

# CS536 Homework 5

Weakest Preconditions 1 & 2; Domain Predicates

Group members: (Group 67)

- Jinit Priyutkumar Parikh (A20517770) (jparikh9@hawk.iit.edu)
- Henilkumar Hareshbhai Patel(A20513297) ([hpatel154@hawk.iit.edu](mailto:hpatel154@hawk.iit.edu))
- Swathi Alapati (A20521251) (salapati3@hawk.iit.edu)

Answers: -

## Question 1:

$IFN \equiv \text{if } B1 \rightarrow S1 \square B2 \rightarrow S2 \text{ fi}$

Then  $wp(IFN, q) \Leftrightarrow ((B1 \wedge W1) \vee (B2 \wedge W2) \vee (B1 \wedge B2 \wedge W1 \wedge W2))$  But for

$p'$  we need to conjunct  $B1 \wedge B2 \wedge W1 \wedge W2$  Otherwise,  $wp(IFN, q) \rightarrow (B1 \wedge$

$W1) \vee (B2 \wedge W2)$  and

This cannot be true  $-wp(IFN, q) \Leftarrow (B1 \wedge W1) \vee (B2 \wedge W2)$

Because  $B1$  and  $B2$  both hold true then we need  $W1$  and  $W2$  to hold the same.

And for the third conjunct it never holds true because it requires  $B \wedge \neg B$ .

## Question 2:

$Wp(S, p \vee q) \rightarrow wp(S, p) \vee wp(S, q)$

Here, if  $\sigma \models wp(S, p \vee q)$ , then  $\models p \vee q$  but it does not imply  $U \models p$  and  $U \models q$

because  $\tau \in U$  does not satisfy all  $p$  and  $q$ . So, this case is not valid.  $Wp(S, p) \vee wp(S, q) \rightarrow wp(S, p \vee q)$

If  $\sigma \models wp(S, p) \vee wp(S, q)$  then it implies that  $\sigma \models wp(S, p)$  or  $\sigma \models wp(S, q)$

therefore  $\tau \in U$  which implies that  $\tau \models p$  or  $\tau \models q$ , so either way it implies  $\tau \models p \vee q$  that implies  $\sigma \models wp(S, p \vee q)$

- $wp(S, p \wedge q) \rightarrow wp(S, p) \wedge wp(S, q)$

In this case, for  $S$  in a state  $\sigma$  satisfies  $p \wedge q$  from  $wp$  therefore  $\tau \in U$  satisfies  $p \wedge q$  which implies that  $\tau \models p$  and  $\tau \models q$  which implies that  $\sigma \models wp(S, p)$  and  $\sigma \models wp(S, q)$

$wp(S, p) \wedge wp(S, q) \rightarrow wp(S, p \wedge q)$

If  $\sigma \models wp(S, p) \wedge wp(S, q)$  then it implies that

$\sigma \models wp(S, p)$  and  $\sigma \models wp(S, q)$

therefore  $\tau \in U$  which implies that  $\tau \models p$  and  $\tau \models q$ , So it implies that  $\tau \models p \wedge q$  that implies  $\sigma \models wp(S, p \wedge q)$

### Question 3:

if  $S$  is deterministic, then  $\models_{\text{tot}} \{b\} S \{q\}$ .

$\models_{\text{tot}} \{b\} S \{q\}$  is true, because  $b$  is stronger than  $wpw (b \rightarrow w)$

- If  $S$  is nondeterministic, then there exists  $\sigma$  such that  $\sigma \models \{\neg c\} S \{\neg q\}$ .

$\sigma \models \{\neg c\} S \{\neg q\}$  is true because  $\sigma \models \{\neg w\} S \{\neg q\}$  and  $w \rightarrow c$  implies  $\neg c \rightarrow \neg w$  that follows  $\sigma \models \{\neg c\} S \{\neg q\}$

- If  $S$  is nondeterministic, then there exist  $\sigma \models \neg c$  and  $\tau \in M(S, \sigma)$  such that  $\tau \models q$ .

True as per the definition of correctness. If  $\tau \in M(S, \sigma)$  such that  $\tau \models q$  then  $\sigma$  is in the  $wp$  and  $wlp$  of  $S$  and  $q$

### Question 4

$wlp(u := u * k; k := u, u > h(k))$

$\equiv wlp(u := u * k, wlp(k := u, u > h(k)))$

$\equiv wlp(u := u * k, u > h(u))$

$\equiv u * k > h(u * k)$

### Question 5

$wlp(\text{if } x < 0 \text{ then } x := -x \text{ fi, } x^2 \geq x)$

$\equiv wlp(\text{if } x < 0 \text{ then } x := -x \text{ else skip fi, } x^2 \geq x)$

$\equiv (x < 0 \rightarrow wlp(x := -x, x^2 \geq x)) \wedge (x \geq 0 \rightarrow wlp(\text{skip}, x^2 \geq x))$

$\equiv (x < 0 \rightarrow -x^2 \geq -x) \wedge (x \geq 0 \rightarrow x^2 \geq x)$

### Question 6

$wp(y := y/x, \text{sqrt}(y) < x)$

- Let  $S \equiv y := y/x$ ,  $q \equiv \text{sqrt}(y) < x$ , and  $w \equiv wlp(S, q)$ .
- We can expand  $w \equiv wlp(S, q) \equiv wlp(y := y/x, \text{sqrt}(y) < x) \equiv \text{sqrt}(y/x) < x$ .
- We can calculate  

$$D(\text{sqrt}(y/x) < x)$$

$$\equiv D(y/x) \wedge y/x \geq 0$$

$$\equiv D(y) \wedge D(x) \wedge x \neq 0 \wedge y/x \geq 0 \equiv x \neq 0 \wedge y/x \geq 0 \quad (D(y) \text{ and } D(x) \text{ is } T)$$
- $wp(S, q) \equiv D(S) \wedge w \wedge D(w)$   

$$\equiv D(y := y/x) \wedge \text{sqrt}(y/x) < x \wedge D(\text{sqrt}(y/x) < x)$$

$$\begin{aligned} &\equiv (D(y) \wedge D(x) \wedge x \neq 0) \wedge \text{sqrt}(y/x) < x \wedge x \neq 0 \wedge y/x \geq 0 \\ &\equiv \text{sqrt}(y/x) < x \wedge x \neq 0 \wedge y/x \geq 0 \end{aligned}$$

### Question 7

Wp (if  $y \geq 0$  then  $x := y/x$  else  $x := -x/y$  fi,  $r < x \leq y$ )

By definition,

$$\text{wp}(S, q) = D(S) \wedge D(w) \wedge w$$

where,

$$q = r < x \leq y$$

So,  $w = \text{wlp}(S, q)$

$$\begin{aligned} &= \text{wlp}(\text{if } y \geq 0 \text{ then } x := y/x \text{ else } x := -x/y, q = r < x \leq y) \\ &= (y \geq 0 \rightarrow \text{wlp}(x := y/x, r < x \leq y)) \wedge (y < 0 \rightarrow \text{wlp}(x := -x/y, q = r < x \leq y)) \\ &= (y \geq 0 \rightarrow r < (y/x) \leq y) \wedge (y < 0 \rightarrow r < (-x/y) \leq y) \end{aligned}$$

Now,

$$\begin{aligned} D(w) &= D((y \geq 0 \rightarrow r < (y/x) \leq y) \wedge (y < 0 \rightarrow r < (-x/y) \leq y)) \\ &= D(y \geq 0 \rightarrow r < (y/x) \leq y) \wedge D(y < 0 \rightarrow r < (-x/y) \leq y) \\ &= x \neq 0 \wedge y \neq 0 \end{aligned}$$

Now,

$$\begin{aligned} D(S) &= D(\text{if } y \geq 0 \text{ then } x := y/x \text{ else } x := -x/y) \\ &= D(y \geq 0) \wedge (y \geq 0 \rightarrow D(x := y/x)) \wedge (y < 0 \rightarrow D(x := -x/y)) \\ &= T \wedge (y \geq 0 \rightarrow x \neq 0) \wedge (y < 0 \rightarrow y \neq 0) \\ &= (y \geq 0 \rightarrow x \neq 0) \wedge (y < 0 \rightarrow y \neq 0) \end{aligned}$$

So,

$$\text{Wp}(S, q) = D(S) \wedge D(w) \wedge w$$

$$(y \geq 0 \rightarrow x \neq 0) \wedge (y < 0 \rightarrow y \neq 0) \wedge x \neq 0 \wedge y \neq 0 \wedge (y \geq 0 \rightarrow r < (y/x) \leq y) \wedge (y < 0 \rightarrow r < (-x/y) \leq y)$$