

Problem 0

a.

States	Single word unigrams; a single letter in the input query
Actions	Insert vowel, insert space
Cost	$F_b(\text{action})$
Start State	"-BEGIN-" keyword so we can start before index 0
End State	Last letter in query

$H(s)$ is consistent because the cost for adding a vowel or space always \leq adding a whole new word with a space.

b. BFS is a special case of UCS, which is a special case of A^* .

c. Writing out potential values of n :

$n = 0$: length of query is 0 \rightarrow states_expanded = 0

$n = 5$: length of query is 5 \rightarrow start, , , , end \rightarrow states_expanded = 5

...

Therefore upper-bound = $O(n)$.

d. The two answers are not similar, as mine was not exponential. This could be because mine did not consider different options for adding vowels or spaces. I only considered either we add a vowel or space, or we don't. But I did not consider if we could have different options for adding vowels and spaces (e.g. adding multiple vowels).

Problem 1

a. $V_{\text{minmax}}(s, d) = \{$

Gamevalue(s);	if isend(s)
Eval(s);	if $d = 0$
$\text{Max}_{a \text{ in } A(s)} V_{\text{minmax}}(\text{next}(s, a), d);$	if player(s) == max
$\text{Min}_{a \text{ in } A(s)} V_{\text{minmax}}(\text{next}(s, a), d-1);$	if player(s) == min
}	

Problem 3

a. $V_{\text{minmax}}(s, d) = \{$

Gamevalue(s);	if isend(s)
Eval(s);	if $d = 0$
$\text{Max}_{a \text{ in } A(s)} V_{\text{minmax}}(\text{next}(s, a), d);$	if player(s) == max
$\sum_{a \text{ in } A(s)} \pi_{\text{opp}}(s, a) V_{\text{minmax}}(\text{next}(s, a), d-1);$	if player(s) == min
}	