# **An Accuracy Argument for Self-Trust**

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**Self-Doubt and Self-Trust** 

#### Rational Self-Doubt

It seems rational to doubt the accuracy of my own beliefs.

- Plenty of evidence that I have been wrong, and that my peers are wrong.
- Preface-like cases: I'm confident that some of my beliefs about biology are *false* (e.g. "Mammals don't lay eggs").
- Cartesian Circle: No non-circular way to rule out the possibility that our beliefs are thoroughly inaccurate.





#### Irrational Self-Doubt

Some cases of extreme self-doubt seem irrational.

E.g. believing a (commissive) Moorean sentence: "It's raining, but I believe it's not raining".



#### Questions

- Why are certain kinds of self-doubt irrational?
- How much may we rationally doubt ourselves?
- What about graded doxastic states?
  - Being very confident in "It's raining, and I'm very confident that it's not raining" seems nearly as bad as believing a Moorean sentence.

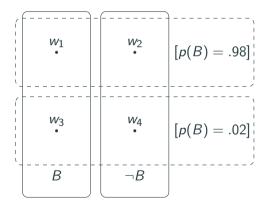
**Goal**: Use accuracy to answer these questions.

#### **Notation**

- $W = \{w_1, ..., w_n\}$  finite set of *possible worlds*.
- Greek letters  $\pi, \gamma$  denote *rigidly designated credence functions*, i.e. vectors in  $\mathbb{R}^n$ .
- Latin letters p, q denote definite descriptions of credence functions.
  - p is a function from possible worlds to credence functions. So  $p(w_i)$  is a credence function for every  $w_i \in \mathcal{W}$ .
  - Can think of them as vector-valued random variables.
  - Abuse notation:  $p_i$  instead of  $p(w_i)$ .
- If  $\phi$  is a property of credence functions,  $[\phi(p)]$  is the proposition  $\{w_i : \phi(p_i)\}$

# **E**xample

- p = My radiologist's credence function.
- B = I have a broken bone.



$$p_1 = p_2 = (.97, .01, .01, .01)$$
  $p_3 = p_4 = (.01, .01, .01, .97)$ 

#### A self-trust requirement

- Let *p* be a definite description of your credence function.
- Let  $\pi$  be your actual credence function (i.e.  $\pi = p_i$  where  $w_i$  is the actual world)

#### **Total Trust**

 $\pi$  Totally Trusts p iff:

$$\mathbb{E}_{\pi}(X|[\mathbb{E}_{\rho}(X) \ge r]) \ge r \tag{1}$$

whenever  $X : \mathcal{W} \to \mathbb{R}$ ,  $r \in \mathbb{R}$ , and the above conditional expectation is defined.

Coherence + Total Trust entails that you cannot be very highly confident of both A and  $[p(A) \le low]$ .

**An Accuracy Argument for Total** 

**Trust** 

### **Measuring Accuracy**

Generalised Strictly Proper (GSP) measures of accuracy can be used to measure the accuracy of a probability function  $\pi$ .

#### An Accuracy Argument for Total Trust

#### Theorem (Dorst et al. 2021, Th.3.2)

 $\pi$  Totally Trusts p iff for every GSP measure of inaccuracy,  $\pi$  expects p to be at least as accurate as  $\pi$ .

- Suppose I don't Totally Trust myself, i.e.  $\pi$  does not Totally Trust p.
- Then there is a rigidly designated credence function  $\pi$  (e.g. (1/3, 1/3, 1/3)) that I think is more accurate than me under some GSP measure S.
- I expect that I would be more accurate, as measured by S, by having credence function  $\pi$  at all possible worlds!

**Problem**: Why care about that measure? Under most reasonable (GSP) measures, I may expect p to be more accurate than  $\pi$ .

# Improving the Argument

### From credences to desirability judgements

Represent opinions via sets of desirable gambles for more expressive power.

**Subtle Point**: We want to show that **rational** agents Totally Trust themselves.

- Rational agents have *coherent* and (for this talk) *precise* doxastic states.
- So we need to show that agents with coherent, precise doxastic states should trust themselves.
- Your beliefs are still representable by a coherent credence function  $\pi$ .
- Added expressive power lets us consider ways your beliefs could be that don't correspond to any coherent credence function.

# From credences to desirability judgements

For any probability function  $\pi$ ,  $D_{\pi} = \{X : \mathbb{E}_{\pi}(X) > 0\}$  is the set of gambles an agent with credence function  $\pi$  finds desirable.

We can express Total Trust in desirability terms.

#### **Total Trust**

 $\pi$  Totally Trusts p iff:

$$X \in D_{\pi(\cdot|[X \in D_p])} \tag{2}$$

whenever  $X: \mathcal{W} \to \mathbb{R}$ ,  $r \in \mathbb{R}$ , and the above conditional expectation is defined.

# From GSP to INSERT NAME HERE measures of inaccuracy

We can use [INSERT NAME HERE] to measure the inaccuracy of an *arbitrary set of desirable gambles* at a world.

- For every *X* which you find desirable, you get a penalty if *X* is not actually desirable.
- For every X which you don't find desirable, you get a penalty if X is actually desirable.

$$S(D, w_i) = \int_{D_{w_i} \sim D} x_i d\mu - \int_{D \sim D_{w_i}} x_i d\mu$$
 (3)

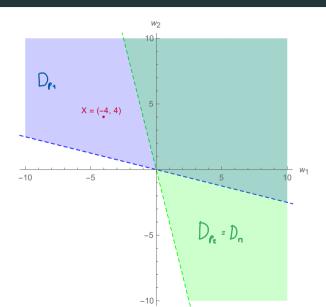
#### A useful fact

#### Fact 1

If Total Trust fails on some gamble, then it fails on some open set of gambles.

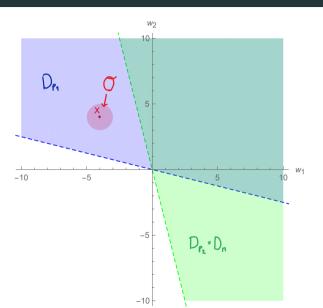
# Example

- $W = \{w_1, w_2\}$
- $w_2$  is the actual world.
- X = (-4, 4)



# Example

- $[X \in D_p] = \{w_1\}.$
- But  $X(w_1) = -4$ .
- So  $X \notin D_{\pi(\cdot|[X \in D_p])}$ , violating Total Trust.
- Similarly for nearby gambles.



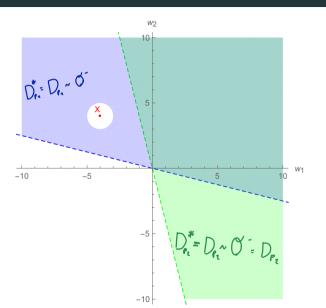
# New accuracy characterisation of Total Trust

- Suppose  $\pi$  does not Totally Trust p.
- Then there is some open set  $\mathcal{O}$  of gambles where Total Trust fails.
- Define:

$$\mathcal{O}^+ = \mathcal{O} \cap D_{\pi}, \quad \mathcal{O}^- = \mathcal{O} \cap D_{\pi}^c$$
 $D_p^* = (D_p \cup \mathcal{O}^+) \sim \mathcal{O}^-$ 

- You actually find the gambles in  $\mathcal{O}^+$  desirable, and those in  $\mathcal{O}^-$  not desirable.
- $D_p^*$  represents the opinions you would have if, at every possible world, you found the gambles in  $\mathcal{O}^+$  desirable and those in  $\mathcal{O}^-$  not desirable.

# **E**xample



# New accuracy characterisation of Total Trust

- **Note**: At some possible worlds,  $D_p^*$  denotes an **incoherent** set of desirable gambles!
- But with INSERT NAME HERE we can measure its inaccuracy at all possible worlds!

#### **Theorem**

- 1. If  $\pi$  does not Totally Trust p, then there are measurable sets of gambles  $\mathcal{O}^+, \mathcal{O}^-$  such that  $\pi$  expects  $D_p^* = (D_p \cup \mathcal{O}^+) \sim \mathcal{O}^-$  to be strictly more accurate than  $D_p$  under every INSERT NAME HERE measure of inaccuracy.
- 2. If  $\pi$  Totally Trusts p, then for any measurable sets of gambles  $\mathcal{O}^+$  and  $\mathcal{O}^-$ ,  $\pi$  expects  $D_p$  to be at least as accurate as  $D_p^*$  under every INSERT NAME HERE measure of inaccuracy.

#### The New Argument

- Suppose  $\pi$  does not Totally Trust p.
- Then there are (rigidly designated!) set of gambles  $\mathcal{O}^+, \mathcal{O}^-$  such that you think you would be more accurate, under **every** INSERT NAME HERE measure of inaccuracy, if you found gambles in  $\mathcal{O}^+$  desirable and gambles in  $\mathcal{O}^-$  not desirable at every possible world.
- There is a way to change your judgements which you expect would make you more accurate, no matter which measure of accuracy you use!

#### Many open questions...

- Is it really bad to expect some *possibly incoherent* definite description to be more accurate than you?
- How do we determine which doxastic states we should compare yours against when evaluating you?
- Self-trust requirements for imprecise probabilities?