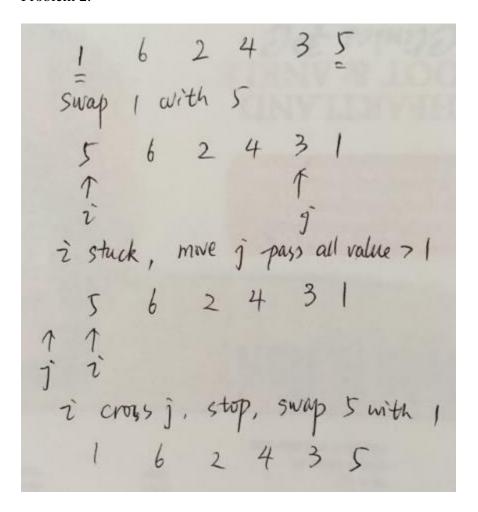
## **Lab 5 Solutions**

# Problem 2.



Problem 3.

a. Good pivots: 2,3,3,4,5

b. Yes: 5/9 of the elements are good pivots.

#### Problem 4.

[22/6] Give an o(n) (that is, better than  $\Theta(n)$ ) algorithm for determining whether a sorted array A of distinct integers contains an element m for which A[m] = m, and then implement as a Java function

```
int findFixedPoint(int[] A)
```

which returns such an m if found, or -1 if no such m is found. You must also provide a proof that your algorithm runs in o(n) time.

```
Step 1: If A[0] = 0, return 0.
Step 2: If A[0] > 0, return -1.
```

Step 3: Do binary search. The base case will examine A[mid] to see if A[mid] = mid, and if so, return A[mid] (it's the a fixed point). If A[mid] > mid, search the left side. If A[mid] < mid, search the right side. If the usual failure signal occurs (lower > upper) return -1. This solves the problem in  $O(\log n)$ .

#### Problem 5.

Review of SubsetSum Problem: Given a set  $S = \{s0, s1, s2, ..., sn-1\}$  of positive integers and a non-negative integer k, find a subset T of S so that the sum of the integers in T equals k or indicate no such subset can be found.

We have already seen a brute force solution to this problem in an earlier lab. In this exercise, you are going to come up with a recursive solution for SubsetSum. Write down pseudo code for your algorithm or Java code.

### Hint:

We are seeking a  $T \subseteq S = \{s_0, s_1, \dots, s_{n-2}, s_{n-1}\}$  whose sum is k. Such a T can be found if and only if one of the following is true:

- (1) A subset  $T_1$  of  $\{s_0, s_1, \dots, s_{n-2}\}$  can be found whose sum is k, OR
- (2) A subset  $T_2$  of  $\{s_0, s_1, \dots, s_{n-2}\}$  can be found whose sum is  $k s_{n-1}$
- If (1) holds, then the desired set T is  $T_1$ . If (2) holds, the desired set T is  $T_2 \cup \{s_{n-1}\}$ .

#### Soln:

```
Algorithm RecSubsetSum(S, k)
    Input: S = \{s_0, s_1, ..., s_{n-1}\} positive integers,
          k nonnegative integer
    Output: T \subseteq S for which sum(T) = k
    //base case
    if S.size() = 1 then
        if k = 0 then return {}
        else if k = s_0 then return \{s_0\}
        else return NULL
    (S, last) \leftarrow S.removeLast()
    T \leftarrow RecSubsetSum(S, k)
    if T not NULL then
        return T
    T \leftarrow RecSubsetSum(S, k - last)
    if T not NULL then
        return T ∪ {last}
    return NULL
```

