4 March 2015

Problem 1 (2 parts, 35 points)

Storage Allocation and Pointers

Part A (15 points) Assuming a **64-bit system with 64-bit memory interface and 64-bit addresses**, show how the following global variables map into static memory. Assume they are allocated starting at address 4000 and are properly aligned. For each variable, draw a box showing its size and position in memory. Label the box with the variable name. Label each element of an array (e.g., M[0]). Note that int and float are still 32-bits.

	$M[] = \{1.44, 3.69, 4.64\};$ Le *P = &G z = 4; *Q = M; 4020 4024 4028	4004	G	
		4008	M[0]	
		4012	M[1]	
		M[2]		
double float		4020	slack	
double			4024	P
float			4028	_
double		4032	z	
		4036	slack	
		4040	Q	
		4044	¥	
		4048	н	
		4052	**	

Part B (20 points) For this part, assume a 32-bit system, such as MIPS-32 that we have worked with.

```
int X = 5;
int V = 7;
char S[] = "hello!";
double Z = 7.25;
double *P = &Z;
int *R = &X;
```

Question:	Answer:
How much space (in bytes) is allocated for R?	4 bytes: pointer = word
How much space (in bytes) is allocated for P?	4 bytes: pointer = word
What is printed by this statement? printf("%d\n", *(R+1));	7
What is printed by this statement? printf("%c\n", *(S+4));	0
What is printed by this statement? printf("%f\n", *P);	7.25

Exam Two Solutions

4 March 2015

Problem 2 (4 parts, 27 points)

Struct and Array Access

Consider the following declarations.

Part A (3 points) Array A contains an element of value **11** (shown in bold). Write a single C statement that overwrites that element with the value **100**.

$$A[2][1] = 100;$$

Part B (3 points) Write a C statement that sets the base field of T to 16, using the variable T.

$$T.base = 16;$$

Part C (3 points) Write a C statement that sets the height field of T to 20, using the variable P.

Part D (18 points) Write the MIPS code implementation of the dynamically allocated array access below in the smallest number of instructions. A pointer to the array **Icons** (declared below) is stored in \$3. Variables **I**, **Y**, **X**, and **Cell** reside in \$4, \$5, \$6, and \$2 respectively. Modify only registers \$1 and \$2. Assume a 32-bit operating system.

Label	Instruction	Comment
	addi \$1, \$0, 192	# Lx*Ly = 192
	mult \$1, \$4	# I*Lx*Ly
	mflo \$1	# \$1: I*Lx*Ly
	sll \$2, \$5, 4	# \$2: Y*Lx = Y*16
	add \$1, \$1, \$2	# \$1: I*Lx*Ly+Y*Lx
	add \$1, \$1, \$6	# \$1: I*Lx*Ly+Y*Lx+X
	sll \$1, \$1, 2	# scale by 4
	add \$1, \$1, \$3	# add base
	lw \$2, 0(\$1)	<pre># read Icons[I][Y][X]</pre>

4 March 2015

Problem 3 (2 parts, 38 points)

Activation Frames

Consider the following C code fragment:

Part A (18 points) Suppose TrapezoidArea has been called with input H=10 so that the state of the stack is as shown below. Describe the state of the stack just before Average deallocates locals and returns to TrapezoidArea. Fill in the unshaded boxes to show TrapezoidArea's (TA's) and Average's activation frames. Include a symbolic description and the actual value (in decimal). For return addresses, show only the symbolic description; do not include a value. Label the frame pointer and stack pointer.

Address 9900 9896 9892 H 10	ck pointer.		
## Pof TA's caller 10 10 10 10 10 10 10 1	address		Value
### 10 RV SP7TrapezoidArea's FP 9888 RV SP 9884 Bases[1] 5 5 5 5 5 5 5 5 5	9900	RA of TA's caller	
## RV 9884 Bases[1] 5 5 9880 Bases[0] 3 3 9876 N 2 2 2 2 2 3 3 3 3 3	9896	FP of TA's caller	
9884 Bases[1] 5 9880 Bases[0] 3 9876 N 2 9872 R 9868 RA 9864 FP 9888 9860 A 9880 9856 S 2 Average's FP 9852 RV 4 9848 X 3 9844 Y 5 SP 9840 avg 4 9836 9832	9892	H	10
9880 Bases[0] 3 9876 N 2 9872 R 9868 RA 9864 FP 9888 9860 A 9880 9856 S 2 Average's FP 9852 RV 4 9848 X 3 9844 Y 5 SP 9840 avg 4 9836 9832	SP, TrapezoidArea's FP 9888	RV	
9876 N 2 9872 R 9868 RA 9864 FP 9888 9860 A 9880 9856 S 2 Average's FP 9852 RV 4 9848 X 3 9844 Y 5 SP 9840 avg 4 9836 9832	9884	Bases[1]	5
9872 R 9868 RA 9864 FP 9888 9860 A 9880 9856 5 2 Average's FP 9852 RV 4 9848 X 3 9844 Y 5 SP 9840 avg 4 9836 9832	9880	Bases[0]	3
9868 RA 9864 FP 9888 9860 A 9880 9856 S 2 Average's FP 9852 RV 4 9848 X 3 9844 Y 5 SP 9840 avg 4 9836 9832	9876	N	2
9864 FP 9888 9860 A 9880 9856 S 2 Average's FP 9852 RV 4 9848 X 3 9844 Y 5 SP 9840 avg 4 9836 9832	9872	R	
9860 A 9880 9856 S 2 Average's FP 9852 RV 4 9848	9868	RA	
9856	9864	FP	9888
Average's FP 9852 RV 4 9848 X 3 9844 Y 5 SP 9840 avg 4 9836 9832	9860	A	9880
9848	9856	5	2
9844	Average's FP 9852	RV	4
SP 9840 avg 4 9836 9832	9848	×	3
9836 9832	9844	У	5
9832	SP 9840	avg	4
	9836		
9828	9832		
	9828		

Exam Two Solutions

4 March 2015

Part B (20 points) Write MIPS code fragments to implement the subroutine Average by following the steps below. Do not use absolute addresses in your code; instead, access variables relative to the frame pointer. Assume no parameters are present in registers (i.e., access all parameters from Average's activation frame). You may not need to use all the blank lines provided. (You may assume that intermediate values your code stores in registers stay in the registers from one part to the next unless your code overwrites them.)

First, write code to properly set Average's frame pointer and to allocate space for Average's local variables and initialize them if necessary.

label	instruction	Comment
Average:	addi \$30, \$29, 0	# set FP base
	addi \$29, \$29, -12	# make room for locals (3 words)

x = A[0];

label	instruction	Comment
	lw \$1, 8(\$30)	# load A (base of array)
	lw \$2, 0(\$1)	# load A[0]
	sw \$2, -4(\$30)	# store it in x

y = A[1];

label	instruction	Comment
	lw \$3, 4(\$1)	# load A[1]
	sw \$3, -8(\$30)	# store it in y

avg = (x+v)/S:

label	instruction	Comment
	add \$1, \$2, \$3	# x+y
	lw \$3, 4(\$30)	# load S
	div \$1, \$3	# (x+y)/S
	mflo \$1	
	sw \$1, -12(\$30)	# store it in avg

return(avg); (store return value, deallocate locals, and return)

label	instruction	Comment
	sw \$1, 0(\$30)	# store return value
	addi \$29, \$30, 0	# deallocate locals
	jr \$31	# return to caller