

# CSC236 Fall 2018

## Assignment #1: induction

### due September 28th, 3 p.m.

The aim of this assignment is to give you some practice with various flavours of induction. For each question below you will present a proof by induction. For full marks you need to make it clear to the reader that the base case(s) is/are verified, that the inductive step follows for each element of the domain (typically the natural numbers), where the inductive hypothesis is used, and that it is used in a valid case.

Your assignment must be **typed** to produce a PDF document **a1.pdf** (hand-written submissions are not acceptable). You may work on the assignment in groups of 1 or 2, and submit a single assignment for the entire group on [MarkUs](#)

1. Recall **bipartite graphs**. Consider the following definitions:

**bipartite graph:** Undirected graph  $G = (V, E)$  is **bipartite** if and only if there exist  $V_1, V_2$  such that  $V = V_1 \cup V_2$ ,  $V_1 \cap V_2 = \emptyset$ , and every edge in  $E$  has one endpoint in  $V_1$  and the other in  $V_2$ .

**P(n):** Every bipartite graph on  $n$  vertices has no more than  $n^2/4$  edges.

- (a) Assume  $P(234)$ . Can you use this<sup>1</sup> to prove that  $P(235)$  follows? Explain why, or why not.
- (b) Assume  $P(235)$ . Can you use this<sup>2</sup> to prove that  $P(236)$  follows? Explain why or why not.
- (c) Use what you've learned from the previous two answers to construct a proof by simple induction that:  $\forall n \in \mathbb{N}, P(n)$ . **Note:** There are proofs of this claim that are not by simple induction, **but** those proofs will receive no marks. **Hint:** You probably need to strengthen the claim in order to devise a successful inductive hypothesis. If this seems mysterious, revisit the previous two answers...

2. Define function  $f$  as follows:

$$f(n) = \begin{cases} 3 & \text{if } n = 0 \\ [f(\lfloor \log_3 n \rfloor)]^2 + f(\lfloor \log_3 n \rfloor) & \text{if } n > 0 \end{cases}$$

Define predicate  $P(n)$  : “ $f(n)$  is a multiple of 4.”

- (a) Assume  $P(3)$ . Can you use this<sup>3</sup> to prove  $P(29)$ ? Explain why or why not.
- (b) Assume  $P(4)$ . Can you use this<sup>4</sup> to prove  $P(29)$ ? Explain why or why not.
- (c) Use complete induction to prove  $\forall n \in \mathbb{N}, n > 0 \Rightarrow P(n)$ .

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<sup>1</sup>If you say yes,  $P(234)$  must be a necessary part of your proof.

<sup>2</sup>If you say yes,  $P(235)$  must be a necessary part of your proof.

<sup>3</sup>If you say yes,  $P(3)$  must be a necessary part of your proof.

<sup>4</sup>If you say yes,  $P(4)$  must be a necessary part of your proof.

3. Use the Principle of Well-Ordering to derive a contradiction that proves there are no positive integers  $x, y, z$  such that:

$$5x^3 + 50y^3 = 3z^3$$

You may assume, without proof, that if a prime number  $p$  divides a perfect cube  $n^3$ , then  $p$  also divides  $n$ .

4. Define  $\mathcal{T}$  as the smallest set of strings that satisfies:

- $"*" \in \mathcal{T}$
- if  $t_1, t_2 \in \mathcal{T}$  then their parenthesized concatenation  $(t_1 t_2) \in \mathcal{T}$ .

Some examples:  $"*" , "(**)" , "(*(**))"$  are all in  $\mathcal{T}$ .

Now read over these four Python functions:

```
def left_count(s: str) -> int:
    """
    Return the number of "(" in s
    """
    return s.count("(")

def double_count(s: str) -> int:
    """
    Return the number of "(" plus number of ")", including possible
    overlaps.
    """
    return (len([s[i:] for i in range(len(s)) if s[i:].startswith("(")])
            + len([s[:i] for i in range(len(s) + 1) if s[:i].endswith(")")])))

def left_surplus(s: str, i: int) -> int:
    """
    Return the number of "(" minus the number of ")"
    in s[:i]
    """
    return s.count("(", 0, i) - s.count(")", 0, i)

def max_left_surplus(s: str) -> int:
    """
    Return the maximum left surplus for all prefixes of s.
    """
    return max([left_surplus(s, i) for i in range(len(s))] + [0])
```

- (a) Use structural induction on  $\mathcal{T}$  to prove:

$$\forall t \in \mathcal{T}, \text{left\_count}(t) \leq 2^{\text{max\_left\_surplus}(t)} - 1$$

- (b) Use structural induction on  $\mathcal{T}$  to prove: **[edit:] error fixed September 9**

$$\forall t \in \mathcal{T}, \text{double\_count}(t) = \begin{cases} 0 & \text{if } t = "*" \\ \text{left\_count}(t) - 1 & \text{otherwise} \end{cases}$$