## UNIVERSITY OF TORONTO

Faculty of Arts and Science

#### Midterm 2 Solutions CSC148H1F

Duration: 50 min. Instructors: Diane Horton, David Liu. Examination Aids: Provided aid sheet

## Name:

# Student Number:

## Please read the following guidelines carefully.

- Please print your name and student number on the front of the exam.
- This examination has 3 questions. There are a total of 12 pages, DOUBLE-SIDED.
- The last page is an aid sheet that may be detached.
- You may always write helper functions/methods unless explicitly asked not to.
- Docstrings are *not* required unless explicitly asked for.

Take a deep breath.

This is your chance to show us How much you've learned.

We **WANT** to give you the credit

That you've earned.

A number does not define you.

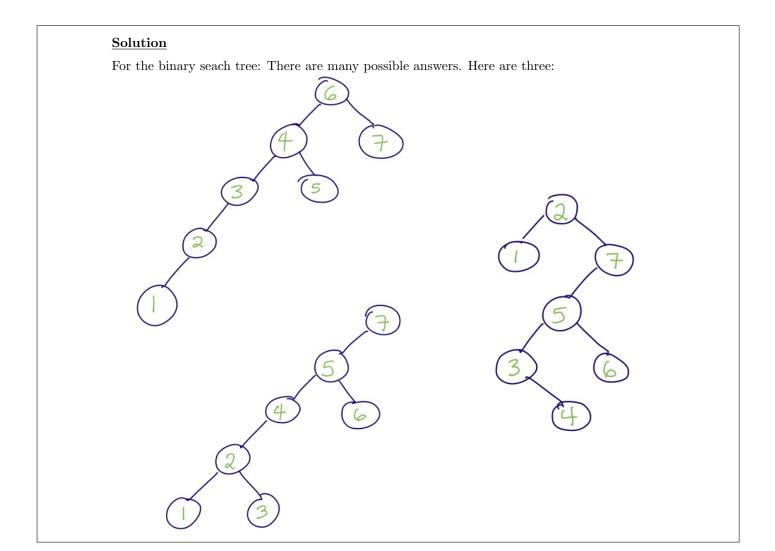
Question	Grade	Out of
Q1		10
Q2		12
Q3		8
Total		30

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priginal question.	won oj

- 1. [10 marks] Point-form responses are acceptable here. You do not need to write a lot for full marks.
  - (a) [3 marks] Recall that we defined the height of a tree in terms of nodes, so that the height of a tree with just a root is 1. Draw a binary search tree of height 5 containing only the numbers 1 through 7 inclusive.

What is the greatest number of nodes we can have in a binary search tree of height 3?

What is the fewest number of nodes we can have in a binary search tree of height 3?



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The easiest way to answer this question is probably to draw the shape and then add the values. Notice that placement of the values is constrained, once you have the shape.

The greatest number of nodes we can have in a binary search tree of height 3 is 7.

The fewest number of nodes we can have in a binary search tree of height 3 is 3.

(b) [1 mark] Suppose we have a sorted LinkedList. Why would binary search be a bad technique for finding a given item?

#### Solution

Every time we want to make a comparison with an item at a particular index, we'd need to iterate through the linked list's nodes to get to that index. We'd also need to iterate through the entire linked list once just to compute the length of the linked list. This wouldn't save any time compared to just doing a linear search on the linked list.

(c) [3 marks] The aid sheet has docstrings for classes LinkedList and Node, which are relevant to this question.

Here is a linked list method you have seen before:

```
class LinkedList:
def remove(self, item: object) -> None:
"""Remove the FIRST occurrence of <item> in this list."""
prev = None
curr = self._first
while curr is not None and curr.item != item:
prev, curr = curr, curr.next
# Rest of method omitted.
```

Whenever we say blah.something, we must know that blah is not None, otherwise the code will raise an error. For each of the following pieces of code, explain how we know that the value to the left of the dot is not None at the moment when it is evaluated.

Line 5: curr = self.\_first

#### Solution

self is a LinkedList object, and so it is not None.

Line 6: curr.item != item

#### Solution

The first part of the condition checks curr is not None, and the expression only evaluates if the first condition is true.

Line 7: prev, curr = curr, curr.next

## **Solution**

This line is inside the while loop, and so the condition must be true; the condition includes the check that curr is not None.

(d) [3 marks] Here is a (possibly incorrect) implementation of the LinkedList remove method.

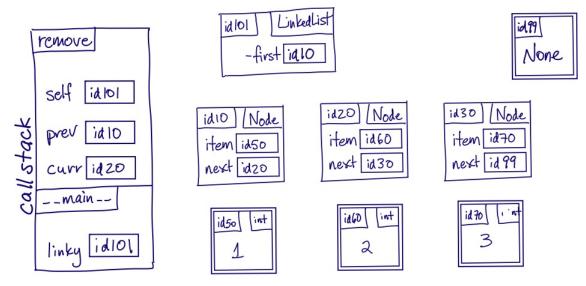
```
def remove(self, item: object) -> None:
    """Remove the FIRST occurrence of <item> in this list."""
    prev = None
    curr = self._first
    while curr is not None and curr.item != item:
        prev, curr = curr, curr.next

# The diagram below shows the state of memory right before this line:
    curr = curr.next

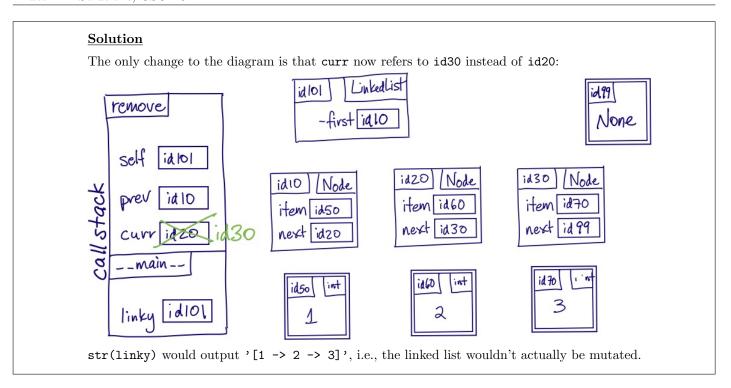
Suppose we run the following code:

>>> linky = LinkedList([1, 2, 3])  # str(linky) == '[1 -> 2 -> 3]'
>>> linky.remove(2)
>>> # What would str(linky) return now?
```

Below, we have drawn a memory model diagram showing the state of the program's memory during the method call linky.remove(2), immediately before the line curr = curr.next.



- (i) Modify the diagram to show the state of memory immediately after the line curr = curr.next is executed.
- (ii) In the space below, write what str(linky) would return if we called it *after* the call to linky.remove(2) is over.



2. [12 marks] Consider the following nested list function. Suggestion: Drawing arcs between matching opening and closing list brackets may help you to see what it is counting.

(a) [3 marks] Suppose we have the following nested list 1st (we've added some extra whitespace for readability):

```
lst = [
    [1, [3, 4]],
    2,
    [],
    [5, 6, [[7]]]
]
```

Note that 1st has length 4, and so we say that it has four sub-nested lists. Complete the table below.

## **Solution**

Sub-nested list of 1st	Correct return value of num_lists on the sub-nested list
[1, [3, 4]]	2
2	0
[]	1
[5, 6, [[7]]]	3

(b) [1 mark] What should num\_lists(1st) return if called with the value for 1st shown in part (a)?

#### Solution

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(c) [2 marks] Explain, in English, how to compute num\_lists on a nested list from the recursive calls on its sub-nested lists.

#### Solution

Add up all of the return values from the recursive calls on each sub-nested list, and then add one. *Note*: of course, this only applies when the input is a list. If it's an integer, then 0 should be returned; no recursive calls are made.

(d) [6 marks] In the space below, implement the num\_lists function using recursion. You may not define any helper functions here.

```
def num_lists(obj: Union[int, List]) -> int:
    """Return the number of list objects in the given nested list.
    If obj is a list itself, include it in the count.

>>> num_lists(4)
0
>>> num_lists([1, 2, 3])
1
>>> num_lists([1, [2], [[3, 4]]])
4 # The four lists are: [1, [2], [[3, 4]]], [2], [[3, 4]], and [3, 4].
"""
```

## Solution

```
if isinstance(obj, int):
    return 0
else:
    num = 1 # For the outermost list
    for lst_i in obj:
        num += num_lists(lst_i)
    return num
```

3. [8 marks] The aid sheet has the docstring for class Tree, which is relevant to this question.

Recall that the *depth* of an item in a tree is equal to the distance between it and the root inclusive, counting items. So the root of a tree has **depth** 1.

Consider the Tree method truncate, with the docstring below:

## class Tree:

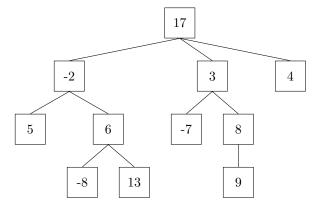
def truncate(self, d: int) -> None:

"""Remove all values in the tree that are at depth <d> or greater.

Precondition:  $d \ge 1$ .

#### Notes:

- 1. Calling truncate with d = 1 always results in an empty tree.
- 2. Calling truncate when d is greater than the tree's height (number of levels) does not change the tree at all.
- (a) [2 marks] Suppose we have a variable t that is a Tree instance representing the following tree:



Here are two calls to t.truncate with our initial tree t, shown above, but with different values of d. Below each call, draw what t would look like after the call. If the tree would be empty, write "empty".

t.truncate(3)
t.truncate(1)



(b) [6 marks] In the space below, implement the truncate method using recursion. You may not use any Tree methods other than is\_empty, but you may access all Tree attributes. You may not define any helper methods here. We have given you some code to help you get started.

## Solution

```
def truncate(self, d: int) -> None:
    if self.is_empty():
        pass # 'return' is also acceptable here
    else:
        if d == 1:
            # Turn the tree into an empty tree.
            self._root = None
            self._subtrees = []
        else:
            # In this branch, we know that self is non-empty and d > 1.
            for subtree in self._subtrees:
                subtree.truncate(d - 1)
            # ALTERNATE VERSION, which removes empty subtrees from self._subtrees.
            # This is *not* required by the representation invariants given on
            # the aid sheet, but useful in practice.
            new_subtrees = []
            for subtree in self._subtrees:
                subtree.truncate(d - 1)
                if not subtree.is_empty():
                    new_subtrees.append(subtree)
            self._subtrees = new_subtrees
```

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Use this page for rough work. If you want work on this page to be marked, please indicate this clearly at the location of the original question.