

## INTEGER PROGRAMMING. EXERCISE SET 2

**E1** Show that the Set Covering Problem can be written as the COP

$$\min_{S \subseteq N} \left\{ \sum_{j \in S} c_j \mid v(S) = v(N) \right\},$$

where  $v(S) = \sum_{i=1}^m \min \{ \sum_{j \in S} a_{ij}, 1 \}$ . What is the meaning of  $v(S)$ ?

**E2** Suppose that an IP or LP has  $n$  constraints and we want to ensure that at least one of them must hold. How can we use binary variables to deal with this case?

**E3** Suppose now that in our IP or LP,  $k$  out of  $n$  constraints must hold. Find a way to express this fact by means of binary variables.

**E4** A set of  $n$  jobs must be carried out on a single machine that can do only one job at a time. Each job  $j$  takes  $p_j$  hours to complete. Given job weights  $w_j$  for  $j = 1, \dots, n$ , in what order should the jobs be carried out so as to minimize the weighted sum of their start times? Formulate this scheduling problem as a mixed integer program.

**E5** Consider the following formulation of the TSP with  $n$  towns where we define  $x_{ij} = 1$  if the salesman goes directly from town  $i$  to town  $j$  and  $x_{ij} = 0$  otherwise:

$$\min \sum_{i=0}^n \sum_{j=0}^n c_{ij} x_{ij}$$

$$\sum_{i=1}^n x_{ij} = 1 \text{ for } j = 1, \dots, n.$$

$$\sum_{j=1}^n x_{ij} = 1 \text{ for } i = 1, \dots, n.$$

$$0 \leq x_{ij} \leq 1 \quad \text{for all } i, j$$

$$x_{ij} \text{ integer for all } i, j$$

Show that we are justified to substitute the second constraint by

$$\sum_{j=1}^n x_{ij} = 1 \text{ for } i = 2, \dots, n.$$