$\begin{array}{ll} termvar, \ x & \text{term variable} \\ variant, \ V & \text{variant} \\ typvar, \ X & \text{type variable} \\ exc, \ Exc & \text{exception} \\ effect, \ E & \text{effect} \\ n & \\ m & \end{array}$

```
program
                 ::=
                        top\ program
                        [x \mapsto t] program
                                                                                        Μ
top
                        \mathbf{let} \, x = t
                        \mathbf{let}\,\mathbf{rec}\,x:\,T=abs
                        type alias X = T
                        \mathbf{type} \ X \ variantArgs = \ V_1 \ tyList_1 | .. | V_n \ tyList_n
                        exception Exc \ tyList
                                                                                              _{\rm term}
                 ::=
                                                                                                  variable
                        \boldsymbol{x}
                         V
                                                                                                  type constructors
                        \lambda(x:T) \to t
                                                                                                  abstraction
                        \lambda(X:K) \to t
                                                                                                  type abstraction
                        \lambda(E:\varphi) \to t
                                                                                                 effect abstraction
                        t t'
                                                                                                  application
                        t[T]
                                                                                                  type application
                        t[[eff]]
                                                                                                 effect application
                        \mathbf{let}\ x = t_1 \, \mathbf{in}\ t_2
                                                                                                 let binding
                        \mathbf{let}\,\mathbf{rec}\,x:\,T=abs\,\mathbf{in}\,t
                                                                                                 recursive let binding
                        match t with p_1 \rightarrow t_1 | ... | p_n \rightarrow t_n end
                                                                                                 pattern matching
                        t : [eff]T
                                                                                                 type annotation
                                                                                        S
                        t:T
                                                                                                  type annotation = (t : [] T)
                        fail [T]Exc t_1 ... t_n
                        try t with pe_1 \rightarrow t_1 | ... | pe_n \rightarrow t_n end
                                                                                                  try
                        t;t'
                                                                                        S
                                                                                                 sequence == let (_{\cdot}: Unit) = t in t'
                                                                                        S
                        (t)
                                                                                        Μ
                        [x_1 \mapsto t_1, \dots, x_n \mapsto t_n]t
                        failure exnval
                                                                                        Μ
                                                                                        Μ
                        TConstr V v_1 ... v_n
                                                                                              value
                 ::=
v
                        TConstr V v_1 ... v_n
                                                                                                  type constructors
                        \lambda(x:T) \to t
                                                                                                 abstraction
                                                                                              lambda abstractions
abs
                 ::=
                        \lambda(x:T) \to t
                                                                                                 abstraction
                        \lambda(X:K) \to abs
                                                                                                  type abstraction
                        \lambda(E:\varphi) \to abs
                                                                                                 effect abstraction
                        abs: [eff] T
                                                                                                 type annotation
                                                                                        S
                        abs: T
                                                                                                  type annotation = (abs : [] T)
                                                                                              pattern
p
                         V p_1 \dots p_n
                                                                                                  variant
```

```
wildcard variable
                     \boldsymbol{x}
                                                                    try pattern
pe
              ::=
                                                                       Exception pattern
                     Exc x_1 \dots x_n
                                                                    runtime value of exceptions
exnval
                     Exc\ v_1 \dots v_n
effelm
                                                                    effects elements
                     E
                                                                       effect
                     \mathbf{IO}
                                                                       IO effect
                     \mathbf{Exn}[exn]
                                                                       exception
                                                                    exceptions
exn
              ::=
               Exc_1 | ... | Exc_n
                                                                    effect
eff
                     effelm_1, ..., effelm_n
                     eff_1 \cup eff_2 \cup ... \cup eff_n
                                                              Μ
                     eff_1 \setminus [exn]
                                                              Μ
                                                              S
                     (eff)
K
              ::=
                                                                    kinds
                                                                       \operatorname{star}
                     K \to K'
                                                                       kind arrow
T
              ::=
                                                                    type
                     X
                                                                       variable
                     Unit
                                                                       Unit type (contained in the module opened by defau
                     T \to \, T'
                                                              S
                                                                       pure function == T - [] - i T'
                     T - [eff] - > T'
                                                                       function
                     \lambda(X:K), T
                                                                       operator abstraction
                     \forall (X:K), T
                                                                       forall
                     \forall (E:\varphi), T
                                                                       {\it effect\ for all}
                     T T'
                                                                       operator application
                     (T)
                                                              S
                     [X \mapsto T]T'
                                                              Μ
                     [E \mapsto \mathit{eff}]T
                                                              Μ
                     T_1 \to \dots \to T_n \to T
                                                              Μ
                     T\ T_1 \dots T_n
                                                              Μ
Γ
                                                                    type environment
                     empty
                                                                       empty
                    \Gamma, x_1: T_1, \ldots, x_n: T_n
                     \Gamma, V : T
                                                              S
                                                                       type constructors (contained in the above values env
                    \Gamma, X_1: K_1, \ldots, X_n: K_n
                                                                       tvars
```

variants

 $\Gamma, X : \{ V_1 \ tyList_1 ... V_n \ tyList_n \}$

```
\Gamma, Exc tyList
                                                                                      exceptions
                                  \Gamma, E
                                                                                      effects
                                  \Gamma_1 \cup ... \cup \Gamma_n
                                                                           Μ
tyList
                          ::=
                                   T_1 \dots T_n
variantArgs
                          ::=
                                  (X_1:K_1)..(X_n:K_n)
terminals
                                  \triangleright
                                   &
formula
                                  judgement
                                  formula_1 .. formula_n
                                  not (formula)
                                   T\equiv\,T'
                                  x:\,T\,\in\,\Gamma
                                   X:K\in\Gamma
                                   V:T\,\in\,\Gamma
                                   T: \{Variant\} \in \Gamma
                                   E \in \Gamma
                                  \mathit{Exc}\;\mathit{tyList}\;\in\;\Gamma
                                  \mathbf{set}\left(\mathit{eff}\right) = \mathbf{set}\left(\mathit{eff}'\right)
                                  \mathbf{set}(exn) = \mathbf{set}(exn')
                                   V \in Variant \triangleright T_1 ... T_n
                                   RetVar = RetVar'
                                  \Gamma = \Gamma'
substs
                          ::=
                                  x_1 \leftarrow v_1, \dots, x_n \leftarrow v_n
                                   substs_1 \cup ... \cup substs_n
                                                                           Μ
```



```
JEffEquiv
                     ::=
                            eff \equiv eff'
                                                                                      Effects equivalence
J\!E\!f\!f\!ElmEquiv
                     ::=
                            effelm \equiv effelm'
                                                                                      Effect element equivalence
JTopOp
                     ::=
                                                                                      Toplevel evaluation
                            program \longrightarrow program'
Jop
                     ::=
                            t \longrightarrow t'
                                                                                      Evaluation
JExnMatches
                     ::=
                            exnval matches pe \rhd \{x_1 \leftarrow v_1, ..., x_n \leftarrow v_n\}
                                                                                      Exception pattern matching with
JMatches
                     ::=
                            v  matches p \rhd \{substs\}
                                                                                      Pattern matching with substitutio
judgement
                     ::=
                            JProgram
                            JTop\,Type
                            JTypeDecl
                            Jtype
                            Jkind
                            JEff
                            JEffElm
                            JP atterns \, Typing
                            JPattern\,Typing
                            {\it JExnPattern\,Typing}
                            Jequiv
                            J\!E\!f\!f\!Equiv
                            JEffElmEquiv
                            JTopOp
                            Jop
                            JExnMatches
                            JMatches
user\_syntax
                     ::=
                            term var
                            variant
                            typvar
                            exc
                            e\!f\!f\!ect
                            n
                            m
```

 $\begin{array}{c} program \\ top \end{array}$

```
abs
                    p
                    pe
                    exnval
                    effelm
                     exn
                    eff
                     K
                     T
                    Γ
                    tyList
                    variantArgs
                    terminals
                    formula
                     substs
                     patterns
                     VArgs
                     Variant
                     RetVar
  \Gamma \vdash program \rhd \Gamma'
                                               Program typing
                                                                                \frac{}{\Gamma \vdash \rhd \Gamma} Prog_Empty
                                                                 \Gamma \vdash top \rhd \Gamma'
                                                                 \Gamma' \vdash \overrightarrow{program} \rhd \Gamma''
                                                                                                               Prog_Program
                                                              \frac{\Gamma \vdash top \ program \rhd \Gamma''}{\Gamma \vdash top \ program \rhd \Gamma''}
     \Gamma \vdash top \rhd \Gamma'
                                       Toplevel typing
                                                                    \frac{\Gamma \vdash t : [\,]\,T}{\Gamma \vdash \mathbf{let}\,x = t \rhd \Gamma, x : T} \quad \text{Top\_Let}
                                                    \frac{\Gamma, x: T \vdash abs: [\,]\,T}{\Gamma \vdash \mathbf{let}\,\mathbf{rec}\,x: T = abs \rhd \Gamma, x: T} \quad \mathsf{TOP\_LETREC}
                                                \frac{\Gamma \vdash T : K}{\Gamma \vdash \mathbf{type} \, \mathbf{alias} \, X = T \rhd \Gamma, X : K} \quad \text{Top\_TypeAlias}
                                            Ret Var = X X_1 ... X_n
                                            \Gamma' = \Gamma, X : K, X_1 : K_1, ..., X_n : K_n
RetVar \& \Gamma' \vdash V_1 \ tyList_1 \rhd \Gamma_1 \quad ... \quad RetVar \& \Gamma' \vdash V_n \ tyList_n \rhd \Gamma_n
\Gamma \vdash \mathbf{type} \ X \ (X_1 : K_1)(X_n : K_n) = V_1 \ tyList_1 | ... | V_n \ tyList_n \rhd \Gamma \cup \Gamma_1 \cup ... \cup \Gamma_n, X : \{ V_1 \ tyList_1 ... \ V_n \ tyList_n \}
                                     \frac{\Gamma \vdash T_1 : K_1 \quad .. \quad \Gamma \vdash T_n : K_n}{\Gamma \vdash \mathbf{exception} \ Exc \ T_1 \ .. \ T_n \rhd \Gamma, Exc \ T_1 \ .. \ T_n} \quad \text{Top\_Exception}
     Ret Var \& \Gamma \vdash V \ tyList \rhd \Gamma'
                                                                        Type declaration
         \frac{\Gamma \vdash T_1 : K_1 \quad .. \quad \Gamma \vdash T_n : K_n}{X \: X_1 \: .. \: X_n \: \& \: \Gamma \vdash V \: T_1 \: .. \: T_n \rhd \mathbf{empty}, \: V : \: T_1 \to \: .. \: \to \: T_n \to X \: X_1 \: .. \: X_n}
                                                                                                                                                                       TyDecl_Decl
     \Gamma \vdash t : [\mathit{eff}] T
                                     Typing
```

t

```
\frac{x:T\in\Gamma}{\Gamma\vdash x:[\,]T}\quad \text{$\mathcal{T}$-Var}
                                                                         \frac{V:T\in\Gamma}{\Gamma\vdash V:[\,]T}\quad \text{$\Tau$\_VARIANT}
                                                                     \Gamma, x_1 : T_1 \vdash t : [eff] T
                                                                    \Gamma \vdash T_1 : *
                                             \frac{1}{\Gamma \vdash \lambda(x_1:T_1) \rightarrow t:[]T_1 - [eff] - > T}
                                                                                                                                                          T_ABS
                                                         \Gamma \vdash t : [eff_1]T_1 - [eff_2] - > T_2
                                                       \frac{\Gamma \vdash t' : [\textit{eff}_3] T_1}{\Gamma \vdash t \ t' : [\textit{eff}_1 \cup \textit{eff}_2 \cup \textit{eff}_3] T_2}
                                             \frac{\Gamma, X : K \vdash t : [\mathit{eff}] \, T}{\Gamma \vdash \lambda(X : K) \rightarrow t : [\mathit{eff}] \forall (X : K), \, T}
                                                                                                                                                      T_{-}TABS
                                                \frac{\Gamma, E \vdash t : [\mathit{eff}] \, T}{\Gamma \vdash \lambda(E : \varphi) \rightarrow t : [\mathit{eff}] \forall \, (E : \varphi), \, T}
                                                                                                                                                    T_EABS
                                                              \Gamma \vdash t : [eff] \forall (X : K), T_2
                                                             \Gamma \vdash T_1 : K
                                                                                                                                          T_TAPP
                                                         \overline{\Gamma \vdash t[T_1] : [\mathit{eff}][X \mapsto T_1]T_2}
                                                               \Gamma \vdash t : [eff] \forall (E : \varphi), T
                                                                \Gamma \vdash eff'
                                                                                                                                            T_EAPP
                                                      \overline{\Gamma \vdash t[[\mathit{eff'}]] : [\mathit{eff}][E \mapsto \mathit{eff'}]T}
                                                                               \Gamma \vdash t : [eff]XX \equiv X'
                                                                             \frac{\Gamma \vdash X' : *}{\Gamma \vdash t : [eff]X'} \quad \text{T-EQ}
                                                                   \Gamma, x: T_1 \vdash t_2: [eff_2] T_2
                                   \overline{\Gamma \vdash \mathbf{let} \ x = (t_1 : [\mathit{eff}_1] \ T_1) \ \mathbf{in} \ t_2 : [\mathit{eff}_1 \ \cup \ \mathit{eff}_2] \ T_2}
                                                                                                                                                                     T_{-}Let
                                                              \Gamma, x: T_1 \vdash abs: []T_1
                                        \frac{\Gamma,x:T_1\vdash t_2:[\mathit{eff}]T_2}{\Gamma\vdash \mathbf{let}\,\mathbf{rec}\,x:T_1=\mathit{abs}\,\mathbf{in}\,t_2:[\mathit{eff}]T_2}
                                                                                                                                               T_LETREC
                           \Gamma \cup \Gamma_1 \vdash t_1 : [\mathit{eff}_1] T_2 \quad .. \quad \Gamma \cup \Gamma_n \vdash t_n : [\mathit{eff}_n] T_2
                            T_1: \{Variant\} \in \Gamma
                            Variant & \Gamma \vdash p_1 ... p_n : T_1 \rhd \Gamma_1 ... \Gamma_n
                           \Gamma \vdash t : [eff] T_1
                                                                                                                                                                                                Т_Матсн
\frac{\Gamma \vdash \mathbf{match} \ t \ \mathbf{with} \ p_1 \to t_1 | ... | p_n \to t_n \ \mathbf{end} : [\mathit{eff} \ \cup \ \mathit{eff}_1 \ \cup \ ... \ \cup \ \mathit{eff}_n] \ T_2}{\Gamma \vdash \mathbf{match} \ t \ \mathbf{with} \ p_1 \to t_1 | ... | p_n \to t_n \ \mathbf{end} : [\mathit{eff} \ \cup \ \mathit{eff}_1 \ \cup \ ... \ \cup \ \mathit{eff}_n] \ T_2}
                                                                           \Gamma \vdash t : [\mathit{eff}] T
                                                                           \Gamma \vdash T : *
                                                                          \Gamma \vdash \mathit{eff}
                                                             \frac{\tilde{\Gamma} \vdash (t : [\mathit{eff}]\,T) : [\mathit{eff}]\,T}{\Gamma} \quad \text{$\Gamma$\_Annot}
                                               \Gamma \vdash t_1 : [eff_1] T_1 \quad \dots \quad \Gamma \vdash t_n : [eff_n] T_n
                                               \Gamma \vdash T : *
                                               Exc\ T_1 ... T_n \in \Gamma
                                                                                                                                                                               T_FAIL
                      \overline{\Gamma \vdash \mathbf{fail} \left[ \mathit{T} \right] \mathit{Exc} \ t_1 \ldots t_n : \left[ \mathbf{Exn} \left[ \mathit{Exc} \right] \ \cup \ \mathit{eff}_1 \ \cup \ \ldots \ \cup \ \mathit{eff}_n \right] \mathit{T}}
```

$$\Gamma_1 \vdash t_1 : [eff_1] T \dots \Gamma_n \vdash t_n : [eff_n] T$$

$$\Gamma \vdash pe_1 \rhd Exc_1 \& \Gamma_1 \dots \Gamma \vdash pe_n \rhd Exc_n \& \Gamma_n$$

$$\Gamma \vdash t : [eff] T$$

 $T_{-}T_{RY}$

 $\overline{\Gamma \vdash \mathbf{try} \, t \, \mathbf{with} \, pe_1 \to t_1 | ... | pe_n \to t_n \, \mathbf{end} : [(eff \setminus [Exc_1| ... | Exc_n]) \cup eff_1 \cup ... \cup eff_n] \, T}$

 $\Gamma \vdash T : K$ Kinding

$$\frac{X:K\in\Gamma}{\Gamma\vdash X:K}\quad \text{K-TVAR}$$

$$\frac{\Gamma,X:K_1\vdash T:K_2}{\Gamma\vdash\lambda(X:K_1),T:K_1\to K_2}\quad \text{K_-Abs}$$

$$\frac{\Gamma\vdash T_1:K_{11}\to K_{12}}{\Gamma\vdash T_2:K_{11}}\quad \text{K_-App}$$

$$\frac{\Gamma\vdash T_1:*}{\Gamma\vdash T_1T_2:K_{12}}\quad \text{K_-App}$$

$$\frac{\Gamma\vdash T_1:*}{\Gamma\vdash eff}$$

$$\frac{\Gamma\vdash T_1:*}{\Gamma\vdash T_1-[eff]->T_2:*}\quad \text{K_-Arrow}$$

$$\frac{\Gamma,X:K_1\vdash T_2:*}{\Gamma\vdash\forall(X:K_1),T_2:*}\quad \text{K_-All}$$

$$\frac{\Gamma,E\vdash T:*}{\Gamma\vdash\forall(E:\varphi),T:*}\quad \text{K_-EAll}$$

 $\Gamma \vdash eff$ Effects typing

$$\frac{\Gamma \vdash \textit{effelm}_1 \quad .. \quad \Gamma \vdash \textit{effelm}_n}{\Gamma \vdash \textit{effelm}_1, \dots, \textit{effelm}_n} \quad \text{Eff_Eff}$$

 $\Gamma \vdash effelm$ Effects elements typing

$$\frac{E \in \Gamma}{\Gamma \vdash E} \quad \text{EffElm_Eff}$$

$$\frac{F}{\Gamma \vdash IO} \quad \text{EffElm_IO}$$

$$\frac{Exc_1 \ tyList_1 \in \Gamma \quad .. \quad Exc_n \ tyList_n \in \Gamma}{\Gamma \vdash \mathbf{Exn} \left[Exc_1 \right] .. \left[Exc_n \right]} \quad \text{EffElm_Exn}$$

Variant & $\Gamma \vdash patterns : T \rhd \Gamma_1 ... \Gamma_n$ Patterns matching typing

Variant & $\Gamma \vdash p : T \rhd \Gamma'$ Pattern matching typing

 $V \in Variant \rhd T_1 ... T_n$ $T_1 : \{Variant_1\} \in \Gamma ... T_n : \{Variant_n\} \in \Gamma$

 $\frac{Variant_1 \& \Gamma \vdash p_1 : T_1 \rhd \Gamma_1 ... Variant_n \& \Gamma \vdash p_n : T_n \rhd \Gamma_n}{Variant \& \Gamma \vdash V p_1 ... p_n : T \rhd \Gamma_1 \cup ... \cup \Gamma_n} \quad \text{PTY_VARIANT}$

$$\overline{\textit{Variant \& } \Gamma \vdash x : T \rhd \mathbf{empty}, x : T} \quad \text{PTY_WILDCARD}$$

 $\Gamma \vdash pe \rhd Exc \& \Gamma'$ Exception pattern matching typing

$$\frac{Exc \ T_1 \dots T_n \in \Gamma}{\Gamma \vdash Exc \ x_1 \dots x_n \rhd Exc \ \& \ \Gamma, x_1 : T_1, \dots, x_n : T_n} \quad \text{PETY_Exc}$$

 $T \equiv T'$ Type equivalence

$$\overline{T \equiv T} \quad \text{Q-Refl}$$

$$\overline{T \equiv T'} \quad \text{Q-Symm}$$

$$T_1 \equiv T_2$$

$$T_2 \equiv T_3$$

$$\overline{T_1 \equiv T_2} \quad \text{Q-Trans}$$

$$T_{11} \equiv T_{21}$$

$$eff_1 \equiv eff_2$$

$$T_{12} \equiv T_{22}$$

$$T_{11} - [eff_1] - > T_{12} \equiv T_{21} - [eff_2] - > T_{22}$$

$$Q-Arrow$$

$$T_1 \equiv T_2$$

$$\forall (X:K), T_1 \equiv \forall (X:K), T_2 \quad Q-All$$

$$T_1 \equiv T_2$$

$$\forall (E:\varphi), T_1 \equiv \forall (E:\varphi), T_2 \quad Q-Ball$$

$$T_1 \equiv T_2$$

$$\forall (X:K), T_1 \equiv \lambda(X:K), T_2 \quad Q-Abs$$

$$T_{11} \equiv T_2$$

$$\lambda(X:K), T_1 \equiv \lambda(X:K), T_2 \quad Q-Abs$$

$$T_{11} \equiv T_{21}$$

$$T_{12} \equiv T_{22}$$

$$T_{11} T_{12} \equiv T_{21} T_{22} \quad Q-App$$

$$(\lambda(X:K), T_{11}) T_{12} \equiv [X \mapsto T_{12}] T_{11} \quad Q-AppAbs$$

 $eff \equiv eff'$ Effects equivalence

$$\frac{eff \equiv eff}{eff \equiv eff} \quad \text{EffEQ_Refl}$$

$$\frac{\text{set } (effelm_1, ..., effelm_n) = \text{set } (effelm'_1, ..., effelm'_n)}{effelm_1, ..., effelm_n \equiv effelm'_1, ..., effelm'_n} \quad \text{EffEQ_EQ}$$

 $effelm \equiv effelm'$ Effect element equivalence

 $program \longrightarrow program'$ Toplevel evaluation

```
\frac{\textbf{not}\left(v \ \textbf{matches} \ p_1 \rhd \{x_1 \leftarrow v_1, \, ..\,, x_n \leftarrow v_n\}\right)}{\textbf{match} \ v \ \textbf{with} \ p_1 \rightarrow t_1 | p_2 \rightarrow t_2 | \, .. \, | p_n \rightarrow t_n \ \textbf{end} \longrightarrow \textbf{match} \ v \ \textbf{with} \ p_2 \rightarrow t_2 | \, .. \, | p_n \rightarrow t_n \ \textbf{end}}
                                                                                                                                                                                                                                      E_MATCHSTEP
                                                                                     \frac{t \longrightarrow t'}{(t : [eff]T) \longrightarrow t'} \quad \text{E-Annot}
                            \frac{t \longrightarrow t'}{\mathbf{fail}\,[T] \mathit{Exc}\,v_1 \ldots v_n\;t\;t_1 \ldots t_n \longrightarrow \mathbf{fail}\,[T] \mathit{Exc}\,v_1 \ldots v_n\;t'\;t_1 \ldots t_n} \quad \text{E_FAILUREARGS}
                                                      \frac{}{\mathbf{fail} [T] Exc \ v_1 \dots v_n \longrightarrow \mathbf{failure} \ Exc \ v_1 \dots v_n} \quad \text{E_FAILURE}
             \frac{t \longrightarrow t'}{\mathbf{try} \ t \ \mathbf{with} \ pe_1 \rightarrow t_1 | ... | pe_n \rightarrow t_n \ \mathbf{end} \longrightarrow \mathbf{try} \ t' \ \mathbf{with} \ pe_1 \rightarrow t_1 | ... | pe_n \rightarrow t_n \ \mathbf{end}}
                                                                                                                                                                                                                       E_{-}T_{RY}
                                          \frac{1}{\mathbf{try} \ v \ \mathbf{with} \ pe_1 \to t_1 | ... | pe_n \to t_n \ \mathbf{end} \longrightarrow v} E_TRYNoFailure
                          \frac{\textbf{not} \left( \textit{exnval} \, \textbf{matches} \, pe_1 \rhd \left\{ x_1 \leftarrow v_1, \, ..., x_n \leftarrow v_n \right\} \right)}{\textbf{try failure} \, \textit{exnval} \, \textbf{with} \, pe_1 \rightarrow \textit{t}_1 \, \textbf{end} \longrightarrow \textbf{failure} \, \textit{exnval}}
                                                                                                                                                                          E_TryNotFound
     \frac{exnval \ \mathbf{matches} \ pe_1 \rhd \{x_1 \leftarrow v_1, \dots, x_n \leftarrow v_n\}}{\mathbf{try \ failure} \ exnval \ \mathbf{with} \ pe_1 \rightarrow t_1 | \dots | pe_n \rightarrow t_n \ \mathbf{end} \longrightarrow [x_1 \mapsto v_1, \dots, x_n \mapsto v_n] t_1} \quad \text{E-TryFound}
\frac{\textbf{not} \ (\textit{exnval} \ \textbf{matches} \ pe_1 \rhd \{x_1 \leftarrow v_1, \, ..\,, x_n \leftarrow v_n\})}{\textbf{try failure} \ \textit{exnval} \ \textbf{with} \ pe_1 \rightarrow t_1 | pe_2 \rightarrow t_2 | \, .. \, | pe_n \rightarrow t_n \ \textbf{end} \longrightarrow \textbf{try failure} \ \textit{exnval} \ \textbf{with} \ pe_2 \rightarrow t_2 | \, .. \, | pe_n \rightarrow t_n \ \textbf{end}}
    exnval matches pe > \{x_1 \leftarrow v_1, ..., x_n \leftarrow v_n\} Exception pattern matching with substitution creation
                                                                                                                                                                  EXNMATCHES_MATCHES
                 \overline{Exc\ v_1 \dots v_n\ \mathbf{matches}\ Exc\ x_1 \dots x_n \rhd \{x_1 \leftarrow v_1, \dots, x_n \leftarrow v_n\}}
      v  matches p \rhd \{substs\}
                                                                          Pattern matching with substitution creation
                                                                      \overline{v \text{ matches } x \rhd \{x \leftarrow v\}} Matches_Any
            \frac{v_1 \text{ matches } p_1 \rhd \{substs_1\} \quad .. \quad v_n \text{ matches } p_n \rhd \{substs_n\}}{\text{TConstr } V \ v_1 \ldots v_n \text{ matches } V \ p_1 \ldots p_n \rhd \{substs_1 \cup \ldots \cup substs_n\}} \quad \text{MATCHES\_MATCHES}
Definition rules:
                                                                           86 good
                                                                                                             0 bad
Definition rule clauses: 174 good
```