### Objective:

### At the end of this lecture, students will understand:

- Time Complexity
- Space Complexity
- Importance of them in Competitive Programming

### Time and Space Complexity:

Sometimes, there are more than one way to solve a problem. We need to learn how to compare the performance of different algorithms and choose the best one to solve a particular problem. While analysing an algorithm, we mostly consider time complexity and space complexity.

### Time and Space Complexity: (Contd.)

- Time complexity of an algorithm quantifies the amount of time taken by an algorithm to run as a function of the length of the input.
- Similarly, Space complexity of an algorithm quantifies the amount of space or memory taken by an algorithm to run as a function of the length of the input.
- Time and space complexity depends on lots of things like hardware, operating system, processors, etc.

### Order of growth:

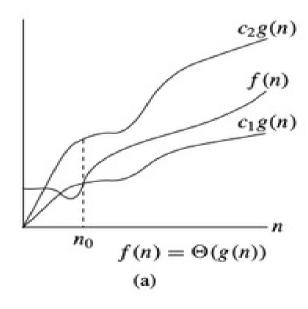
- Order of growth of an algorithm is a way of saying /predicting how execution time of a program and the space/memory occupied by it changes with the input size.
- The most famous way is the Big-Oh notation. It gives the worst case possibility for an algorithm.

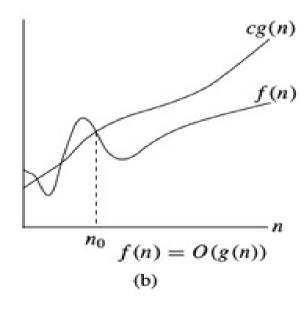
• E.g. a program takes O(n) time means it takes at most n operations to compute the answer. Now if your input for n is 10<sup>6</sup>, it takes 10<sup>6</sup> operations, while a O(n<sup>2</sup>) algorithm takes 10<sup>12</sup> operations for the same input.

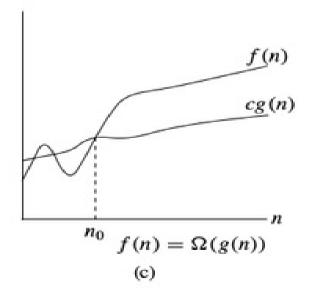
#### O-notation:

To denote asymptotic upper bound, we use O-notation. For a given function g(n), we denote by O(g(n)) (pronounced "big-oh of g of n") the set of functions:

 $O(g(n)) = \{ f(n) : \text{there exist positive constants } c \text{ and } n_0 \text{ such that } 0 \leq f(n) \leq c * g(n) \text{ for all } n \geq n_0 \}$ 



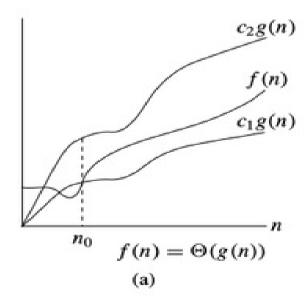


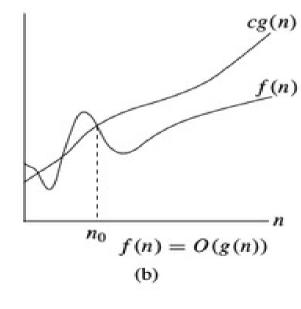


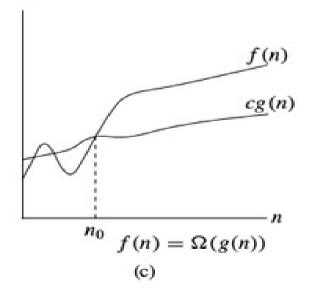
#### $\Omega$ -notation:

To denote asymptotic lower bound, we use  $\Omega$ -notation. For a given function g(n), we denote by  $\Omega(g(n))$  (pronounced "big-omega of g of n") the set of functions:

 $\Omega(g(n)) = \{ f(n) : \text{there exist positive constants } c \text{ and } n_0 \text{ such that } 0 \leq c * g(n) \leq f(n) \text{ for all } n \geq n_0 \}$ 



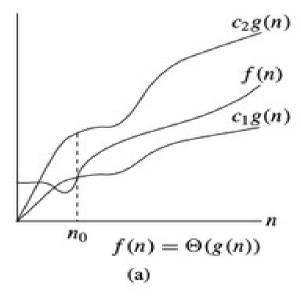


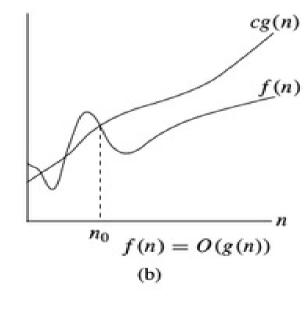


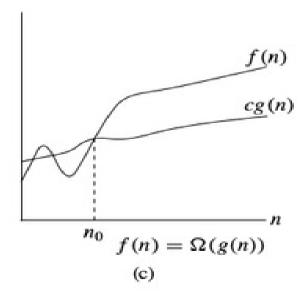
#### $\Theta$ -notation:

To denote asymptotic tight bound, we use  $\Theta$ -notation. For a given function g(n), we denote by  $\Theta(g(n))$  (pronounced "big-theta of g of n") the set of functions:

 $\Theta(g(n)) = \{ f(n) : \text{there exist positive constants } c_1, \ c_2 \ \text{and} \ n_0 \ \text{such that} \ 0 \le c_1 * g(n) \le f(n) \le c_2 * g(n) \ \text{for all} \ n > n_0 \ \}$ 

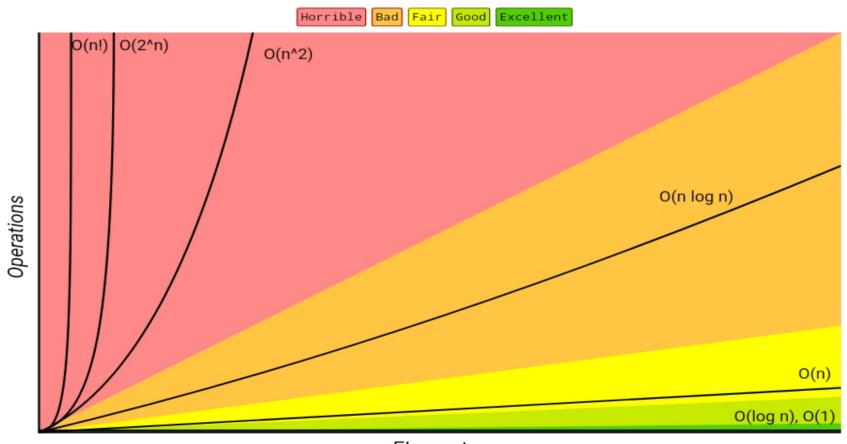






## N vs. Growth of Complexity:

#### **Big-O Complexity Chart**



Elements

### Complexity O(1):

Time complexity of a function (or set of statements) is considered as O(1) if it doesn't contain loop, recursion and call to any other non-constant time function.

- For example swap() function has O(1) time complexity.
- A loop or recursion that runs a constant number of times is also considered as O(1). E.g. the following loop is O(1).

```
for (int i = 1; i <= c; i++) { // Here c is a constant
   // some O(1) expressions
}</pre>
```

### Complexity O(n):

Time Complexity of a loop is considered as O(n) if the loop variable is incremented / decremented by a constant amount. For example following functions have O(n) time complexity.

```
for (int i = 1; i <= n; i += c) {  // Here c is a +ve integer constant
  // some O(1) expressions
}
for (int i = n; i > 0; i -= c) {
  // some O(1) expressions
```

### Complexity $O(n^c)$ :

Time complexity of nested loops is equal to the number of times the innermost statement is executed. For example the following sample loops have  $O(n^2)$  time complexity

```
for (int i = 1; i <=n; i += c) {
    for (int i = n; i > 0; i -= c) {
      for (int j = 1; j <=n; j += c) {
            // some O(1) expressions
      }
    }
}</pre>
```

For example Selection sort and Insertion Sort have O(n²) time complexity.

### Complexity O(logn):

Time Complexity of a loop is considered as O(Log<sub>c</sub>n) if the loop variables is divided / multiplied by a constant amount.

```
for (int i = 1; i <=n; i *= c) {
    // some O(1) expressions
}
for (int i = n; i > 0; i /= c) {
    // some O(1) expressions
}
```

For example Binary Search has O(Log<sub>2</sub>n) time complexity.

### Complexity O(LogLogn):

Time Complexity of a loop is considered as O(LogLog<sub>c</sub>n) if the loop variables is reduced / increased exponentially by a const amount.

```
for (int i = 2; i <=n; i = pow(i, c)) {  // c is a const greater than 1
   // some O(1) expressions
}
// Here fun is sqrt or cuberoot or any other constant root
for (int i = n; i > 1; i = fun(i)) {
   // some O(1) expressions
}
```

### Time complexities of consecutive loops:

Time Complexity When there are consecutive loops, we calculate time complexity as sum of time complexities of individual loops.

```
for (int i = 1; i <=m; i += c) {
    // some O(1) expressions
}
for (int i = 1; i <=n; i += c) {
    // some O(1) expressions
}</pre>
```

```
Time complexity of the code is O(m) + O(n) which is O(m + n)
```

```
If m == n, the time complexity becomes O(2n) which is O(n).
```

# Quiz!

```
for (int i = 0; i < n; i += 1) {
    for (int j = 0; j < m; j +=1) {
        print "Hello"
    }
    }
A. O(n)
B. O(m)
C. O(n+m)</pre>
```

D. O(nm)

```
for (int i = 0; i < n; i += 1) {
    for (int j = 0; j < m; j +=1) {
        print "Hello"
     }
    }
A. O(n)
B. O(m)
C. O(n+m)</pre>
```

```
for (int j = 0; j < n; j += 1) {
    for (int i = 0; i < n; i += 1)
        print "First loop"
    }
for (int i = 0; i < n; i += 1) {
        print "Second loop"
    }
for (int i = 0; i < n; i += 1) {
        print "Third loop"
    }
}</pre>
```

- A. O(n)
- B.  $O(n^2)$
- C. Càn't be determined
- D. None of these

```
for (int j = 0; j < n; j += 1) {
    for (int i = 0; i < n; i += 1)
        print "First loop"
    }
for (int i = 0; i < n; i += 1) {
        print "Second loop"
    }
for (int i = 0; i < n; i += 1) {
        print "Third loop"
    }
}</pre>
```

- A O(n) F. O(n^2)
- C. Càn't be determined
- D. None of these

```
while (var < n){
    cout << pie << endl;
    for (int j = 1; j < n; j+=2){
        sum += 1
    }
    var *= 3
    }

A. O(n)
B. O(logn)
C. O(nlog<sub>3</sub>n)
D. O(n<sup>2</sup>)
```

```
while (var < n){
    cout << pie << endl;
    for (int j = 1; j < n; j+=2){
        sum += 1
    }
    var *= 3
    }

A. O(n)
B. O(logn)
C. O(nlog<sub>3</sub>n)
```

In the outer loop we have var \*= 3, therefore outer loop will have log<sub>3</sub>(n) complexity but the inner loop proceeds 1 to n, therefore ultimately the whole code will have complexity nlog<sub>3</sub>n.

```
for (int i = 0; i < n; i += 3) {
    for (int j = 0; j < m; j += 2) {
        print "Hello"
    }
}
A. O(n)
```

B. O(m) C. O(n+m) D. O(nm)

```
for (int i = 0; i < n; i += 3) {
    for (int j = 0; j < m; j += 2) {
        print "Hello"
    }
}
A. O(n)
B. O(m)
C. O(n+m)</pre>
```

```
for (int i = n/2; i \le n; i + = 1) {
   for (int j = 1; j < m; j *= 2) \{
      print "Hello"
```

- A. O(nlog<sub>2</sub>m) B. O(nm)
- C.  $O(\log_2 n^* \log_2 m)$
- D. None of these

```
for (int i = n/2; i <= n; i += 1) {
  for (int j = 1; j < m; j *= 2) {
          print "Hello"
```

- ⟨⟨⟨⟨ O(nlog₂m⟩⟩ B. O(nm)
- C.  $O(\log_2 n^* \log_2 m)$
- D. None of these

```
int i = 1, j = 1
for (; i < n; i += 1) {
    while(j < i) {
        print "Hello"
```

- A. O(log<sub>2</sub>n) B. O(n)
- C. O(n $\log_2$ n)
- D. None of the above

```
int i = 1, j = 1
for (; i < n; i += 1) {
    while(j < i) {
        print "Hello"
```

A. O(log<sub>2</sub>n)
B. O(n)
C. O(nlog<sub>2</sub>n)
D. None of the above

```
for (int i = 1; i < n; i += i) {
    for (int j = 0; j < n; j += 1) {
        print "Hello"
```

A. O(logn)
B. O(nlogn)
C. O(n)

```
for (int i = 1; i < n; i += i) {
    for (int j = 0; j < n; j += 1) {
        print "Hello"
    }
}</pre>
```

A. O(logn)
B. O(nlogn)
C. O(n)

D. O(n^2)

```
for (int i = 1; i < n; i *= 3) {
        for (int j = 1; j < n; j += 2) {
            print "Hello"
        }
      }

A. O(log<sub>3</sub>n)
B. O(nlog<sub>3</sub>n)
C. O(n)
D. O(n^2)
```

```
for (int i = 1; i < n; i *= 3) {
    for (int j = 1; j < n; j += 2) {
        print "Hello"
    }
    }

A. O(log<sub>3</sub>n)
5. O(nlog<sub>3</sub>n)
C. O(n)
```

```
int i = 1, j = 1
for (; i < n; i +=1) {
    while(j < i) {
        print "Hello"
```

- A. O(logn) B. O(n)
- C. O(nlogn)
- D. None of the above

```
int i = 1, j = 1
for (; i < n; i +=1) {
    while(j < i) {
        print "Hello"
```

- A, O(logn)
  F. O(n)
  C. O(nlogn)
- D. None of the above

```
for (int i = 1; i < n; i += 2) {
      for (int j = 1; j < n; j *= 3) {
          print "Hello"
      }
    }
A. O(log<sub>3</sub>n)
B. O(nlog<sub>3</sub>n)
C. O(n)
D. O(n^2)
```

```
for (int i = 1; i < n; i += 2) {
     for (int j = 1; j < n; j *= 3) {
        print "Hello"
     }
    }

A. O(log<sub>3</sub>n)
5. O(nlog<sub>3</sub>n)
C. O(n)
```