Nano-RK:

an Energy-aware Resource-centric RTOS for Sensor Networks

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Presented by Anthony Rowe RTSS 2005





Secure Sensor Networks for Physical Infrastructures



Develop a secure software platform for infrastructure sensor networks to be deployed in physical structures such as factories, buildings, homes, bridges, campuses, vehicles, ships and planes.

- Provide continual monitoring of operational health and safety,
- Report malfunctions (nearly) instantly, and
- Support mobile nodes and track people.









What is nanoRK?

- A Real-time Operating System for sensor nodes for use in wireless sensor networks
 - Priority-Based Preemptive Multitasking
 - Resource Reservations
 - Built-in Multi-hop Networking Support





Related Work

- Tiny OS (Berkeley)
 - Large public following and support
 - Compact
 - Cooperative Multitasking
 - No timing guarantees for tasks
- Mantis OS (University of Colorado)
 - No real-time scheduling
 - No reservation support
- uCos, Emerald, OSEK
 - Still too large
 - Not network centric



NanoRK Motivation

Quicker Development Cycles

 Traditional OS abstractions allow for quicker learning curves and help abstract away low level details

Scaling Technology

- Sensor Nodes will become increasingly complex as they begin to manage multiple tasks
- Local Processing is much cheaper than using the network

Resource Reservations (Energy Budget)

 Reservations can enforce energy and communication budgets to minimize negative impact on network lifetimes from unintentional errors or malicious behavior on some nodes



NanoRK Motivation

Why Priority-based scheduling?

	Period	Execution Time	
Network Radio	Sporadic	10ms	
Audio Sensor	200 hz	10us	
Light Sensor	166 hz	10us	
Smart Camera	1 hz	300ms	
Global Positioning	5 hz	10ms	

Time-triggered task interleaving can become daunting...

Furthermore, what if a new sensor is added or a period changes?



Goals of a Sensor RTOS (1 of 2)

- Multitasking with Priority-Based Preemption
 - Respond on a timely basis to critical events
 - Support predictability and schedulability
- Built-in Network Support
 - Decrease coding effort to implement custom communication protocols
- Enforce Energy Usage Limits
 - Meet Battery Lifetime Requirements



Goals of a Sensor RTOS (2 of 2)

Unified Sensor / Network API

 Sensors and actuators support in a device driver manner that is abstracted away from the user

Small Footprint

 Current trends in microcontrollers show larger ROM sizes (64Kb to 128KB) and smaller RAM sizes (2KB to 8KB)



FireFly Hardware

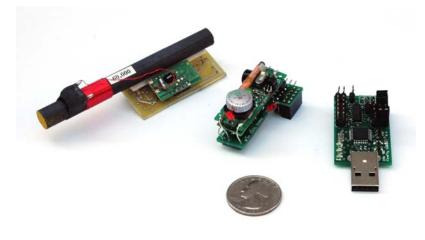
Atmega128L

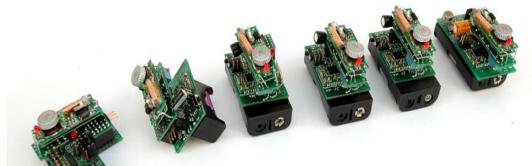
CC2420 (802.15.4)

Light, Temp, Audio, PIR, Acceleration, Ultrasound

Global AM band Synchronization



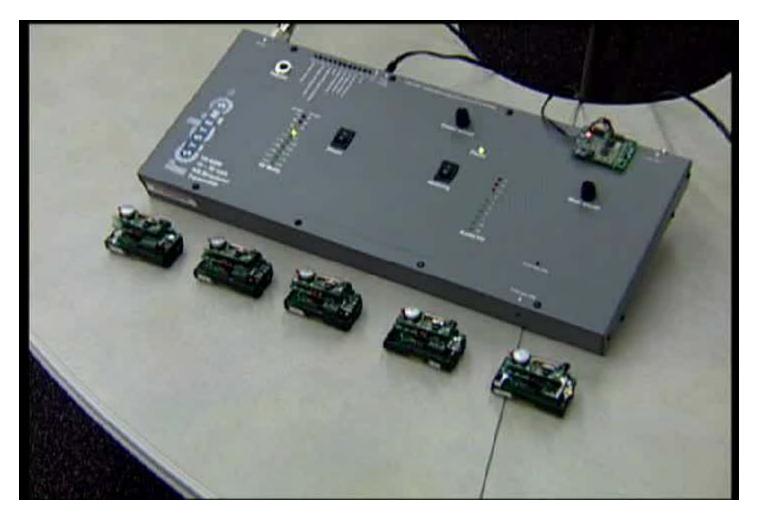








FireFly Synchronization Hardware





NanoRK Resources

Component	Resource	
Context Swap Time	45 μS	
Mutex Structure Overhead	5 Bytes per Resource	
Stack Size Per Task	32→128 bytes (64 bytes by default)	
OS Struct Overhead	50 bytes Per Task	
Network Overhead	164 bytes	
Total Typical Configuration: (8 tasks, 8 mutexes, 4 16 byte network buffers)	2KB RAM, 10KB ROM	

Current Hardware Platform:

4Mhz Atmega128L with Chipcon CC2420 802.15.4 transceiver



NanoRK Reservations

CPU

 Each Task can be given a budget of how long it is allowed to execute per a given period

Network

- Per Task Budget on Network Usage
- Transmit and Receive Packets and/or Bytes

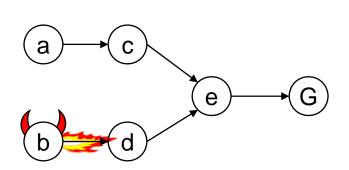
Sensors / Actuators

 Number of System Calls to a particular peripheral per a given period



Virtual Energy Reservations

- {CPU, Network, Sensors}
 - Together comprise the total energy usage of the node
- Static Offline Budget Enforcement
 - It is possible to calculate a node lifetime given a certain energy budget



Node	Reserve [TX, RX]	TX Rate	Lifetime w/ out Reserve	Lifetime w/ Reserve
а	[1,2]	1	8 years	8 years
b	X	300	3.5 days	3.5 days
С	[1,2]	1	5 years	5 years
d	[1,2]	1	3.9 days	5 years
е	[1,2]	1	4 days	2.9 years



Real-Time and Multimedia Systems Laboratory

NanoRK Reservations

- What if a reservation is violated?
- Hard Reservation
 - For CPU reservations, the task is preempted
 - For Sensor/Network reservations, an error is returned to the task making the system call
- Firm Reservations
 - The task is allowed to consume slack resources until they have been depleted



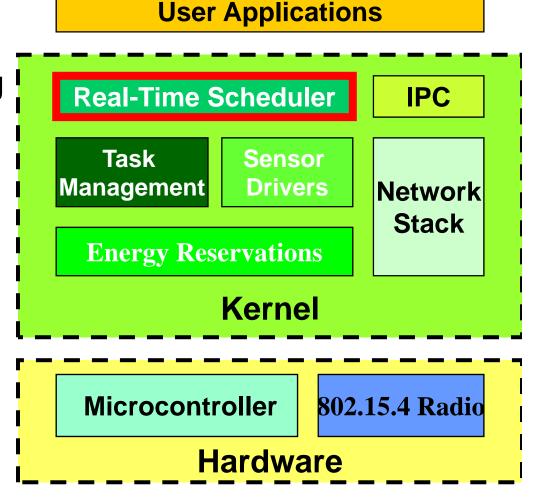
User Applications Real-Time Scheduler IPC Task Sensor Management **Drivers Network Stack Energy Reservations Kernel** Microcontroller 802.15.4 Radio

Hardware



Real-Time Scheduler

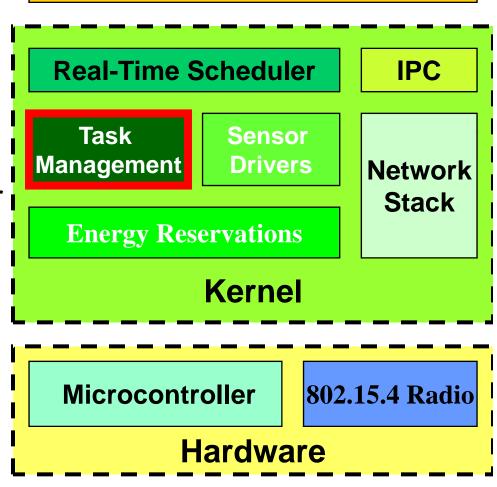
- Priority-based Scheduling
- Static Design-time Framework
- Priority Ceiling Protocol





Task Management

- Two 32 bit counters
 - { Seconds, Nanoseconds }
- Periodic Task Switching
 - Triggered by a One Shot Timer or an Event
- Enables Fine Grained Timing
 Control



User Applications



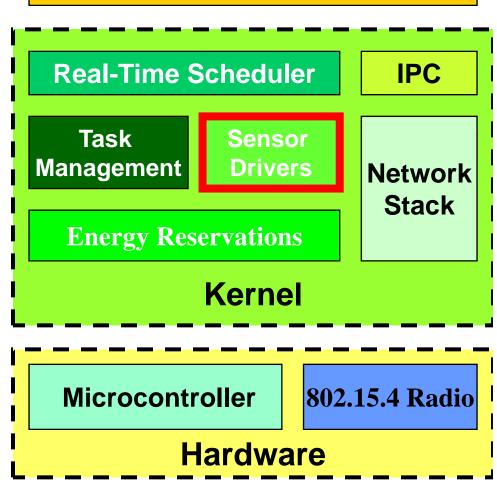
Nano-RK Task Specs

```
nrk task type Task1;
Task1.task = Sound Task;
Task1.Ptos = (void *) &Stack1[STACKSIZE - 1];
Task1.TaskID = 1;
Task1.priority = 3;
Task1.Period = 1000; // ms
Task1.set reserve[CPU] = 50; //ms
Task1.set_reserve[NETWORK_PACKETS] = 3;
Task1.set reserve[SENSE ACTUATE] = 12;
nrk activate task(Task1);
```



Sensor Drivers

- System Calls
 - Allows for arbitration if two tasks read from a non-atomic sensor
- Raw ADC Values
 - Temp = 67
- Real World Values
 - Temp = 24 (°C)

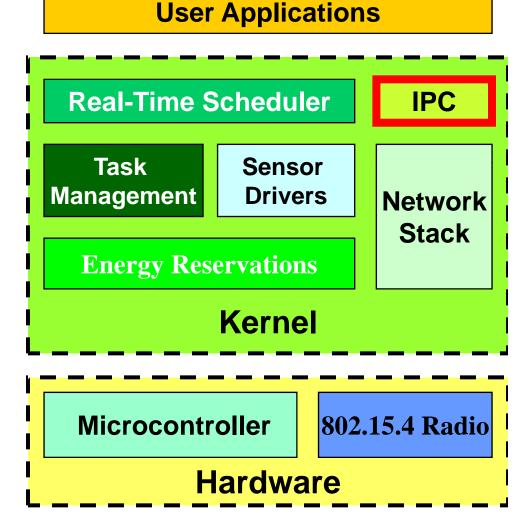


User Applications



Inter Process Communication

- Use Semaphores to arbitrate shared memory communication
- "Message Boxes"





NanoRK Network Stack

Zerocopy Buffering Mechanism

- TX and RX buffers are in application space
- Kernel Manages network data through pointer manipulation into application space

Port Abstraction

A port allows applications to direct their data to individual tasks on each node

Network Task

Handles Forwarding Packets, Routing, etc.





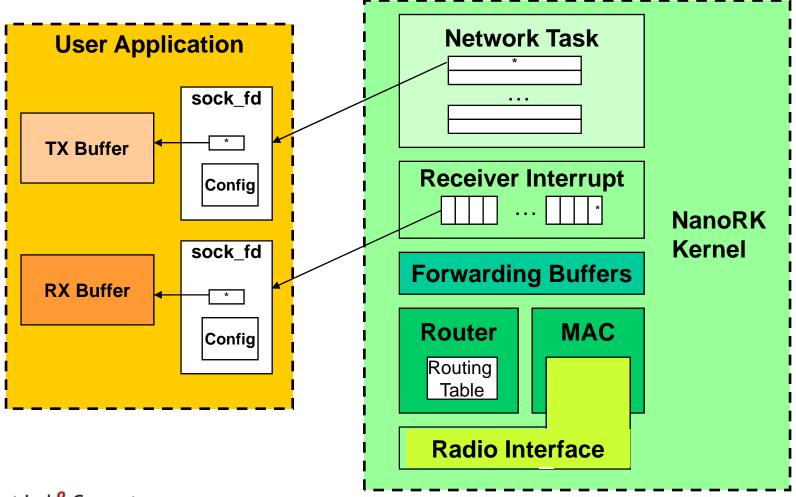
NanoRK Network Stack

- CSMA MAC Support
 - Low Power Listen (LPL) MAC support
- TDMA MAC Support
 - RT-Link: Globally Time Synchronized MAC protocol
- Ad-Hoc Routing Support
 - Tasks can access low level routing table
 - Ex: An AODV or DSR <u>TASK</u> can be responsible for establishing and managing routes





Detailed NanoRK Network Stack





NanoRK Limitations

- No User / Kernel Space Boundary
 - Single Memory Space with no MMU
- Somewhat higher memory overhead
 - Compared to non-multitasking operating systems



Simulation / Modeling Tool

Represent Network as a Graph

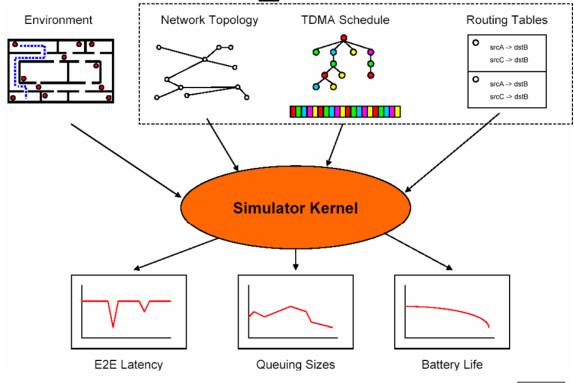
Easily evaluate different scheduling algorithms

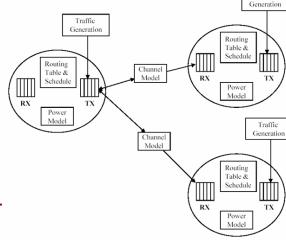
Cycle-Accurate Power and Channel Modeling

 Analyze how different schedules effect power, throughput, and latency

Hybrid simulator

- Allow virtual expansion of our real network
- Works with real nodes





Traffic



Future Work

- A public domain release
- Dynamic upgrades
- Support for multiple microcontrollers
- Family of applications
- Complete set of tools
- Field deployment



Conclusions

- NanoRK: a sensor network Operating System
 - priority-based preemptive multitasking
 - task synchronization
 - multi-hop communications support
- Reservations can enforce system-wide energy and communication usage
- Easy-to-use high level Networking and Sensor Abstractions



Questions?

