

MIPS Instruction Set (last revision: 11/05/2002)

Instruction	Example	Meaning	Comments
Arithmetic			
add	add \$s1,\$s2,\$s3	\$s1 = \$s2 + \$s3	overflow detected
subtract	sub \$s1,\$s2,\$s3	\$s1 = \$s2 - \$s3	overflow detected
add imm	addi \$s1,\$s2,100	\$s1 = \$s2 + 100	overflow detected
add uns	addu \$s1,\$s2,\$s3	\$s1 = \$s2 + \$s3	overflow <i>not</i> detected
subtract uns	subu \$s1,\$s2,\$s3	\$s1 = \$s2 - \$s3	overflow <i>not</i> detected
add uns imm	addiu \$s1,\$s2,\$s3	\$s1 = \$s2 + 100	overflow <i>not</i> detected
multiply	mult \$s2,\$s3	Hi,Lo = \$s2 × \$s3	64-bit signed product
multiply uns	mult \$s2,\$s3	Hi,Lo = \$s2 × \$s3	64-bit unsigned product
divide	div \$s2,\$s3	Lo = \$s2/\$s3; Hi = \$s2 mod \$s3	Lo,Hi = quotient,remainder
divide uns	divu \$s2,\$s3	Lo = \$s2/\$s3; Hi = \$s2 mod \$s3	unsigned Lo,Hi
move from Hi	mfhi \$s1	\$s1 = Hi	copy Hi
move from Lo	mflo \$s1	\$s1 = Lo	copy Lo
Logical			
and	and \$s1,\$s2,\$s3	\$s1 = \$s2 & \$s3	logical AND
or	or \$s1,\$s2,\$s3	\$s1 = \$s2 \$s3	logical OR
and imm	andi \$s1,\$s2,100	\$s1 = \$s2 & 100	logical AND
or imm	ori \$s1,\$s2,100	\$s1 = \$s2 100	logical OR
shift left logical	sll \$s1,\$s2,10	\$s1 = \$s2 << 10	shift left by const
shift right logical	srl \$s1,\$s2,10	\$s1 = \$s2 >> 10	shift right by const
Data transfer			
load word	lw \$s1,100(\$s2)	\$s1 = Mem[\$s2+100]	word from mem to reg
store word	sw \$s1,100(\$s2)	Mem[\$s2+100] = \$s1	word from reg to mem
load byte	lb \$s1,100(\$s2)	\$s1 = Mem[\$s2+100]	sign extended
load byte uns	lbu \$s1,100(\$s2)	\$s1 = Mem[\$s2+100]	not sign extended
store byte	sb \$s1,100(\$s2)	Mem[\$s2+100] = \$s1	low byte from reg to mem
load upper imm	lui \$s1,100	\$s1 = 100 * 2 ¹⁶	load const, upper 16 bits
load imm†	li \$s1,100	\$s1 = 100	load const
load address†	la \$s1,L1	\$s1 = address L1	load address
move registers†	move \$s1,\$s2	\$s1 = \$s2	
Conditional branch			
branch on equal	beq \$s1,\$s2,25	if (\$s1 == \$s2) goto PC + 4 + 4*25	PC-relative branch
branch on not equal	bne \$s1,\$s2,25	if (\$s1 != \$s2) goto PC + 4 + 4*25	PC-relative branch
branch on greater than†	bgt \$s1,\$s2,25	if (\$s1 > \$s2) goto PC + 4 + 4*25	PC-relative branch
branch on less than†	blt \$s1,\$s2,25	if (\$s1 < \$s2) goto PC + 4 + 4*25	PC-relative branch
branch on greater or eq†	bge \$s1,\$s2,25	if (\$s1 >= \$s2) goto PC + 4 + 4*25	PC-relative branch
branch on less than or eq†	ble \$s1,\$s2,25	if (\$s1 <= \$s2) goto PC + 4 + 4*25	PC-relative branch
set on less than	slt \$s1,\$s2,\$s3	if (\$s2<\$s3) \$s1=1; else \$s1 = 0	two's complement
set on less than imm	slti \$s1,\$s2,100	if (\$s2<100) \$s1=1; else \$s1 = 0	two's complement
set on less than uns	sltu \$s1,\$s2,\$s3	if (\$s2<\$s3) \$s1=1; else \$s1 = 0	natural numbers
set on less than imm uns	sltiu \$s1,\$s2,100	if (\$s2<100) \$s1=1; else \$s1 = 0	natural numbers
Unconditional jump			
jump	j 2500	go to 4*2500	jump to target address
jump register	jr \$ra	go to addr in \$ra	for switch, proc return
jump and link	jal 2500	\$ra = PC + 4; go to 2500*4	for proc call
Operating system			
syscall	syscall	system call	see A-48, A-49

uns = Unsigned; *imm* = Immediate; † = Pseudo-instruction

Floating-point

Instruction	Example	Meaning	Comments
Arithmetic			
FP add single	add.s \$f2,\$f4,\$f6	\$f2 = \$f4 + \$f6	
FP add double	add.d \$f2,\$f4,\$f6	\$f2 = \$f4 + \$f6	
FP subtract single	sub.s \$f2,\$f4,\$f6	\$f2 = \$f4 - \$f6	
FP subtract double	sub.d \$f2,\$f4,\$f6	\$f2 = \$f4 - \$f6	
FP multiply single	mul.s \$f2,\$f4,\$f6	\$f2 = \$f4 × \$f6	
FP multiply double	mul.d \$f2,\$f4,\$f6	\$f2 = \$f4 × \$f6	
FP divide single	div.s \$f2,\$f4,\$f6	\$f2 = \$f4 / \$f6	
FP divide double	div.d \$f2,\$f4,\$f6	\$f2 = \$f4 / \$f6	
Data transfer and conversion			
load single†	l.s \$f0,addr	\$f0 = Mem[addr]	memory (32 bits) to FP reg
load double†	l.d \$f0,addr	\$f0 = Mem[addr]	memory (64 bits) to FP reg
load word coproc 1	lwc1 \$f1,100(\$s2)	\$f1 = Mem[\$s2 + 100]	32-bit data to FP reg
store word coproc 1	swc1 \$f1,100(\$s2)	Mem[\$s2 + 100] = \$s2	32-bit data to memory
convert single to double	cvt.d.s \$f2,\$f4	\$f2 = (double) \$f4	also .d.s, .w.d, etc.; w=int
move FP regs, single	mov.s \$f0,\$f2	\$f0 = \$f2	copy FP reg, single
move FP regs, double	mov.d \$f0,\$f2	\$f0 = \$f2	copy FP reg, double
move to FP, single	mtc1 \$s0,\$f0	\$f0 = \$s0	32-bit value, not converted
move from FP, single	mfc1 \$s0,\$f0	\$s0 = \$f0	32-bit value, not converted
Conditional branch			
branch on FP true	bc1t 25	if (cond == 1) goto PC + 4 + 4*25	PC relative branch
branch on FP false	bc1f 25	if (cond == 0) goto PC + 4 + 4*25	PC relative branch
FP compare single	c.lt.s \$f2,\$f4	if (\$f2<\$f4) cond=1; else cond=0	also eq,ne,lt,le,gt,ge; also .d
FP compare double	c.lt.d \$f2,\$f4	if (\$f2<\$f4) cond=1; else cond = 0	also eq,ne,lt,le,gt,ge; also .d

Registers	Usage
\$zero	constant 0 (enforced by hardware)
\$v0, \$v1	expression evaluation, function return value
\$a0 ... \$a3	first four function arguments (rest on stack)
\$t0 ... \$t9	temporary; caller-saved
\$s0 ... \$s7	callee-saved
\$sp	stack pointer; last location on stack; grow by subtracting
\$fp	frame pointer; first word in current stack frame
\$ra	return address, written by jal instruction
\$f0 ... \$f31	floating-point registers; for double precision even numbers refer to pairs, i.e., \$f0 uses \$f0 and \$f1.

Compiler directive	Description
.asciiz "string"	Store <i>string</i> in memory, add null terminator
.data	Subsequent items stored in data segment
.double 3.1415926535898	Store a double-precision floating-point number in memory
.float 3.141593	Store a single-precision floating-point number in memory
.globl <i>symbol</i>	Declare label <i>symbol</i> to be global
.space <i>n</i>	Allocate <i>n</i> bytes of space (must be in data segment)
.text	Subsequent items stored in text segment
.word <i>q</i>	Store 32-bit quantity <i>q</i> in memory

Based on Computer Organization and Design (Hennessy and Patterson), p. 274, p. 291 and Appendix A.