

CSCE 420 - Fall 2025

Homework 3 (HW3)

due: Tues, Oct 28, 11:59 pm

Turn-in answers as a PDF (HW3.pdf) and commit/push it to your class github repo.

All homeworks must be typed, *not* hand-written and scanned as photos.

1a. Prove that $(A \wedge B \rightarrow C \wedge D) \vdash (A \wedge B \rightarrow C)$ ("conjunctive rule splitting") is a **sound rule-of-inference** using a **truth table**.

A	B	C	D	$A \wedge B$	$C \wedge D$	$A \wedge B \rightarrow C \wedge D$	$A \wedge B \rightarrow C$
T	T	T	T	T	T	T	T
T	T	T	F	T	F	F	T
T	T	F	T	T	F	F	F
T	T	F	F	T	F	F	F
T	F	T	T	F	T	T	T
T	F	T	F	F	F	T	T
T	F	F	T	F	F	T	T
T	F	F	F	F	F	T	T
F	T	T	T	F	T	T	T
F	T	T	F	F	F	T	T
F	T	F	T	F	F	T	T
F	T	F	F	F	F	T	T
F	F	T	T	F	T	T	T
F	F	T	F	F	F	T	T
F	F	F	T	F	F	T	T
F	F	F	F	F	F	T	T

Our table proves by modus ponens rule of inference that all models that satisfy the premises $A \wedge B \rightarrow C \wedge D$, also satisfy the derived sentences $A \wedge B \rightarrow C$. I've highlighted all the matching truths in yellow.

1b. Also prove $(A \wedge B \rightarrow C \wedge D) \models (A \wedge B \rightarrow C)$ using **Natural Deduction**.

(Hint: it might help to use a ROI for "Implication Introduction". If you have a Horn clause, with 1 positive literal and $n-1$ negative literals, like $(\neg X \vee \neg Z \vee \neg Y)$, you can transform it into a conjunctive rule by collecting the negative literals as positive antecedents, e.g. $X \wedge Y \rightarrow Z$. This is a truth-preserving operation (hence sound), which you could prove to yourself using a truth table.)

1. Start with $A \wedge B \rightarrow C \wedge D$
2. Modify into $A \wedge B \rightarrow C$ using And Elimination (if $C \wedge D$ is true, then C must be true, so D can be removed)

1c. Also prove $(A \wedge B \rightarrow C \wedge D) \models (A \wedge B \rightarrow C)$ using **Resolution**.

Using the hint above, let's convert this sentence to CNF and resolve so we get a contradiction

Premises:

1. $A \wedge B \rightarrow C \wedge D$ turns into $\neg A \vee (C \wedge D) \vee \neg B$ by Horn Clause implication elimination (used the hint)
2. $\neg A \vee (C \wedge D) \vee \neg B$ turns into $\neg A \vee C \vee \neg B$ and $\neg A \vee D \vee \neg B$ by distribution

Conclusion negation:

1. $A \wedge B \rightarrow C$ turns into $\neg A \vee C \vee \neg B$ via Horn Clause implication elimination (used the hint)
2. $\neg A \vee C \vee \neg B$ turns into $A \wedge \neg C \wedge B$ via negation

Total Clauses:

1. $(\neg A \vee C \vee \neg B)$
2. $(\neg A \vee D \vee \neg B)$
3. A
4. $\neg C$
5. B

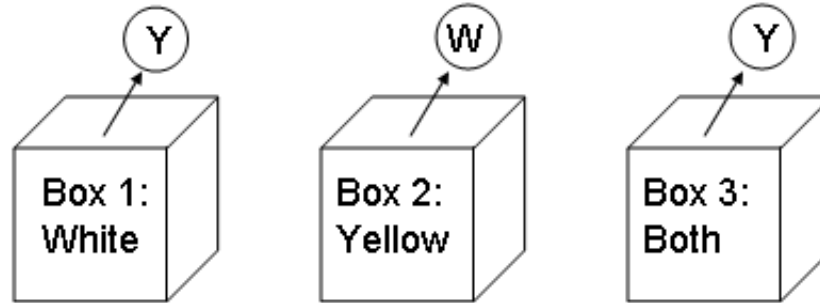
Steps:

6. Resolve 1 and 4 and turn it into $\neg A \vee \neg B$
7. Resolve 5 and 6 and turn it into $\neg A$
8. Resolve 3 and 7 and see that A and $\neg A = \emptyset$ empty set which is a **CONTRADICTION**
9. Resolve 3 and 6 and turn it into $\neg B$
10. Resolve 5 and 9 and see that B and $\neg B = \emptyset$ empty set which is a **CONTRADICTION**
11. Resolve 1 and 3 to get $C \vee \neg B$
12. Resolve 11 and 5 to get C
13. Resolve 12 and 4 to get C and $\neg C = \emptyset$ empty set which is a **CONTRADICTION**

We've proved entailment through resolution for all 3 conclusion variables :D

2. Sammy's Sport Shop

You are the proprietor of *Sammy's Sport Shop*. You have just received a shipment of three boxes filled with tennis balls. One box contains only yellow tennis balls, one box contains only white tennis balls, and one contains both yellow and white tennis balls. You would like to stock the tennis balls in appropriate places on your shelves. Unfortunately, the boxes have been labeled incorrectly; the manufacturer tells you that you have exactly one box of each, but that **each box is definitely labeled wrong**. You draw one ball from each box and observe its color. Given the initial (incorrect) labeling of the boxes above, and the three observations, use Propositional Logic to infer the correct contents of the middle box.



Use propositional symbols in the following form: O1Y means a yellow ball was drawn (observed) from box 1, L1W means box 1 was initially labeled white, C1W means box 1 contains (only) white balls, and C1B means box 1 actually contains both types of tennis balls. Note, there is no 'O1B', etc, because you can't directly "observe both". When you draw a tennis ball, it will either be white or yellow.

The initial facts describing this particular situation are: {O1Y, L1W, O2W, L2Y, O3Y, L3B}

2a. Using these propositional symbols, write a propositional knowledge base (sammy.kb) that captures the knowledge in this domain (i.e. implications of what different observations or labels mean, as well as constraints inherent in this problem, such as that all boxes have different contents). *Do it in a complete and general way*, writing down *all* the rules and constraints, not just the ones needed to make the specific inference about the middle box. *Do not include derived knowledge* that depends on the particular labeling of this instance shown above; stick to what is stated in the problem description above. Your KB should be general enough to reason about any alternative scenario, not just the one given above (e.g. with different observations and labels and box contents).

Knowledge Base KB

Direct Observation:

1. O1Y
2. L1W
3. O2W
4. L2Y
5. O3Y
6. L3B

What we know from the labels being wrong:

7. $(L1W \rightarrow \neg C1W)$ since the label on box 1 is white, the color cannot be white because the label is wrong
8. $(L2Y \rightarrow \neg C2Y)$ same reasoning as above
9. $(L3B \rightarrow \neg C3B)$ same reasoning as above
10. $(C1Y \vee C1B)$ box 1 has to be yellow or both
11. $(C2W \vee C2B)$ same reasoning logic as above
12. $(C3W \vee C3Y)$ same reasoning logic as above

Constraints that the boxes have on each other:

13. $(C1Y \rightarrow \neg C3Y)$ If box 1 is yellow then no other boxes can be yellow (in this case it only affects box 3 because box 2 is labeled wrong and isn't yellow balls)
14. $(C1B \rightarrow \neg C2B)$ If box 1 is both then no other boxes can be both

15. $(C2W \rightarrow \neg C3W)$ If box 2 is white then box 3 cannot be white

General Constraints (one color exists at a time):

16. $(C1Y \vee C2Y \vee C3Y)$

17. $(C1W \vee C2W \vee C3W)$

18. $(C1B \vee C2B \vee C3B)$

19. $\neg (C1Y \wedge C3Y)$ Only one box c1 or c3 can be yellow at one time

20. $\neg (C1B \wedge C2B)$ same reasoning logic as above

21. $\neg (C2W \wedge C3W)$ same reasoning logic as above

What we know from the observed balls:

22. $(O1Y \rightarrow (C1Y \vee C1B))$ since the observed ball in box 1 is yellow then the resulting color must either be yellow or both

23. $(O2W \rightarrow (C2W \vee C1B))$

24. $(O3Y \rightarrow (C1W \vee C1Y))$

25. $O1Y \rightarrow \neg C1W$

26. $O2W \rightarrow \neg C2Y$

27. $O3Y \rightarrow \neg C3W$

2b. Prove that box 2 must contain white balls (**C2W**) using **Natural Deduction**.

Premises: (I tried my best.)

1. $O2W \rightarrow \neg C2Y$ (rule 26)

2. $O2W$ (rule 3)

3. Through modus ponens, we know that $\neg C2Y$ is true

4. $O3Y \rightarrow \neg C3W$ (rule 27)

5. $O3Y$ (rule 5)

6. Through modus ponens, we know that $\neg C3W$ is true

7. $L3B \rightarrow \neg C3B$ (rule 9)

8. $L3B$ (rule 6)

9. Through modus ponens, we know that $\neg C3B$ is true

10. $C3W \vee C3Y$ can be combined with premise 6 $\neg C3W$ through disjunctive syllogism to conclude $C3Y$

11. Use rule 19 $\neg (C1Y \wedge C3Y)$, since $C3Y$ is true, we infer that $C1Y$ must be false in order for the clause to be true

12. We can infer $\neg C1Y$ through premise 11

13. $(L1W \rightarrow \neg C1W)$ rule 7

14. $L1W$ rule 2

15. Through modus ponens, we know that $\neg C1W$ is true

16. We can infer $C1B$ is true through premise 12 premise 15

17. Using rule 20 $\neg (C1B \wedge C2B)$, this clause can only be true if $C2B$ is false

18. We can infer $\neg C2B$ through premise 17

19. $(C2W \vee C2B)$ rule 11

20. Using disjunctive syllogism of premise 18 and 19, we can deduce that **C2W is true.**

2c. Convert your KB to CNF.

Direct Observation:

1. $O1Y$

2. $L1W$

3. $O2W$

4. $L2Y$

5. $O3Y$

6. $L3B$

What we know from the labels being wrong:

7. $(L1W \rightarrow \neg C1W)$ turns into $(\neg L1W \vee \neg C1W)$ via implication elimination

8. $(L2Y \rightarrow \neg C2Y)$ turns into $(\neg L2Y \vee \neg C2Y)$ via implication elimination

9. $(L3B \rightarrow \neg C3B)$ turns into $(\neg L3B \vee \neg C3B)$ via implication elimination

10. $(C1Y \vee C1B)$
11. $(C2W \vee C2B)$
12. $(C3W \vee C3Y)$

Constraints that the boxes have on each other:

13. $(C1Y \rightarrow \neg C3Y)$ turns into $(\neg C1Y \vee \neg C3Y)$ via implication elimination
14. $(C1B \rightarrow \neg C2B)$ turns into $(\neg C1B \vee \neg C2B)$ via implication elimination
15. $(C2W \rightarrow \neg C3W)$ turns into $(\neg C2W \vee \neg C3W)$ via implication elimination

General Constraints (one color exists at a time):

16. $(C1Y \vee C2Y \vee C3Y)$
17. $(C1W \vee C2W \vee C3W)$
18. $(C1B \vee C2B \vee C3B)$
19. $\neg (C1Y \wedge C3Y)$ turns into $(\neg C1Y \vee \neg C3Y)$ via DeMorgan's Laws
20. $\neg (C1B \wedge C2B)$ turns into $(\neg C1B \vee \neg C2B)$ via DeMorgan's Laws
21. $\neg (C2W \wedge C3W)$ turns into $(\neg C2W \vee \neg C3W)$ via DeMorgan's Laws

What we know from the observed balls:

22. $(O1Y \rightarrow (C1Y \vee C1B))$ turns into $\neg O1Y \vee C1Y \vee C1B$ via implication elimination
23. $(O2W \rightarrow (C2W \vee C1B))$ turns into $\neg O2W \vee C2W \vee C1B$ via implication elimination
24. $(O3Y \rightarrow (C1W \vee C1Y))$ turns into $\neg O3Y \vee C1W \vee C1Y$ via implication elimination
25. $O1Y \rightarrow \neg C1W$ turns into $(\neg O1Y \vee \neg C1W)$ via implication elimination
26. $O2W \rightarrow \neg C2Y$ turns into $(\neg O2W \vee \neg C2Y)$ via implication elimination
27. $O3Y \rightarrow \neg C3W$ turns into $(\neg O3Y \vee \neg C3W)$ via implication elimination

2d. Prove $C2W$ using **Resolution**.

We can use all of the premises from above AND an additional clause/rule:

28. $C2W$ conclusion needs to be negated into $\neg C2W$ during resolution refutation

Steps:

29. Resolve rule 3 and rule 11 to get $C2W \vee C1B$
30. Resolve rule 11 and rule 16 to get $C2B$
31. Resolve rule 14 and rule 30 to get $\neg C1B$
32. Resolve rule 29 and rule 31 to get $C2W$
33. Resolve rule 32 $C2W$ and rule 28 $\neg C2W$ to get \emptyset empty set which is a Contradiction! Which proves $C2W$ is true.

3. Do **Forward Chaining** for the *CanGetToWork* KB below.

You don't need to follow the formal FC algorithm (with agenda/queue and counts array). Just indicate which rules are triggered (in any order), and keep going until all consequences are generated.

Show the final list of all inferred propositions at the end. *Is CanGetToWork among them?*

```
KB = { a. CanBikeToWork → CanGetToWork
      b. CanDriveToWork → CanGetToWork
      c. CanWalkToWork → CanGetToWork
      d. HaveBike ∧ WorkCloseToHome ∧ Sunny → CanBikeToWork
      e. HaveMountainBike → HaveBike
      f. HaveTenSpeed → HaveBike
      g. OwnCar → CanDriveToWork
      h. OwnCar → MustGetAnnualInspection
      i. OwnCar → MustHaveValidLicense
      j. CanRentCar → CanDriveToWork
      k. HaveMoney ∧ CarRentalOpen → CanRentCar
      l. HertzOpen → CarRentalOpen
      m. AvisOpen → CarRentalOpen
      n. EnterpriseOpen → CarRentalOpen
      o. CarRentalOpen → IsNotAHoliday
      p. HaveMoney ∧ TaxiAvailable → CanDriveToWork
      q. Sunny ∧ WorkCloseToHome → CanWalkToWork
      r. HaveUmbrella ∧ WorkCloseToHome → CanWalkToWork
      s. Sunny → StreetsDry }
```

Facts: { Rainy, HaveMountainBike, EnjoyPlayingSoccer, WorkForUniversity, WorkCloseToHome, HaveMoney, HertzClosed, AvisOpen, McDonaldsOpen }

We'll start with going through the facts to generate all consequences:

1. HaveMountainBike fires rule e HaveBike
2. AvisOpen fact fires rule m -> CarRentalOpen
3. CarRentalOpen fires rule o-> IsNotAHoliday
4. HaveMoney and CarRentalOpen fires rule k->CanRentCar
5. CanRentCar fires rule j ->CanDriveToWork
6. CanDriveToWork fires rule b-> CanGetToWork **(SOLUTION end)**

4. Do **Backward Chaining** for the *CanGetToWork* KB.

In this case, you should follow the BC algorithm closely (the pseudocode for the propositional version of Back-chaining is given in the lecture slides).

Important: when you pop a subgoal (proposition) from the goal stack, you should systematically go through all rules that can be used to prove it **IN THE ORDER THEY APPEAR IN THE KB**.

In some cases, this will lead to *back-tracking*, which you should show.

Also, the sequence of results depends on order in which antecedents are pushed onto the stack. If you have a rule like $A \wedge B \rightarrow C$, and you pop C off the stack, push the antecedents in reverse order, so B goes in first, then A; in the next iteration, A would be the next subgoal popped off the stack.

Lets track backward chaining from CanGetToWork:

Steps:

1. Push **CanGetToWork** onto stack
2. Pop **CanGetToWork** and push CanBikeToWork (first rule in knowledge base)
3. Pop CanBikeToWork and push HaveBike, WorkCloseToHome, Sunny (antecedents are checked in reverse order)
4. Pop Sunny-> it fails->backtrack to CanGetToWork since none of the other conditions on stack need to be checked
5. Push CanDriveToWork
6. Pop CanDriveToWork and push OwnCar
7. Pop OwnCar->it fails-> backtrack to CanDriveToWork
8. Pop CanDriveToWork and push CanRentCar
9. Pop CanRentCar and push HaveMoney and CarRentalOpen (check antecedents in reverse order)
10. Pop CarRentalOpen and push HertzOpen
11. Pop HertzOpen-> it fails-> backtrack to CarRentalOpen
12. Pop CarRentalOpen and push AvisOpen
13. Pop AvisOpen-> it succeeds/is a fact-> move to HaveMoney
14. Pop HaveMoney-> it succeeds/is a fact
15. We then have proved all the variables for CanRentCar
16. We then have proved all the variables for CanDriveToWork
17. We then have proved all the variables for CanGetToWork