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complete induction

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use Introduction to the Theory of Computation, Section 1.3



Outline

Principle of complete induction

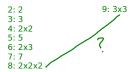
Examples of complete induction

Complete Induction

another flavour needed

Every natural number greater than 1 has a prime factorization write them as products of 1 or more primes.

Try some examples



How does the factorization of 8 help with the factorization of 9?

notational convenience...

I will use (though you don't have to) the following:

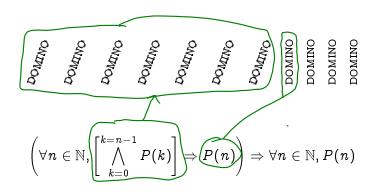
$$\bigwedge_{k=0}^{k=n-1} P(k)$$

... as equivalent to

$$\forall k \in \mathbb{N}, k < n \Rightarrow P(k)$$



More dominoes



If all the previous cases always imply the current case then all cases are true





complete induction outline

inductive step: state inductive hypothesis H(n)

derive conclusion C(n): show that C(n) follows from H(n), indicating where you use H(n) and why that is valid

verify base case(s): verify that the claim is true for any cases not covered in the inductive step

sometimes I embed the base(s) in the inductive step. Also, some people go from assuming from start to n, then show this implies P(n+1)

Wait! isn't that the same outline as simple induction?

Yes, we just modify the inductive hypothesis, H(n) so that it assumes the main claim for every natural number from the starting point up to n-1, and the conclusion, C(n), is now the main claim for n.

watch the base cases, part 1

$$f(n) = egin{cases} 1 & n \leq 1 \ \left[f(\lfloor \sqrt{n}
floor)
ight]^2 + 2 f(\lfloor \sqrt{n}
floor) & n > 1 \end{cases}$$

Check a few cases, and make a conjecture

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\begin{array}{lll} f(0) = 1 \\ f(1) = 1 \\ f(2) = 3 \\ f(3) = 3 \\ f(4) = 15 \\ f(5) = 15 \\ f(6) = 15 \\ f(7) = 15 \\ f(8) = 15 \end{array}
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f(9) = 15

For all natural numbers n > 1, f(n) is a multiple of 3x

use the complete induction outline

P(n): f(n) is a multiple of 3.

Let $n \in P(n-1)$. I wish to show that P(n) follows, that is that P(n) is also multiple of 3.

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We know that f(n) = f(floor(sqrt(n)))^2 + 2f(floor(sqrt(n))) # since n > 1
= f(floor(sqrt(n))) \times [f(floor(sqrt(n))) + 2]
= 3k \times [f(floor(sqrt(n))) + 2], \text{ for some } k \text{ in } N \text{ # by hypothesis since sqrt(n) } >= 2
# since n >= 4
# also n > 1 implies n > 2 > n
# which implies n > 3 sqrt(n > 3)
```

so f(n) is also a multiple of 3.

Base case(s). f(2) = 3 = f(3), as shown previously, which verifies our base cases 2, and 3.

zero pair free binary strings, zpfbs...

Denote by zpfbs(n) the number of binary strings of length n that contain no pairs of adjacent zeros. What is zpfbs(n) for the first few natural numbers n?

$$f(0) = 1$$

 $f(1) = 2$
 $f(2) = 3$
 $f(3) = 5$
 $f(4) = 8$
 $f(5) = 13$
 $f(6) = 21$

$$f(n) = 1 \text{ if } n \text{ is } 0$$

= 2 if n is 1
= $f(n-2) + f(n-1) \text{ if}$

n > 1

b1 b2 1 0

break up the counting into 2 parts...

f(4) counts binary strings of length 4 that have no adjacent zeros.

binary strings of length 4 that end in 1 and have no adjacent zeros are simply a bs of length 3 with no adjacent zeros with a 1 appended.

binary strings of length 4 that end in 0 and have no adjacent zeros are simply the binary strings of length 2 with 10 appended.



what is zpfbs(n)?

use the complete induction outline

P(n): zpfbs(n) = 1 if n is 0, 2 if n is 1, andis zpfbs(n-2) + zpfbs(n-1) if n > 1

Let $n \in P(0)$ and P(1) and ... P(n-1). I want to show that P(n) follows, that is zpfbs(n) is what the formula on the previous page gives.

- case n >= 2. Then note that zpfbs(n) counts those binary strings of length n that end in a 1 and those that end in 0. Ending in 1 doesn't introduce, or take away any pairs of zeros, so there are zpfbs(n-1) of those, by the hypthesis # since n-1 < n and n-1 >= 0. Ending in 0 is a problem, so we will only count the binary strings of length n that end in 10, and there are zpfbs(n-2) of those, by the hypothesis # since n-2 < n and n-2 >= 0.

 Thus, altogether, we have zpfbs(n-2) + zpfbs(n-1) strings with no adjacent 0s of length n, as the formula claims (although it's called f rather than zpfbs).
- case n < 2: If n were 0, then the only possible binary string is the empty string, and it has no pairs of zeros, so zpfbs(0) = 1, as claimed. If n were 1, then the only possible binary strings are 0 and 1, neither of which have adjacent 0s, so zpfbs(1) = 2, as claimed. These verify the base case.

Every natural number greater than 1 has a prime factorization

use the complete induction outline

P(n): n has a representation as product of 1 or more primes.

prime factorization: representation as product of 1 or more primes

Let n be an arbitrary, fixed, natural number that is no smaller than 2. Assume $P(2) \dots P(n-1)$. We will show P(n), that n has a representation as a product of 1 or more primes.

case n is prime: we're done, n's prime factorization is itself.

case n is composite. Then by definition being composite n is the product of some f1 and f2 \in n such that 1 < f1, f2 < n. By P(f1) and P(f2), f1 and f2 have representations as products of primes, so n's representation is simply the product of these.

* Some students worry that you must know that there exists at least one prime (alternatively that 2 is prime) independently of this proof before it is valid. Not so. The worry is that by applying the induction step enough times you will eventually encounter a number that is not composite. Yes, so a *consequence* of this proof, not a *prerequisite* of it, is that for every integer n >= 2 there is at least one prime less than or equal to n.



After a certain natural number n, every postage can be made up by combining 3- and 5- cent stamps what is the "certain natural number"?

After a certain natural number n, every postage can be made up by combining 3- and 5- cent stamps use the complete induction outline

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notes...