

## Comments on “A Linear Time Algorithm for the Optimal Discrete IRS Beamforming”

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**Abstract**—Comments on [2][3] are provided. Updated necessary and sufficient conditions for its benign I are given. Consequently, Consolidated Algorithm is provided with full specification. Simulation results with improved performance over the implementation of Algorithm are provided. 1 are provided.

**Index Terms**—Intelligent reflective surface (IRS), reconfigurable intelligent surface (RIS), discrete beamforming for IRS/RIS.

## I. I. INTRODUCTION

Reference [3] presented an algorithm to solve the problem of finding the values  $\theta_1, \theta_2, \dots, \theta_N$  to maximize  $|h_0 + \sum_{n=1}^N h_n e^{j\theta_n}|$  where  $\theta_n \in \Phi_K$  and  $\Phi_K = \{\omega, 2\omega, \dots, K\omega\}$  with  $\omega = \frac{2\pi}{N}$  and  $j = \sqrt{-1}$ . The set  $\Phi_K$  can equivalently be described as  $\{0, \omega, 2\omega, \dots, (K-1)\omega\}$ . In [3], the values  $h_n \in \mathbb{C}$ ,  $n = 1, 2, \dots, N$  are the channel coefficients and  $\theta_n$  are the phase values added to the corresponding  $h_n$  by an intelligent reflecting surface (IRS), also known as a reconfigurable intelligent surface (RIS).

## II. TWO STATEMENTS FROM [?]1

Towards achieving its goal, [?] introduced the following lemma.

**Lemma 1:** For an optimal solution  $(\theta_1^*, \dots, \theta_n^*)$  to problem (8), each  $\theta_n^*$  must satisfy

$$\theta_n^* \equiv \arg \min_{\theta_n} \min_{\theta_n \in \mathbb{S}_K} |(\theta_n + \alpha_n - \underline{\mu}) \bmod 2\pi| \quad (11)$$

where  $\mu$  stands for the phase of  $\mu$  in (10)<sup>1</sup>.

In [?], problem (8) is defined as

$$\begin{aligned} \text{maximize } f(\theta) & \quad (8a) \\ \text{maximize } f(\theta) & \quad (8a) \end{aligned}$$

$$\text{subject to } \theta_n \in \Phi_K \text{ for } n = 1, 2, \dots, N \quad (8b)$$

where  
where

$$\frac{f(\theta)}{f(\theta_0)} = \frac{1}{\beta_0^2} \left| \beta_0 e^{j\alpha_0} + \sum_{n=1}^N \beta_n e^{j(\alpha_n + \theta_n)} \right|^2, \quad (7b)$$

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$h_n = \beta_n e^{j\alpha_n}$  for  $n = 0, 1, \dots, N$ , and  $\theta = (\theta_1, \theta_2, \dots, \theta_N)$ . Also,  $q$  is defined as

$$g \equiv h_0 + \sum_{n=1}^N h_n e^{j\theta_n^*} \quad (9)$$

and  $\mu$  as  
and  $\mu$  as

$$\mu = \frac{g}{|g|}. \quad (10)$$

Lemma 1 does not hold. This can be seen by numerical examples. We give one such example in Table ??.

We define

$$g_0(\theta_1, \theta_2) = h_0 + \sum_{n=1}^2 h_n e^{j\theta_n} \quad (19)$$

$$g_0(\theta_1, \theta_2) = h_0 + \sum_{n=1}^N h_n e^{j\theta_n} \quad (19)$$

and list the values of  $g_0(\theta_1, \theta_2)$  for all possible  $\theta_1, \theta_2 \in \{0, \pi\}$ . There are four such values and they are listed in rows 5–8 of Table ???. The set of values for  $\theta_1$  and  $\theta_2$  that maximize  $g_0$ , or equivalently, that achieve  $g$  in (9), are  $\theta_1 = \theta_2 = \pi$  as in row 8 of Table ???. Note that this operation results in  $\underline{\mu} = 2.3719$  radians as shown in column 5 of row 8 of Table ???.

At this point, we would like to emphasize that [?] uses a particular convention for the phases of complex numbers. They are defined to be in  $[0, 2\pi]$ , see the text that follows. At this point, we would like to emphasize that [?] uses a particular convention for the phases of complex numbers. They are defined to be in  $[0, 2\pi]$ , see the text that follows. We use the same convention in generating Table ??, see its column 5, as well as in generating Table ?. With this convention, we list  $\theta_n + \alpha_n - \frac{\mu}{n}$  and  $(\theta_n + \alpha_n - \frac{\mu}{n}) \bmod 2\pi$  for possibilities of  $\frac{\mu}{n} = 0$  and  $\theta_n = \pi$  and  $n = 1, 2$  in rows 1–8 of Table ??. It can be seen from rows 1–4 of Table ? that the method results in  $\theta_1 = \pi$  as the potential  $\theta_1^*$ , which we know from the discussion in the previous paragraph to be correct. When we carry out the calculation  $(\theta_2 + \alpha_2 - \frac{\mu}{n}) \bmod 2\pi$  in rows 5–8 of Table ?, we find that the method suggests  $\theta_2 = 0$  should be  $\theta_2^* = 0$ . However, we know from the exhaustive search in rows 5–8 of Table ? that  $\theta_2^* = \pi$ . Thus, Lemma 1 is not correct.

It is possible to come up with a correct lemma similar to Lemma 1. We specify this lemma below.

<sup>2</sup>Note that absolute values  $\text{sign}(l)$  in  $\text{arg}(1)$  are not needed since  $\text{arg}(1)$  is in  $[0, 2\pi)$ .





Since, Proposition 1 is compatible with Lemma 2. To see this, assume  $\mu$  satisfies (12). Then,

$$\alpha_n - \mu \in \left( (-2k-1)\frac{\pi}{K}, (-2k+1)\frac{\pi}{K} \right) \quad (26)$$

$$\mu \in \left( \alpha_n + \left( k - \frac{1}{2} \right) \omega, \alpha_n + \left( k + \frac{1}{2} \right) \omega \right) \quad (25)$$

considering the reversal of order due to the subtraction of  $\mu$ . Now let  $\theta_n = \frac{\theta_n - \mu}{K}$ ,  $k\omega = 2k\frac{\pi}{K}$ . Then

$$\alpha_n - \frac{\theta_n - \mu}{K} \in \left( \frac{\pi}{K}, \left( \frac{\pi}{K} - \frac{\pi}{K} \right) \frac{\pi}{K} \right) \quad (27)$$

and thus  $\cos(\theta_n + \alpha_n - \frac{\theta_n - \mu}{K})$  is the largest among all other possibilities for  $\theta_n$  because the slice  $(\frac{\pi}{K}, \frac{\pi}{K})$  corresponds to the largest values of the cosine function among all slices corresponding to different values of  $\theta_n \in (\frac{\pi}{K}, \frac{\pi}{K})$   $k = 1, 2, \dots$  (27)

and thus  $\cos(\theta_n + \alpha_n - \frac{\theta_n - \mu}{K})$  is the largest among all other possibilities for  $\theta_n$ . Algorithm 2 is the proposed algorithm to find the largest value of the cosine function among all slices corresponding to different values of  $\theta_n$  and  $\alpha_n$  present in Algorithm 1 of [?].

### Algorithm 2 Update for Algorithm 1

We now specify Algorithm 2 to replace Algorithm 1 in [?]. In doing so, not only do we incorporate Lemma 2 into the algorithm but also we eliminate the many uncertainties present in Algorithm 1 of [?].

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1: Initialization: Compute  $s_{nk} = e^{j(\alpha_n + (k-0.5)\omega)}$  for  $n = 1, 2, \dots, N$  and  $k = 1, 2, \dots, K$ .
2: Eliminate duplicates among  $s_{nk}$  and sort to get  $0 \leq \lambda_1 < \lambda_2 < \dots < \lambda_L < 2\pi$ .
3: Let, for  $l = 1, 2, \dots, L$ ,  $\mathcal{N}(\lambda_l) = \{n | s_{nk} = \lambda_l\}$ .
4: Set  $\mu = 0$ . For  $n = 1, 2, \dots, N$ , calculate  $\theta_n = \arg \max_{\theta_n \in \Phi_K} \cos(\theta_n + \alpha_n - \frac{\theta_n - \mu}{K})$ .
5: Set  $g_1 = h_0 + \sum_{n \in \mathcal{N}(\lambda_1)} h_n e^{j\theta_n}$ ,  $\text{absmax} = 0$ .
6: for  $l = 2, 3, \dots, L$  then
7:   For each  $n \in \mathcal{N}(\lambda_l)$  let  $(\theta_n + \omega \leftarrow \theta_n) \bmod \Phi_K$ .
8:   Let  $\theta_n = \theta_n + \omega$ .
9:   for  $l = 2, 3, \dots, L$  then
10:    For each  $n \in \mathcal{N}(\lambda_l)$  let  $(\theta_n + \omega \leftarrow \theta_n) \bmod \Phi_K$ .
11:    Let  $\theta_n = \theta_n + \omega$ .
12:   end if
13: end for
14: Read out  $\theta_n^*$  as the stored  $\theta_n$ ,  $n = 1, 2, \dots, N$ .
15: if  $|g_l| > \text{absmax}$  then
16:   Let  $\text{absmax} = |g_l|$ 
17:   Store  $\theta_n^*$  for  $n = 1, 2, \dots, N$ 
18: end if
19: end for
20: Read out  $\theta_n^*$  as the stored  $\theta_n$ ,  $n = 1, 2, \dots, N$ .

```

Because its description is based on Lemma 1, which does not provide an equivalency condition for finding  $\theta_n^*$ , the performance of Algorithm 2 will in general not achieve the optimum result for SNR Boost [?].

We have implemented Algorithm 1 to the best of our interpretation. We have also implemented Algorithm 2. We present the CDF results for SNR Boost in Fig. 3, which does not provide a clear indication of the average of 1000 realizations of the channel performance. Algorithm 1 is not optimal. Algorithm 2 performs better than Algorithm 1 for SNR Boost. The gains decrease with  $N$ . Plots for  $K = 4$  show smaller gains as compared to  $K = 2$ . Algorithm 2 always performs better than Algorithm 1. We have also implemented Algorithm 2. We

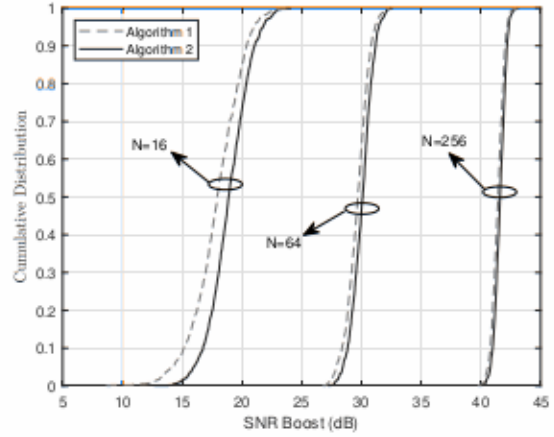


Fig. 1: CDF plots for SNR Boost. Algorithm 1 and Algorithm 2,  $K = 2$ .

We note that CDF is possible for SNR Boost. The maximum value of  $K$  is 2 and  $N$  is 16. The mid-256th percentile of the sample is 1000. For example, of the maximization of  $\cos(x)$ , Algorithm 2 is not the same as Algorithm 1. Clearly, Algorithm 2 is not the same as Algorithm 1. However, this is different. Plots for  $K = 2$  and  $N = 16$  show smaller gains proposed in [?] and [?]. Algorithm 2 is shown by plotting these functions against  $x$  and  $N$ . In addition to being periodic with period  $2\pi$ , have even symmetry around odd multiples of the maximization of  $\cos(x)$ . For example, minimization of  $f_1(x) = \pi - |(x \bmod 2\pi) - \pi|$  is the same as maximization of  $\cos(x)$  within the context of Lemma 2. However, this is different than minimization of  $|x \bmod 2\pi|$  proposed in Lemma 1 of [?]. The reason can be seen by plotting these functions against  $x$ . While  $f_1(x)$  and  $\cos(x)$ , in addition to being periodic with period  $2\pi$ , have even symmetry around odd multiples of  $\pi$ ,  $|x \bmod 2\pi|$  (or equivalently,  $(x \bmod 2\pi)$ ) does not have this symmetry.