Theorem 2.1.1 — Logical Equivalences (Epp page 35)

Given any statement variables p, q, and r, a tautology \mathbf{t} , and a contradiction \mathbf{c} , the following logical equivalences hold:

| $p \wedge q \equiv q \wedge p$ | $p \lor q \equiv q \lor p$ |
|--|--|
| $(p \land q) \land r \equiv p \land (q \land r)$ | $(p \lor q) \lor r \equiv p \lor (q \lor r)$ |
| $p \land (q \lor r) \equiv (p \land q) \lor (p \land r)$ | $p \lor (q \land r) \equiv (p \lor q) \land (p \lor r)$ |
| $p \wedge \mathbf{t} \equiv p$ | $p \lor \mathbf{c} \equiv p$ |
| $p \lor \sim p \equiv \mathbf{t}$ | $p \wedge \sim p \equiv \mathbf{c}$ |
| $\sim (\sim p) \equiv p$ | |
| $p \wedge p \equiv p$ | $p \lor p \equiv p$ |
| $p \lor \mathbf{t} \equiv \mathbf{t}$ | $p \wedge \mathbf{c} \equiv \mathbf{c}$ |
| $\sim (p \land q) \equiv \sim p \lor \sim q$ | $\sim (p \vee q) \equiv \sim p \wedge \sim q$ |
| $p \lor (p \land q) \equiv p$ | $p \land (p \lor q) \equiv p$ |
| \sim t \equiv c | \sim c \equiv t |
| | $(p \land q) \land r \equiv p \land (q \land r)$ $p \land (q \lor r) \equiv (p \land q) \lor (p \land r)$ $p \land \mathbf{t} \equiv p$ $p \lor \sim p \equiv \mathbf{t}$ $\sim (\sim p) \equiv p$ $p \land p \equiv p$ $p \lor \mathbf{t} \equiv \mathbf{t}$ $\sim (p \land q) \equiv \sim p \lor \sim q$ $p \lor (p \land q) \equiv p$ |

Table 2.3.1 — Rules of Inference (Epp, page 60)

| Modus ponens | $p \rightarrow q$ | | Disjunctive syllogism | $p \lor q$ $p \lor q$ |
|----------------------------|------------------------|-----------------------|------------------------|--|
| | p | | | $ \sim q \sim p$ |
| | $\therefore q$ | | | $ \therefore p \qquad \therefore q$ |
| Modus tollens | $p \rightarrow q$ | | Hypothetical syllogism | $p \rightarrow q$ |
| | $\sim q$ | | | $q \rightarrow r$ |
| | $\therefore \sim p$ | | | $\therefore p \to r$ |
| Disjunctive addition | p | q | Dilemma, or | $p \lor q$ |
| | $\therefore p \lor q$ | $\therefore p \lor q$ | Proof by division | $p \rightarrow r$ |
| Conjunctive simplification | $p \wedge q$ | $p \wedge q$ | into cases | $q \rightarrow r$ |
| | $\therefore p$ | $\therefore q$ | | $ \therefore r$ |
| Conjunctive addition | p | | Contradiction rule | $\sim p \rightarrow \mathbf{c}$ |
| | q | | | $ \therefore p$ |
| | $\therefore p \land q$ | | | |
| Closing conditional | p (a | ssumed) | Closing conditional | p (assumed) |
| world without | q (d | lerived) | world with | $ q \wedge \sim q \text{ (derived)}$ |
| contradiction | $\therefore p \to q$ | | contradiction | .∴~p |

Other Logical Equivalences and Rules of Inference

| Definition of imp | olication: | $\mid p \to q \equiv \sim p \vee$ | q | $ \mid \sim (p \to q) \equiv p \land \sim q $ | |
|--------------------|--|---|-----------------------|---|--|
| Contrapositive r | ule: | $p \to q \equiv \sim q -$ | $\rightarrow \sim p$ | | |
| Definition of bice | onditional: | $p \leftrightarrow q \equiv (p \rightarrow$ | $q) \wedge (q \to p)$ | | |
| Negation of quar | ntifiers: | $\sim (\forall x \ P(x)) \equiv$ | $\exists x \sim P(x)$ | | |
| Universal | $\forall x \in D, P(x) \to Q(x)$ | | Universal | $\forall x \in D, P(x) \to Q(x)$ | |
| modus ponens: | $P(a)$ where $a \in D$ | | modus tollens: | $\sim Q(a)$ where $a \in D$ | |
| | $\therefore Q(a)$ | | | $\therefore \sim P(a)$ | |
| Universal | $\forall x \in D, P(x)$ | | Existential | $\exists x \in D, P(x)$ | |
| instantiation: | $\therefore P(a) \text{ where } a \in D$ | | instantiation*: | $\therefore P(a) \text{ where } a \in D$ | |
| Universal | $P(a)$ where $a \in D$ | | Existential | $P(a)$ where $a \in D$ | |
| generalization*: | $\therefore \forall x \in D, P(x)$ | | generalization: | $\therefore \exists x \in D, P(x)$ | |

^{*}Remember the special circumstances required for the rules marked by the stars.