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ME-HPTs: Memory-Efficient Hashed Page Tables

HPCA 2023

Jovan Stojkovic, Namrata Mantri, Dimitrios Skarlatos*, Tianyin Xu, Josep Torrellas

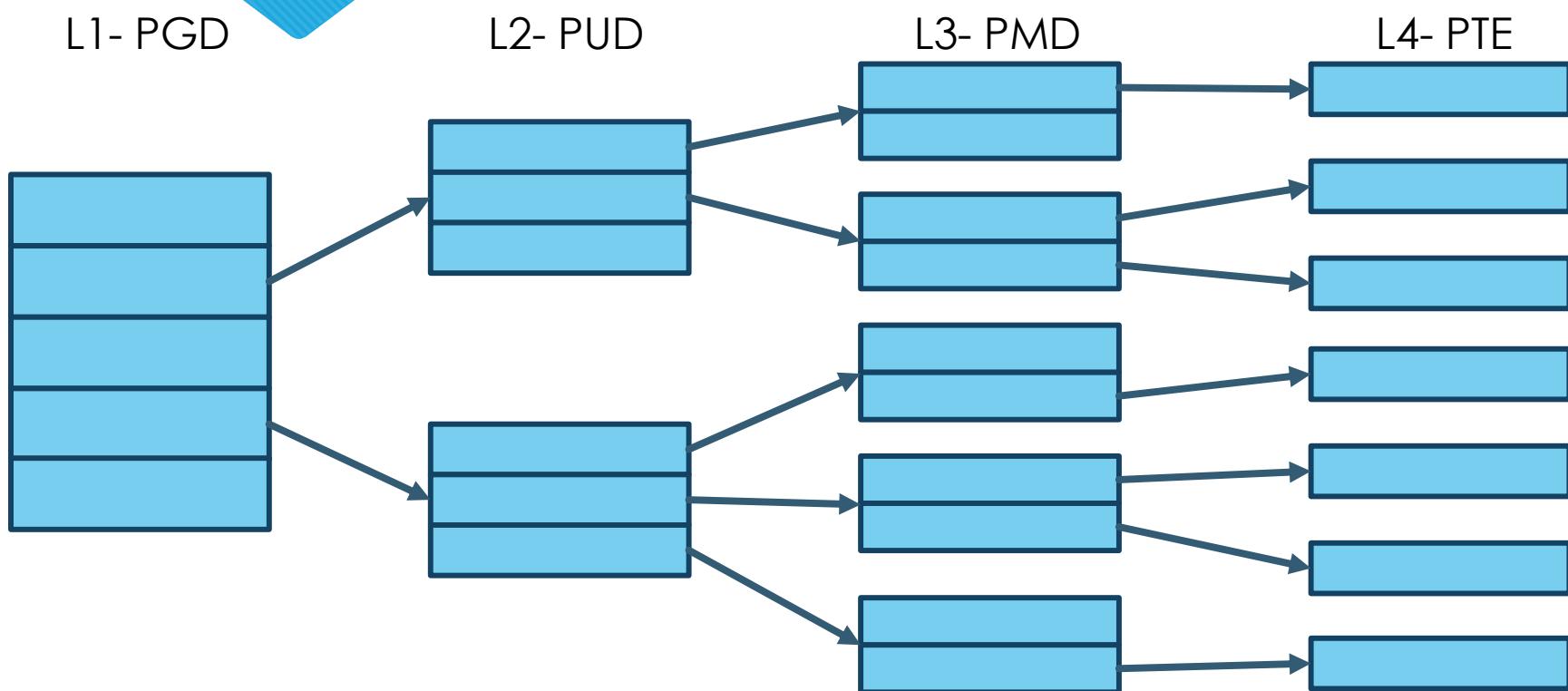
University of Illinois at Urbana-Champaign

*Carnegie Mellon University

Virtual Memory and Page Tables

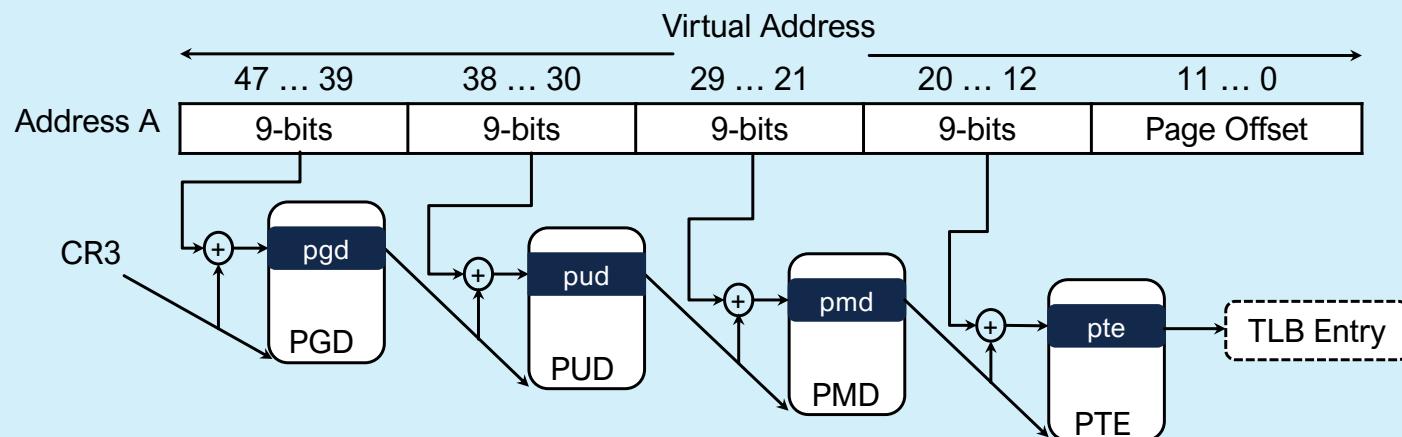
- Virtual memory is an essential technique in modern computing systems
 - Memory virtualization
 - Process isolation
- Virtual memory performance depends on the page table organization
 - Radix page tables – slow and not scalable
 - Hashed page tables – memory inefficient

Radix Page Tables: Memory-Efficient Multi-Level Trees



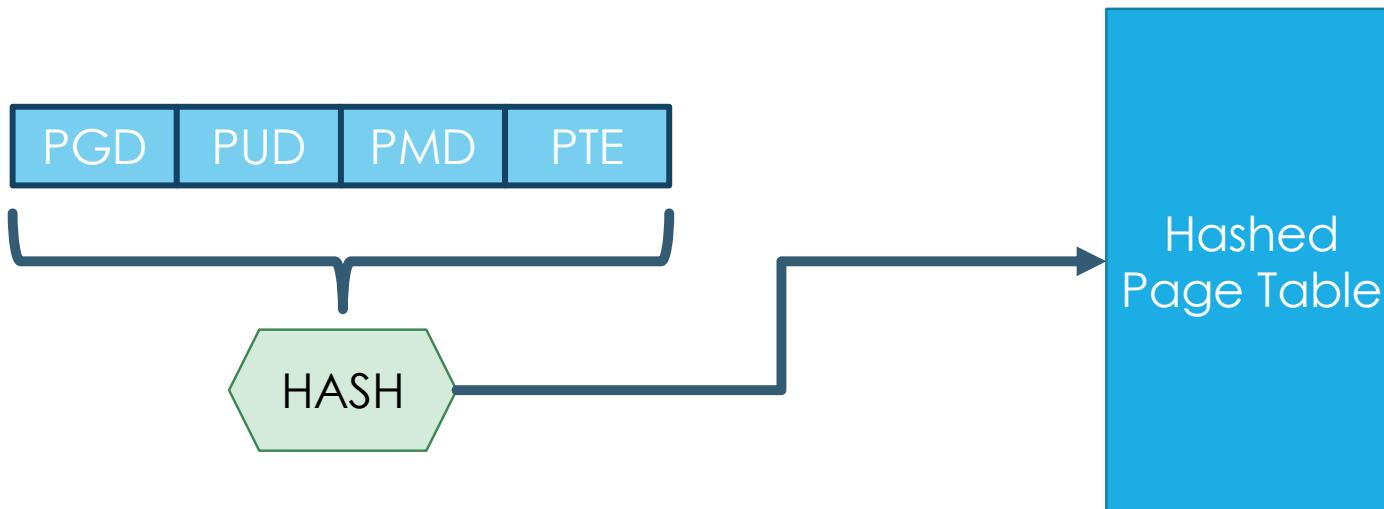
Radix Page Walk: Expensive Pointer Chase

x86-64 Radix Page Tables



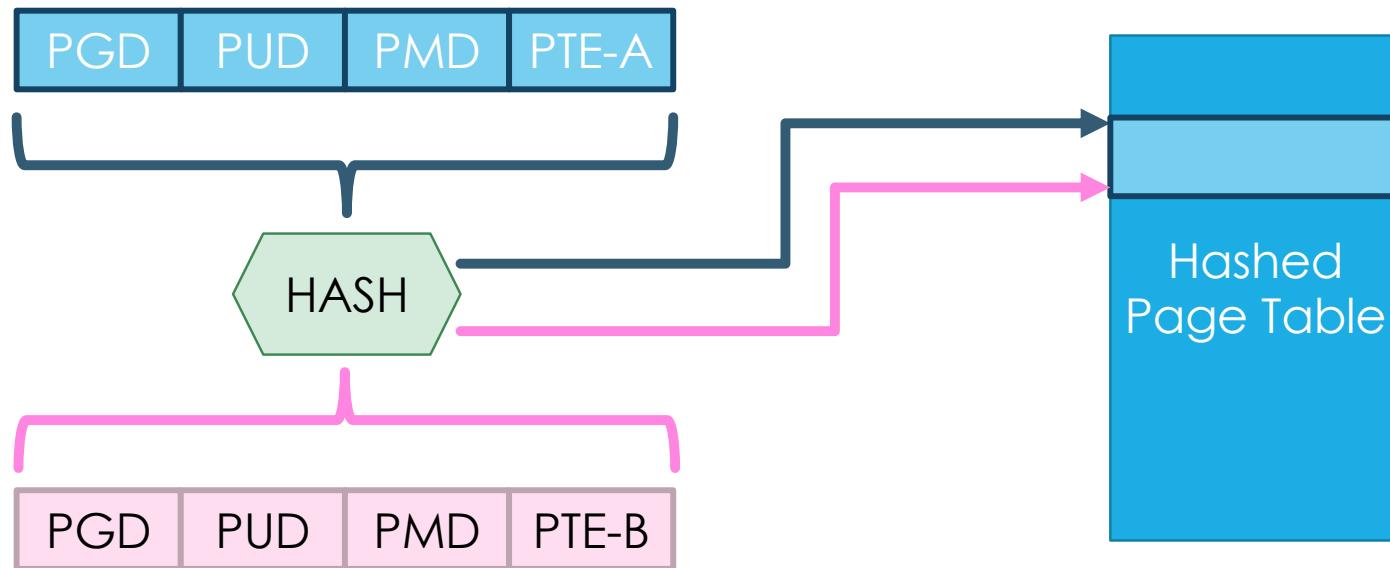
Hashed Page Tables

- 😊 Page walk requires a single memory access

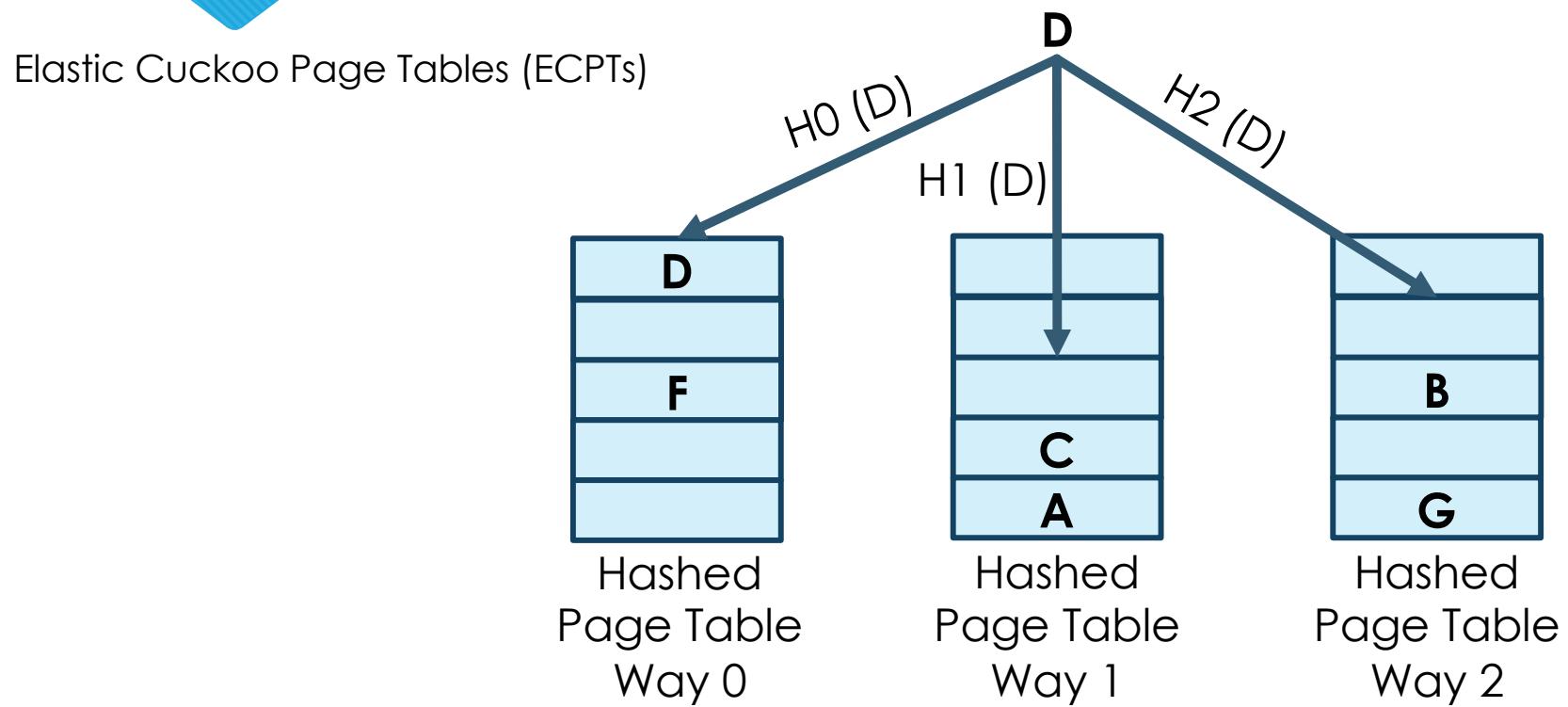


Hashed Page Tables

😢 Hash collisions



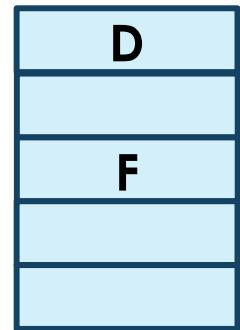
Hashed Page Tables: Recent Advances Make Them Compelling



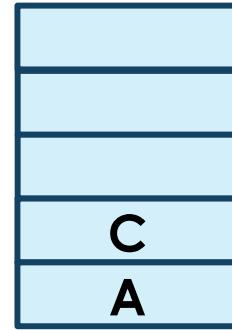
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Cuckoo Hashing

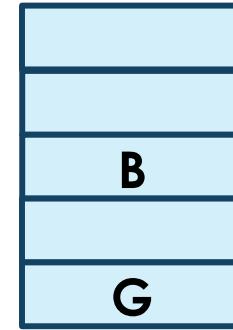
Insert E



Hashed
Page Table
Way 0

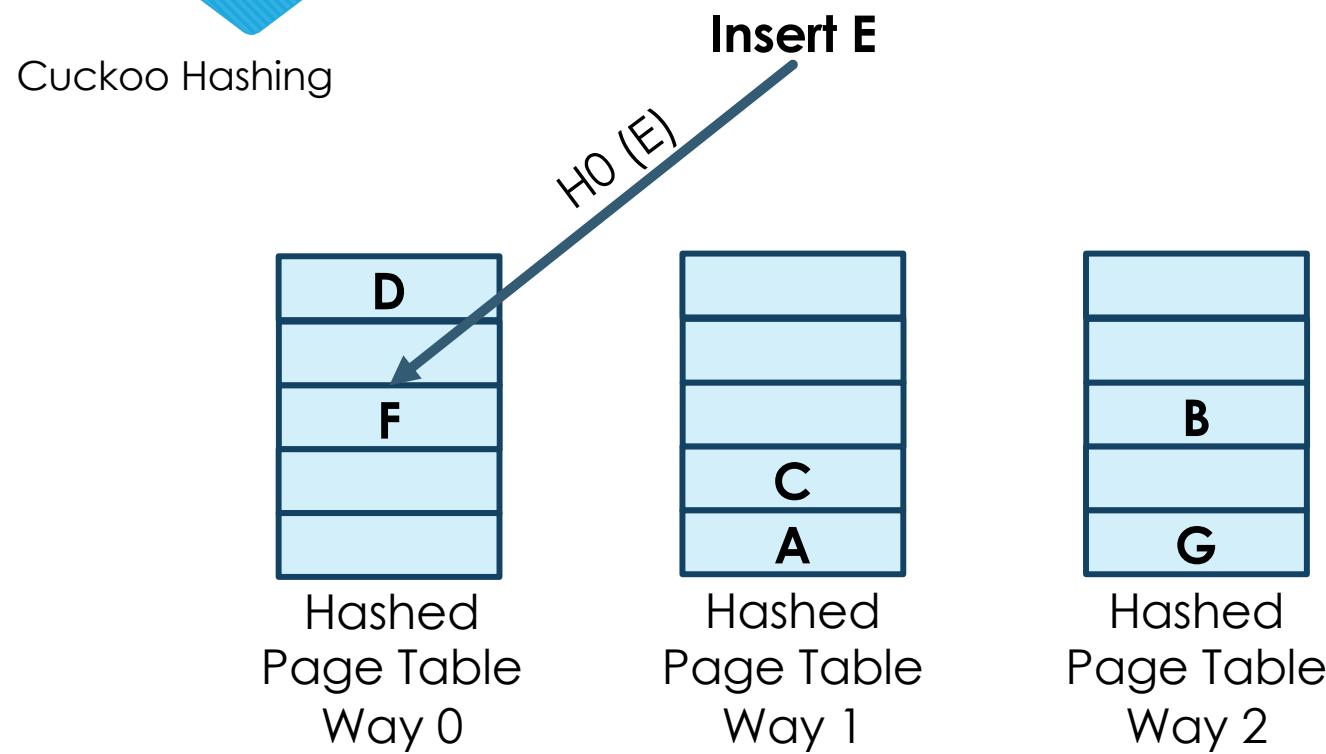


Hashed
Page Table
Way 1

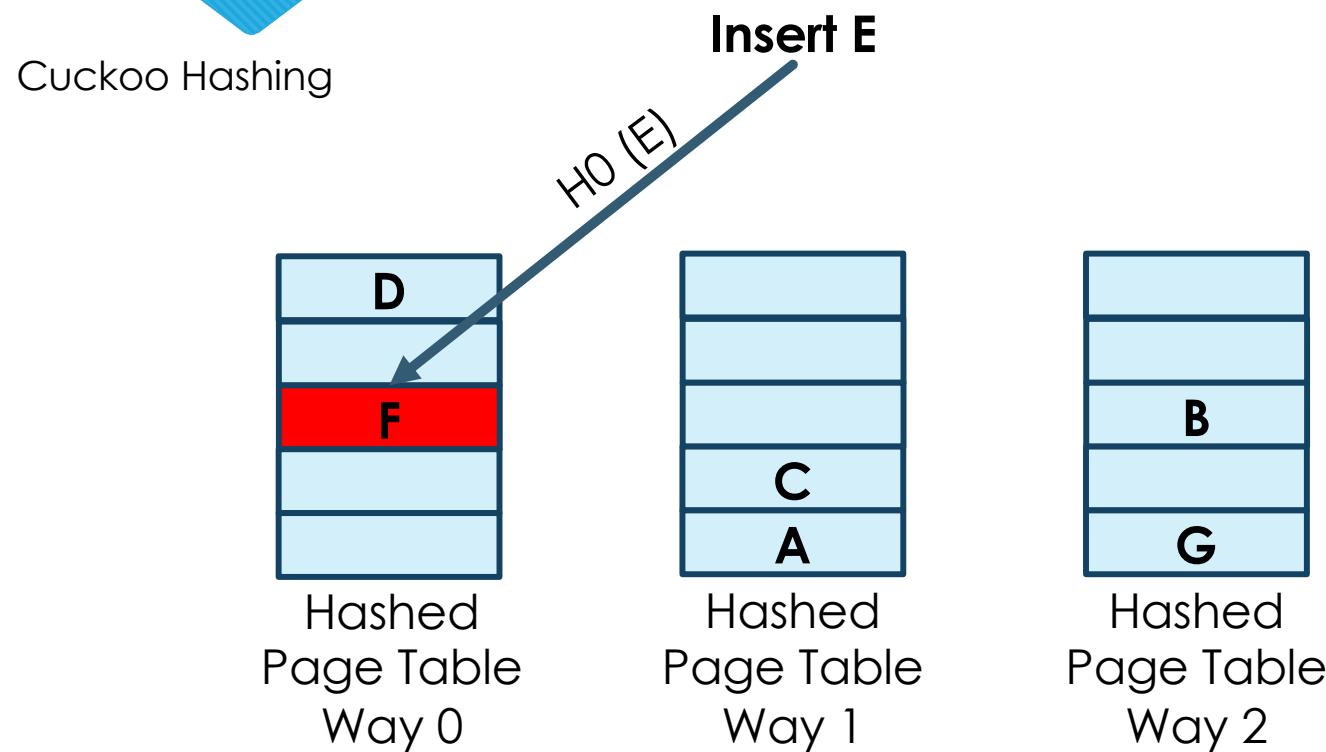


Hashed
Page Table
Way 2

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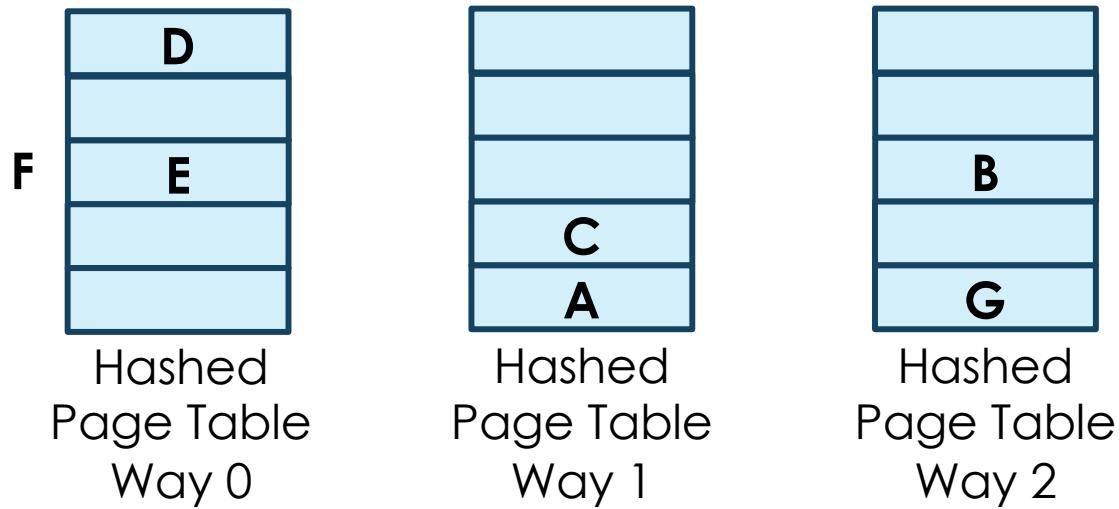


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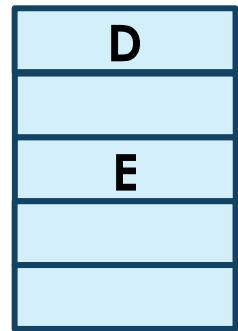
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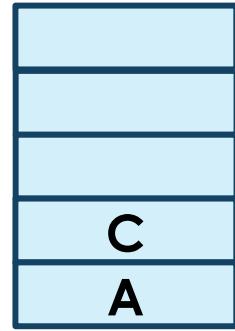
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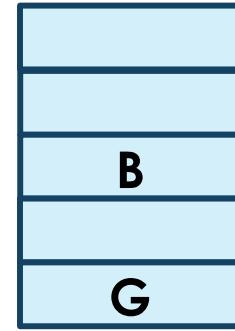
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Hashed
Page Table
Way 0



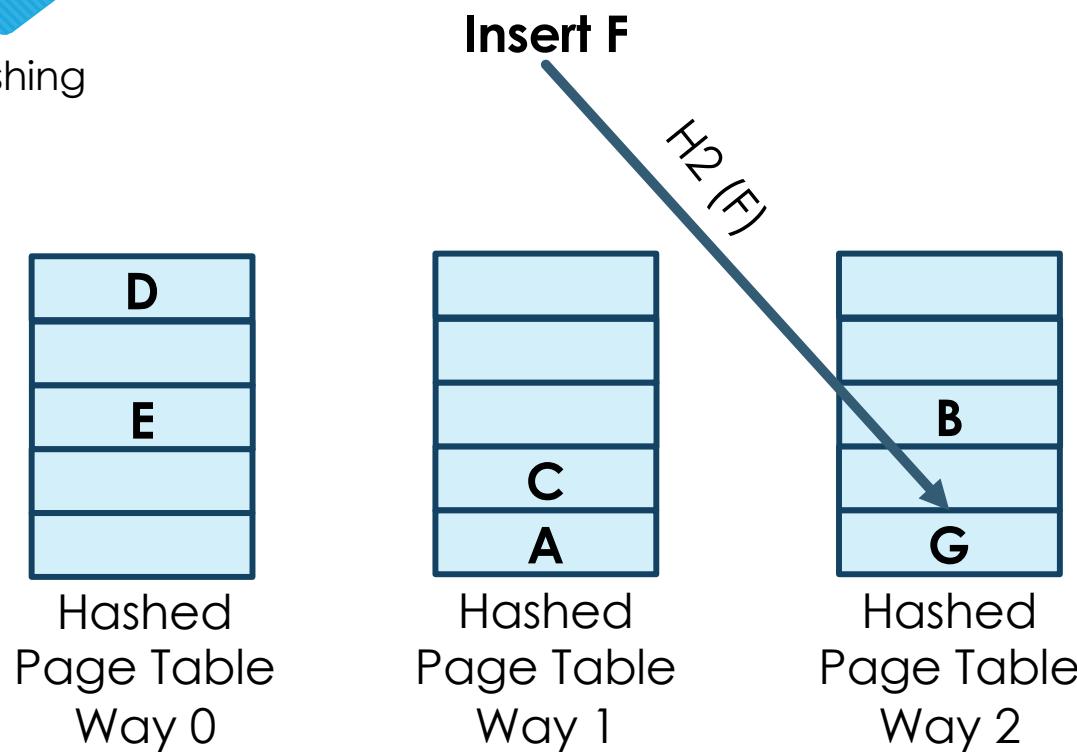
Hashed
Page Table
Way 1



Hashed
Page Table
Way 2

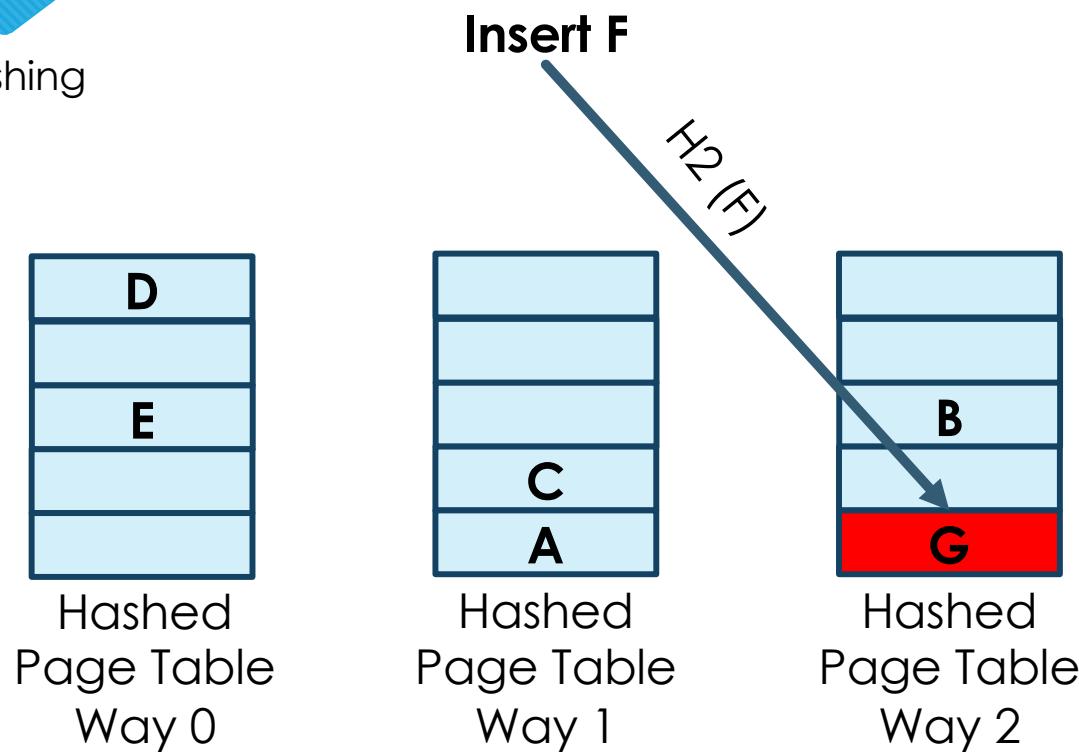
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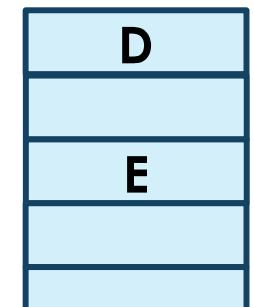
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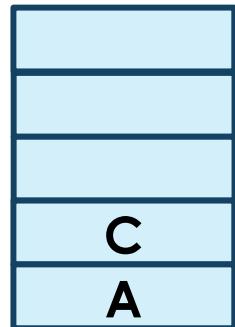


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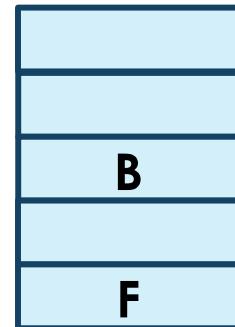
Cuckoo Hashing



Hashed
Page Table
Way 0



Hashed
Page Table
Way 1



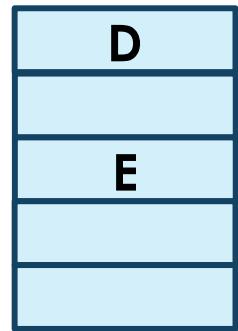
Hashed
Page Table
Way 2

G

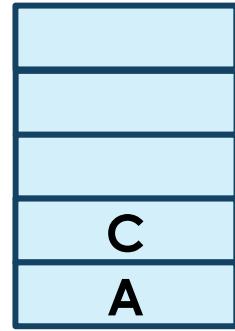
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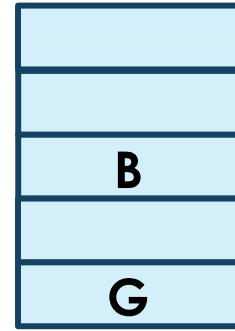
Insert G



Hashed
Page Table
Way 0

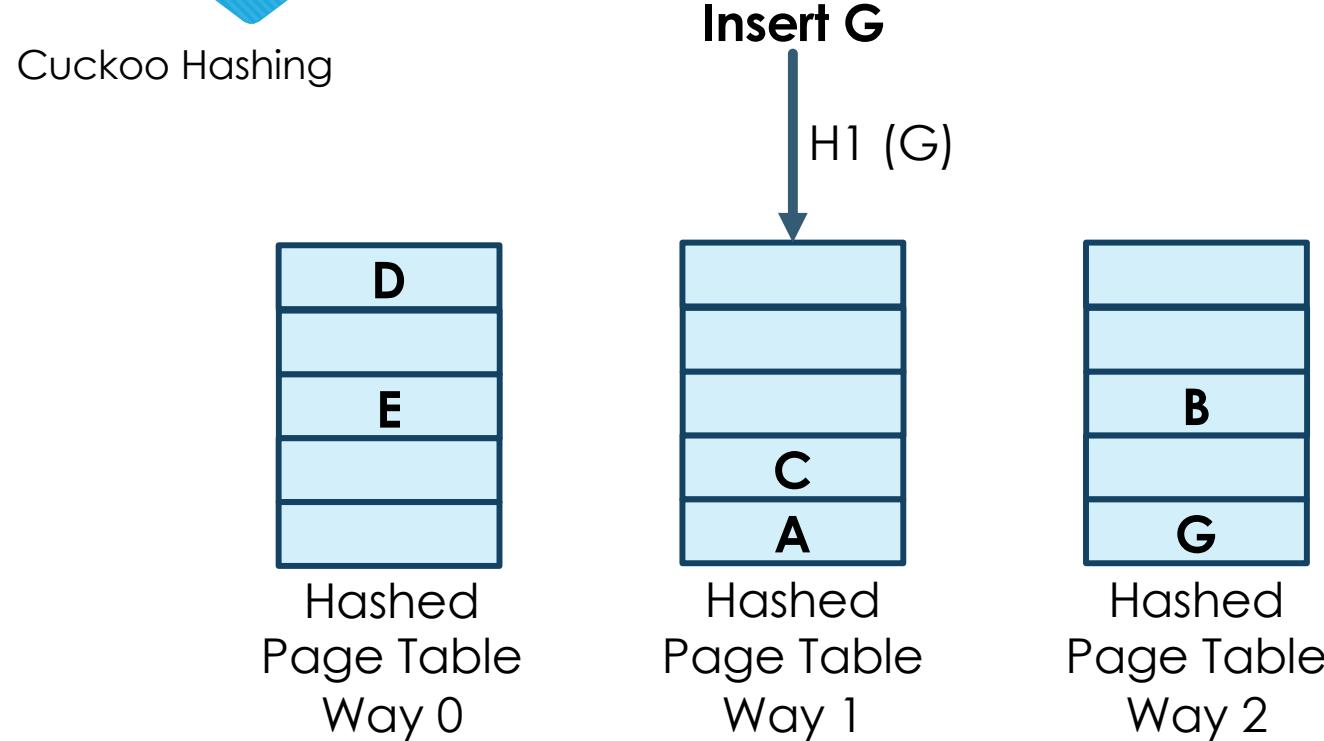


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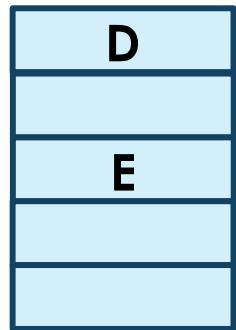
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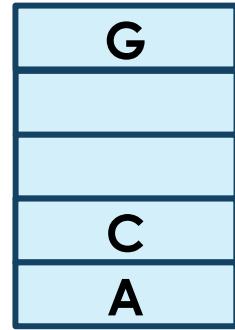


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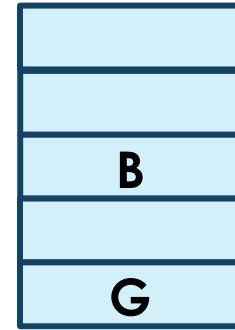
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Hashed
Page Table
Way 0



Hashed
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Hashed
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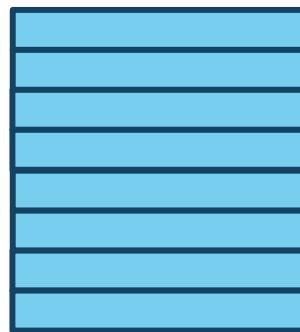
Outline of this talk

- **Problem: Contiguous Memory Requirements of Hashed Page Tables**
- ME-HPTs: Memory-Efficient Hashed Page Tables
 - ME-HPTs Design
 - ME-HPTs Key Results
- Conclusion

Hashed Page Tables: Large Contiguous Memory Chunks

- With hashed page tables – unity of allocation is one way of the page table

Hashed
Page Table
Way 0



Hashed
Page Table
Way 1

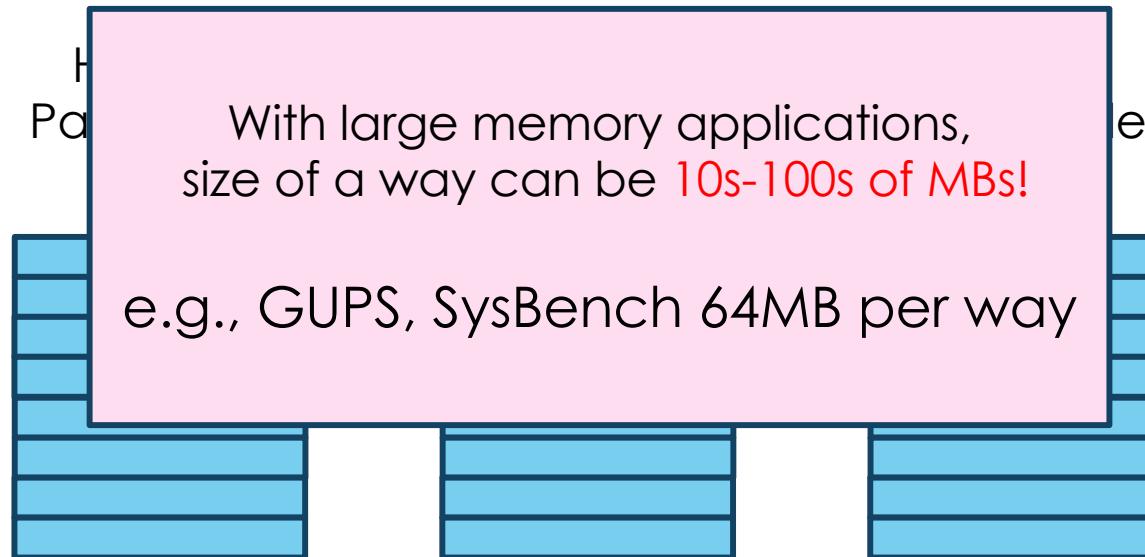


Hashed
Page Table
Way 2



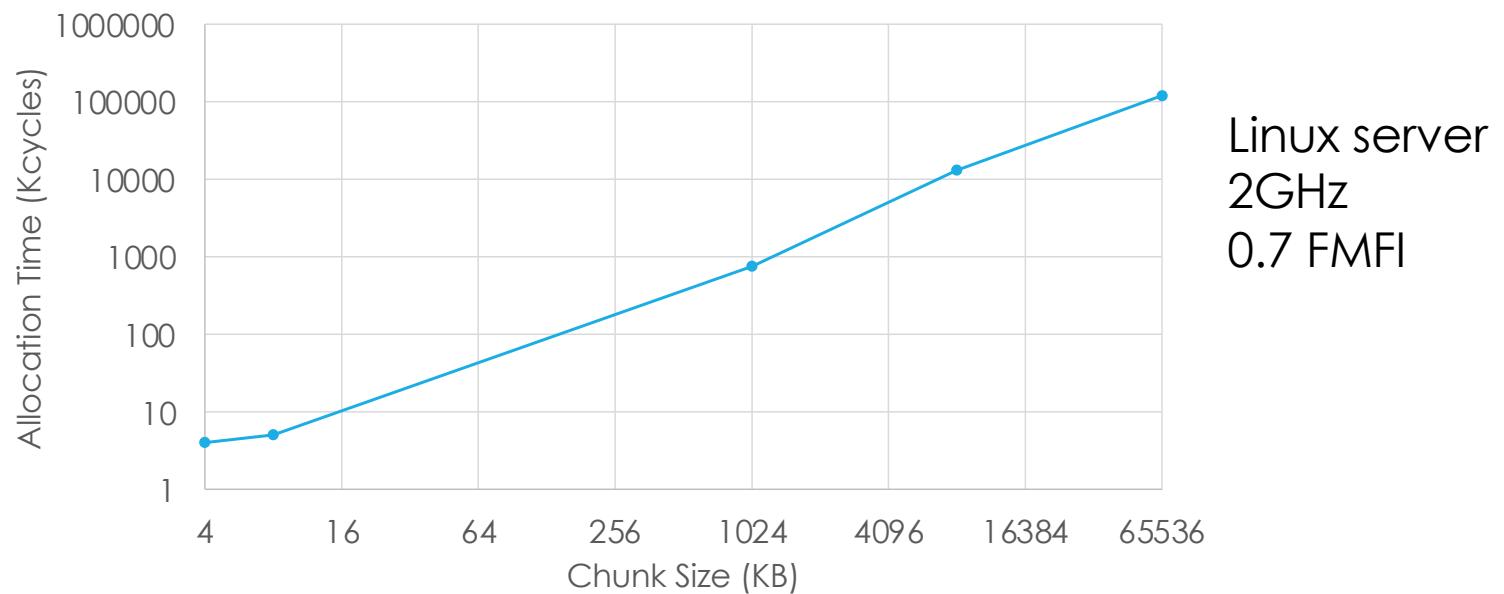
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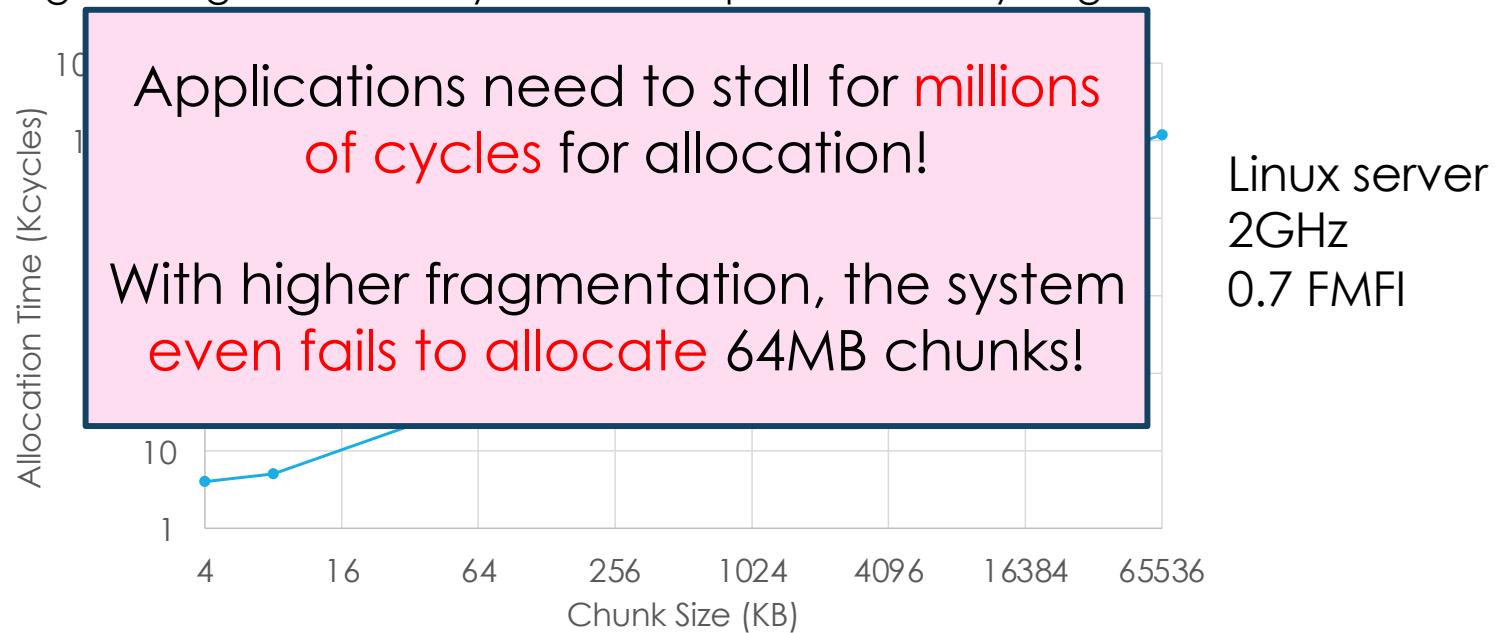
Hashed Page Tables: Contiguity is Expensive!

- Finding large contiguous memory chunks is expensive in busy fragmented servers



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Contributions

- Four novel architectural techniques to provide Memory-Efficient Hashed Page Tables (**ME-HPTs**)
- Reduced memory contiguity requirement by 92%
- Sped-up applications by 9% on average
- Allow large-memory applications to run at high performance on highly fragmented servers

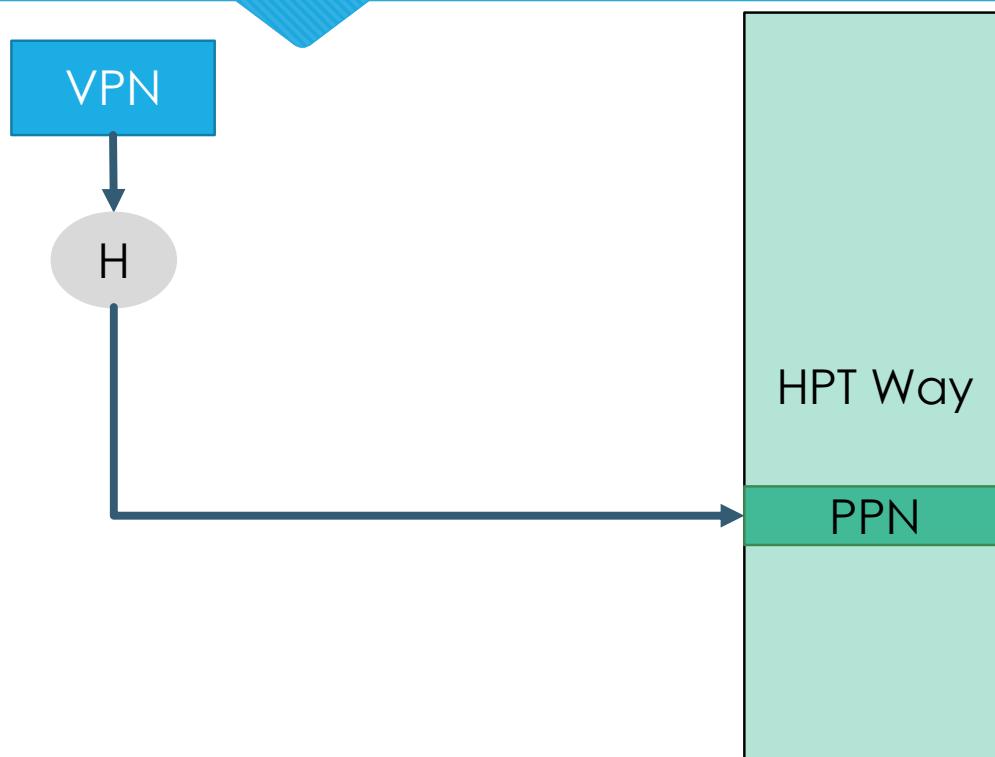
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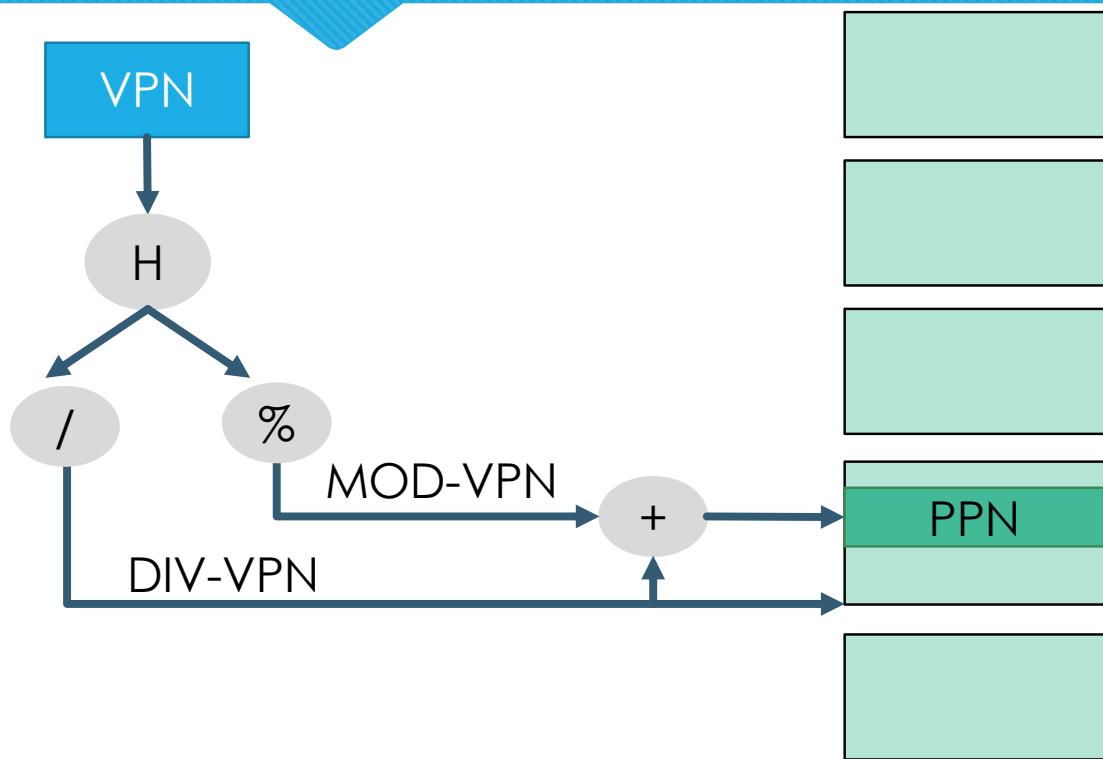
Memory-Efficient Hashed Page Tables: ME-HPTs Design Overview

- Memory-Efficient Hashed Page Tables (ME-HPTs): Four novel architectural techniques
- Directly minimizing contiguity requirements
 - Logical-to-Physical (L2P) Table
 - Dynamically Changing Chunk Size
- Indirectly minimizing contiguity requirements by minimizing memory consumption
 - In-place Page Table Resizing
 - Per-way Page Table Resizing

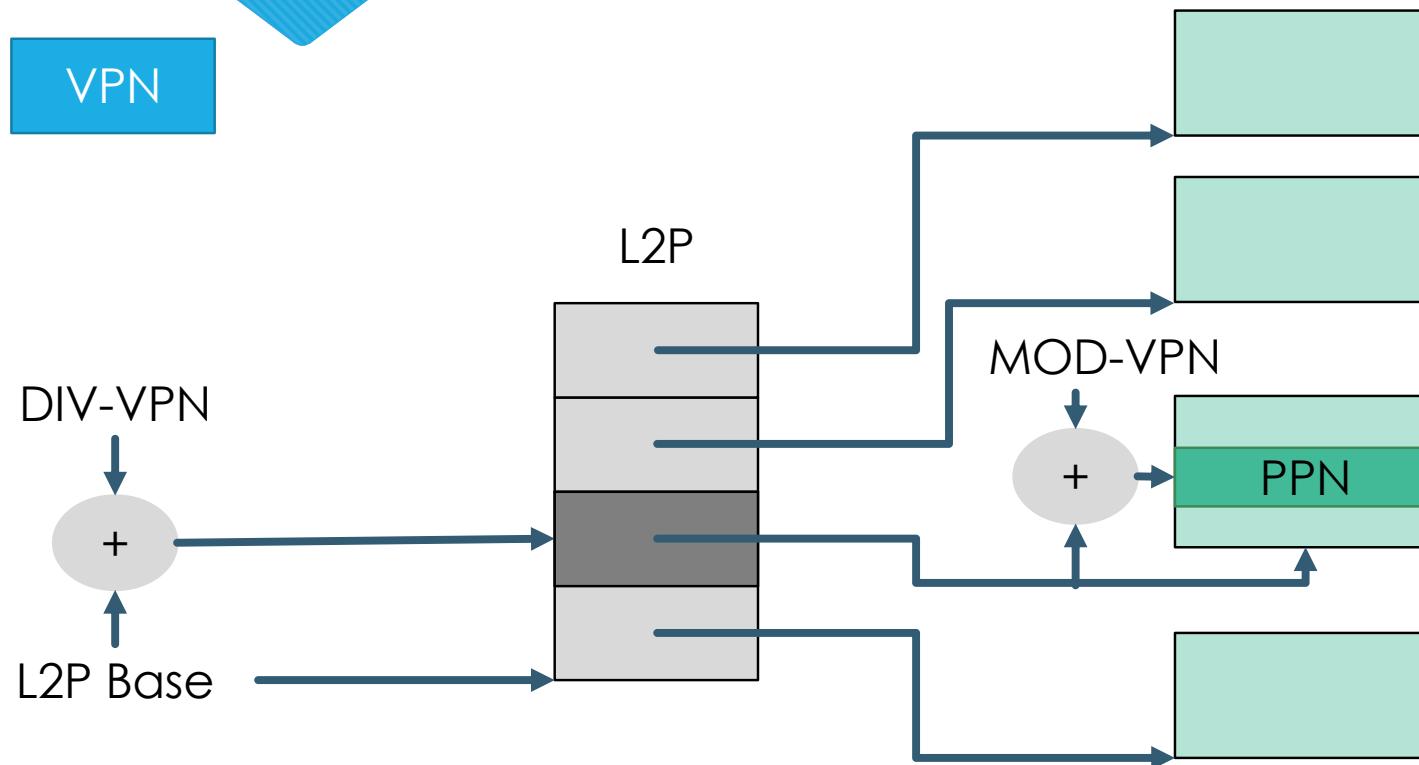
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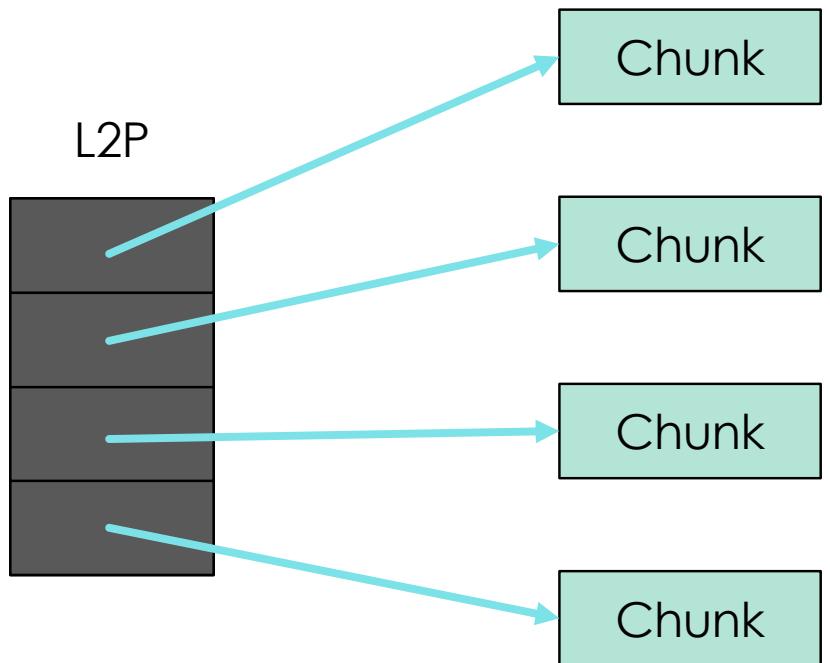
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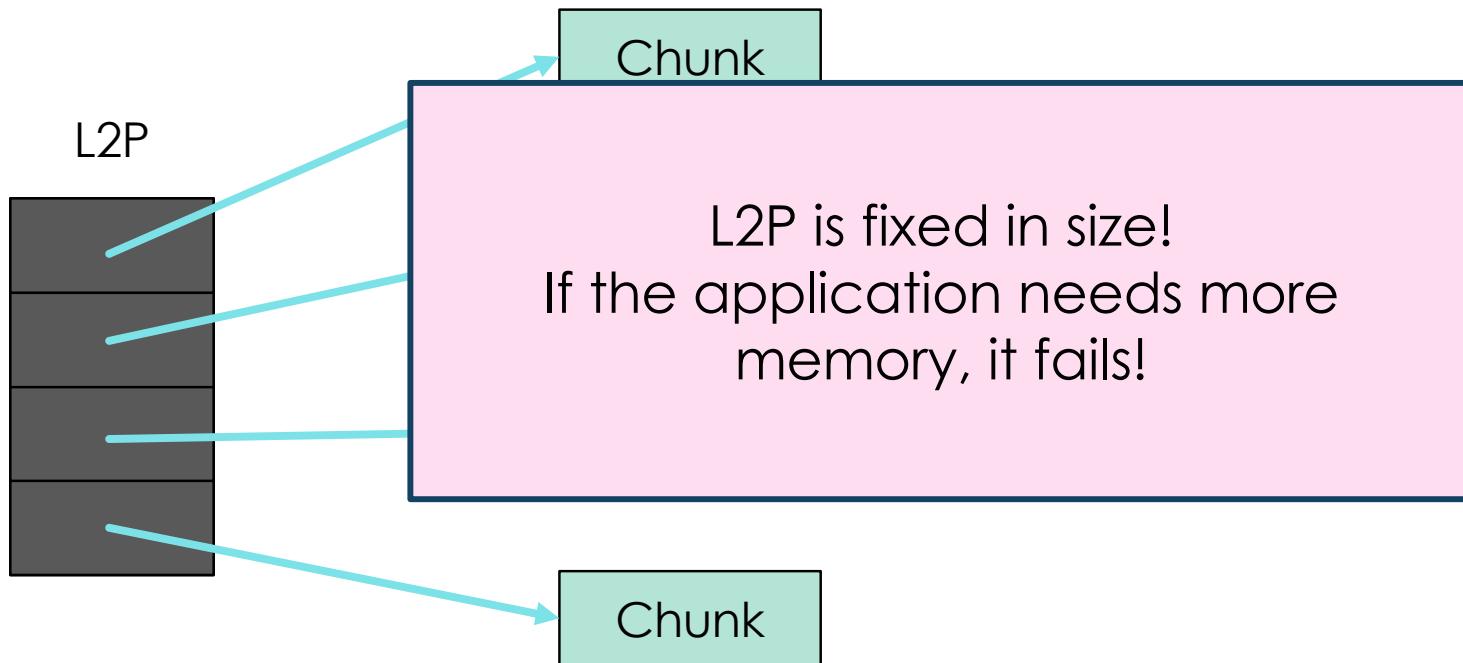
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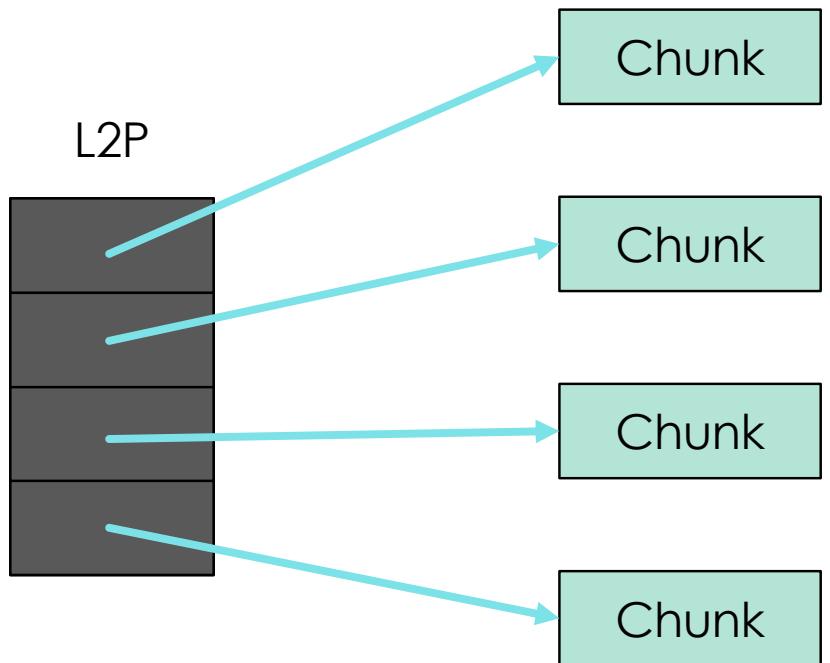
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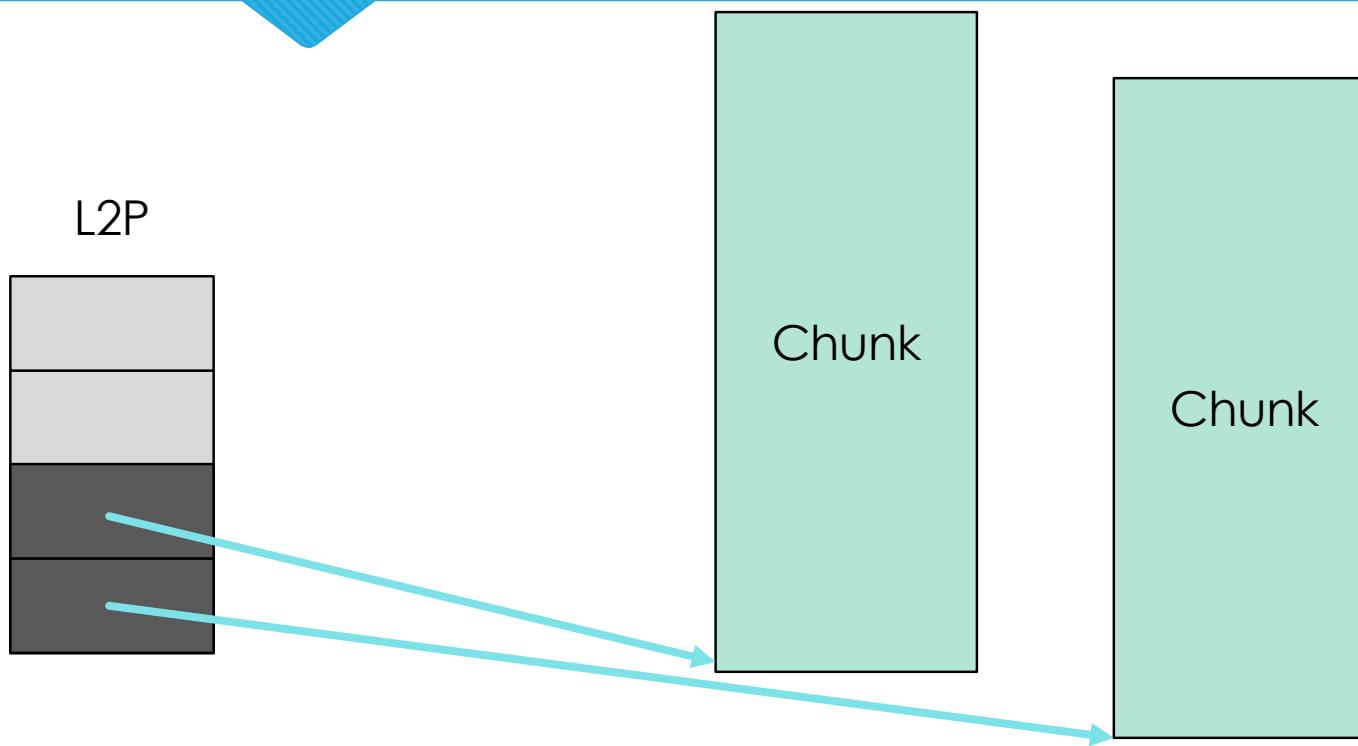
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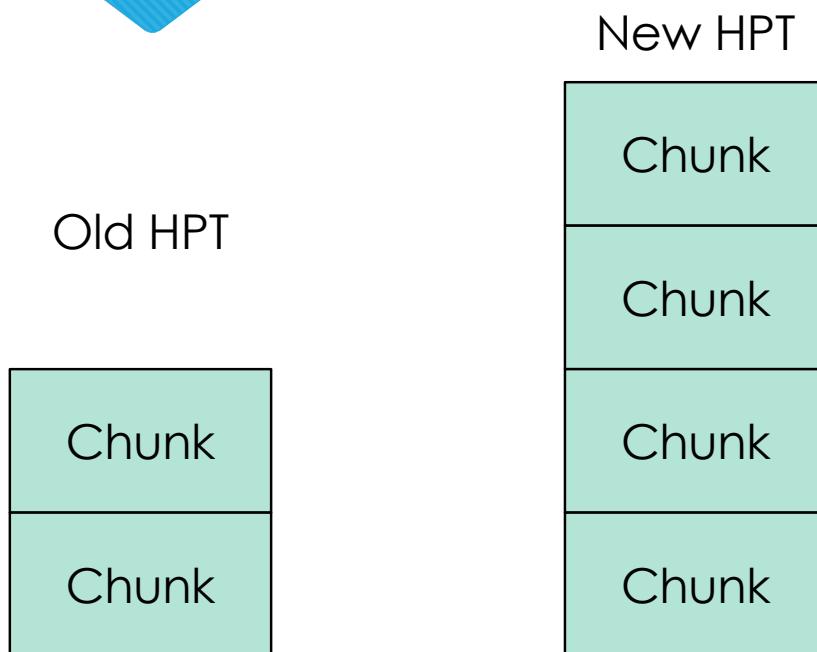
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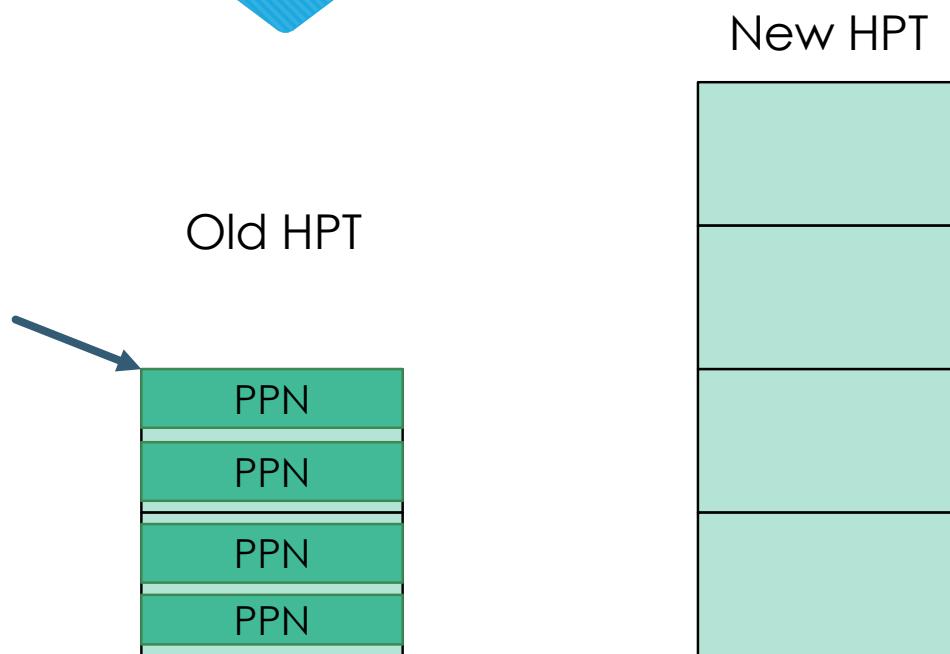
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Memory Efficient Hashed Page Tables: In-Place Page Table Resizing

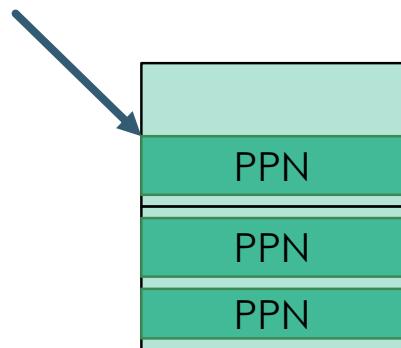


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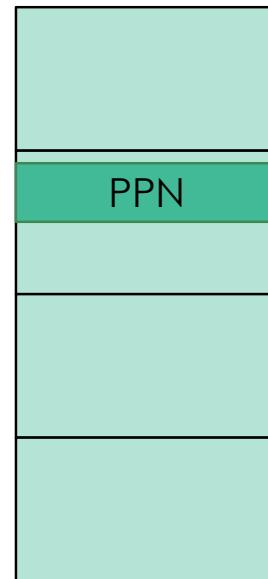


Memory Efficient Hashed Page Tables: In-Place Page Table Resizing

Old HPT

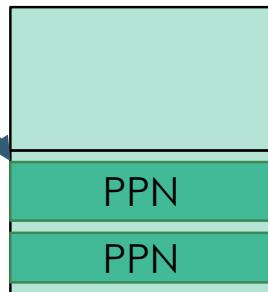


New HPT

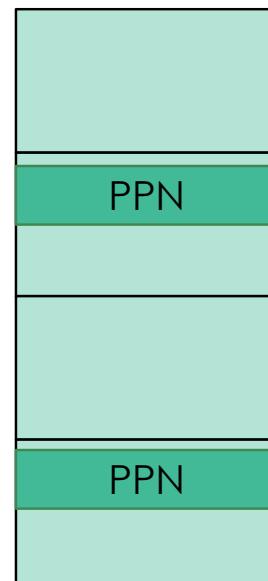


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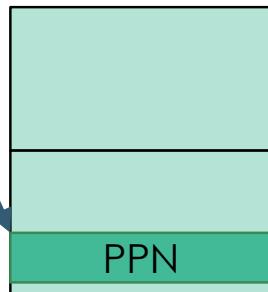


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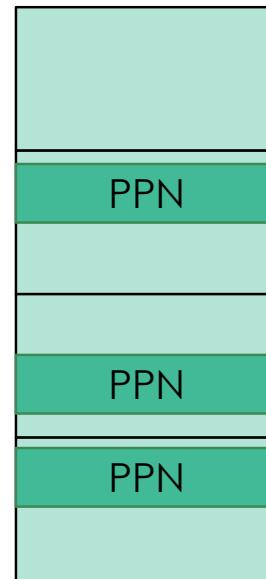


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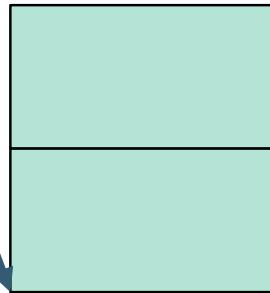


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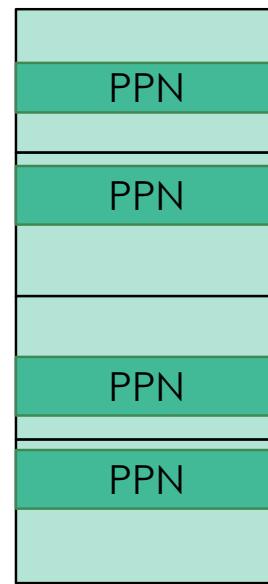


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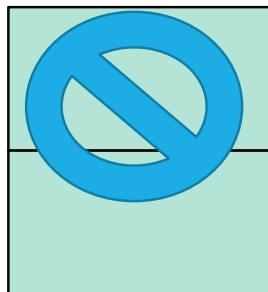
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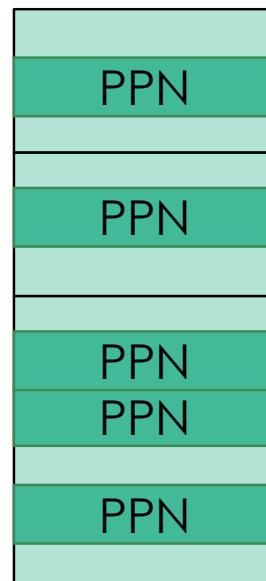
New HPT



Memory Efficient Hashed Page Tables: In-Place Page Table Resizing



Deallocate
old table!



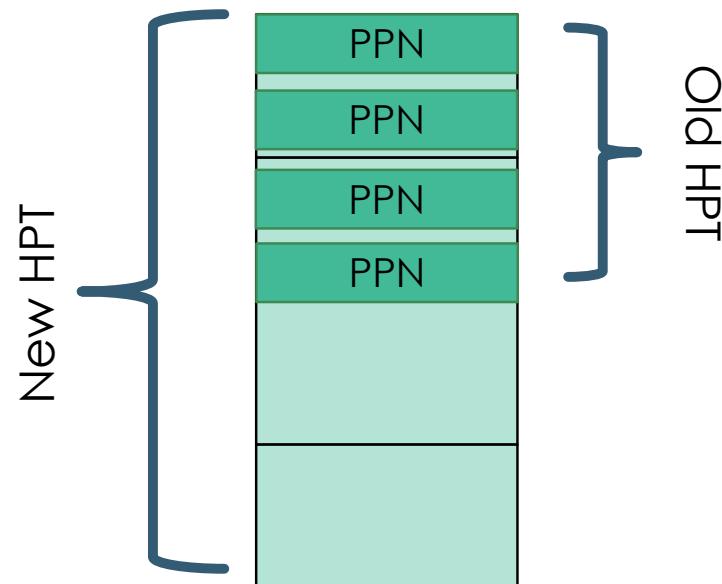
Memory Efficient Hashed Page Tables: In-Place Page Table Resizing

Until the old table is deallocated, we
keep **both tables** in memory!



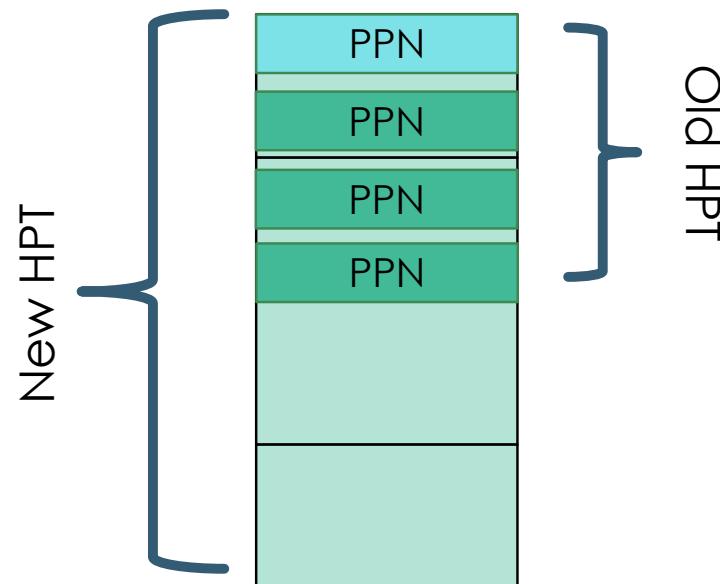
Memory Efficient Hashed Page Tables: In-Place Page Table Resizing

- Keep both tables in shared memory space
- Same hash function for both tables
- On rehash, some entries stay in the same chunk, others move to new chunks



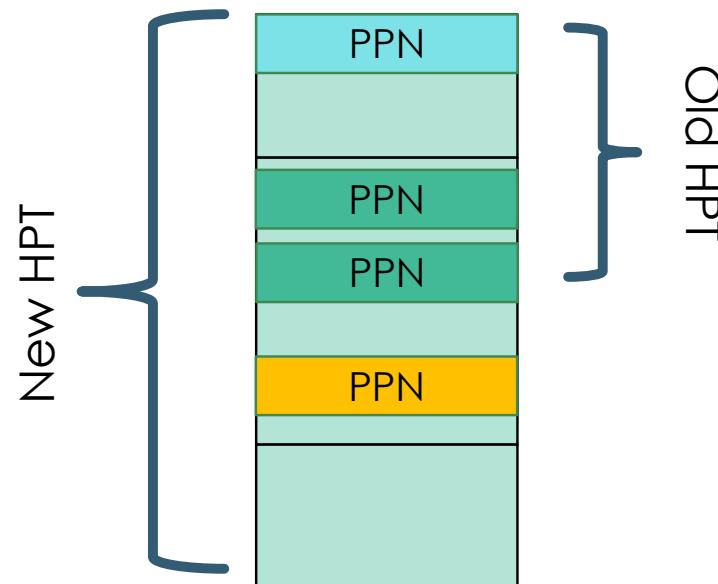
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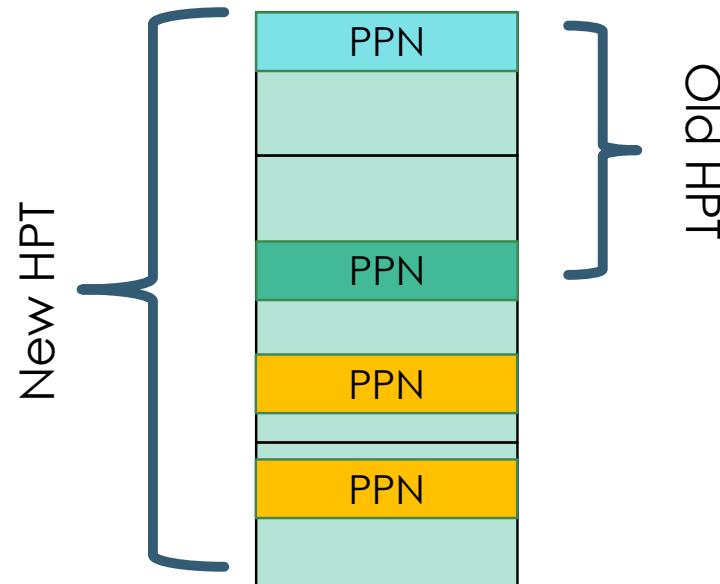
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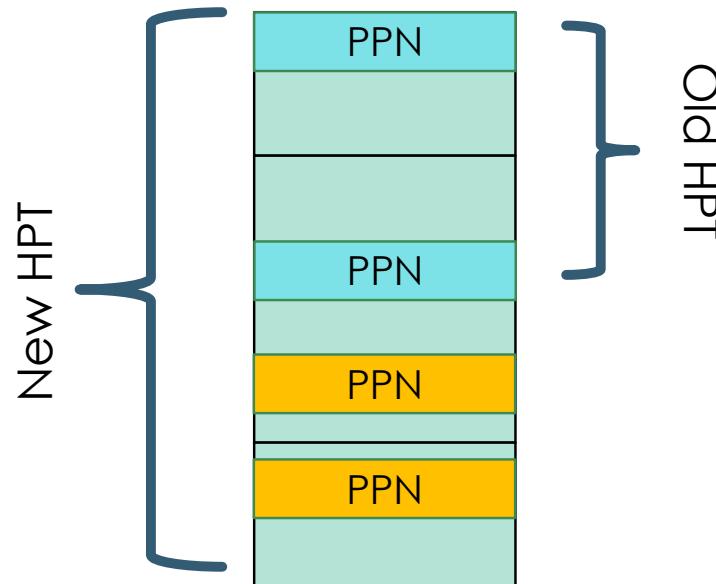
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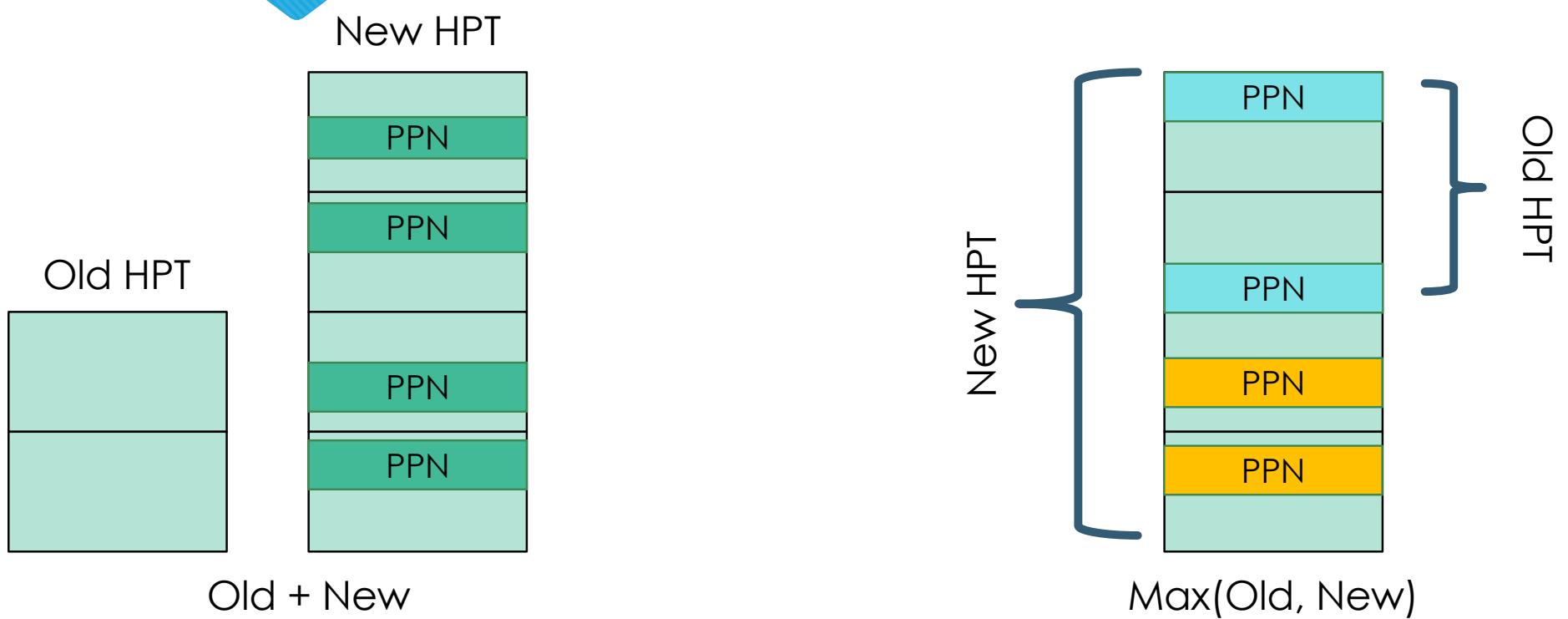


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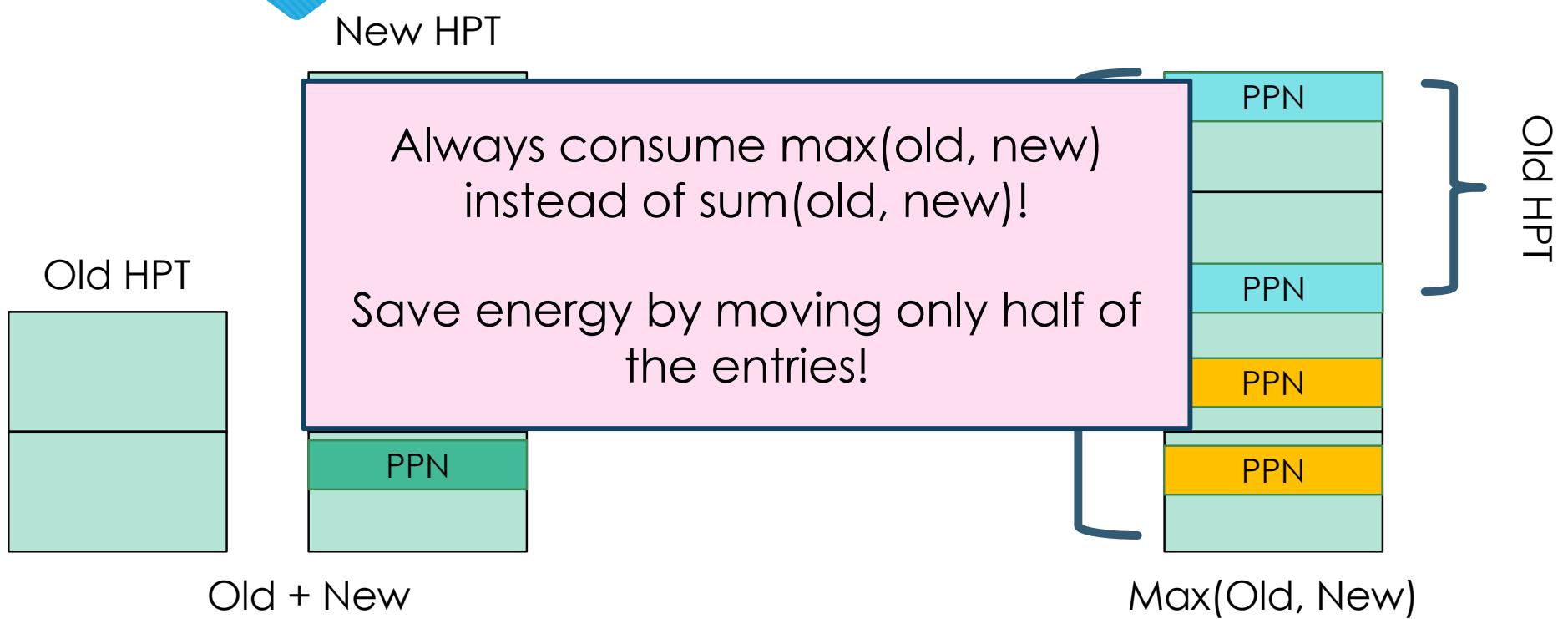
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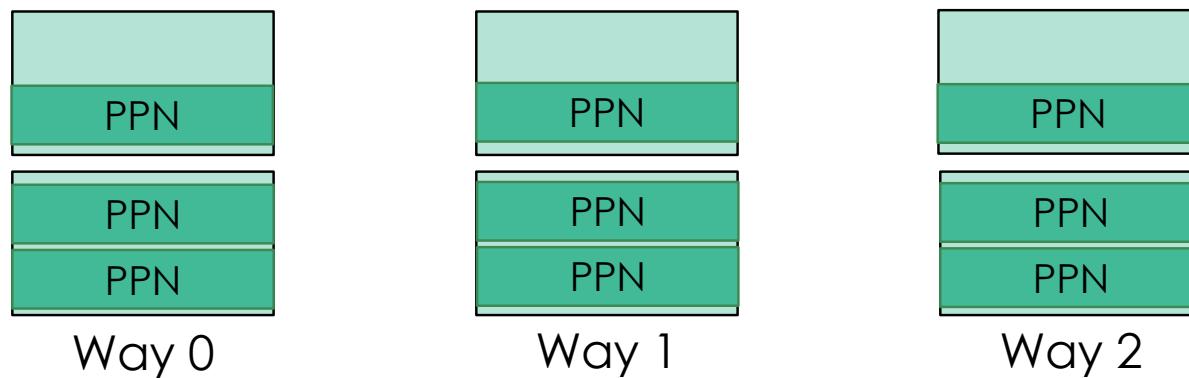
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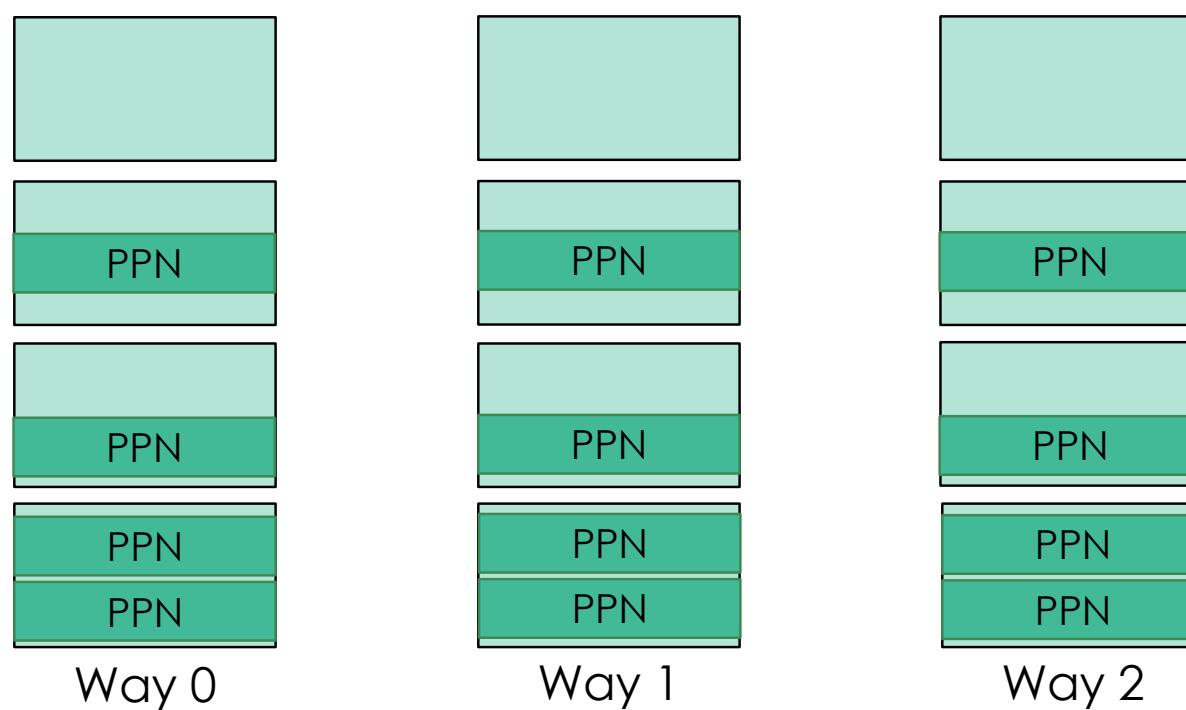
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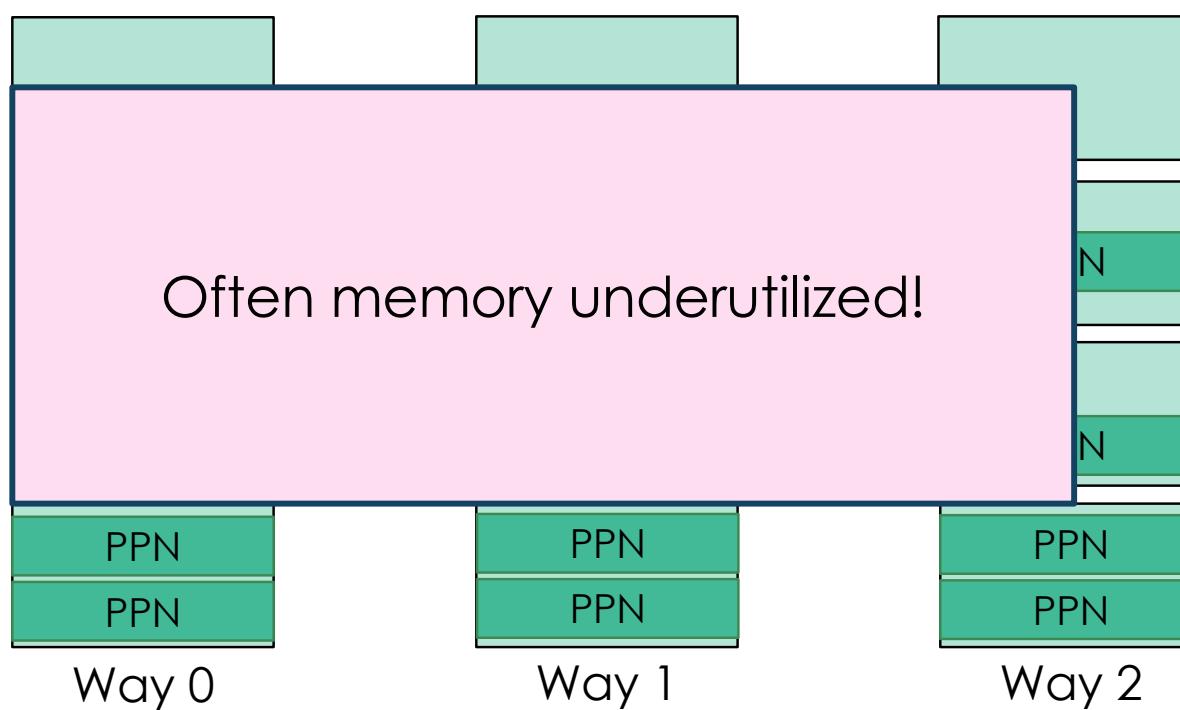
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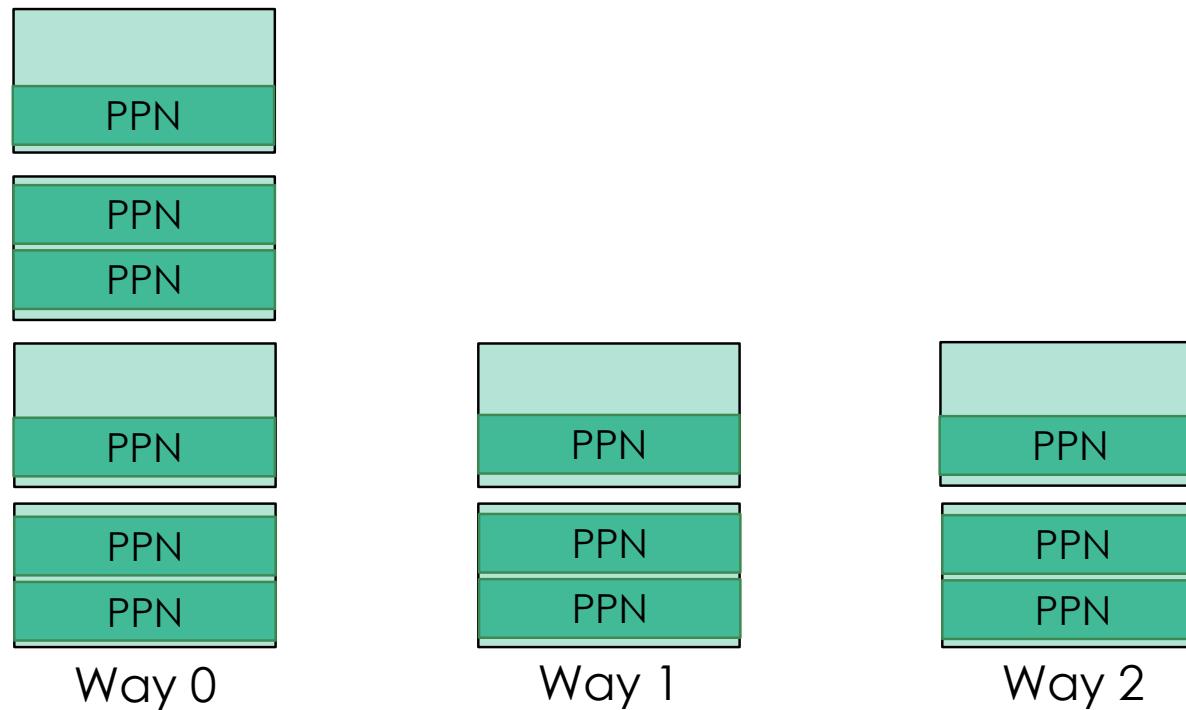
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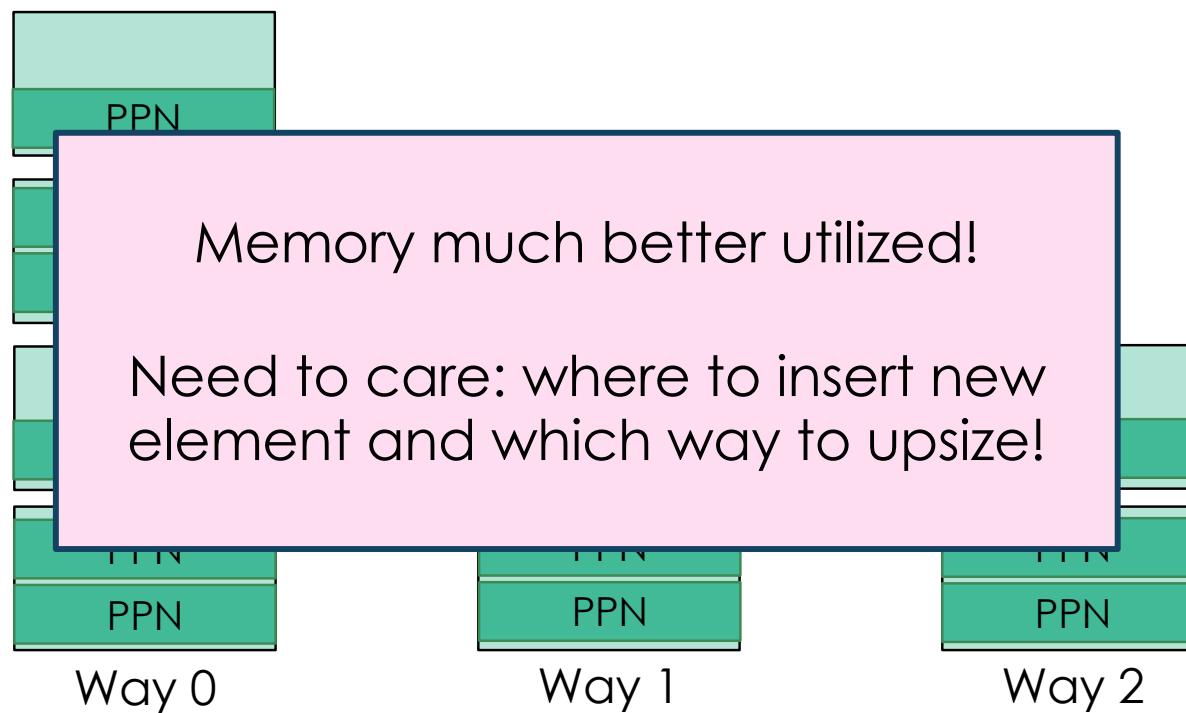
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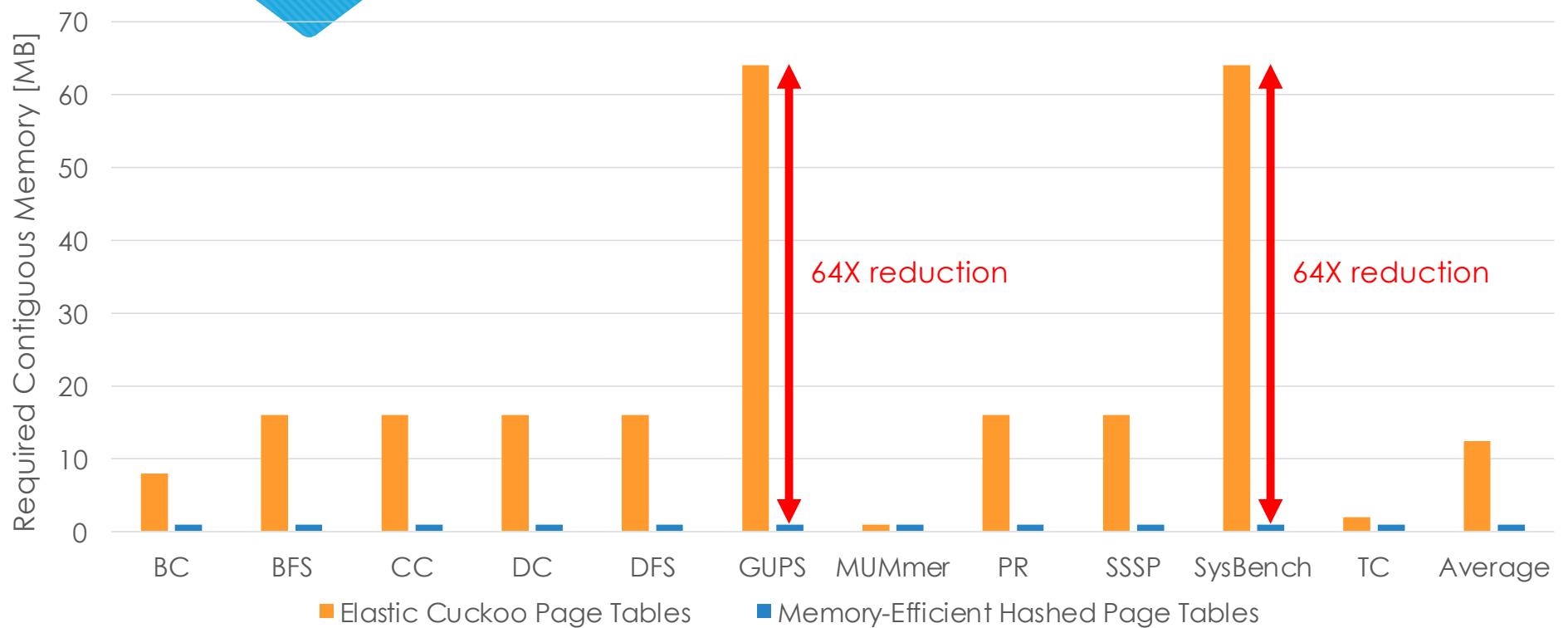
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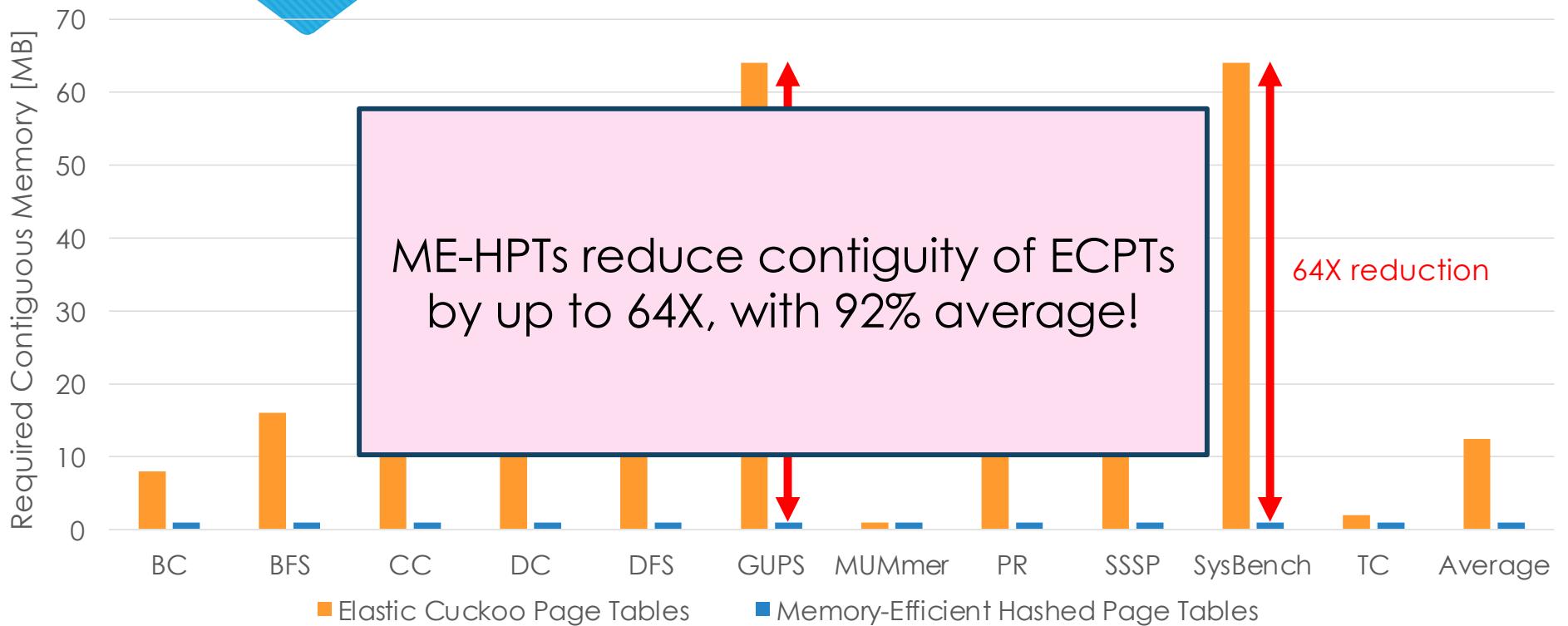
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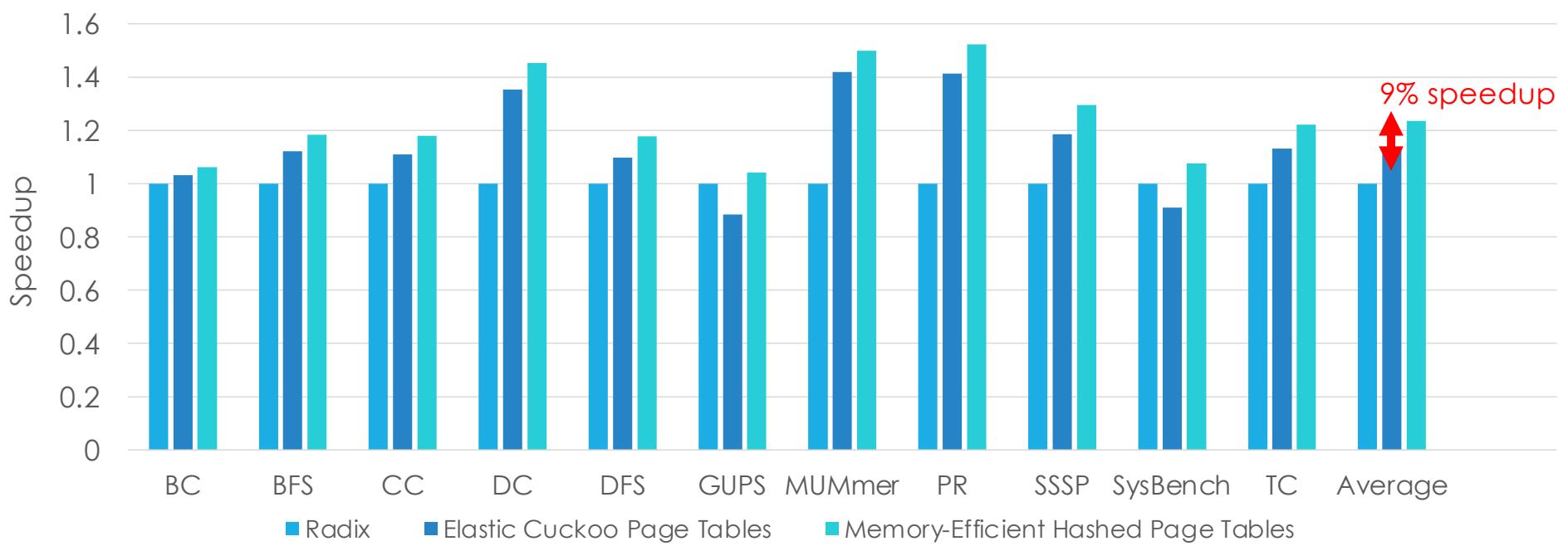
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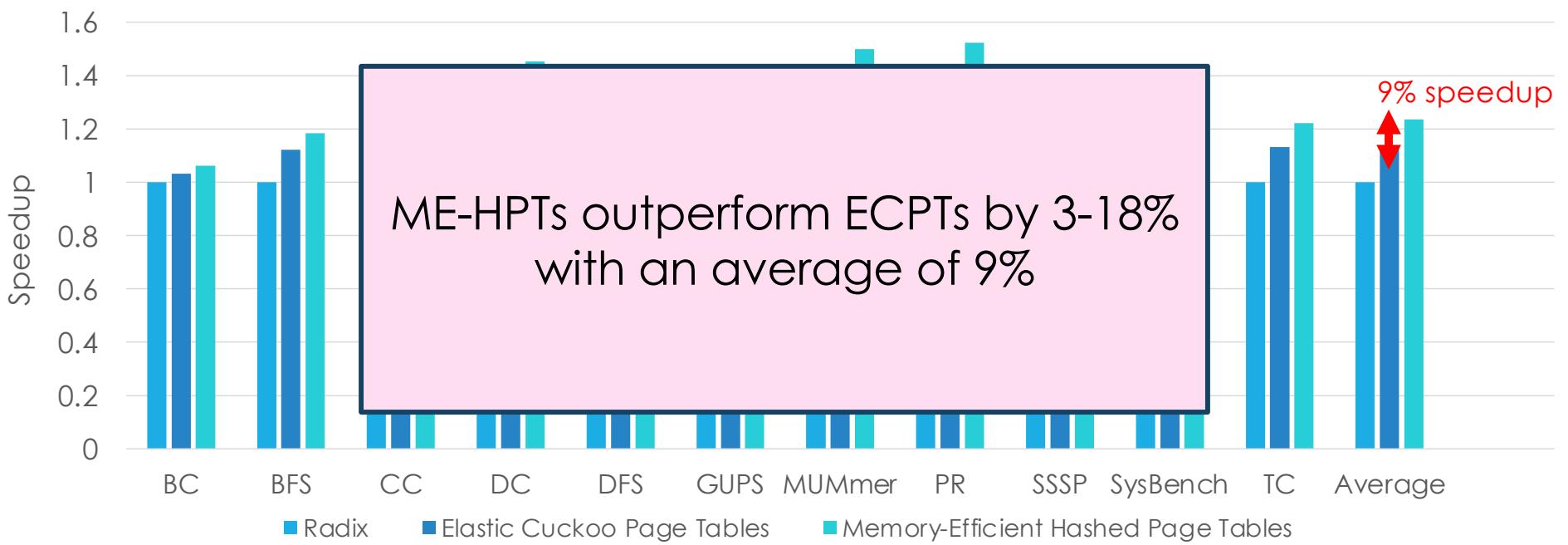
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Improved Application Performance



Improved Application Performance



Conclusion

- Four novel architectural techniques to provide Memory-Efficient Hashed Page Tables
 - L2P Table
 - Dynamically Changing Chunk Sizes
 - In-Place Page Table Resizing
 - Per-Way Page Table Resizing
- Reduced memory contiguity requirement by 92%
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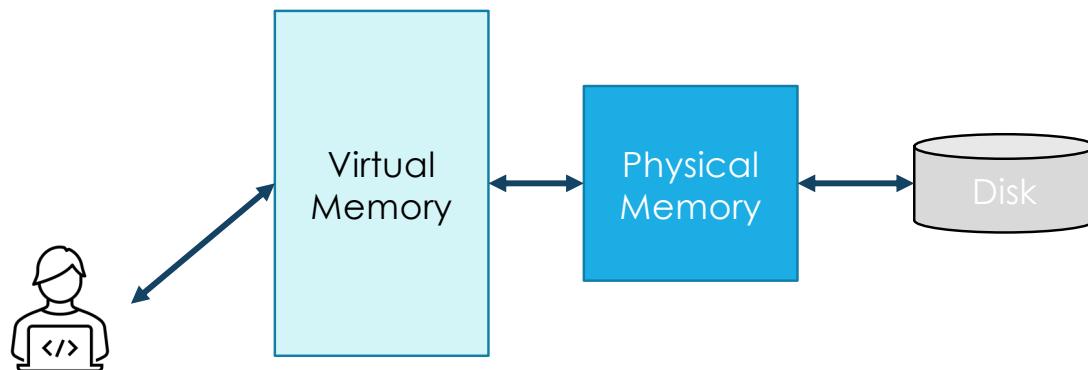
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Questions?

Backup Slides

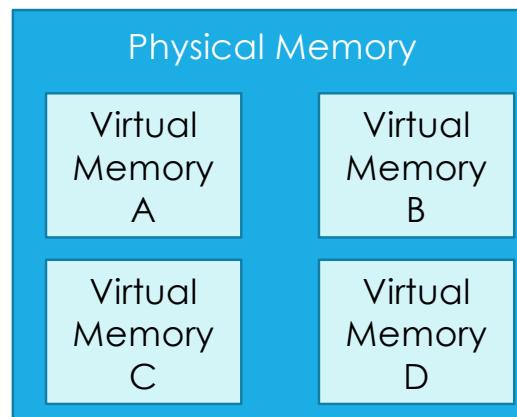
Virtual Memory Needs Memory-Efficient Hashed Page Tables

- Virtual memory is essential technique in modern computing systems
 - Memory virtualization



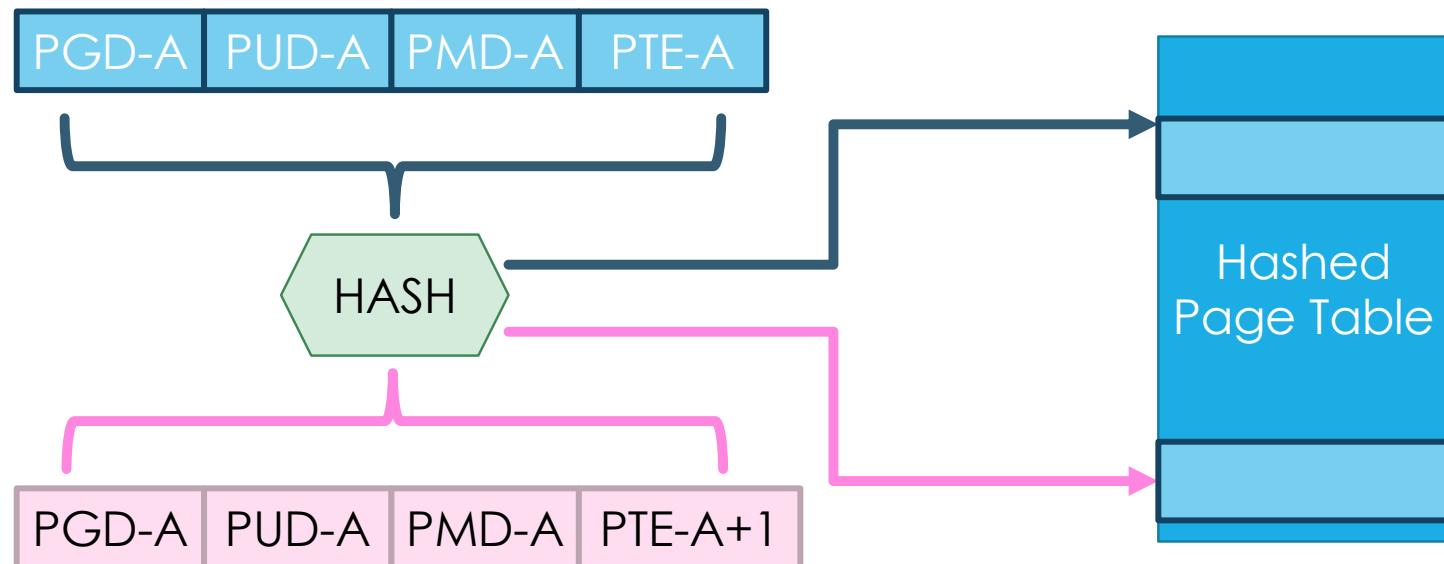
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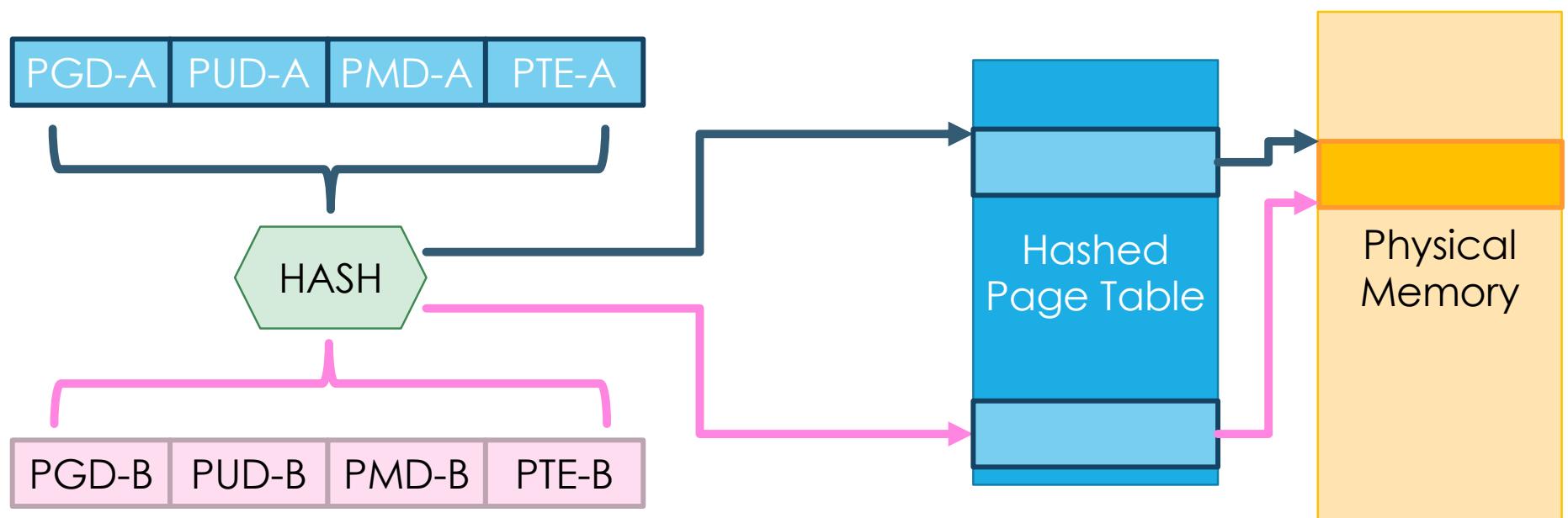
Hashed Page Tables

Loss of spatial locality



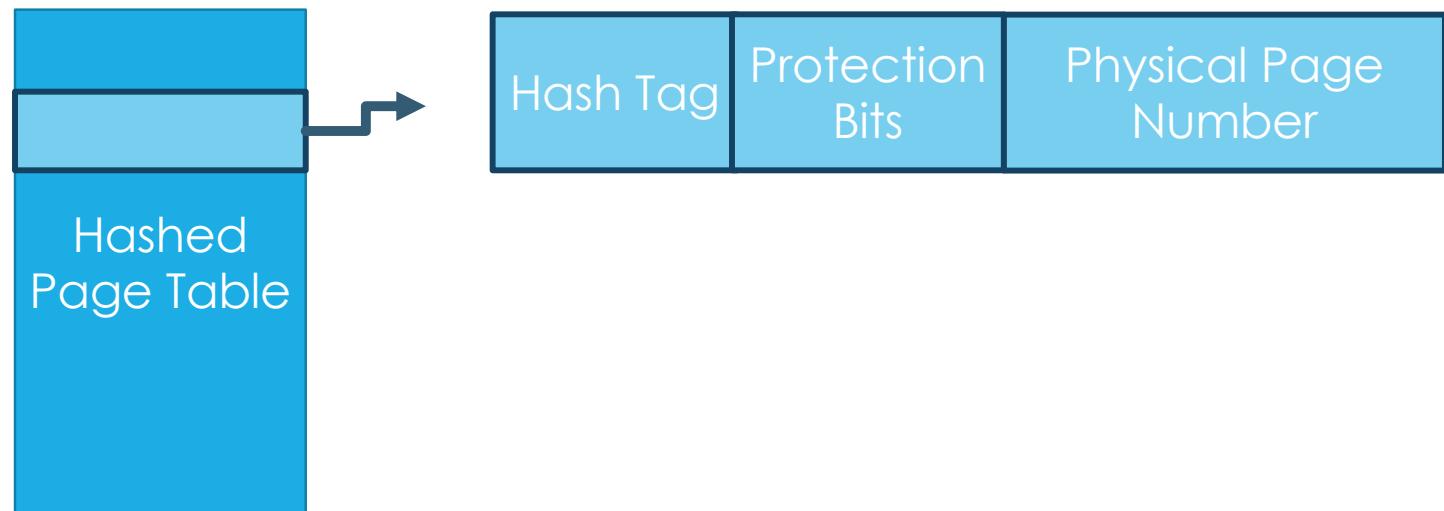
Hashed Page Tables

■ Page sharing and multiple page sizes not easy to support with global hashed table



Hashed Page Tables: One Step Forward, Two Steps Back

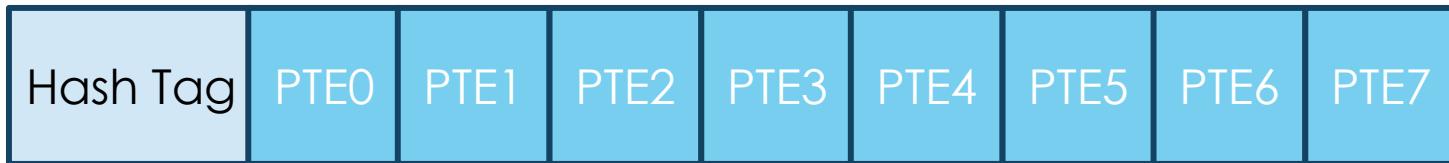
Extra space for hash tags



Hashed Page Tables: Recent Advances Make Them Compelling



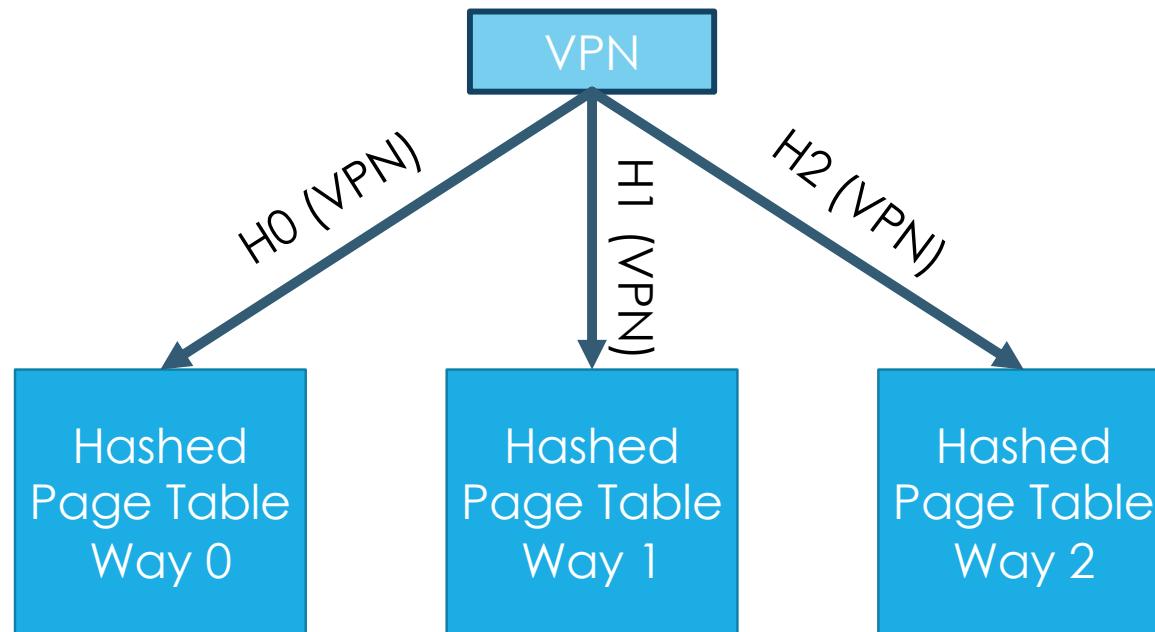
PTE Compaction and Clustering



Hashed Page Tables: Recent Advances Make Them Compelling

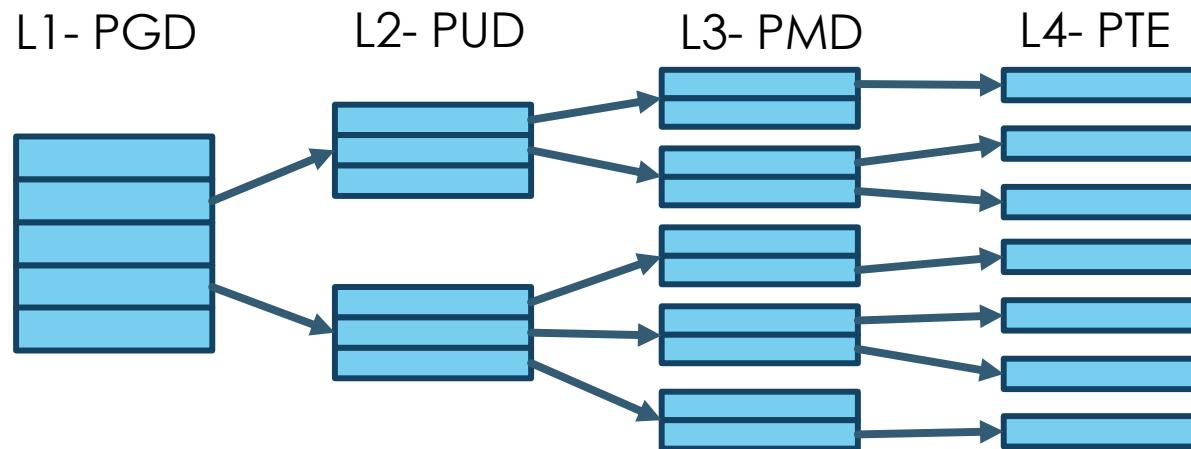


Cuckoo Hashing for Collision Handling

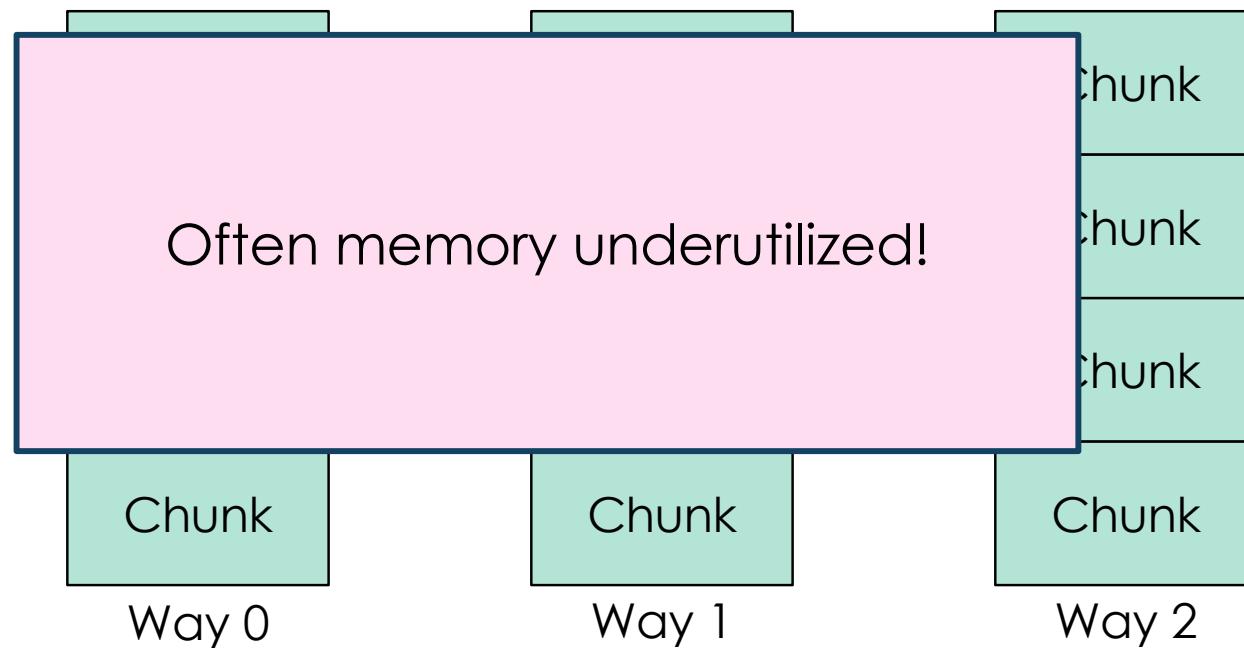


Hashed Page Tables: Large Contiguous Memory Chunks

- With radix page tables – unity of allocation is a 4KB page
 - L1 and each L2, L3 and L4 page tables are allocated independently



ME-HPTs: Per Way Page Table Resizing



ME-HPTs: Per Way Page Table Resizing

Memory much better utilized!

Need to care: where to insert new element and which way to upsize!

Chunk

Way 0

Chunk

Way 1

chunk

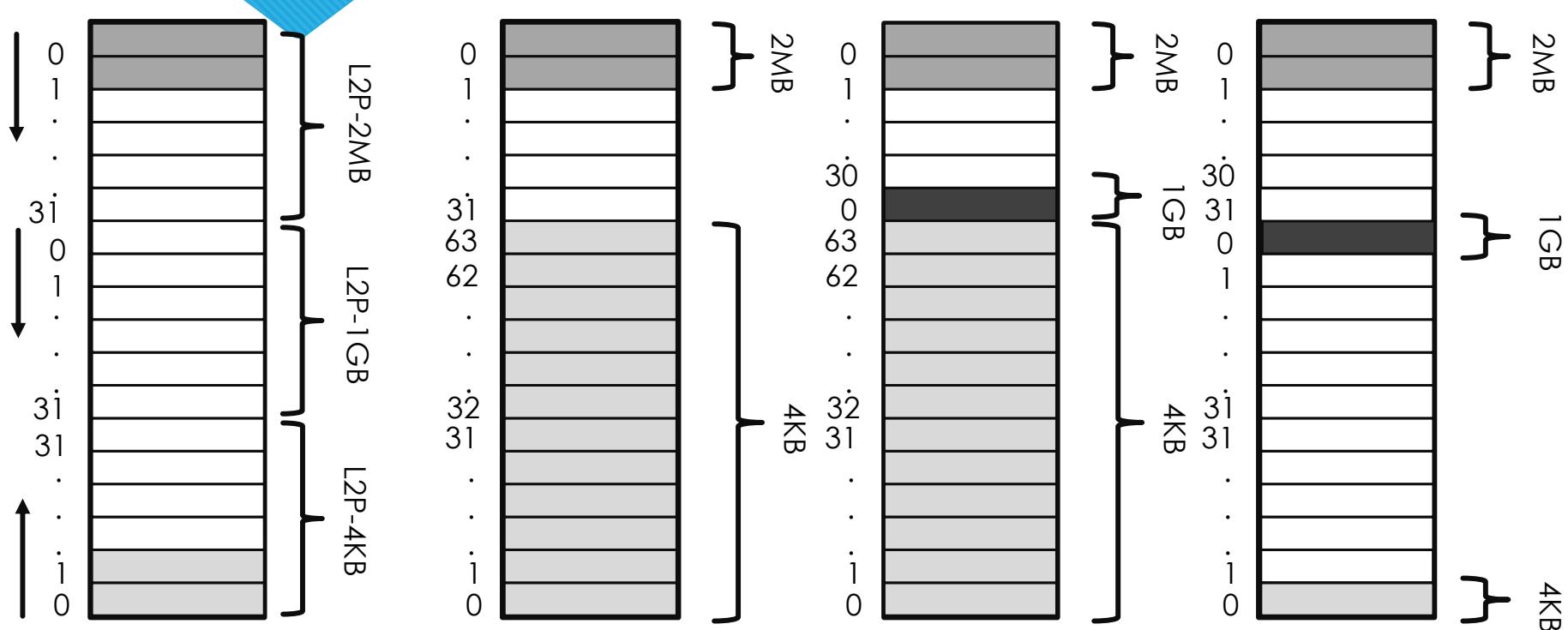
Chunk

Way 2

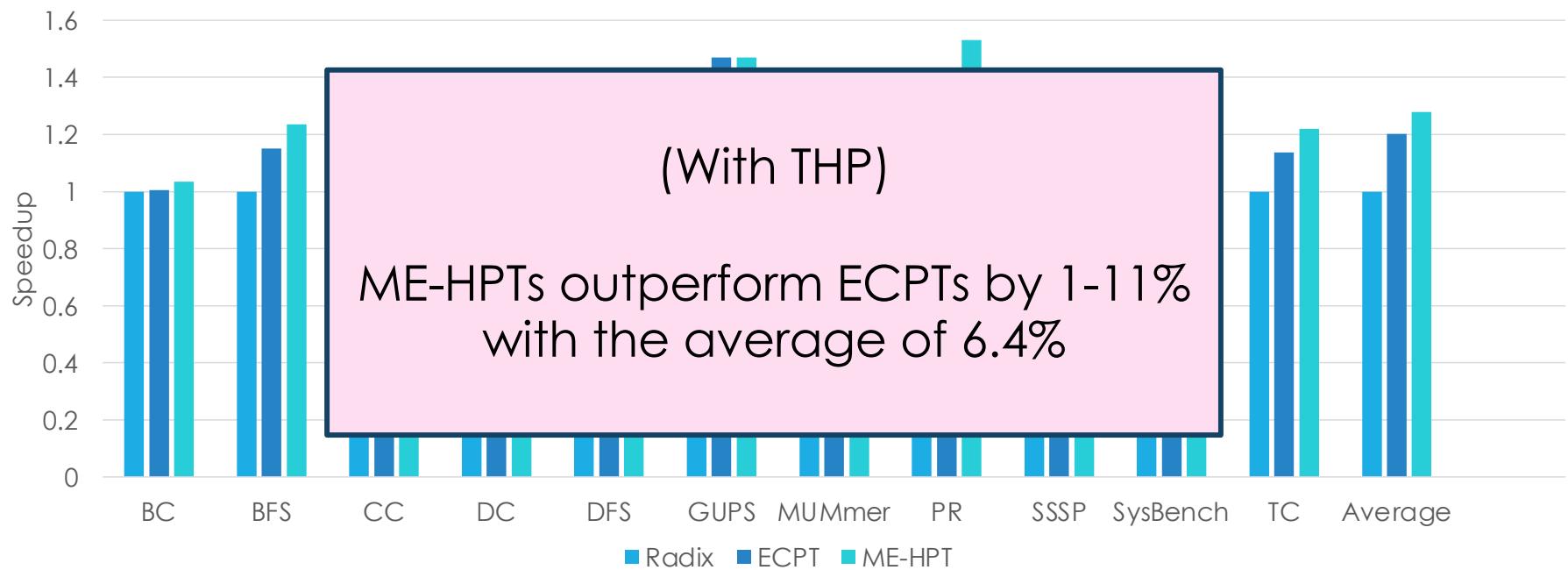
ME-HPTs Implementation: L2P Table Entry Stealing

- L2P Table is per page size of each page table way and its size is fixed
- Applications rarely use all page tables extensively
 - Some L2P tables will be less used than the others
- **Steal L2P entries from one L2P table and give them to another L2P table**

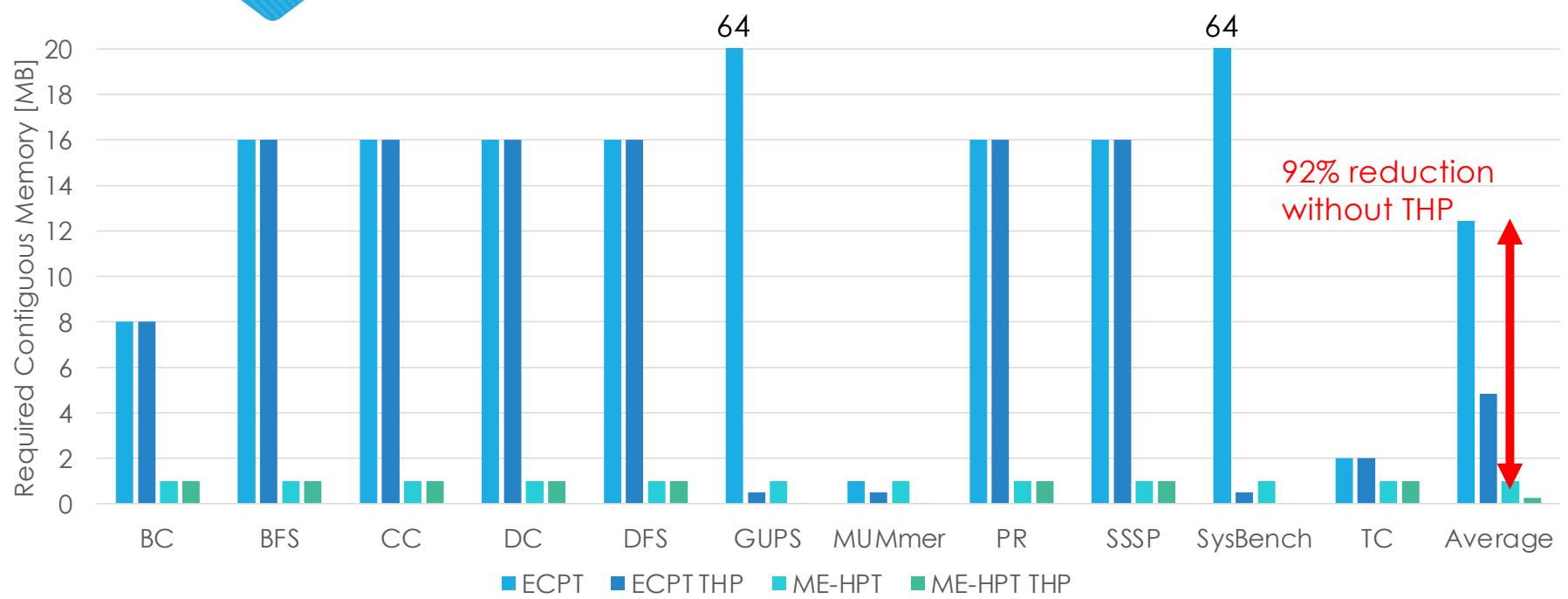
ME-HPTs Implementation: L2P Table Entry Stealing



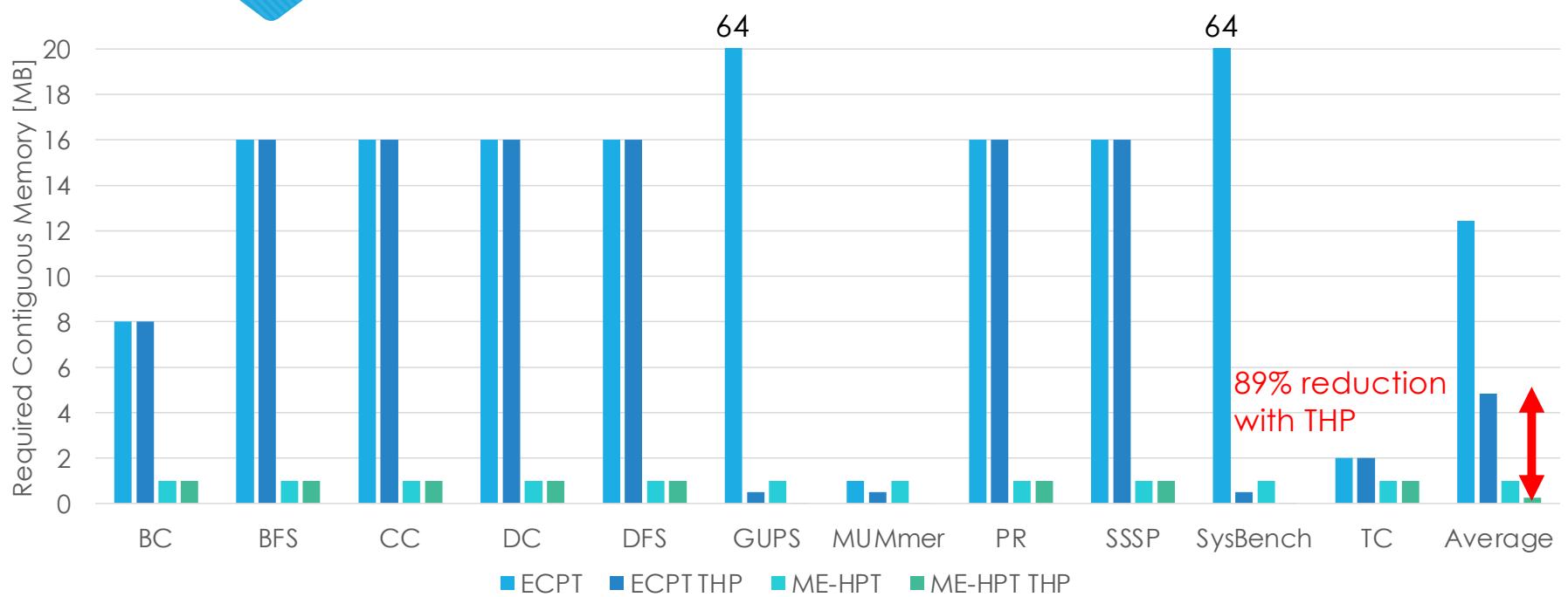
ME-HPTs Key Results: Improved Application Performance



ME-HPTs Key Results: Significant Memory Contiguity Savings



ME-HPTs Key Results: Significant Memory Contiguity Savings



ME-HPTs Key Results: Memory Consumption Reduction

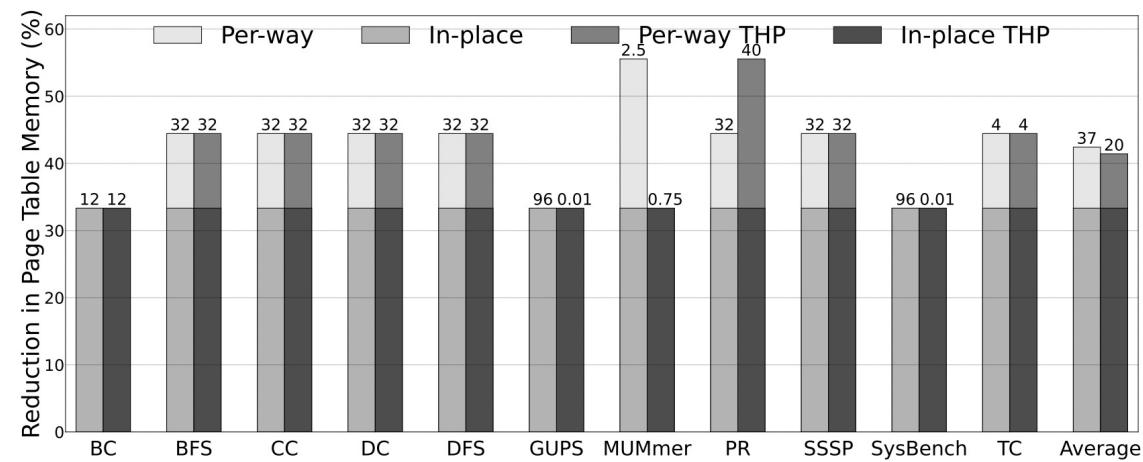


Fig. 10: Reduction in page table memory attained by ME-HPT over the ECPT baseline. The number on top of each bar is the absolute reduction in Mbytes.

ME-HPTs Key Results: Number of L2P Table Entries Used per App

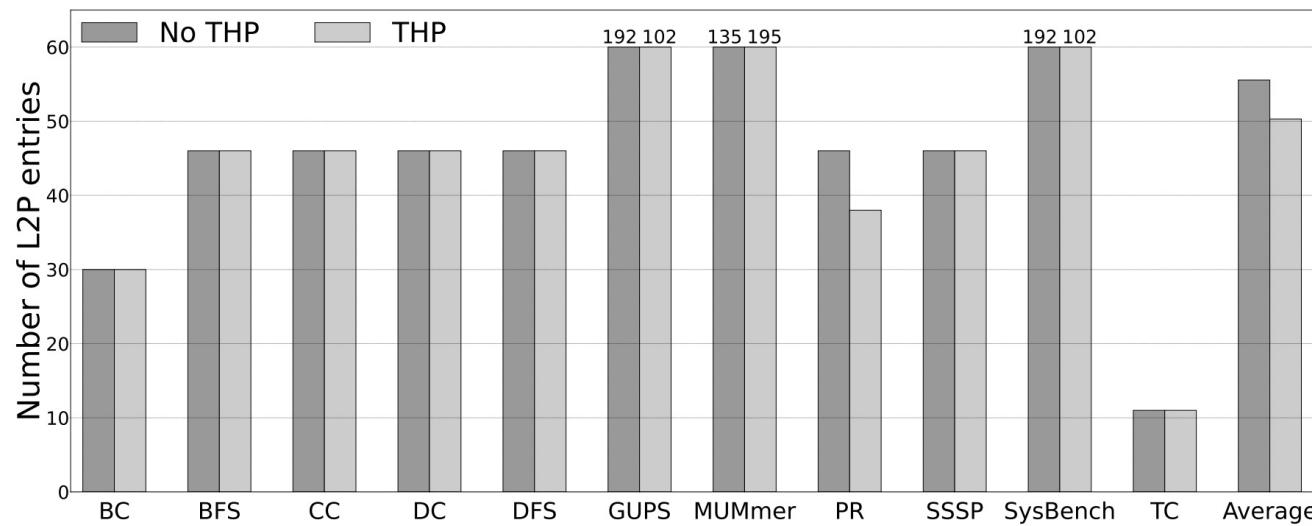


Fig. 14: Number of L2P table entries used per application.

ME-HPTs Other Use Cases

- Techniques applicable to various hash table designs beyond HPTs
- **Scalable Secure Directories**
 - Directories as set-associative structures
 - Efficient resizing required
- **Memory Indexing**
 - Hash tables commonly used to implement memory indices of databases, file systems...
 - Dynamic resizing key operation: in-place resizing useful
- **Key-value Stores**
 - Dynamic structures whose size is unknown ahead of time