

This assignment is based on lecture 11 (chapter 23 – Query Processing)

- Submit your *own work* on time. No credit will be given if the assignment is submitted after the due date.
 - Note that the completed assignment should be submitted in .doc, .docx, .rtf or .pdf format only.
 - In MCQs, if you think that your answer needs more explanation to get credit then please write it down.
 - You are encouraged to discuss these questions in the Sakai forum.
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(1) What are the objectives of query processing?

The aims of query processing are to transform a query written in a high-level language, typically SQL, into a correct and efficient execution strategy expressed in a low-level language (implementing relational algebra), and to execute the strategy to retrieve the required data.

(2) What are the typical phases of query processing?

Query decomposition,

Query optimization,

Code generation

Compile time Runtime query execution

Run time

(3) State the heuristics that should be applied to improve the processing of a query.

Perform Selection operations as early as possible.

Combine the Cartesian product with a subsequent Selection operation whose predicate represents a join condition into a Join operation.

Use associativity of binary operations to rearrange leaf nodes so that the leaf nodes with the most restrictive Selection operations are executed first.

Perform Projection operations as early as possible.

Compute common expressions once.

(4) What types of statistics should a DBMS hold to be able to derive estimates of relational algebra operations?

For each base relation R

$nTuples(R)$ – the number of tuples (records) in relation R (that is, its cardinality).

$n bFactor(R)$ – the blocking factor of R (that is, the number of tuples of R that fit into one block).

$nBlocks(R)$ – the number of blocks required to store R. If the tuples of R are stored physically together, then: $nBlocks(R) = \lceil nTuples(R)/bFactor(R) \rceil$

We use $\lceil x \rceil$ to indicate that the result of the calculation is rounded to the smallest integer that is greater than or equal to x.

For each attribute A of base relation R

$nDistinctA(R)$ – the number of distinct values that appear for attribute A in relation R.

$minA(R), maxA(R)$ – the minimum and maximum possible values for the attribute A in relation R.

$SCA(R)$ – the selection cardinality of attribute A in relation R. This is the average number of tuples that satisfy an equality condition on attribute A.

(5) What are the differences between materialization and pipelining?

Materialization: is the process of temporarily writing the results of intermediate relational algebra operations to disk: the output of one operation is stored in a temporary relation for processing by the next operation.

Pipelining: sometimes known as stream-based processing or on-the-fly processing. And is an alternative approach is to pipeline the results of one operation to another operation without creating a temporary relation to hold the intermediate result. Clearly, if we can use pipelining we can save on the cost of creating temporary relations and reading the results back in again.