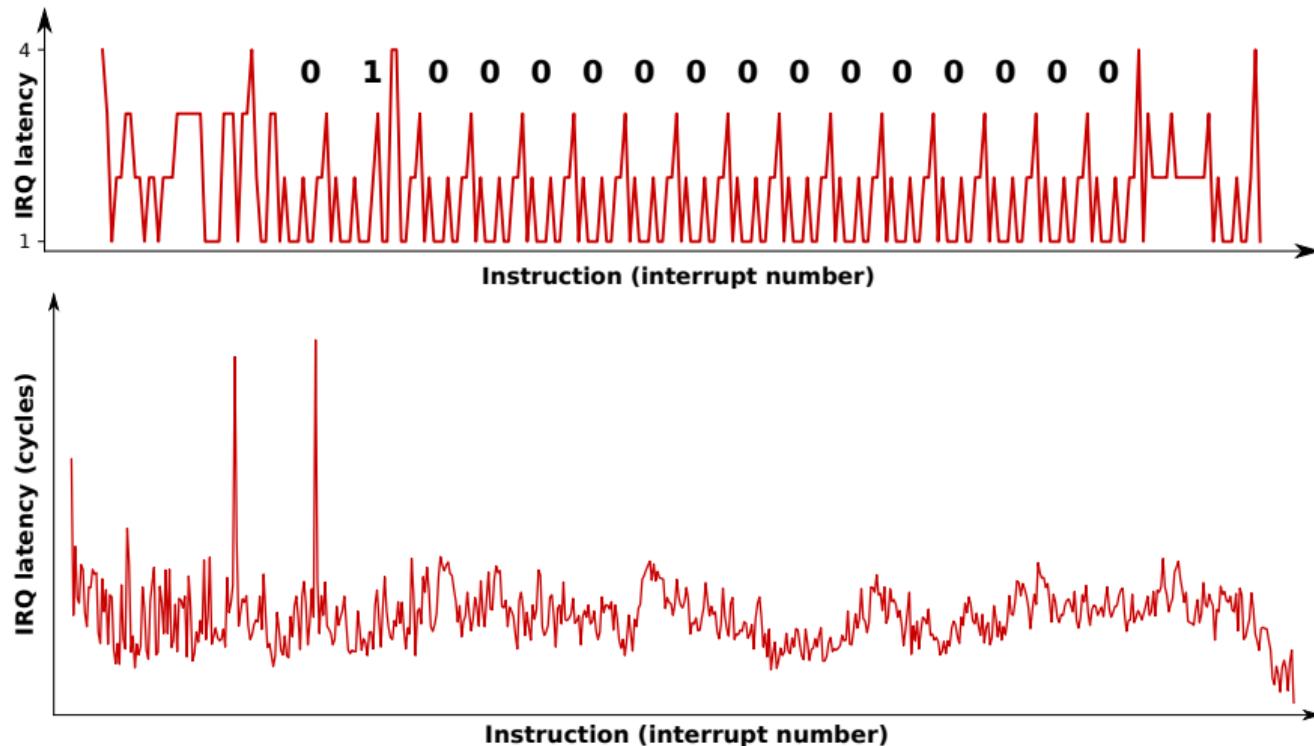
THICKNESS: 22 Inches

STEEL: 11 Layers of Special  
Cutting and Drill Resistant

LOCKS: 4 Hamilton Watch  
Movements for Time Locks



# Microarchitectural timing leaks in practice



## Example: CPU cache timing side channel

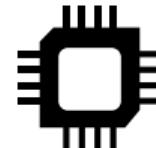


**Cache principle:** CPU speed  $\gg$  DRAM  $\rightarrow$  *cache code/data*

```
while true do  
    maccess(&a);  
endwh
```



**CPU + cache**



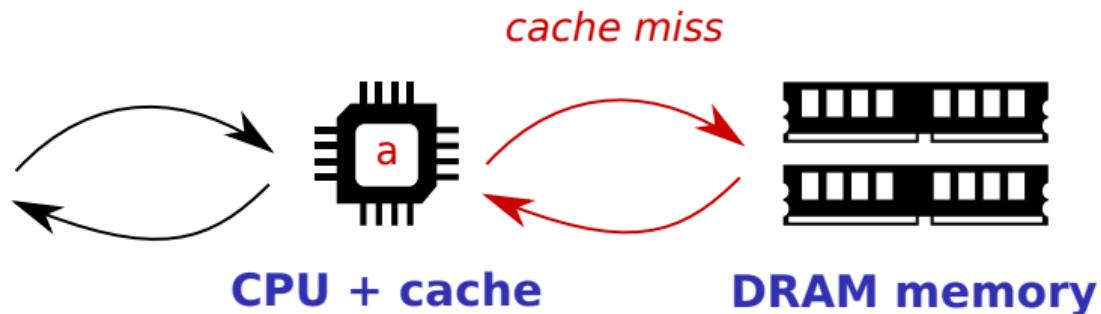
**DRAM memory**

# Example: CPU cache timing side channel



**Cache miss:** Request data from (slow) DRAM upon first use

```
while true do  
    maccess(&a);  
endwh
```



# Example: CPU cache timing side channel



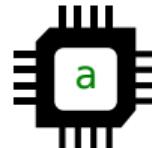
**Cache hit:** No DRAM access required for subsequent uses

```
while true do  
    maccess(&a);  
endwh
```

*cache hit*



**CPU + cache**



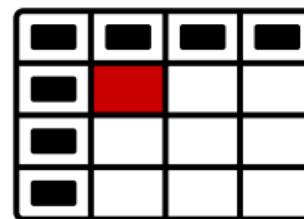
**DRAM memory**

# Cache timing attacks in practice: Flush+Reload

```
if secret do  
    maccess(&a);  
else  
    maccess(&b);  
endif
```



```
flush(&a);  
start_timer  
    maccess(&a);  
end_timer
```



CPU cache

*flush 'a' to memory*



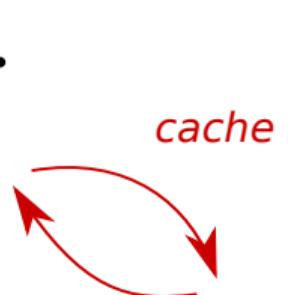
DRAM memory

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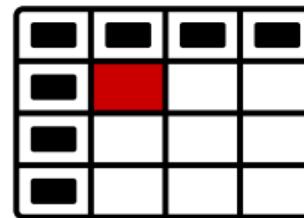
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*cache miss*

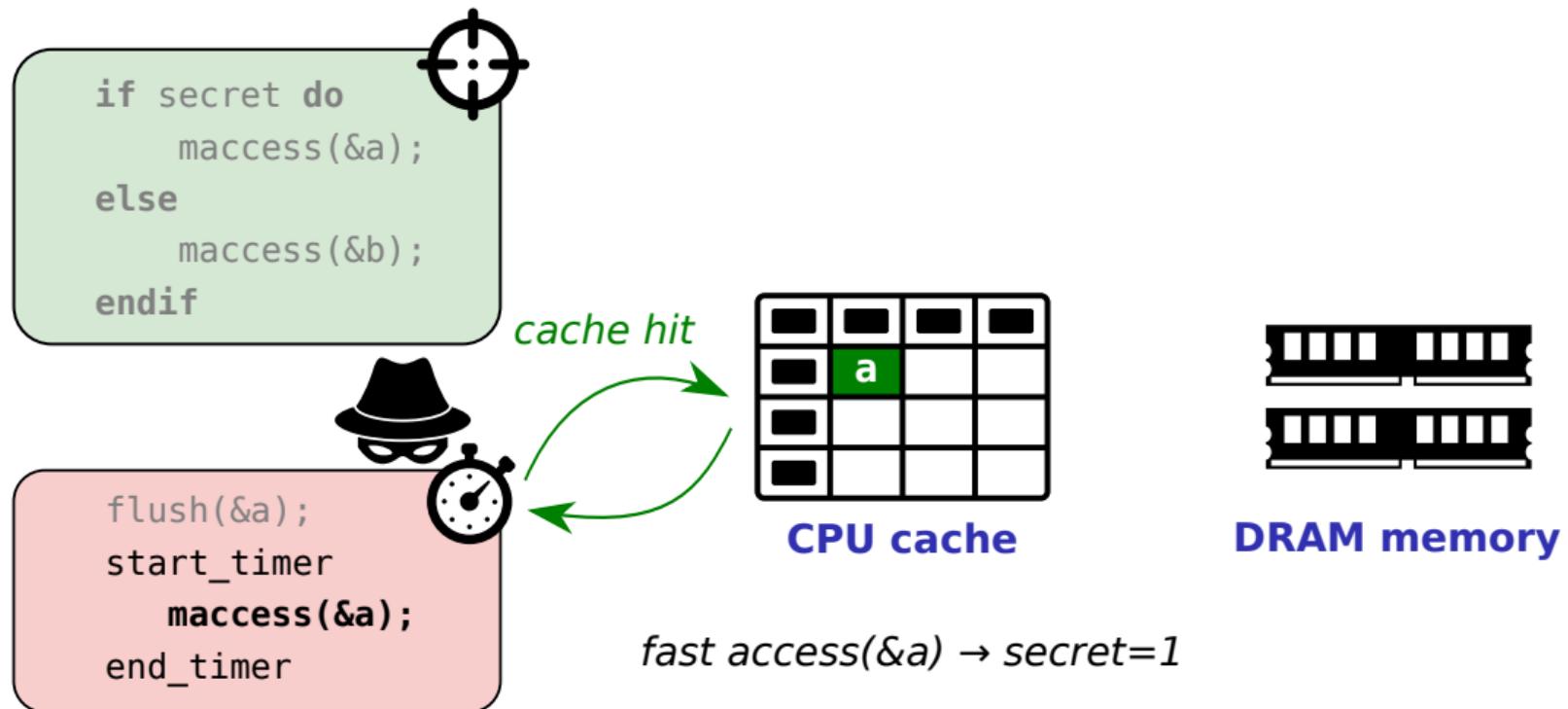


**CPU cache**



**DRAM memory**

# Cache timing attacks in practice: Flush+Reload





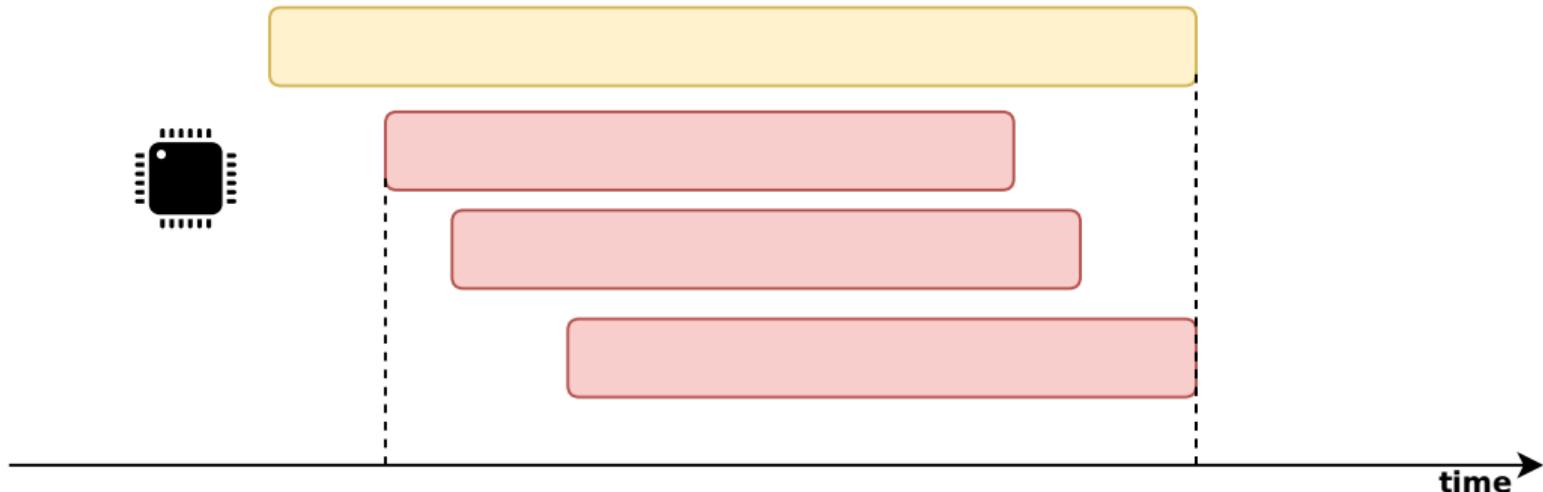
We can communicate across protection walls  
using microarchitectural side channels!

# WHAT IF I TOLD YOU

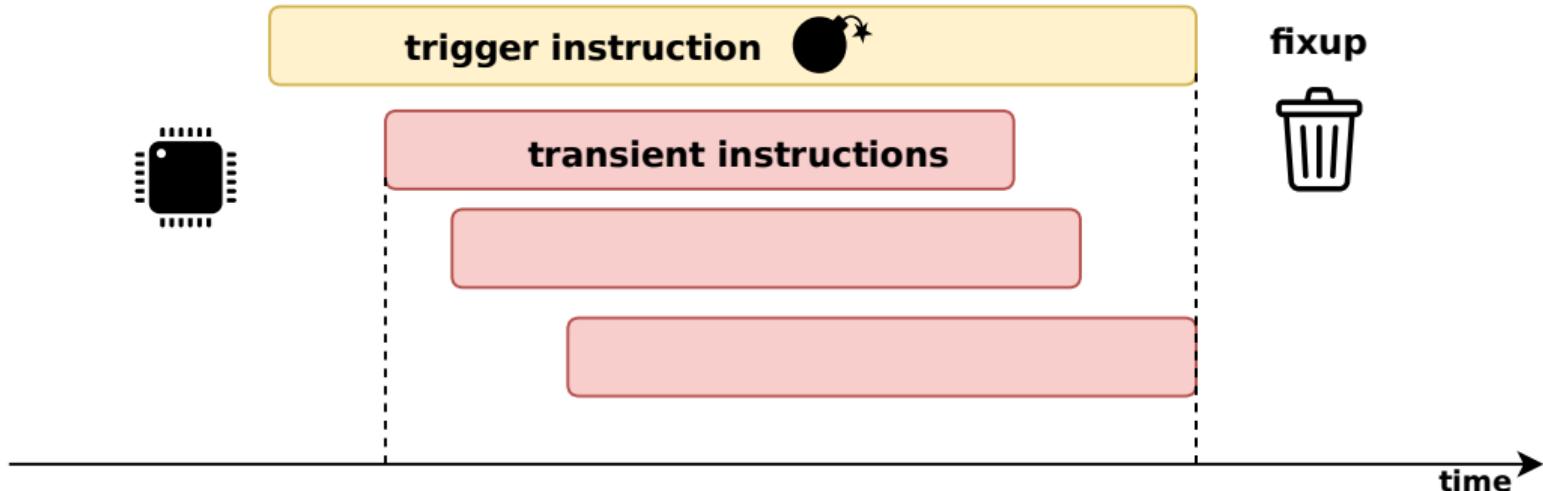


# YOU CAN CHANGE RULES MID-GAME

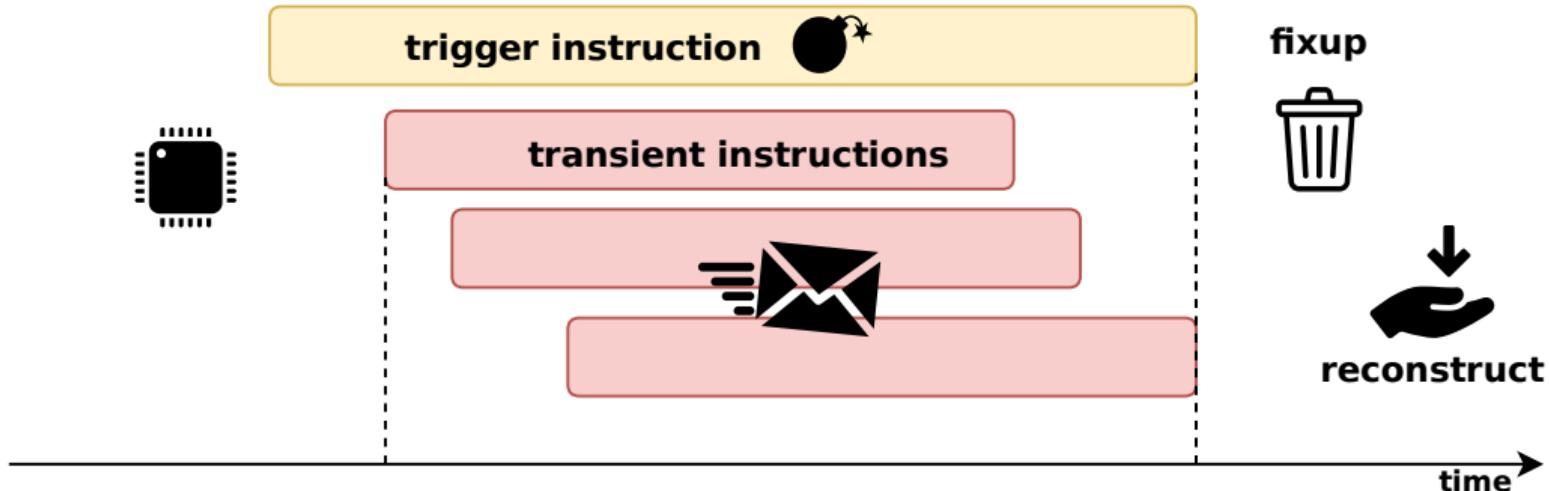
# Abusing out-of-order and speculative execution



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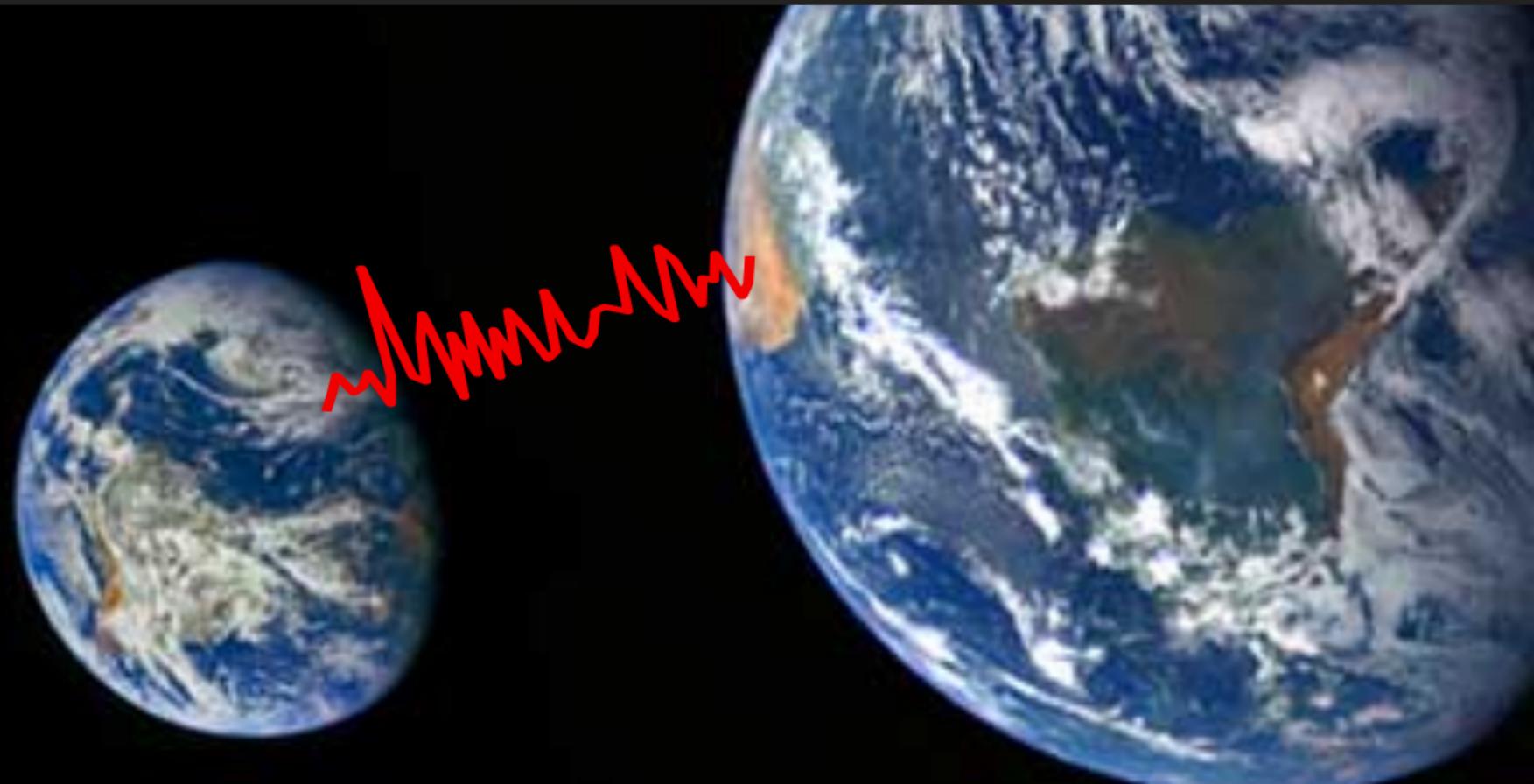


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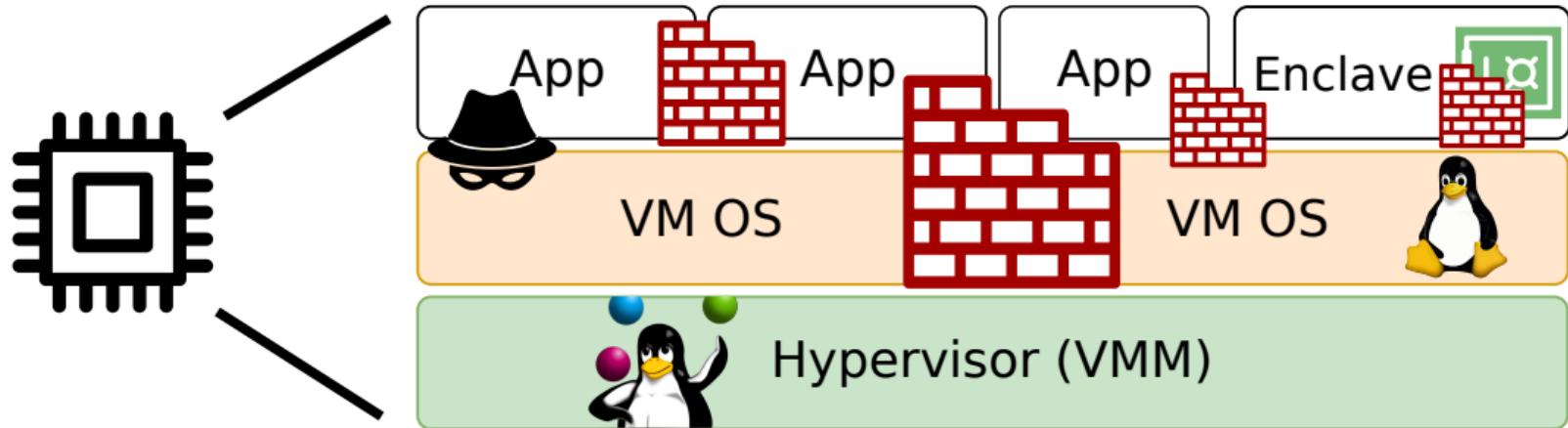




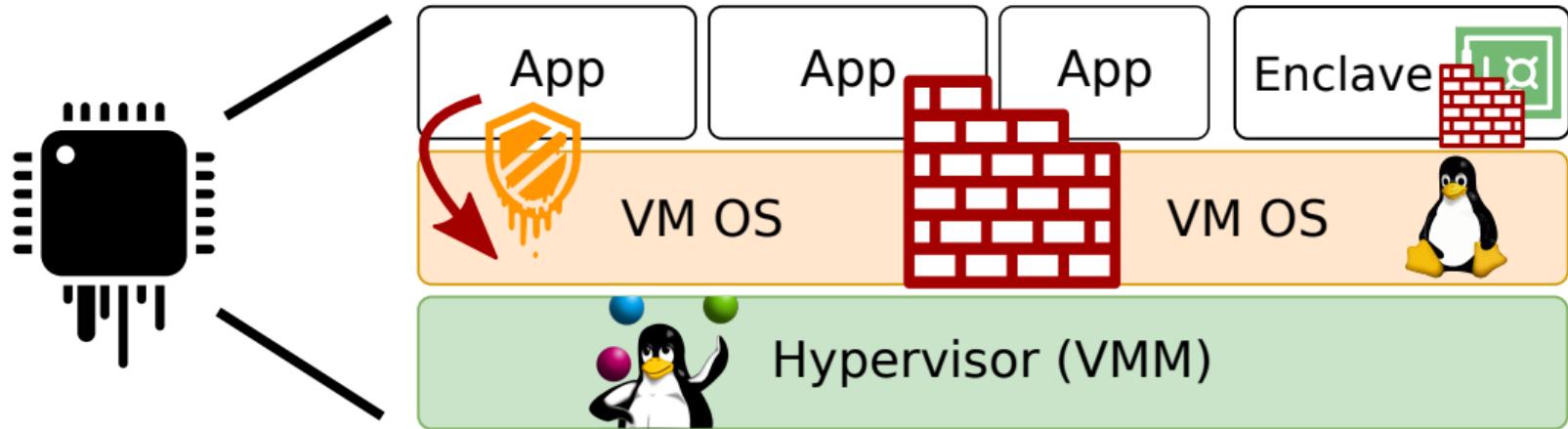
Transient-execution attacks: Welcome to the world of fun!



# Leaky processors: Breaking isolation mechanisms

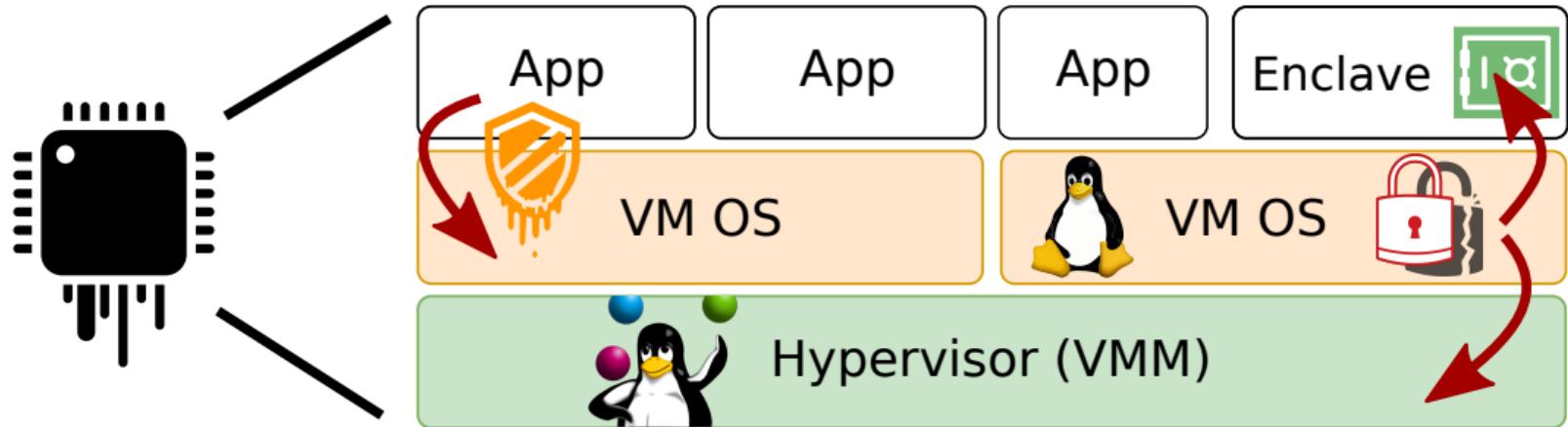


# Leaky processors: Breaking isolation mechanisms



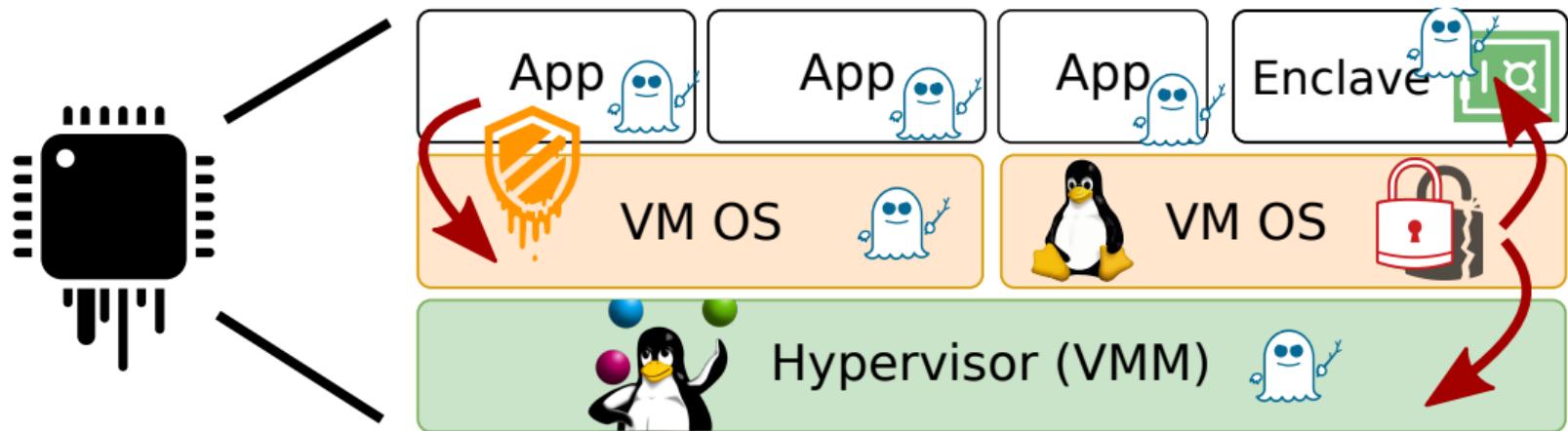
- **Meltdown** breaks user/kernel isolation

# Leaky processors: Breaking isolation mechanisms



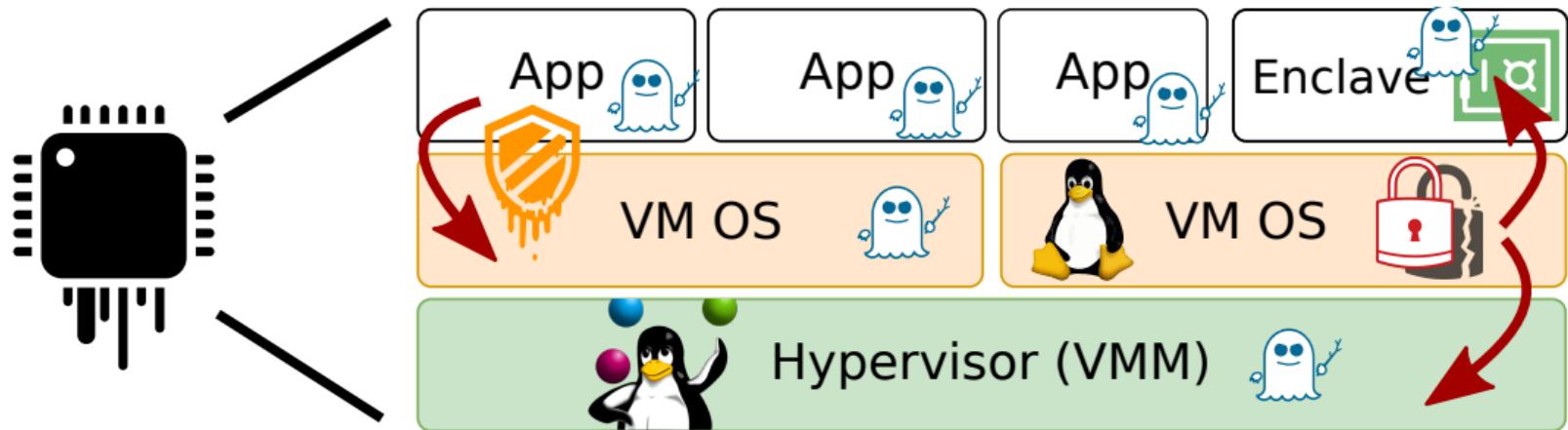
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- **Foreshadow** breaks SGX enclave and virtual machine isolation

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- **Meltdown** breaks user/kernel isolation
- **Foreshadow** breaks SGX enclave and virtual machine isolation
- **Spectre** breaks software-defined isolation on various levels
- ... many more – but all exploit the same underlying insights!

# Spectre v1: Speculative buffer over-read



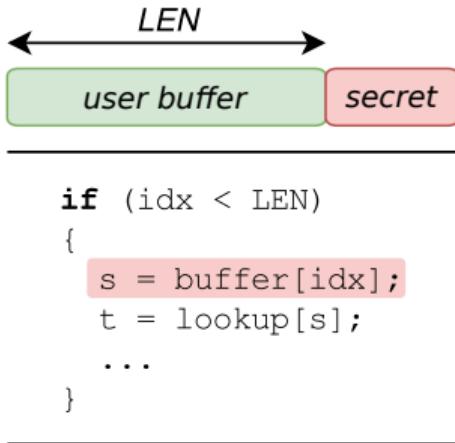
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```
if (idx < LEN)
{
    s = buffer[idx];
    t = lookup[s];
    ...
}
```

---

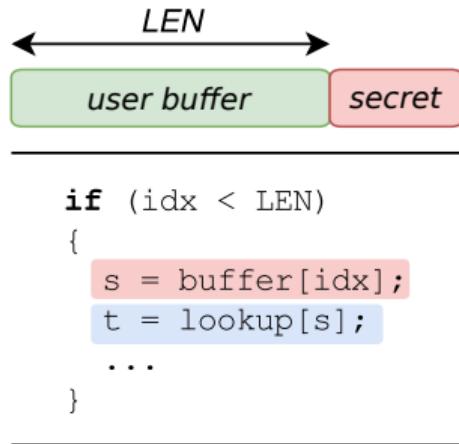
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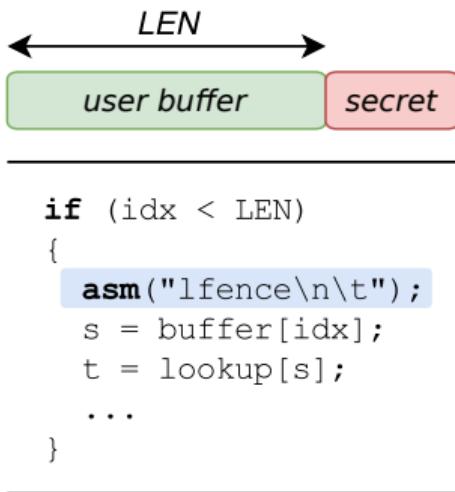
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# Spectre v1: Speculative buffer over-read



- Programmer *intention*: no out-of-bounds accesses
- **Mistrain gadget** to **speculatively** “ahead of time” execute with  $idx \geq LEN$  in the transient world
- **Side channels** may leave traces after roll-back!
- Insert explicit **speculation barriers** to tell the CPU to halt the transient world...

Terminal

Foreshadow Demo

(100%) 12:27 AM

SGX enclave: secret string at 0x7f19ee646000

Press enter to naively read enclave memory at address 0x7f19ee646000...

Segment 0: 0x7f19ee646000 - 0x7f19ee646317

Victim address = 0x7f19ee646316... 0xFF

Actual success rate = 0/791 = 0.00 %

Press enter to use Foreshadow to read enclave memory at address 0x7f19ee646000 ...

Segment 0: 0x7f19ee646000 - 0x7f19ee646317

Victim address = 0x7f19ee6460dd... 0x69

Extracted Bytes-----

49	74	20	77	61	73	20	6F	6E	65	20	6F	66	20	74	68	6F	73	65	20	70	69	63	74	75	72	65	73	20	77	68	69	63	68
20	61	72	65	20	73	6F	20	63	6F	6E	74	72	69	76	65	64	20	74	68	61	74	20	74	68	65	20	65	79	65	73	20	66	6F
6C	6C	6F	77	20	79	6F	75	20	61	62	6F	75	20	77	68	65	6F	79	6F	75	20	6D	6F	76	65	2F	20	42	49	47	20		
42	2	F	71	48	15	57	20	40	57	20	7	41	1	4c	4	5	4	5	C	0	'4	88	5	?	!	61	70	74	69	6C	61	74	20
61	0	6	7	L	U	9	3	0	6	63	5	26	0	6	1	6	6	6	6	67	0	6	6	6	6	6	20	61	20	6C			
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However FORESHADOW  
can read the actual  
enclave memory

It was one of those pictures which are so contrived that the eyes follow you about when you move. BIG BROTHER IS WATCHING YOU, the caption beneath it ran. Inside the flat a fruity voice was reading out a list of figures w.....



**SHARING IS NOT CARING**

**SHARING IS LOSING YOUR STUFF TO OTHERS**

# A new golden age for computer architecture?



## Conclusions and take-away



**Hardware + software patches:** Update your systems!

# Conclusions and take-away



**Hardware + software patches:** Update your systems!

- ⇒ New emerging and powerful class of **transient-execution** attacks
- ⇒ Importance of fundamental **side-channel research**; no silver-bullet defenses
- ⇒ Security **cross-cuts** the system stack: hardware, OS, VMM, compiler, app

