# Algorithmic Game Theory Homework 6 - Cooperative Games

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## 1 Problem 10

## 1.1 Question a

#### 1.1.1 Monotone

In order to probe if this game is *monotone* we need to probe that  $v(C) \leq v(D)$  for any C, D such that  $C \subseteq D$ . Since we have a case function we are going to analyze the different parts.

Let have a graph  $G_x = (V, X)$ . Since by definition the players are the Edges, we can add player to X, so we are adding edges Y to the graph  $G_x$ , such that  $X' = X \cup Y$ . Lets call this new graph  $G'_x = (V', X')$ .

- If  $G'_x$  is connected then  $v(X') = 2|X'| \operatorname{diam}(G'_x)$ . Since we are adding more edges |X'| > |X|, and the diameter is at least greater or equal  $\operatorname{diam}(G'_x) \geq \operatorname{diam}(G_x)$  because it is connected and by definition of diameter if we are adding edges the greatest length of the shortest path cannot be smaller with more edges.
- If  $G'_x$  is not connected also  $\frac{|X'|}{2} > \frac{|X|}{2}$ .

Therefore for any  $G_x \subseteq G'_x$ ,  $v(X) \le v(X')$ , so it is **Monotone**.

#### 1.1.2 Superadditive

A game is Superadditive if  $v(C \cup D) \ge v(C) + v(D)$  for any 2 disjoint coalitions C and D.

Lets analyze each case. Lets take if  $G_{X \cup Y}$  is connected, then

$$v(X \cup Y) = 2(|X| + |Y|) - diam(G_{X \cup Y})$$
(1a)

$$= 2|X| + 2|Y| - diam(G_{X \cup Y})$$
 (1b)

And  $2|X| + 2|Y| - diam(G_{X \cup Y}) \ge 2|X| - diam(G_X) + 2|Y| - diam(G_Y)$  because  $diam(G_{X \cup Y}) \le diam(G_X) + diam(G_Y)$  since it is included in one of the 2.

Lets take if  $G_{X \cup Y}$  is not connected, then

$$v(X \cup Y) = \frac{|X| + |Y|}{2} \tag{2a}$$

$$=\frac{|X|}{2} + \frac{|Y|}{2} \tag{2b}$$

$$= v(X) + v(Y) \tag{2c}$$

Therefore it is superadditive.

#### 1.1.3 Supermodular

A game is Supermodular if  $v(C \cup D) + v(C \cap D) \ge v(C) + v(D)$ .

By the previous item we have probe that it is superadditive, so if we sum  $v(C \cap D)$  to the union, we are still going to have something greater or equal of the sum of the individual valuation functions, therefore it is also supermodular.

### 1.2 Question b

Yes, a graph of only 1 edge because. Let suppose that that graph is A = (Z, Y) where Z = z, v and Y = z, v. Lets assume that we take X = Y, then

- $v(X) = 2|X| diam(A_X) = 2 \times 1 1 = 1$
- Also we have that x(X) = 1 because we have 1 edge.
- Therefore the core is not empty.