x	term variable
α	type variable
T	type constructor
c	coercion variable

t	$::=$		term
	$\mid$	x	variable
	$\mid$	$\lambda(x : \tau) \Rightarrow t$	abstraction
	$\mid$	$\lambda\{\alpha : \kappa\} \Rightarrow t$	type abstraction
	$\mid$	$\lambda\{c : \tau_1 \sim \tau_2\} \Rightarrow t$	coercion abstraction
	$\mid$	$t_1\, t_2$	application
	$\mid$	$t\ [\tau]$	type application
	$\mid$	$t\ \sim[\gamma]$	coercion application
	$\mid$	$t : \tau$	type annotation
	$\mid$	$t \blacktriangleright \gamma$	coercion
κ	$::=$		kind
	$\mid$	$*$	star
	$\mid$	$\kappa_1 \to \kappa_2$	kind arrow
τ	$::=$		type
	$\mid$	α	type variable
	$\mid$	T	type constructor
	$\mid$	$\tau_1 \to \tau_2$	$\equiv (\to)\ \tau_1\ \tau_2$
	$\mid$	$\{\tau_1 \sim \tau_2\} \to \tau_3$	coercion arrow
	$\mid$	$\lambda(\alpha : \kappa), \tau$	operator abstraction
	$\mid$	$\forall\,(\alpha : \kappa), \tau$	universal quantification
	$\mid$	$\tau_1\, \tau_2$	operator application
γ	$::=$		coercion proof term
	$\mid$	c	variable
	$\mid$	$\mathbf{refl}\, \tau$	reflexivity
	$\mid$	$\mathbf{sym}\, \gamma$	symmetry
	$\mid$	$\gamma_1 \circ \gamma_2$	composition
	$\mid$	$\gamma_1 \to \gamma_2$	$\equiv (\to)\ \gamma_1\ \gamma_2$ where $(\to) :$ refl $(\to)$
	$\mid$	$\{\gamma_1 \sim \gamma_2\} \to \gamma_3$	coercion arrow introduction
	$\mid$	$\lambda(\alpha : \kappa), \gamma$	operator abstraction introduction
	$\mid$	$\forall\,(\alpha : \kappa), \gamma$	universal quantification introduction
	$\mid$	$\gamma_1\, \gamma_2$	application introduction
	$\mid$	$\mathbf{left}\, \gamma$	left elimination
	$\mid$	$\mathbf{right}\, \gamma$	right elimination
Γ	$::=$		typing environment
	$\mid$	$\emptyset$	empty
	$\mid$	$\Gamma, x : \tau$	variable
	$\mid$	$\Gamma, T : \kappa$	type constructor
	$\mid$	$\Gamma, \alpha : \kappa$	type variable
	$\mid$	$\Gamma, c : \tau_1 \sim \tau_2$	coercion variable
tv	$::=$		typed value
	$\mid$	$\lambda(x : \tau) \Rightarrow t$	abstraction
	$\mid$	$\lambda\{\alpha : \kappa\} \Rightarrow tv$	type abstraction

	$\mid$	$\lambda\{c : \tau_1 \sim \tau_2\} \Rightarrow tv$	coercion abstraction
	$\mid$	$tv\ [\tau]$	type application
	$\mid$	$tv\ \sim[\gamma]$	coercion application
	$\mid$	$tv : \tau$	type annotation
	$\mid$	$tv \blacktriangleright \gamma$	coercion
v	$::=$		value
	$\mid$	$\lambda x \Rightarrow t$	abstraction

Initial environment: $\Gamma = \emptyset, (\to) : * \to * \to *$

$\boxed{\Gamma \vdash t : \tau}$ Typing rules

$$\frac{x : \tau \in \Gamma}{\Gamma \vdash x : \tau}\ \text{T\_VAR}$$

$$\frac{\Gamma, x : \tau_1 \vdash t : \tau_2 \qquad \Gamma \vdash \tau_1 : *}{\Gamma \vdash (\lambda(x : \tau_1) \Rightarrow t) : \tau_1 \to \tau_2}\ \text{T\_ABS}$$

$$\frac{\alpha \notin \Gamma \qquad \Gamma, \alpha : \kappa \vdash t : \tau}{\Gamma \vdash (\lambda\{\alpha : \kappa\} \Rightarrow t) : \forall (\alpha : \kappa), \tau}\ \text{T\_TYABS}$$

$$\frac{\Gamma, c : \tau_1 \sim \tau_2 \vdash t : \tau_3 \qquad \Gamma \vdash \tau_1 : \kappa \qquad \Gamma \vdash \tau_2 : \kappa}{\Gamma \vdash \lambda\{c : \tau_1 \sim \tau_2\} \Rightarrow t : \{\tau_1 \sim \tau_2\} \to \tau_3}\ \text{T\_CABS}$$

$$\frac{\Gamma \vdash t_1 : \tau_2 \to \tau_1 \qquad \tau_2 \equiv \tau_2' \qquad \Gamma \vdash t_2 : \tau_2'}{\Gamma \vdash t_1\ t_2 : \tau_1}\ \text{T\_APP}$$

$$\frac{\Gamma \vdash t : \forall (\alpha : \kappa), \tau_2 \qquad \Gamma \vdash \tau_1 : \kappa \qquad [\alpha \mapsto \tau_1]\ \tau_2 \rhd \tau_2'}{\Gamma \vdash t\ [\tau_1] : \tau_2'}\ \text{T\_TYAPP}$$

$$\frac{\Gamma \vdash t : \{\tau_1 \sim \tau_2\} \to \tau_3 \qquad \Gamma \vdash \gamma : \tau_1' \sim \tau_2' \qquad \tau_1 \equiv \tau_1' \qquad \tau_2 \equiv \tau_2'}{\Gamma \vdash t\ \sim[\gamma] : \tau_3}\ \text{T\_CAPP}$$

$$\frac{\Gamma \vdash t : \tau_2 \qquad \tau_1 \equiv \tau_2}{\Gamma \vdash (t : \tau_1) : \tau_1}\ \text{T\_ANNOT}$$

$$\frac{\Gamma \vdash \gamma : \tau_1 \sim \tau_2 \qquad \Gamma \vdash t : \tau_1' \qquad \tau_1 \equiv \tau_1'}{\Gamma \vdash (t \blacktriangleright \gamma) : \tau_2}\ \text{T\_COERCE}$$

$\boxed{\Gamma \vdash \gamma : \tau_1 \sim \tau_2}$ Coercion typing

$$\frac{c : \tau_1 \sim \tau_2 \in \Gamma}{\Gamma \vdash c : \tau_1 \sim \tau_2}\ \text{C\_VAR}$$

$$\frac{}{\Gamma \vdash \mathbf{refl}\,\tau : \tau \sim \tau}\ \text{C\_REFL}$$

$$\frac{\Gamma \vdash \gamma : \tau_2 \sim \tau_1}{\Gamma \vdash \mathbf{sym}\,\gamma : \tau_1 \sim \tau_2}\ \text{C\_SYM}$$

$$\frac{\Gamma \vdash \gamma_1 : \tau_1 \sim \tau_2 \qquad \tau_2 \equiv \tau_2' \qquad \Gamma \vdash \gamma_2 : \tau_2' \sim \tau_3}{\Gamma \vdash \gamma_1 \circ \gamma_2 : \tau_1 \sim \tau_3}\ \text{C\_COMP}$$

$$\frac{\Gamma \vdash \gamma_1 : \tau_1 \sim \tau_1' \qquad \Gamma \vdash \gamma_2 : \tau_2 \sim \tau_2' \qquad \Gamma \vdash \gamma_3 : \tau_3 \sim \tau_3' \qquad \Gamma \vdash \{\tau_1 \sim \tau_2\} \to \tau_3 : *}{\Gamma \vdash \{\gamma_1 \sim \gamma_2\} \to \gamma_3 : \{\tau_1 \sim \tau_2\} \to \tau_3 \sim \{\tau_1' \sim \tau_2'\} \to \tau_3'} \quad \text{C\_CARROW}$$

$$\frac{\alpha \notin \Gamma \qquad \Gamma, \alpha : \kappa \vdash \gamma : \tau_1 \sim \tau_2}{\Gamma \vdash \lambda(\alpha : \kappa), \gamma : \lambda(\alpha : \kappa), \tau_1 \sim \lambda(\alpha : \kappa), \tau_2} \quad \text{C\_ABS}$$

$$\frac{\alpha \notin \Gamma \qquad \Gamma, \alpha : \kappa \vdash \gamma : \tau_1 \sim \tau_2 \qquad \Gamma \vdash \forall(\alpha : \kappa), \tau_1 : *}{\Gamma \vdash \forall(\alpha : \kappa), \gamma : \forall(\alpha : \kappa), \tau_1 \sim \forall(\alpha : \kappa), \tau_2} \quad \text{C\_FORALL}$$

$$\frac{\Gamma \vdash \gamma_1 : \tau_1 \sim \tau_1' \qquad \Gamma \vdash \gamma_2 : \tau_2 \sim \tau_2' \qquad \Gamma \vdash \tau_1\, \tau_2 : \kappa}{\Gamma \vdash \gamma_1\, \gamma_2 : \tau_1\, \tau_2 \sim \tau_1'\, \tau_2'} \quad \text{C\_APP}$$

$$\frac{\Gamma \vdash \gamma : \tau_1 \to \tau_2 \sim \tau_1' \to \tau_2'}{\Gamma \vdash \mathbf{left}\, \gamma : \tau_1 \sim \tau_1'} \quad \text{C\_LEFT1}$$

$$\frac{\Gamma \vdash \gamma : \tau_1\, \tau_2 \sim \tau_1'\, \tau_2'}{\Gamma \vdash \mathbf{left}\, \gamma : \tau_1 \sim \tau_1'} \quad \text{C\_LEFT2}$$

$$\frac{\Gamma \vdash \gamma : \tau_1 \to \tau_2 \sim \tau_1' \to \tau_2'}{\Gamma \vdash \mathbf{right}\, \gamma : \tau_2 \sim \tau_2'} \quad \text{C\_RIGHT1}$$

$$\frac{\Gamma \vdash \gamma : \{\tau_1 \sim \tau_2\} \to \tau_3 \sim \{\tau_1' \sim \tau_2'\} \to \tau_3'}{\Gamma \vdash \mathbf{right}\, \gamma : \tau_3 \sim \tau_3'} \quad \text{C\_RIGHT2}$$

$$\frac{\Gamma \vdash \gamma : \tau_1\, \tau_2 \sim \tau_1'\, \tau_2'}{\Gamma \vdash \mathbf{right}\, \gamma : \tau_2 \sim \tau_2'} \quad \text{C\_RIGHT3}$$

$\boxed{\Gamma \vdash \tau : \kappa}$ Kinding rules

$$\frac{\alpha : \kappa \in \Gamma}{\Gamma \vdash \alpha : \kappa} \quad \text{K\_VAR}$$

$$\frac{T : \kappa \in \Gamma}{\Gamma \vdash T : \kappa} \quad \text{K\_TYPECONSTR}$$

$$\frac{\alpha \notin \Gamma \qquad \Gamma, \alpha : \kappa_1 \vdash \tau : \kappa_2}{\Gamma \vdash (\lambda(\alpha : \kappa_1), \tau) : \kappa_1 \to \kappa_2} \quad \text{K\_ABS}$$

$$\frac{\Gamma \vdash \tau_1 : \kappa_2 \to \kappa_1 \qquad \Gamma \vdash \tau_2 : \kappa_2}{\Gamma \vdash \tau_1\, \tau_2 : \kappa_1} \quad \text{K\_APP}$$

$$\frac{\Gamma \vdash \tau_1 : \kappa \qquad \Gamma \vdash \tau_2 : \kappa \qquad \Gamma \vdash \tau_3 : *}{\Gamma \vdash (\{\tau_1 \sim \tau_2\} \to \tau_3) : *} \quad \text{K\_CARROW}$$

$$\frac{\alpha \notin \Gamma \qquad \Gamma, \alpha : \kappa \vdash \tau : *}{\Gamma \vdash (\forall(\alpha : \kappa), \tau) : *} \quad \text{K\_FORALL}$$

$\boxed{\tau_1 \equiv \tau_2}$ Type equivalence

$$\frac{}{\tau \equiv \tau} \quad \text{EQ\_REFL}$$

$$\frac{\tau_2 \equiv \tau_1}{\tau_1 \equiv \tau_2} \quad \text{EQ\_SYM}$$

$$\frac{\tau_1 \equiv \tau_2 \qquad \tau_2 \equiv \tau_3}{\tau_1 \equiv \tau_3} \quad \text{EQ\_TRANS}$$

$$\frac{}{\alpha \equiv \alpha} \quad \text{EQ\_VAR}$$

$$\frac{}{T \equiv T}\ \textsc{Eq\_TypeConstr}$$

$$\frac{\tau_1 \equiv \tau_1' \qquad \tau_2 \equiv \tau_2' \qquad \tau_3 \equiv \tau_3'}{(\{\tau_1 \sim \tau_2\} \to \tau_3) \equiv (\{\tau_1' \sim \tau_2'\} \to \tau_3')}\ \textsc{Eq\_CArrow}$$

$$\frac{\tau_1 \equiv \tau_2}{\forall\,(\alpha : \kappa), \tau_1 \equiv \forall\,(\alpha : \kappa), \tau_2}\ \textsc{Eq\_Forall}$$

$$\frac{\tau_1 \equiv \tau_2}{\lambda(\alpha : \kappa), \tau_1 \equiv \lambda(\alpha : \kappa), \tau_2}\ \textsc{Eq\_Abs}$$

$$\frac{\tau_1 \equiv \tau_1' \qquad \tau_2 \equiv \tau_2'}{\tau_1\ \tau_2 \equiv \tau_1'\ \tau_2'}\ \textsc{Eq\_App}$$

$$\frac{[\alpha \mapsto \tau_2]\ \tau_1 \rhd \tau_1'}{(\lambda(\alpha : \kappa), \tau_1)\ \tau_2 \equiv \tau_1'}\ \textsc{Eq\_AppAbs}$$

$\boxed{t \longrightarrow tv}$ Operational semantics

$$\frac{t_2 \longrightarrow tv_2 \qquad t_1\ tv_2 \longrightarrow tv_3}{t_1\ t_2 \longrightarrow tv_3}\ \textsc{E\_App1}$$

$$\frac{t \longrightarrow tv_1 \qquad tv_1\ tv_2 \longrightarrow tv_3}{t\ tv_2 \longrightarrow tv_3}\ \textsc{E\_App2}$$

$$\frac{[x \mapsto tv_1]\ t \rhd t' \qquad t' \longrightarrow tv_2}{(\lambda(x : \tau) \Rightarrow t)\ tv_1 \longrightarrow tv_2}\ \textsc{E\_AppAbs}$$

$$\frac{t \longrightarrow tv}{(\lambda\{\alpha : \kappa\} \Rightarrow t) \longrightarrow (\lambda\{\alpha : \kappa\} \Rightarrow tv)}\ \textsc{E\_TAbs}$$

$$\frac{t \longrightarrow tv}{(\lambda\{c : \tau_1 \sim \tau_2\} \Rightarrow t) \longrightarrow (\lambda\{c : \tau_1 \sim \tau_2\} \Rightarrow tv)}\ \textsc{E\_CAbs}$$

$$\frac{t \longrightarrow tv}{t\ [\tau] \longrightarrow tv\ [\tau]}\ \textsc{E\_TApp}$$

$$\frac{[\alpha \mapsto \tau]\ t \rhd t' \qquad t' \longrightarrow tv}{(\lambda\{\alpha : \kappa\} \Rightarrow t)\ [\tau] \longrightarrow tv}\ \textsc{E\_TAppAbs}$$

$$\frac{t \longrightarrow tv}{(t \sim\![\gamma]) \longrightarrow (tv \sim\![\gamma])}\ \textsc{E\_CApp}$$

$$\frac{\texttt{<<no parses (char 6): [c*** |-> C] @ t |> t' \& t' --> tv >>}}{(\lambda\{c : \tau_1 \sim \tau_2\} \Rightarrow t) \sim\![\gamma] \longrightarrow tv}\ \textsc{E\_CAppAbs}$$

$$\frac{t \longrightarrow tv}{(t : \tau) \longrightarrow (tv : \tau)}\ \textsc{E\_Annot}$$

$$\frac{t \longrightarrow tv}{(t \blacktriangleright \gamma) \longrightarrow (tv \blacktriangleright \gamma)}\ \textsc{E\_Coerce}$$

$\boxed{tv \longrightarrow v}$ type erasure

$$\frac{}{(\lambda(x : \tau) \Rightarrow t) \longrightarrow (\lambda x \Rightarrow t)}\ \textsc{Erase\_Abs}$$

$$\frac{tv \longrightarrow v}{(\lambda\{\alpha : \kappa\} \Rightarrow tv) \longrightarrow v}\ \textsc{Erase\_TAbs}$$

$$\frac{tv \longrightarrow v}{(\lambda\{c : \tau_1 \sim \tau_2\} \Rightarrow tv) \longrightarrow v} \quad \textsc{Erase\_CAbs}$$

$$\frac{tv \longrightarrow v}{(tv\ [\tau]) \longrightarrow v} \quad \textsc{Erase\_TApp}$$

$$\frac{tv \longrightarrow v}{(tv\ \sim[\gamma]) \longrightarrow v} \quad \textsc{Erase\_CApp}$$

$$\frac{tv \longrightarrow v}{(tv : \tau) \longrightarrow v} \quad \textsc{Erase\_Annot}$$

$$\frac{tv \longrightarrow v}{(tv \blacktriangleright \gamma) \longrightarrow v} \quad \textsc{Erase\_Coerce}$$

$\boxed{[\alpha \mapsto \tau_1]\ \tau_2 \rhd \tau_3}$ Type substitution

$$\frac{}{[\alpha \mapsto \tau]\ \alpha \rhd \tau} \quad \textsc{SubstT\_Var1}$$

$$\frac{\alpha_1 \neq \alpha_2}{[\alpha_1 \mapsto \tau]\ \alpha_2 \rhd \alpha_2} \quad \textsc{SubstT\_Var2}$$

$$\frac{}{[\alpha \mapsto \tau]\ T \rhd T} \quad \textsc{SubstT\_Type}$$

$$\frac{[\alpha \mapsto \tau_1]\ \tau_2 \rhd \tau_2' \qquad [\alpha \mapsto \tau_1]\ \tau_3 \rhd \tau_3' \qquad [\alpha \mapsto \tau_1]\ \tau_4 \rhd \tau_4'}{[\alpha \mapsto \tau_1]\ \{\tau_2 \sim \tau_3\} \to \tau_4 \rhd \{\tau_2' \sim \tau_3'\} \to \tau_4'} \quad \textsc{SubstT\_CArrow}$$

$$\frac{\alpha_1 \neq \alpha_2 \qquad [\alpha_1 \mapsto \tau_1]\ \tau_2 \rhd \tau_2'}{[\alpha_1 \mapsto \tau_1]\ (\lambda(\alpha_2 : \kappa), \tau_2) \rhd (\lambda(\alpha_2 : \kappa), \tau_2')} \quad \textsc{SubstT\_Abs}$$

$$\frac{\alpha_1 \neq \alpha_2 \qquad [\alpha_1 \mapsto \tau_1]\ \tau_2 \rhd \tau_2'}{[\alpha_1 \mapsto \tau_1]\ (\forall\,(\alpha_2 : \kappa), \tau_2) \rhd (\forall\,(\alpha_2 : \kappa), \tau_2')} \quad \textsc{SubstT\_Forall}$$

$$\frac{[\alpha \mapsto \tau_1]\ \tau_2 \rhd \tau_2' \qquad [\alpha \mapsto \tau_1]\ \tau_3 \rhd \tau_3'}{[\alpha \mapsto \tau_1]\ \tau_2\,\tau_3 \rhd \tau_2'\,\tau_3'} \quad \textsc{SubstT\_App}$$

$\boxed{[x \mapsto tv]\ t_1 \rhd t_2}$ substitution

$$\frac{}{[x \mapsto tv]\ x \rhd tv} \quad \textsc{Subst\_Var1}$$

$$\frac{x_1 \neq x_2}{[x_1 \mapsto tv]\ x_2 \rhd x_2} \quad \textsc{Subst\_Var2}$$

$$\frac{}{[x \mapsto tv]\ (\lambda(x : \tau) \Rightarrow t) \rhd (\lambda(x : \tau) \Rightarrow t)} \quad \textsc{Subst\_Abs1}$$

$$\frac{x_1 \neq x_2 \qquad [x_1 \mapsto tv]\ t_1 \rhd t_2}{[x_1 \mapsto tv]\ (\lambda(x_2 : \tau) \Rightarrow t_1) \rhd (\lambda(x_2 : \tau) \Rightarrow t_2)} \quad \textsc{Subst\_Abs2}$$

$$\frac{[x \mapsto tv]\ t_1 \rhd t_2}{[x \mapsto tv]\ (\lambda\{\alpha : \kappa\} \Rightarrow t_1) \rhd (\lambda\{\alpha : \kappa\} \Rightarrow t_2)} \quad \textsc{Subst\_TAbs}$$

$$\frac{[x \mapsto tv]\ t_1 \rhd t_2}{[x \mapsto tv]\ (\lambda\{c : \tau_1 \sim \tau_2\} \Rightarrow t_1) \rhd (\lambda\{c : \tau_1 \sim \tau_2\} \Rightarrow t_2)} \quad \textsc{Subst\_CAbs}$$

$$\frac{[x \mapsto tv]\ t_1 \rhd t_1' \qquad [x \mapsto tv]\ t_2 \rhd t_2'}{[x \mapsto tv]\ t_1\,t_2 \rhd t_1'\,t_2'} \quad \textsc{Subst\_App}$$

$$\frac{[x \mapsto tv]\ t_1 \rhd t_2}{[x \mapsto tv]\ (t_1\ [\tau]) \rhd (t_2\ [\tau])} \quad \textsc{Subst\_TApp}$$

$$\frac{[x \mapsto tv]\ t_1 \rhd t_2}{[x \mapsto tv]\ (t_1\ {\sim}[\gamma]) \rhd (t_2\ {\sim}[\gamma])} \quad \textsc{Subst\_CApp}$$

$$\frac{[x \mapsto tv]\ t_1 \rhd t_2}{[x \mapsto tv]\ (t_1 : \tau) \rhd (t_2 : \tau)} \quad \textsc{Subst\_Annot}$$

$$\frac{[x \mapsto tv]\ t_1 \rhd t_2}{[x \mapsto tv]\ (t_1 \blacktriangleright \gamma) \rhd (t_2 \blacktriangleright \gamma)} \quad \textsc{Subst\_Coerce}$$

$\boxed{[\alpha \mapsto \tau]\ t_1 \rhd t_2}$ substitution of type variable in term

$$\frac{}{[\alpha \mapsto \tau]\ x \rhd x} \quad \textsc{TtSubst\_Var}$$

$$\frac{[\alpha \mapsto \tau_1]\ \tau_2 \rhd \tau_2' \qquad [\alpha \mapsto \tau_1]\ t_1 \rhd t_2}{[\alpha \mapsto \tau_1]\ (\lambda(x : \tau_2) \Rightarrow t_1) \rhd (\lambda(x : \tau_2') \Rightarrow t_2)} \quad \textsc{TtSubst\_Abs}$$

$$\frac{\alpha_1 \neq \alpha_2 \qquad [\alpha_1 \mapsto \tau]\ t_1 \rhd t_2}{[\alpha_1 \mapsto \tau]\ (\lambda\{\alpha_2 : \kappa\} \Rightarrow t_1) \rhd (\lambda\{\alpha_2 : \kappa\} \Rightarrow t_2)} \quad \textsc{TtSubst\_TAbs}$$

$$\frac{[\alpha \mapsto \tau_1]\ \tau_2 \rhd \tau_2' \qquad [\alpha \mapsto \tau_1]\ \tau_3 \rhd \tau_3' \qquad [\alpha \mapsto \tau_1]\ t_1 \rhd t_2}{[\alpha \mapsto \tau_1]\ (\lambda\{c : \tau_2 \sim \tau_3\} \Rightarrow t_1) \rhd (\lambda\{c : \tau_2' \sim \tau_3'\} \Rightarrow t_2)} \quad \textsc{TtSubst\_CAbs}$$

$$\frac{[\alpha \mapsto \tau]\ t_1 \rhd t_1' \qquad [\alpha \mapsto \tau]\ t_2 \rhd t_2'}{[\alpha \mapsto \tau]\ t_1\ t_2 \rhd t_1'\ t_2'} \quad \textsc{TtSubst\_App}$$

$$\frac{[\alpha \mapsto \tau_1]\ t_1 \rhd t_2 \qquad [\alpha \mapsto \tau_1]\ \tau_2 \rhd \tau_2'}{[\alpha \mapsto \tau_1]\ (t_1\ [\tau_2]) \rhd (t_2\ [\tau_2'])} \quad \textsc{TtSubst\_TApp}$$

$$\frac{\texttt{<<no parses (char 46): [a |-> T] @ t1 |> t2 \& [a |-> T] @ C***1 |> C2 >>}}{[\alpha \mapsto \tau]\ (t_1\ {\sim}[\gamma_1]) \rhd (t_2\ {\sim}[\gamma_2])} \quad \textsc{TtSubst\_CApp}$$

$$\frac{[\alpha \mapsto \tau_1]\ t_1 \rhd t_2 \qquad [\alpha \mapsto \tau_1]\ \tau_2 \rhd \tau_2'}{[\alpha \mapsto \tau_1]\ (t_1 : \tau_2) \rhd (t_2 : \tau_2')} \quad \textsc{TtSubst\_Annot}$$

$$\frac{\texttt{<<no parses (char 46): [a |-> T] @ t1 |> t2 \& [a |-> T] @ C***1 |> C2 >>}}{[\alpha \mapsto \tau]\ (t_1 \blacktriangleright \gamma_1) \rhd (t_2 \blacktriangleright \gamma_2)} \quad \textsc{TtSubst\_Coerce}$$

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Definition rules:        80 good    3 bad
Definition rule clauses: 153 good    3 bad
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