Sampling in networks

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Please go to http://www.cs.upc.edu/~csn for all course's material, schedule, lab work, etc.

The "problem" of analyzing networks

Sampling comes to our rescue

A few possible scenarios:

- 1. We have collected a *large* graph that fits into memory, but want to run an expensive algorithm that may take too long. How can we speed up the computation?
- 2. We have collected a *huge* graph that fits into disk but not main memory. How can we analyze it in reasonable time?
- 3. It is extremely costly or impossible to collect the entire graph (think Facebook, WWW, Twitter, etc.), we only have access to subgraphs via *crawling*, and yet we want to infer properties of the underlying graph.

In all of these scenarios, sampling (implicitly or explicitly) is used!



Understanding sampling is important!

A little story of not so long ago..

- ▶ 1999-2000: several acclaimed reports on power-law degree distribution of various networks
 - ▶ Internet: [Faloutsos et al., 1999]
 - WWW: [Albert et al., 1999]
 - Metabolic networks: [Jeong et al., 2000]
- ▶ 2003: it is shown empirically that the sampling procedure may induce a power-law, even if the underlying graph is not scale-free! [Lakhina et al., 2003]
- ➤ 2005: further empirical and theoretical studies support this [Achlioptas et al., 2005, Clauset and Moore, 2005]

Conclusion: it is very important to understand how biases in sampling affect results



In today's lecture

Sampling strategies

Biases of sampling strategies

Sampling General Goals

How do we measure the goodness of a sample, as well as the method of sampling?

Depends on what do we compare against:

Scale-down goal: We want the sample graph S to have similar properties as the original G

Back-in-time goal: We want the sample graph S to be similar to what G looked like back in the time when it had same size of S

Goals

- 1. Sample a representative subgraph (scale-down goal)
 - that is, obtain a subgraph that has similar properties, for a set of representative properties simultaneously (e.g.: degree distribution, clustering coefficient, community structure, etc.)
- 2. Estimation of a network parameter (back-in-time goal)
 - ► E.g.: average degree of nodes, diameter, ...
- 3. Estimate node attributes (back-in-time goal)
 - ▶ E.g.: age of users in a social network
- 4. Estimate edge attributes (back-in-time goal)
 - ▶ E.g.: relationship type of friends in a social network

Different sampling strategies will work for certain goals better than others



Overview of sampling strategies

From [Leskovec and Faloutsos, 2006, Maiya and Berger-Wolf, 2011, Ahmed et al., 2014]

- Random node selection
 - Only possible when access to entire graph is given
- Random edge selection
 - Only possible when access to entire graph is given
- Crawling-based
 - Snowball sampling: BFS, DFS, Forest Fire, ...
 - Random walks

[A spoiler note: For scale-down sampling goal best performers are based on random walks, since these are biased towards high degree nodes and guarantee connectivity. For back-in-time goal: Forest-fire, PageRank sampling of nodes; these mimic the temporal evolution of the graph]



Random node selection

Several possibilities

- Uniform node sampling
- Degree-based sampling [Adamic et al., 2001]
 - Probability of visiting node proportional to its degree (assumed known)
 - Originally used for searching [Adamic et al., 2001]
- Pagerank-based sampling [Leskovec and Faloutsos, 2006]
 - Probability of visiting node proportional to its pagerank (assumed known)

Random edge selection

Several possibilities

- Uniform edge sampling
 - sample edges and then include incident nodes
- Random node-edge sampling
 - select node uniformly at random, then select incident edge uniformly at random
- ▶ Hybrid sampling [Krishnamurthy et al., 2005]
 - ▶ With probability 0.8, perform random node-edge sampling
 - ▶ With probability 0.2, perform uniform edge sampling
- Induced edge sampling [Ahmed et al., 2014]
 - ▶ Uniformly sample edges
 - Complete graph sample with edges between nodes incident on sampled edges



Crawling I

a.k.a. "sampling by exploration"

- Breadth-First search (BFS)
 - explore neighbors of least recently visited nodes
- Depth-First search (DFS)
 - explore neighbors of most recently visited nodes
- ► Random walk (RW) [Gjoka et al., 2010]
 - explore neighbors of most recently visited nodes uniformly at random (no queue)
- ▶ Forest Fire sampling (FFS) [Leskovec et al., 2005]
 - probabilistic version of BFS
 - with probability p (typically 0.7), visit neighbor



Crawling II

a.k.a. "sampling by exploration"

- Expansion sampling (XS)
 [Maiya and Berger-Wolf, 2010, Maiya and Berger-Wolf, 2011]
 - greedily add node maximizing expansion $\frac{|N(S)|}{|S|}$
- Random walk with jump (RJ) [Ribeiro and Towsley, 2010]
 - same as random walk, but jump to random node with probaility p

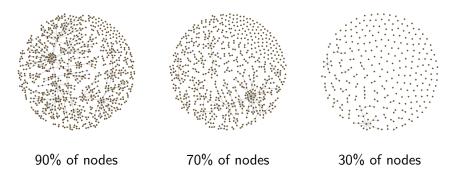
In today's lecture

Sampling strategies

Biases of sampling strategies

Uniform node sampling

- ► Induced subgraphs of scale-free networks are not scale-free [Stumpf et al., 2005]
- Induced subgraphs of connected scale-free networks are sparse



Crawled subsets of ER graphs are scale-free

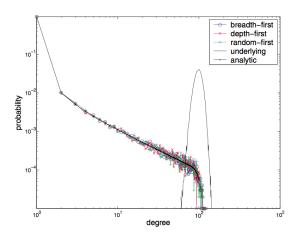
[Lakhina et al., 2003][Clauset and Moore, 2005]

[Lakhina et al., 2003] observe this empirically by sampling ER-graphs with trace-route routine (a minimum spanning tree)

[Clauset and Moore, 2005] Give a general proof of this fact (worth reading!). Basic argument is that traceroutes from single source can be modelled as a spanning tree. Then show that building a spanning tree in Erdos-Renyi graph gives subgraph with degree distribution following a power law of the form $P(k) \approx k^{-1}$

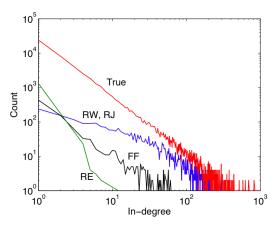
Crawled subsets of ER graphs are scale-free

[Clauset and Moore, 2005]



More crawling biases,

In general, random walks, DFS, and BFS lead to over-sampling of high-degree nodes

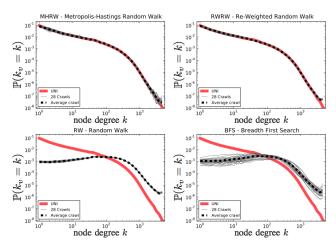


Compensating for RW bias

- Random Walk (RW)
 - Nodes with high degree are over-represented since probability of visiting a node $v \propto k_v$
- Re-Weighted random walk (RWRW)
 - Hansen-Hurwitz estimator for non-uniform selection probabilities
 - After the walk, re-weight $\hat{p}(k) = rac{\sum_{v:k_v=k} 1/k_v}{\sum_v 1/k_v}$
- Metropolis-Hastings random walk (MHRW)
 - ▶ Walk with new transition probabilities $P_{v \to w} = \frac{1}{k_v} min(1, \frac{k_v}{k_w})$
 - i.e. select random neighbor, and move with probability $min(1, \frac{k_v}{k_w})$
 - i.e. always accept moves to nodes of lower degree, reject some moves to nodes of higher degree
 - results in uniform probabilities of visiting nodes

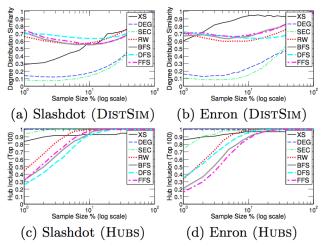


Uniform sampling of Facebook users using random walks [Gjoka et al., 2010]



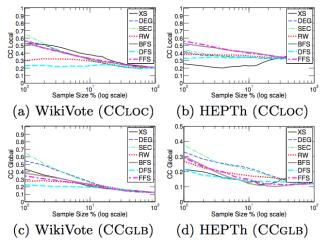
Results from [Maiya and Berger-Wolf, 2011]

Degree distribution



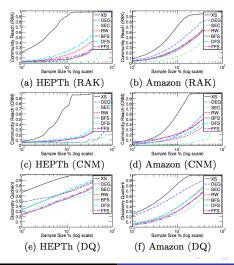
Results from [Maiya and Berger-Wolf, 2011]

Clustering coefficient



Results from [Maiya and Berger-Wolf, 2011]

Network reach



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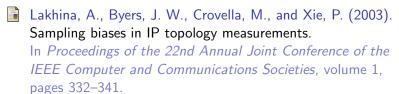


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