

# Comments on “A Linear Time Algorithm for the Optimal Discrete IRS Beamforming”

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Abstract—Comments on [?] are provided. Updated necessary and sufficient conditions for its Lemma 1 are given. Consequently, an updated Algorithm 1 is provided with full specification. Simulation results with improved performance over the implementation of Algorithm 1 are provided.

Index Terms—Intelligent reflective surface (IRS), reconfigurable intelligent surface (RIS), discrete beamforming for IRS/RIS.

## I. Introduction

Reference [?] presented an algorithm to solve the problem of finding the values  $\theta_1, \theta_2, \dots, \theta_N$  to maximize  $|h_0 + \sum_{n=1}^N h_n e^{j\theta_n}|$  where  $\theta_n \in \Phi_K$  and  $\Phi_K = \{\omega, 2\omega, \dots, K\omega\}$  with  $\omega = \frac{2\pi}{K}$  and  $j = \sqrt{-1}$ . The set  $\Phi_K$  can equivalently be described as  $\{0, \omega, 2\omega, \dots, (K-1)\omega\}$ . In [?], the values  $h_n \in \mathbb{C}$ ,  $n = 1, 2, \dots, N$  are the channel coefficients and  $\theta_n$  are the phase values added to the corresponding  $h_n$  by an intelligent reflective surface (IRS), also known as reconfigurable intelligent surface (RIS).

## II. Two Statements from [?]

Towards achieving its goal, [?] introduced the following lemma.

Lemma 1: For an optimal solution  $(\theta_1^*, \dots, \theta_N^*)$  to problem (8), each  $\theta_n^*$  must satisfy

$$\theta_n^* = \arg \min_{\theta_n \in \Phi_K} |(\theta_n + \alpha_n - \underline{\mu}) \bmod 2\pi| \quad (11)$$

where  $\underline{\mu}$  stands for the phase of  $\mu$  in (10)<sup>1</sup>.

In [?], problem (8) is defined as

$$\underset{\boldsymbol{\theta}}{\text{maximize}} \quad f(\boldsymbol{\theta}) \quad (8a)$$

$$\text{subject to } \theta_n \in \Phi_K \quad \text{for } n = 1, 2, \dots, N \quad (8b)$$

where

$$f(\boldsymbol{\theta}) = \frac{1}{\beta_0^2} \left| \beta_0 e^{j\alpha_0} + \sum_{n=1}^N \beta_n e^{j(\alpha_n + \theta_n)} \right|^2, \quad (7b)$$

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<sup>1</sup>To prevent confusion, we will use the same equation numbers (7)–(13) in [?] in our equation numbers as available in [?], will be (19)–(24) and will be incremented from that number. Similarly, we will not use Lemma 2 and Algorithm 2 in [?] as Lemma 1 and Algorithm 1 in [?]. Note that a lemma or an algorithm with number 2 does not exist in [?].

$h_n = \beta_n e^{j\alpha_n}$  for  $n = 0, 1, \dots, N$ , and  $\boldsymbol{\theta} = (\theta_1, \theta_2, \dots, \theta_N)$ . Also,  $g$  is defined as

$$g = h_0 + \sum_{n=1}^N h_n e^{j\theta_n} \quad (9)$$

and  $\mu$  as

$$\mu = \frac{g}{|g|}. \quad (10)$$

Lemma 1 does not hold. This can be seen by numerical examples. We give one such example in Table ?? In this table, we look at the simple case of  $K = 2$ ,  $N = 2$ . According to Lemma 1 in [?], the condition in (11) should satisfy (8) for this simple case. We draw values of  $h_n$  according to the first paragraph of Sec. IV in [?]. We list these values in rows 2–4 of Table ?. We define

$$g_0(\theta_1, \theta_2) = h_0 + \sum_{n=1}^2 h_n e^{j\theta_n} \quad (19)$$

and list the values of  $g_0(\theta_1, \theta_2)$  for all possible  $\theta_1, \theta_2 \in \{0, \pi\}$ . There are four such values and they are listed in rows 5–8 of Table ?. The set of values for  $\theta_1$  and  $\theta_2$  that maximize  $|g_0|$ , or equivalently, that achieve  $g$  in (9), are  $\theta_1 = \theta_2 = \pi$  as in row 8 of Table ?. Note that this operation results in  $\underline{\mu} = 2.3719$  radians as shown in column 5 of row 8 of Table ?.

At this point, we would like to emphasize that [?] uses a particular convention for the phases of complex numbers. They are defined to be in  $[0, 2\pi)$ , see the text that follows (2) in [?]. We use the same convention in generating Table ?, see its column 5, as well as in generating Table ?. With this convention, we list  $\theta_n + \alpha_n - \underline{\mu}$  and  $(\theta_n + \alpha_n - \underline{\mu}) \bmod 2\pi$  for possibilities of  $\theta_n = 0$  and  $\theta_n = \pi$  and  $n = 1, 2$  in rows 1–8 of Table ?.<sup>2</sup> It can be seen from rows 1–4 of Table ? that the method results in  $\theta_1 = \pi$  as the potential  $\theta_1^*$ , which we know from the discussion in the previous paragraph to be correct. When we carry out the calculation  $(\theta_2 + \alpha_2 - \underline{\mu}) \bmod 2\pi$  in rows 5–8 of Table ?, we find that the method suggests  $\theta_2 = 0$  should be  $\theta_2^*$ . However, we know from the exhaustive search in rows 5–8 of Table ? that  $\theta_2^* = \pi$ . Thus, Lemma 1 is not correct.

It is possible to come up with a correct lemma similar to Lemma 1. We specify this lemma below.

<sup>2</sup>Note that absolute value signs in (11) are not needed since the argument of the minimum operation in (11) is in  $[0, 2\pi)$ .

hand, Proposition 1 is compatible with Lemma 2. To see this, assume  $\mu$  satisfies (12). Then,

$$\underline{\mu} \in \left( \alpha_n + \left( k - \frac{1}{2} \right) \omega, \alpha_n + \left( k + \frac{1}{2} \right) \omega \right). \quad (25)$$

Since  $\omega = \frac{2\pi}{K}$ ,

$$\alpha_n - \underline{\mu} \in \left( (-2k-1)\frac{\pi}{K}, (-2k+1)\frac{\pi}{K} \right) \quad (26)$$

considering the reversal of order due to the subtraction of  $\underline{\mu}$ . Now, let  $\theta_n = k\omega = 2k\frac{\pi}{K}$ . Then

$$\theta_n + \alpha_n - \underline{\mu} \in \left( -\frac{\pi}{K}, \frac{\pi}{K} \right) \quad (27)$$

and thus  $\cos(\theta_n + \alpha_n - \underline{\mu})$  is the largest among all other possibilities for  $\theta_n$  because the slice  $(-\frac{\pi}{K}, \frac{\pi}{K})$  corresponds to the largest values of the cosine function among all slices corresponding to different values of  $\theta_k \in \Phi_K$  for  $k = 1, 2, \dots, K$ .

### III. New Algorithm

We now specify Algorithm 2 to replace Algorithm 1 in [?]. In doing so, not only do we incorporate Lemma 2 instead of Lemma 1 but also we eliminate the many uncertainties present in Algorithm 1 of [?].

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#### Algorithm 2 Update for Algorithm 1

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- 1: Initialization: Compute  $s_{nk} = e^{j(\alpha_n + (k-0.5)\omega)}$  for  $n = 1, 2, \dots, N$  and  $k = 1, 2, \dots, K$ .
- 2: Eliminate duplicates among  $s_{nk}$  and sort to get  $0 \leq \lambda_1 < \lambda_2 < \dots < \lambda_L < 2\pi$ .
- 3: Let, for  $l = 1, 2, \dots, L$ ,  $\mathcal{N}(\lambda_l) = \{n | s_{nk} = \lambda_l\}$ .
- 4: Set  $\underline{\mu} = 0$ . For  $n = 1, 2, \dots, N$ , calculate  $\theta_n = \arg \max_{\theta_n \in \Phi_K} \cos(\theta_n + \alpha_n - \underline{\mu})$ .
- 5: Set  $g_1 = h_0 + \sum_{n=1}^N h_n e^{j\theta_n}$ , **absgmax** = 0.
- 6: for  $l = 2, 3, \dots, L$  do
- 7:   For each  $n \in \mathcal{N}(\lambda_l)$ , let  $(\theta_n + \omega \leftarrow \theta_n) \bmod \Phi_K$ .
- 8:   Let

$$g_l = g_{l-1} + \sum_{n \in \mathcal{N}(\lambda_l)} h_n (e^{j\theta_n} - e^{j(\theta_n - \omega) \bmod \Phi_K})$$

- 9:   if  $|g_l| > \mathbf{absgmax}$  then
  - 10:     Let **absgmax** =  $|g_l|$
  - 11:     Store  $\theta_n$  for  $n = 1, 2, \dots, N$
  - 12:   end if
  - 13: end for
  - 14: Read out  $\theta_n^*$  as the stored  $\theta_n$ ,  $n = 1, 2, \dots, N$ .
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### IV. Results and Remarks

Because its description is based on Lemma 1, which does not provide an equivalency condition for finding  $\theta_1^*, \theta_2^*, \dots, \theta_N^*$ , the performance of Algorithm 1 will in general not achieve the optimum result for SNR Boost [?].

We have implemented Algorithm 1 to the best of our interpretation. We have also implemented Algorithm 2. We

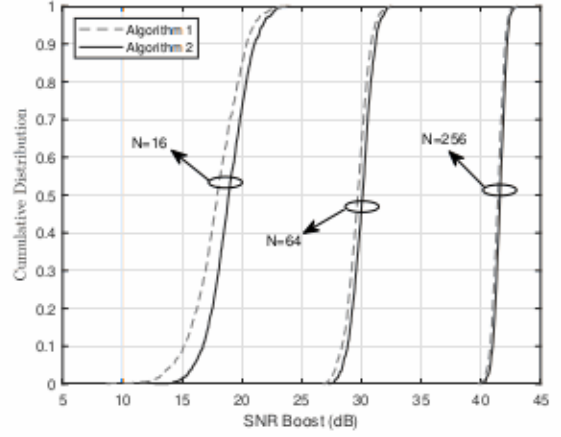


Fig. 1: CDF plots for SNR Boost with Algorithm 1 and Algorithm 2,  $K = 2$ .

present the CDF results for SNR Boost [?] in Fig. ?? for  $K = 2$  and  $N = 16, 64$ , and  $256$ , using the average of 1,000 realizations of the channel. Clearly, Algorithm 1 is not optimal. Algorithm 2 performs better than Algorithm 2 although the gains decrease with  $N$ . Plots for  $K = 4$  show smaller gains as compared to  $K = 2$ , but still, Algorithm 2 always performs better than Algorithm 1 for the same  $K$  and  $N$ .

We note that it is possible to convert the maximization of  $\cos(\theta_n + \alpha_n - \underline{\mu})$  to the minimization of a simple expression. For example, minimization of  $f_1(x) = \pi - |(x \bmod 2\pi) - \pi|$  is the same as maximization of  $\cos(x)$  within the context of Lemma 2. However, this is different than minimization of  $|x \bmod 2\pi|$  proposed in Lemma 1 of [?]. The reason can be seen by plotting these functions against  $x$ . While  $f_1(x)$  and  $\cos(x)$ , in addition to being periodic with period  $2\pi$ , have even symmetry around odd multiples of  $\pi$ ,  $|x \bmod 2\pi|$  (or equivalently,  $(x \bmod 2\pi)$ ) does not have this symmetry.