x	term variable
α	type variable
T	type
c	coercion variable

t	$::=$			term
	$\mid$	x		variable
	$\mid$	$\lambda(x : \tau) \Rightarrow t$		abstraction
	$\mid$	$\lambda\{\alpha : \kappa\} \Rightarrow t$		type abstraction
	$\mid$	$\lambda\{c : \tau_1 \sim \tau_2\} \Rightarrow t$		coercion abstraction
	$\mid$	$t_1\ t_2$		application
	$\mid$	$t\ [\tau]$		type application
	$\mid$	$t\ {\sim}[\gamma]$		coercion application
	$\mid$	$t : \tau$		type annotation
	$\mid$	$t \blacktriangleright \gamma$		coercion
	$\mid$	(t)	S	parenthesis
v	$::=$			value
	$\mid$	$\lambda(x : \tau) \Rightarrow t$		abstraction
κ	$::=$			kind
	$\mid$	$*$		star
	$\mid$	$\kappa_1 \to \kappa_2$		kind arrow
	$\mid$	(κ)	S	parenthesis
τ	$::=$			type
	$\mid$	α		type variable
	$\mid$	T		type
	$\mid$	$\tau_1 \to \tau_2$		arrow
	$\mid$	$\{\tau_1 \sim \tau_2\} \rightsquigarrow \tau_3$		coercion arrow
	$\mid$	$\lambda(\alpha : \kappa), \tau$		operator abstraction
	$\mid$	$\forall\,(\alpha : \kappa), \tau$		forall
	$\mid$	$\tau_1\,\tau_2$		operator application
	$\mid$	(τ)	S	parenthesis
γ	$::=$			coercion proof term
	$\mid$	c		variable
	$\mid$	$\mathbf{refl}\,\tau$		reflexivity
	$\mid$	$\mathbf{sym}\,\gamma$		symmetry
	$\mid$	$\gamma_1 \circ \gamma_2$		composition
	$\mid$	$\gamma_1 \to \gamma_2$		arrow introduction
	$\mid$	$\{\tau_1 \sim \tau_2\} \rightsquigarrow \gamma$		coercion arrow introduction
	$\mid$	$\lambda(\alpha : \kappa), \gamma$		operator abstraction introduction
	$\mid$	$\forall\,(\alpha : \kappa), \gamma$		forall introduction
	$\mid$	$\gamma_1\,\gamma_2$		application introduction
	$\mid$	$\mathbf{left}\,\gamma$		left elimination
	$\mid$	$\mathbf{right}\,\gamma$		right elimination
	$\mid$	(γ)	S	parenthesis
Γ	$::=$			typing environment
	$\mid$	$\emptyset$		empty
	$\mid$	$\Gamma, x : \tau$		variable

$$
\begin{array}{lll}
| & \Gamma, T : \kappa & \text{abstract type} \\
| & \Gamma, \alpha : \kappa & \text{type variable} \\
| & \Gamma, c : c_1 \sim c_2 & \text{coercion variable}
\end{array}
$$

$\boxed{[\alpha \mapsto \tau_1]\tau_2 \rhd \tau_3}$ Type substitution

$$\frac{}{[\alpha \mapsto \tau]\alpha \rhd \tau} \quad \textsc{SubstT\_Var1}$$

$$\frac{}{[\alpha_1 \mapsto \tau]\alpha_2 \rhd \alpha_2} \quad \textsc{SubstT\_Var2}$$

$$\frac{}{[\alpha \mapsto \tau]T \rhd T} \quad \textsc{SubstT\_Type}$$

$$\frac{\begin{array}{l}[\alpha \mapsto \tau_1]\tau_2 \rhd \tau_2' \\ [\alpha \mapsto \tau_1]\tau_3 \rhd \tau_3'\end{array}}{[\alpha \mapsto \tau_1]\tau_2 \to \tau_3 \rhd \tau_2' \to \tau_3'} \quad \textsc{SubstT\_Arrow}$$

$$\frac{}{[\alpha \mapsto \tau_1]\lambda(\alpha : \kappa), \tau_2 \rhd \lambda(\alpha : \kappa), \tau_2} \quad \textsc{SubstT\_Abs1}$$

$$\frac{[\alpha_1 \mapsto \tau_1]\tau_2 \rhd \tau_2'}{[\alpha_1 \mapsto \tau_1]\lambda(\alpha_2 : \kappa), \tau_2 \rhd \lambda(\alpha_2 : \kappa), \tau_2'} \quad \textsc{SubstT\_Abs2}$$

$$\frac{}{[\alpha \mapsto \tau_1]\forall\,(\alpha : \kappa), \tau_2 \rhd \forall\,(\alpha : \kappa), \tau_2} \quad \textsc{SubstT\_Forall1}$$

$$\frac{[\alpha_1 \mapsto \tau_1]\tau_2 \rhd \tau_2'}{[\alpha_1 \mapsto \tau_1]\forall\,(\alpha_2 : \kappa), \tau_2 \rhd \forall\,(\alpha_2 : \kappa), \tau_2'} \quad \textsc{SubstT\_Forall2}$$

$$\frac{\begin{array}{l}[\alpha \mapsto \tau_1]\tau_2 \rhd \tau_2' \\ [\alpha \mapsto \tau_1]\tau_3 \rhd \tau_3'\end{array}}{[\alpha \mapsto \tau_1]\tau_2\,\tau_3 \rhd \tau_2'\,\tau_3'} \quad \textsc{SubstT\_App}$$

$\boxed{\Gamma \vdash t : \tau}$ Typing rules

$$\frac{x : \tau \in \Gamma}{\Gamma \vdash x : \tau} \quad \textsc{T\_Var}$$

$$\frac{\begin{array}{l}\Gamma, x : \tau_1 \vdash t : \tau_2 \\ \Gamma \vdash \tau_1 : *\end{array}}{\Gamma \vdash \lambda(x : \tau_1) \Rightarrow t : \tau_1 \to \tau_2} \quad \textsc{T\_Abs}$$

$$\frac{\begin{array}{l}\Gamma \vdash t_1 : \tau_2 \to \tau_1 \\ \Gamma \vdash t_2 : \tau_2\end{array}}{\Gamma \vdash t_1\,t_2 : \tau_1} \quad \textsc{T\_App}$$

$$\frac{\begin{array}{l}\Gamma, \alpha : \kappa \vdash t : \tau \\ \alpha \notin \Gamma\end{array}}{\Gamma \vdash \lambda\{\alpha : \kappa\} \Rightarrow t : \forall\,(\alpha : \kappa), \tau} \quad \textsc{T\_TyAbs}$$

$$\frac{\begin{array}{l}\Gamma \vdash t : \forall\,(\alpha : \kappa), \tau_2 \\ \Gamma \vdash \tau_1 : \kappa \\ [\alpha \mapsto \tau_1]\tau_2 \rhd \tau_2'\end{array}}{\Gamma \vdash t\ [\tau_1] : \tau_2'} \quad \textsc{T\_TyApp}$$

$$\frac{\begin{array}{l}\Gamma \vdash t : \tau_2 \\ \tau_1 \equiv \tau_2\end{array}}{\Gamma \vdash (t : \tau_1) : \tau_1} \quad \textsc{T\_Annot}$$

$\boxed{\Gamma \vdash \tau : \kappa}$ Kinding

$$\frac{\alpha : \kappa \in \Gamma}{\Gamma \vdash \alpha : \kappa} \quad \textsc{K\_Var}$$

$$\frac{T : \kappa \in \Gamma}{\Gamma \vdash T : \kappa} \quad \textsc{K\_AbsType}$$

$$\frac{\Gamma, \alpha : \kappa_1 \vdash \tau : \kappa_2}{\Gamma \vdash \lambda(\alpha : \kappa_1), \tau : \kappa_1 \to \kappa_2} \quad \textsc{K\_Abs}$$

$$\frac{\begin{array}{c}\Gamma \vdash \tau_1 : \kappa_2 \to \kappa_1 \\ \Gamma \vdash \tau_2 : \kappa_2\end{array}}{\Gamma \vdash \tau_1\, \tau_2 : \kappa_1} \quad \textsc{K\_App}$$

$$\frac{\begin{array}{c}\Gamma \vdash \tau_1 : * \\ \Gamma \vdash \tau_2 : *\end{array}}{\Gamma \vdash \tau_1 \to \tau_2 : *} \quad \textsc{K\_Arrow}$$

$$\frac{\Gamma, \alpha : \kappa \vdash \tau : *}{\Gamma \vdash \forall\, (\alpha : \kappa), \tau : *} \quad \textsc{K\_Forall}$$

$\boxed{\tau_1 \equiv \tau_2}$ Type equivalence

$$\frac{}{\tau \equiv \tau} \quad \textsc{EQ\_Refl}$$

$$\frac{\tau_2 \equiv \tau_1}{\tau_1 \equiv \tau_2} \quad \textsc{EQ\_Symm}$$

$$\frac{\begin{array}{c}\tau_1 \equiv \tau_2 \\ \tau_2 \equiv \tau_3\end{array}}{\tau_1 \equiv \tau_3} \quad \textsc{EQ\_Trans}$$

$$\frac{}{\alpha \equiv \alpha} \quad \textsc{EQ\_Var}$$

$$\frac{}{T \equiv T} \quad \textsc{EQ\_AbsType}$$

$$\frac{\begin{array}{c}\tau_1 \equiv \tau_1' \\ \tau_2 \equiv \tau_2'\end{array}}{\tau_1 \to \tau_2 \equiv \tau_1' \to \tau_2'} \quad \textsc{EQ\_Arrow}$$

$$\frac{\tau_1 \equiv \tau_2}{\forall\, (\alpha : \kappa), \tau_1 \equiv \forall\, (\alpha : \kappa), \tau_2} \quad \textsc{EQ\_Forall}$$

$$\frac{\tau_1 \equiv \tau_2}{\lambda(\alpha : \kappa), \tau_1 \equiv \lambda(\alpha : \kappa), \tau_2} \quad \textsc{EQ\_Abs}$$

$$\frac{\begin{array}{c}\tau_1 \equiv \tau_1' \\ \tau_2 \equiv \tau_2'\end{array}}{\tau_1\, \tau_2 \equiv \tau_1'\, \tau_2'} \quad \textsc{EQ\_App}$$

$$\frac{[\alpha \mapsto \tau_2]\tau_1 \rhd \tau_1'}{(\lambda(\alpha : \kappa), \tau_1)\, \tau_2 \equiv \tau_1'} \quad \textsc{EQ\_AppAbs}$$

$\boxed{[x \mapsto v]t_1 \rhd t_2}$ substitution

$$\frac{}{[x \mapsto v]x \rhd v} \quad \textsc{Subst\_Var1}$$

$$\frac{}{[x_1 \mapsto v] x_2 \rhd x_2} \quad \textsc{Subst\_Var2}$$

$$\frac{}{[x \mapsto v] \lambda (x : \tau) \Rightarrow t \rhd \lambda (x : \tau) \Rightarrow t} \quad \textsc{Subst\_Abs1}$$

$$\frac{[x_1 \mapsto v] t_1 \rhd t_2}{[x_1 \mapsto v] \lambda (x_2 : \tau) \Rightarrow t_1 \rhd \lambda (x_2 : \tau) \Rightarrow t_2} \quad \textsc{Subst\_Abs2}$$

$$\frac{[x \mapsto v] t_1 \rhd t_2}{[x \mapsto v] \lambda \{\alpha : \kappa\} \Rightarrow t_1 \rhd \lambda \{\alpha : \kappa\} \Rightarrow t_2} \quad \textsc{Subst\_TAbs}$$

$$\frac{[x \mapsto v] t_1 \rhd t_2}{[x \mapsto v] \lambda \{c : \tau_1 \sim \tau_2\} \Rightarrow t_1 \rhd \lambda \{c : \tau_1 \sim \tau_2\} \Rightarrow t_2} \quad \textsc{Subst\_CAbs}$$

$$\frac{\begin{array}{c}[x \mapsto v] t_1 \rhd t_1' \\ [x \mapsto v] t_2 \rhd t_2'\end{array}}{[x \mapsto v] t_1\ t_2 \rhd t_1'\ t_2'} \quad \textsc{Subst\_App}$$

$$\frac{[x \mapsto v] t_1 \rhd t_2}{[x \mapsto v] t_1\ [\tau] \rhd t_2\ [\tau]} \quad \textsc{Subst\_TApp}$$

$$\frac{[x \mapsto v] t_1 \rhd t_2}{[x \mapsto v] t_1\ \sim[\gamma] \rhd t_2\ \sim[\gamma]} \quad \textsc{Subst\_CApp}$$

$$\frac{[x \mapsto v] t_1 \rhd t_2}{[x \mapsto v] t_1 : \tau \rhd t_2 : \tau} \quad \textsc{Subst\_Annot}$$

$$\frac{[x \mapsto v] t_1 \rhd t_2}{[x \mapsto v] t_1 \blacktriangleright \gamma \rhd t_2 \blacktriangleright \gamma} \quad \textsc{Subst\_Coerce}$$

$\boxed{t_1 \longrightarrow t_2}$ Evaluation

$$\frac{t_2 \longrightarrow t_2'}{t_1\ t_2 \longrightarrow t_1\ t_2'} \quad \textsc{E\_App1}$$

$$\frac{t_1 \longrightarrow t_1'}{t_1\ v \longrightarrow t_1'\ v} \quad \textsc{E\_App2}$$

$$\frac{\begin{array}{c}[x \mapsto v] t \rhd t' \\ t' \longrightarrow t''\end{array}}{(\lambda (x : \tau) \Rightarrow t)\ v \longrightarrow t''} \quad \textsc{E\_AppAbs}$$

$$\frac{t \longrightarrow t'}{\lambda \{\alpha : \kappa\} \Rightarrow t \longrightarrow t'} \quad \textsc{E\_TAbs}$$

$$\frac{t \longrightarrow t'}{\lambda \{c : \tau_1 \sim \tau_2\} \Rightarrow t \longrightarrow t'} \quad \textsc{E\_CAbs}$$

$$\frac{t \longrightarrow t'}{t\ [\tau] \longrightarrow t'} \quad \textsc{E\_TApp}$$

$$\frac{t \longrightarrow t'}{t\ \sim[\gamma] \longrightarrow t'} \quad \textsc{E\_CApp}$$

$$\frac{t \longrightarrow t'}{t : \tau \longrightarrow t'} \quad \textsc{E\_Annot}$$

$$\frac{t \longrightarrow t'}{t \blacktriangleright \gamma \longrightarrow t'} \quad \textsc{E\_Coerce}$$

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Definition rules:        51 good    0 bad
Definition rule clauses: 106 good    0 bad
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