Quantum Hackenbush

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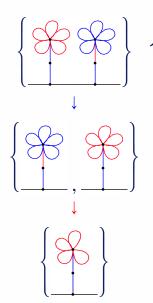


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Introduction

QUANTUM HACKENBUSH

- "HACKENBUSH?"
- QUANTUM(-inspired) HACKENBUSH
- Results
- Future research



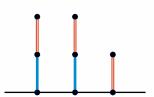
Combinatorial game, just like tic-tac-toe or chess.



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 - Two-player
 - Deterministic
 - Perfect-information
 - Short (finite and loopfree)



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 - Two-player
 - Deterministic
 - Perfect-information
 - Short (finite and loopfree)
- Played on graph with blue and red edges.





Combinatorial games

Definition

A *short game* is an ordered pair

$$G = \left\{ G_1^L, \ldots, G_m^L \mid G_1^R, \ldots, G_n^R \right\}$$

of Left options G^L and Right options G^R .

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of Left options G^L and Right options G^R .

$$G = \left\{ \begin{array}{c|c} & & & \\ & & & \\ \end{array} \right\}$$

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Given game state (position) G, assuming both players play optimally:

What do we want to study?

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• Winnability: which player is able to win?

	$\exists \mathbf{G}^R \in \mathscr{P} \cup \mathscr{R}$	$\forall \mathbf{G}^{R} \in \mathscr{N} \cup \mathscr{L}$
$\exists \mathbf{G}^{L} \in \mathscr{P} \cup \mathscr{L}$	$G \in \mathscr{N}$	$oldsymbol{G} \in \mathscr{L}$
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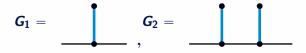
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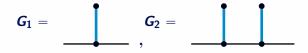
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• Advantage: by how much does a player win?

$$G_1 = G_2 =$$



We know $G_1, G_2 \in \mathcal{L}$.



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But: $G_1 = 1$ and $G_2 = 2$.

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Alternating turns: $G \rightarrow G^L \rightarrow G^{LR} \rightarrow G^{LRL} \rightarrow \cdots$ Why give Left two turns in a row?

Definition

In the disjunctive sum

$$G + H$$

the player must move in exactly one of the components.

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$$G + H \rightarrow G^{L} + H \rightarrow G^{L} + H^{R} \rightarrow G^{LL} + H^{R}$$

Disjunctive sums in HACKENBUSH

Simplest number theorem

Theorem

A short game G, where $G^L < G^R$ for all G^L , G^R , is:

- Integer $G^L < n < G^R$ smallest in absolute value, if it exists.
- Fraction $G^L < \frac{i}{2^j} < G^R$ with smallest j, otherwise.

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Examples:

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$$\{ | \} = 0$$

•
$$\{-999, -21 \mid 42, 181\} = 0$$

•
$$\{1 \mid 2\} = \frac{3}{2} = 1\frac{1}{2}$$

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Not all games are numbers! Take $\{0 \mid 0\} = *$.

"Half" a move

Quantum-inspired combinatorial games

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Altered variant of a ruleset, in which *superposed moves* are allowed, leading to *superposed game states*.

Preliminaries

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A labelled short game $\hat{\boldsymbol{G}}$:

Preliminaries

11

A labelled short game $\hat{\mathbf{G}}$:

Definition

A *ruleset* is a pair of move functions $\mathcal{R} = (\rho_L, \rho_R)$.

$$\rho_L \left(\begin{array}{c|c} 3 & 4 & \\ 1 & 2 & 5 \end{array} \right) \equiv \begin{array}{c} 4 & \\ 2 & 5 \end{array}$$

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How do we label moves?

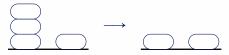
• Hackenbush: $move \iff edge_id$.

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- Hackenbush: move ←⇒ edge_id.
- Not so easy for other games, such as NIM:

Definition

NIM is played on a number of heaps, each containing at least one token. On their turn, the player selects one heap to remove a (nonzero) number of tokens from.



Definition 1

NIM-Subtract: the player selects \underline{a} heap and removes a (nonzero) number of tokens off the top.

Definition 2

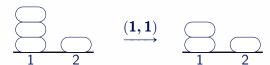
NIM-DECREASE: the player selects <u>a heap</u> and cuts it at some height, leaving a smaller heap.

NIM-SUBTRACT labelling

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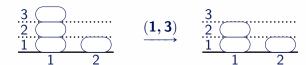
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Main idea 16

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Main idea

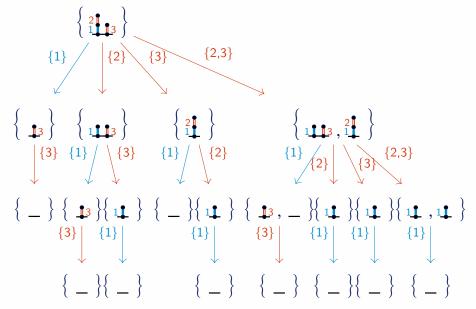
$$\left\{\begin{array}{c} \mathbf{3} \, | \, \mathbf{4} \, | \\ \mathbf{1} \, | \, \mathbf{2} \, | \, \mathbf{5} \, | \end{array}\right\} \stackrel{\left\{\mathbf{1},\,\mathbf{2}\right\}}{\longrightarrow} \left\{\begin{array}{c} \mathbf{4} \, | & \mathbf{3} \, | \\ \mathbf{2} \, | \, \mathbf{5} \, | \end{array}\right\}$$

Main idea

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$$\left\{\begin{array}{c} 3 \ | \ 4 \ | \\ 1 \ | \ 2 \ | \ 5 \ | \end{array}\right\} \xrightarrow{\left\{\begin{array}{c} 1,2 \right\}} \left\{\begin{array}{c} 4 \ | \\ 2 \ | \ 5 \ | \end{array}\right\} \xrightarrow{\left\{\begin{array}{c} 4 \ | \ 5 \ | \\ 2 \ | \ 5 \ | \end{array}\right\}} \left\{\begin{array}{c} 4 \ | \ 5 \ | \\ 2 \ | \ 5 \ | \end{array}\right\} \xrightarrow{\left\{\begin{array}{c} 4 \ | \ 5 \ | \\ 2 \ | \ 5 \ | \end{array}\right\}} \left\{\begin{array}{c} 4 \ | \ 3 \ | \\ 2 \ | \ 5 \ | \end{array}\right\}$$

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The flavours 18

When is an unsuperposed move allowed?

- **A**: Never.
- B: Never, except if the only option.
- ullet C: Only if legal in all realisations.
- C': Only if legal in all realisations in which the player still has at least one legal move.
- \mathcal{D} : Always.

Some properties of HACKENBUSH

• **Dead-ending**: moves cannot "come back".

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- Non-repeating: any move can only be played once.

Some properties of HACKENBUSH

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Some properties of HACKENBUSH

- Dead-ending: moves cannot "come back".
- Non-repeating: any move can only be played once.
- Persistent: interpretation of move does not change.
- Consistent: any move always has the same effect.

Useful theorem

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Theorem

Under flavour $f \in \{A, B, D\}$, if a realisation is *weakly covered*, then it can be left out of the superposed game state without changing its legal superposed move options.

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$$moves \left(\begin{array}{c} 1 \\ 1 \\ \end{array} \right)$$

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$$\left\{ \begin{array}{c} 1 \\ 1 \\ \end{array} \right\}^{f} \equiv \left\{ \begin{array}{c} 1 \\ 1 \\ \end{array} \right\}^{f}$$

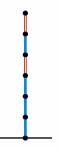
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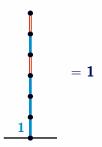




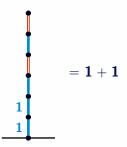
Theorem by Thea van Roode

Starting from the ground, count number of edges until first colour change, sign determined by colour of grounded edge. Then, divide by 2 for every next edge, add for blue and subtract for red.

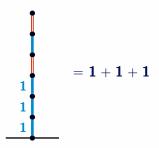
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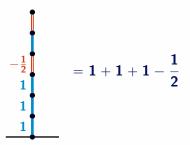
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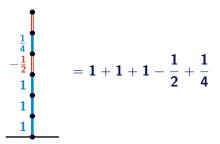
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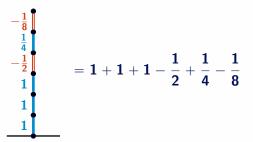
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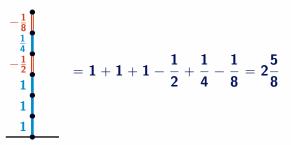
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Theorem

Under flavour $f \in \{\mathcal{B}, \mathcal{D}\}$, any HACKENBUSH stalk has the same value as in the classical case.

Proof

The realisation containing the tallest stalk weakly covers all other realisations.

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Theorem

Under flavour \mathcal{A} , any Hackenbush stalk has the same value as in the classical case, ignoring the first edge of either colour.

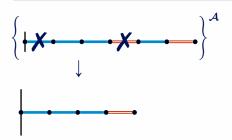


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Theorem

Under flavour ${\cal A}$, any Hackenbush stalk has the same value as in the classical case, ignoring the first edge of either colour.

$$\left\{\begin{array}{c|c} X & X & \end{array}\right\}^{A}$$

$$\downarrow \frac{1}{2} & \frac{1}{2} & \frac{1}{2} & = 1 + 1 + 1 - \frac{1}{2} = 2\frac{1}{2}$$

	4	5	6
000100	{1 1}	{ 2 2 }	{3 3 }
001000		{ 2 1 }	{3 2 }
001001		$\{1 \mid 1, \{1 \mid 1\}\}$	$\{2 \mid 2, \{2 \mid 2\}\}$
001011		$\{\frac{1}{2} \mid \frac{1}{2}\}$	$\{1\frac{1}{2} \mid 1\frac{1}{2}\}$
001100		$ \{\frac{1}{2} \mid \frac{1}{2}\} \\ \{\frac{1}{2} \mid \frac{1}{2}\} $	$\{1\frac{f}{2} \mid 1\frac{f}{2}\}$
010000			$\{3 \mid 1\}$
010001			$\{2 \mid 1, \{2 \mid 1\}\}$
010010			$\{1\frac{1}{2} \mid 1, \{1 \mid 1\}\}$
010011			$\{1 \mid 1, \{1 \mid 1, \{1 \mid 1\}\}\}$
010100			$\left\{\frac{3}{4} \mid \frac{3}{4}\right\}$
010110			$\{\frac{1}{2}, \{\frac{1}{2} \mid \frac{1}{2}\} \mid \frac{1}{2}\}$
010111			$\{\frac{1}{2} \mid \frac{1}{4}\}$
011000			$ \left\{ \frac{1}{2} \mid \frac{1}{4} \right\} \\ \left\{ \frac{3}{4} \mid \frac{1}{2} \right\} $
011001			
011011			$\left\{\frac{1}{4} \mid \frac{1}{4}\right\}$
011100			$ \left\{ \frac{1}{2} \mid \frac{1}{2}, \left\{ \frac{1}{2} \mid \frac{1}{2} \right\} \right\} \\ \left\{ \frac{1}{4} \mid \frac{1}{4} \right\} \\ \left\{ \frac{1}{4} \mid \frac{1}{4} \right\} $

Values of binary-encoded stalks under flavour C.

\mathcal{A}	0	1	2	3	4	5	\mathcal{B}	0	1	2	3	4	5	C	0	1	2	3	4	5
0	P	P	\mathscr{R}	R	R	R	0	P	R	R	\mathscr{R}	R	\mathscr{R}	0	Đ	R	R	R	R	\mathscr{R}
1	\mathscr{P}	\mathscr{R}	\mathscr{R}	\mathscr{R}	\mathscr{R}	\mathscr{R}	1	\mathscr{L}	\mathscr{R}	\mathscr{R}	\mathscr{R}	\mathscr{R}	\mathscr{R}	1	\mathscr{L}	\mathscr{R}	\mathscr{R}	\mathscr{R}	\mathscr{R}	\mathscr{R}
2	$\frac{\mathscr{P}}{\mathscr{L}}$	\mathscr{L}	\mathscr{R}	\mathscr{R}	\mathscr{R}	\mathscr{R}	2	\mathscr{L}	\mathscr{R}	\mathscr{R}	\mathscr{R}	\mathscr{R}	\mathscr{R}	2	\mathscr{L}	Đ	\mathscr{R}	\mathscr{R}	\mathscr{R}	\mathscr{R}
3	\mathscr{L}	9 R L L	\mathscr{R}	\mathscr{R}	\mathscr{R}	\mathscr{R}	3	\mathscr{L}	9	\mathscr{R}	\mathscr{R}	\mathscr{R}	\mathscr{R}	3	\mathscr{L}	\mathscr{L}	\mathscr{R}	\mathscr{R}	\mathscr{R}	\mathscr{R}
4	\mathscr{L}	\mathscr{L}	\mathscr{R}	\mathscr{R}	\mathscr{R}	\mathscr{R}	4	\mathscr{L}	<u>R</u> <u>P</u> L	\mathscr{R}	\mathscr{R}	\mathscr{R}	\mathscr{R}	4	\mathscr{L}	\mathscr{L}	\mathscr{R}	\mathscr{R}	\mathscr{R}	\mathscr{R}
5	\mathscr{L}	\mathscr{L}	$\frac{\mathscr{R}}{\mathscr{L}}$	\mathscr{R}	\mathscr{R}	\mathscr{R}	5	\mathscr{L}	\mathscr{L}	$\frac{\mathscr{R}}{\mathscr{R}}$ \mathscr{L}	\mathscr{R}	\mathscr{R}	\mathscr{R}	5	\mathscr{L}	\mathscr{L}	$\frac{\mathscr{R}}{\mathscr{L}}$	\mathscr{R}	\mathscr{R}	\mathscr{R}
6	\mathscr{L}	$_{\mathscr{L}}$	\mathscr{L}	\mathscr{R}	\mathscr{R}	\mathscr{R}	6	\mathscr{L}	\mathscr{L}	\mathscr{L}	\mathscr{R}	\mathscr{R}	\mathscr{R}	6	\mathscr{L}	$_{\mathscr{L}}$	\mathscr{L}	\mathscr{R}	\mathscr{R}	\mathscr{R}
7	\mathscr{L}	\mathscr{L}	\mathscr{L}				7	\mathscr{L}	\mathscr{L}	\mathscr{L}				7	\mathscr{L}	\mathscr{L}	\mathscr{L}			
8							8							8						
c'	0	1	2	3	4	5	\mathcal{D}	0	1	2	3	4	5		1					
$\frac{c'}{0}$	0 <i>P</i>	$\frac{1}{\mathscr{R}}$	2 R	3 R	4 R	<u>5</u>	D	P	1 R	2 R	3 R	4 R	5 R		ı					
	P								${\mathscr R}$											
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Outcome classes for different \mathbf{n} (vertical) and \mathbf{m} (horizontal). Max. move width $\mathbf{2}$.

$$\mathcal{N}\text{-positions}$$

$$\begin{cases}
3 & 4 & \\
1 & 2 & 5
\end{cases}
\end{cases}$$

$$\begin{cases}
1,2 \\
2 & 5
\end{cases}$$

$$\begin{cases}
3,5 \\
2 & 5
\end{cases}$$

$$\begin{cases}
4 & 3 \\
2 & 5
\end{cases}$$

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Future research

- Proving conjectures.
- Different classes of positions.
- Efficient representation of superposed game states.
- Quantum circuit to play Quantum Hackenbush.

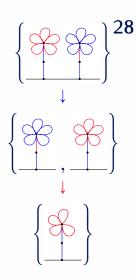
Finishing up

Summary and conclusion

- Values.
- Defining quantum-inspired combinatorial games.
- Quantum Hackenbush stalks.
- Outcome classes, $\mathcal N$ -positions.

Think (and taste) carefully...





References

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