

# Notes on Stochastic Network Modeling (SNM)

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		<b>Preface</b>	
		These notes were prepared for the subject <i>Stochastic Network Mod-</i>	
		<i>eling (SNM)</i> from the <i>Master in Innovation and Research in Infor-</i>	
		<i>matics (MIRI)</i> , Universitat Politècnica de Catalunya.	
		Further reading: The book of Mor [3], Nelson [7] or Trivedi [8]	
		are excellent books on Markov Chains and cover most of the course.	
		The book of Kemeny [4] covers absorbing DTMC.	



## Part I

### Introduction

## Chapter 1

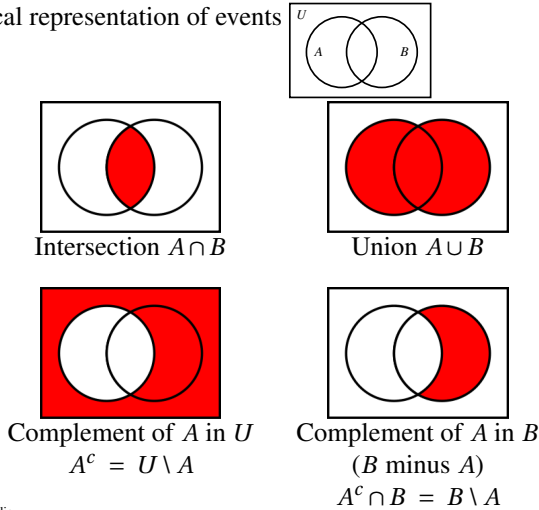
### Probability Review

#### 1.1 Ingredients of Probability

- **Random experiment**, e.g. toss a die.
- **Outcome**,  $\omega$ , e.g. tossing a die can be  $\omega = 2$ , choosing a fruit can be  $\omega = \text{orange}$ .
- **Sample space or Universal set**,  $U$ , set of all possible outcomes. E.g. tossing a die  $U = \{1, 2, 3, 4, 5, 6\}$ .
- **Event**,  $A$ , any subset of  $U$  (e.g. tossing a die  $A = \{1, 2, 3\}$ ). We say the event  $A$  occurs if the outcome of the experiment  $\omega \in A$ .  $U$  is the **sure event**, and we represent by the empty set  $\emptyset$  an **impossible outcome**.

#### Venn Diagrams

Graphical representation of events



source: wikipedia

#### Random Variable

- For simplicity it is defined a **random variable (RV)**,  $X$  as a function that assigns a real number to each outcome in the sample space  $U$ , i.e.:

$$X : U \rightarrow \mathbb{R} \quad (1.1)$$

- We will represent the experiment by a RV,  $X$ , and the possible outcomes by its values.  $X = x_i$  is the **outcome**  $X(\omega_i) = x_i$ .
- Using RVs the sample space is mapped in a subset of  $\mathbb{R}$ . So, in terms of  $X$ ,  $U$  is a set of points of  $\mathbb{R}$ . The same for any event.
- Normally the definition of  $X$  comes naturally from the experiment, e.g. tossing a die:  $X = \{\text{number in the toss}\}$ .
- RVs can be **discrete** (e.g. tossing a die) or **continuous** (e.g. waiting time of a packet in a queue).

#### Probability Measure

- If the sample space  $U$  of the RV  $X$  is **finite (discrete RV)**,  $U = \{x_1, \dots, x_n\}$ , a **probability measure** is an assignment of numbers  $P(x_i)$ , referred to as **probabilities**, to each **outcome**  $x_i$  such that:

$$0 \leq P(x_i) \leq 1$$

$$P(A) = \sum_{x_i \in A} P(x_i) \quad (1.2)$$

$$P(U) = 1$$

E.g. tossing a fair die,

$$P(x_i) = 1/6$$

$$P(X \in \{2, 4, 6\}) = \frac{1}{6} + \frac{1}{6} + \frac{1}{6} = \frac{1}{2} \quad (1.3)$$

- If the sample space of the RV  $X$  is continuous (**continuous RV**), the events are intervals of  $\mathbb{R}$ . The probability measure is defined by means of the **cumulative distribution function, CDF**:

$$F(x) = P(X \in (-\infty, x]) = P(X \leq x) \quad (1.4)$$

- $X$  is called absolutely continuous<sup>1</sup> if there exists the **probability density function, PDF**, such that for any interval  $I = \{x \mid a \leq x \leq b\}$ :

$$\int_a^b f(x) dx = P(X \in I) = F(b) - F(a) \quad (1.5)$$

#### Conditional Probability and Bayes Formula

- Given the the sample space  $U$  and the **events**  $A, B \in U$  with  $P(B) > 0$  the **probability of A conditioned by B** is defined as:

$$P(A|B) = \frac{P(A \cap B)}{P(B)} \quad (1.6)$$

NOTE: It's common to use commas to denote set intersection, and write  $P(A \cap B)$  as  $P(A, B)$ .

- **Bayes Formula**

$$P(A|B) P(B) = P(B|A) P(A) \Rightarrow P(A|B) = \frac{P(B|A) P(A)}{P(B)} \quad (1.7)$$

#### Law of total probability

- Let  $B_i$  a **partition** of the sample space  $U$  ( $\cup_i B_i = U$ ,  $B_i \cap B_j = \emptyset, \forall i \neq j$ ), then

$$P(A) = \sum_i P(A|B_i) P(B_i) \quad (1.8)$$

- For **conditional probabilities**:

$$P(A|C) = \sum_i P(A|C \cap B_i) P(B_i|C) \quad (1.9)$$

- If  $C$  is **independent** of any of the  $B_i$

$$P(A|C) = \sum_i P(A|C \cap B_i) P(B_i) \quad (1.10)$$

#### 1.2 Expected value

- Given the discrete  $N \in \mathbb{Z}$ , respectively continuous  $X \in \mathbb{R}$  RV, the **expected value** is:

$$E[N] = \sum_{k=-\infty}^{\infty} k P(N = k)$$

$$E[X] = \int_{-\infty}^{\infty} x f(x) dx \quad (1.11)$$

<sup>1</sup>Some special distributions, called singular, do not have a PDF. One example is the Cantor distribution (see Wikipedia).

**Example** A number  $X_1 \in \{1, 2, \dots, 6\}$  is obtained tossing a dice. Then, a number  $X_2 \in [0, \infty]$  is obtained exponentially distributed with parameter  $X_1$ . Compute  $f(x_1, x_2)$ ,  $f(x_2)$  and  $E[X_2]$ .  
Note: Exponential distribution with parameter  $\alpha$ :

$$f(x) = \alpha e^{-\alpha x}, x \in [0, \infty], E[X] = \frac{1}{\alpha}. \quad (1.12)$$

**Solution:**

$$f(x_1, x_2) = f(x_2|x_1) P(x_1) = x_1 e^{-x_1 x_2} \frac{1}{6}, \begin{cases} x_1 \in \{1, 2, \dots, 6\} \\ x_2 \in [0, \infty] \end{cases}$$

$$f(x_2) = \sum_{x_1} f(x_2|x_1) P(x_1) = \frac{1}{6} \sum_{n=1}^6 n e^{-n x_2}, x_2 \in [0, \infty]$$

$$E[X_2] = \frac{1}{6} \sum_{n=1}^6 \int_{x_2=0}^{\infty} x_2 n e^{-n x_2} = \frac{1}{6} \sum_{n=1}^6 \frac{1}{n} = \frac{49}{120}$$

### 1.3 Variance

- The amount of dispersion of a RV  $X$  with expected value  $\mu = E[X]$  is measured by the **Variance**:

$$\text{Var}(X) = \sigma^2 = E[(X - \mu)^2] = E[X^2] - \mu^2 \quad (1.13)$$

- Often it is used the **standard deviation**  $\sigma = \sqrt{\text{Var}(X)}$ .

### 1.4 Indicator Function

$$I(A) = \begin{cases} 1, & \text{if } A \text{ occurs,} \\ 0, & \text{otherwise.} \end{cases} \quad (1.14)$$

Therefore:

$$E[I(A)] = 0 \times P(I(A) = 0) + 1 \times P(I(A) = 1) = P(A) \quad (1.15)$$

### 1.5 Expected value of non negative RVs

- For **non negative** RVs,  $N \geq 0$  discrete and  $X \geq 0$  continuous:

$$E[N] = \sum_{k=0}^{\infty} k P(N = k) = \sum_{k=0}^{\infty} P(N > k)$$

$$E[X] = \int_0^{\infty} x f(x) dx = \int_0^{\infty} P(X > x) dx = \int_0^{\infty} (1 - F(x)) dx \quad (1.16)$$

$$N = \sum_{k=0}^{N-1} 1 = \sum_{k=0}^{\infty} I(N > k) \quad (1.17)$$

$$X = \int_0^X dx = \int_0^{\infty} I(X > x) dx$$

and take expectations.

### 1.6 Wald's Equation

- Definition:** An positive integer RV  $N > 0$  is a **stopping time** of a sequence  $X_1, X_2, \dots$  if the event  $N = n$  is independent of  $X_{n+1}, X_{n+2}, \dots$ .

E.g. toss a die until you get 6. Let  $N$  be the number of tosses.  $N$  does not depend on the values obtained after getting 6.

- Wald's Equation** If  $X_1, X_2, \dots$  are independent and identically distributed and  $N$  is a stopping time:

$$E \left[ \sum_{n=1}^N X_n \right] = E[X] E[N] \quad (1.18)$$

*Proof.*

$$\begin{aligned} E \left[ \sum_{n=1}^N X_n \right] &= E \left[ \sum_{n=1}^{\infty} X_n I(n \leq N) \right] = \\ &= \sum_{n=1}^{\infty} E[X_n] E[I(n \leq N)] = \\ &= E[X] \sum_{n=1}^{\infty} P(n \leq N) = \\ &= E[X] \sum_{n=0}^{\infty} P(N > n) = E[X] E[N] \quad \square \end{aligned}$$

### 1.7 Probability in $\mathbb{R}^k$

If we have a set of  $k$  RV  $\mathbf{X} = (X_1, \dots, X_k)$  taking values in  $\mathbb{R}^k$  ( $\mathbf{X} \in \mathbb{R}^k$ ), we define the **joint distribution**:

- Discrete RV

$$P(\mathbf{x}) = P(x_1, \dots, x_k) = P(X_1 = x_1, \dots, X_k = x_k) \quad (1.19)$$

- Continuous RV:

- **cumulative distribution function, CDF:**

$$F(\mathbf{x}) = F(x_1, \dots, x_k) = P(X_1 \in (-\infty, x_1], \dots, X_k \in (-\infty, x_k]) \quad (1.20)$$

- with **joint density** function  $f(\mathbf{x}) = f(x_1, \dots, x_k)$  (if exists):

$$\begin{aligned} F(\mathbf{x}) &= F(x_1, \dots, x_k) = \int_{-\infty}^{x_1} \dots \int_{-\infty}^{x_k} f(x_1, \dots, x_k) dx_k \dots dx_1 \\ f(\mathbf{x}) &= f(x_1, \dots, x_k) = \frac{\partial^k F(x_1, \dots, x_k)}{\partial x_1 \dots \partial x_k} \end{aligned} \quad (1.21)$$

**Marginal distributions** Let  $\mathbf{X} = (X_1, X_2)$ , where  $\mathbf{X} \in \mathbb{R}^k, X_1 \in \mathbb{R}^r, X_2 \in \mathbb{R}^{k-r}, 1 \leq r < k$ :

- Discrete RV

$$P(\mathbf{x}_2) = \sum_{x_1} \dots \sum_{x_r} P(\mathbf{x}_1, \mathbf{x}_2) \quad (1.22)$$

- Continuous RV

$$f(\mathbf{x}_2) = \int_{x_1} \dots \int_{x_r} f(\mathbf{x}_1, \mathbf{x}_2) dx_1 \dots dx_r \quad (1.23)$$

### Independent RV

- Discrete RV

$$\begin{aligned} P(\mathbf{x}) &= P(x_1, \dots, x_k) = \\ &= P(X_1 = x_1, \dots, X_k = x_k) = \\ &= P(X_1 = x_1) \dots P(X_k = x_k) \end{aligned} \quad (1.24)$$

- Continuous RV

$$\begin{aligned} F(\mathbf{x}) &= F(x_1, \dots, x_k) = F_{X_1}(x_1) \dots F_{X_k}(x_k) \\ f(\mathbf{x}) &= f(x_1, \dots, x_k) = f_{X_1}(x_1) \dots f_{X_k}(x_k) \end{aligned} \quad (1.25)$$

### Conditional Distribution

- Let  $\mathbf{X} = (X_1, X_2)$ , where  $\mathbf{X} \in \mathbb{R}^k$ ,  $X_1 \in \mathbb{R}^r$ ,  $X_2 \in \mathbb{R}^{k-r}$ , the  $r$ -dimensional distribution of  $X_1$  **conditioned by**  $X_2 = \mathbf{x}_2$ ,  $P(\{X_2 = \mathbf{x}_2\}) > 0$  is:

$$F(X_1|X_2) = P(X_1 \leq \mathbf{x}_1 | X_2 = \mathbf{x}_2) = \frac{P(X_1 \leq \mathbf{x}_1, X_2 = \mathbf{x}_2)}{P(X_2 = \mathbf{x}_2)}.$$

If  $\mathbf{X}$  is **discrete** with probability  $P(\mathbf{x}_1, \mathbf{x}_2)$  or absolutely **continuous** with density  $f(\mathbf{x}_1, \mathbf{x}_2)$ :

$$\begin{aligned} P(\mathbf{x}_1|\mathbf{x}_2) &= \frac{P(\mathbf{x}_1, \mathbf{x}_2)}{P(\mathbf{x}_2)} \\ f(\mathbf{x}_1|\mathbf{x}_2) &= \frac{f(\mathbf{x}_1, \mathbf{x}_2)}{f(\mathbf{x}_2)} \end{aligned} \quad (1.26)$$

**Composition of marginals and conditionals** Using the previous formulas we can compute ( $\mathbf{X}$  can be a mixture of discrete and continuous RV):

#### Law of total probability

- If  $\mathbf{x}_1, \mathbf{x}_2$  are **discrete** RV:  $P(\mathbf{x}_2) = \sum_{\mathbf{x}_1} P(\mathbf{x}_2|\mathbf{x}_1) P(\mathbf{x}_1)$
- If  $\mathbf{x}_1$  is **discrete** and  $\mathbf{x}_2$  is **cont.**:  $f(\mathbf{x}_2) = \sum_{\mathbf{x}_1} f(\mathbf{x}_2|\mathbf{x}_1) P(\mathbf{x}_1)$
- If  $\mathbf{x}_1, \mathbf{x}_2$  are **cont.**:  $f(\mathbf{x}_2) = \int_{\mathbf{x}_1} f(\mathbf{x}_2|\mathbf{x}_1) f(\mathbf{x}_1) d\mathbf{x}_1$
- If  $\mathbf{x}_1$  is **cont.** and  $\mathbf{x}_2$  is **discrete**:  $P(\mathbf{x}_2) = \int_{\mathbf{x}_1} P(\mathbf{x}_2|\mathbf{x}_1) f(\mathbf{x}_1) d\mathbf{x}_1$

### Conditional expected value

- Given  $X \in \mathbb{R}$ ,  $Y \in \mathbb{R}^k$  with density  $f(x, y)$ :

$$\begin{aligned} E[X | Y = \mathbf{y}] &= \int_{\mathbb{R}} x f(x|\mathbf{y}) dx \\ E[X] &= \int_{\mathbb{R}^k} E[X | Y = \mathbf{y}] f(\mathbf{y}) d\mathbf{y} \end{aligned} \quad (1.27)$$

where the **marginal**  $f(\mathbf{y}) = \int_{x=-\infty}^{\infty} f(x, y) dx$  and the **conditional**  $f(x|\mathbf{y}) = f(x, y)/f(\mathbf{y})$ .

Thus, the law of total probability also applies to expected value, and it is known as **law of total expectation**.

## Chapter 2

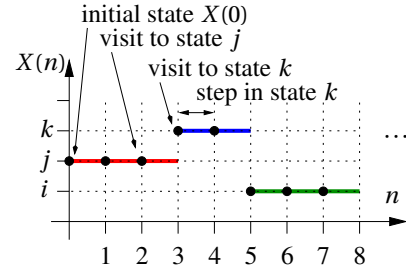
### Stochastic Process (SP)

#### 2.1 Introduction

- Sequence of RVs**  $\{X(t)\}_{t \geq 0}$ .
- $X(t)$  is the **state** at time  $t$ .
- The **state**  $X(t)$  can be **continuous** or **discrete**.
- The **index** can be **continuous** or **discrete**. We shall use  $n$  for the **index**, and refer to it as **steps** when it is **discrete**, and  $t$  and refer to it as **time** when it is **continuous**.
- We call a possible sequence of states of the SP the **sample function** (or sample path) of the SP.

#### Sample Path

- Possible evolution (**sample path**) of a **discrete state, discrete time** SP  $\{X(n)\}_{n \geq 0}$ :



- To characterize the stochastic process we would need the distribution and **joint probabilities** of the  $\{X(n)\}_{n \geq 0}$  RVs:

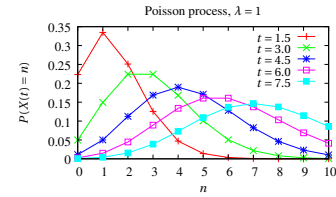
$$P(X(n) = i, X(n-1) = k, \dots, X(0) = j) \quad (2.1)$$

#### Example 2: Poisson Process

- It is a discrete state continuous time SP.
- It counts the number of events occurred in a time interval.
- Often used to build models of other stochastic processes.
- Definition:** The number of “events” in any interval of length  $t$ ,  $X(t)$ , is **Poisson distributed** with mean  $\lambda t$ , i.e.

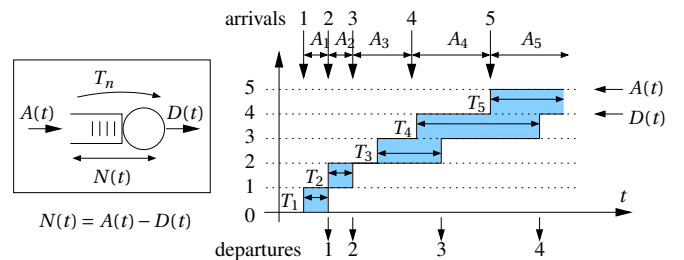
$$\begin{aligned} P(X(t+s) - X(s) = n) &= P(X(t) - X(0) = n) = \\ P(X(t) = n) &= e^{-\lambda t} \frac{(\lambda t)^n}{n!} \end{aligned} \quad (2.2)$$

where we assume  $X(0) = 0$ .



#### Example 3: Queue with Poisson Arrivals

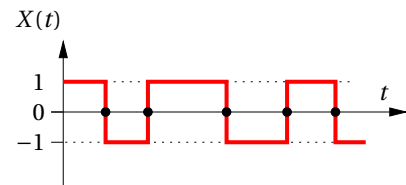
- The queue arrivals,  $A(t)$ , are modeled as a **Poisson process** with mean  $\lambda t$ . Each event model an arrival.



#### Example 4: Telegraph signal

- The signal is modeled as a **Poisson process** with mean  $\lambda t$  such that  $X(0) = 1$  or  $X(0) = -1$  with equal probability of  $1/2$  and:

$$X(t) = \begin{cases} 1 & \text{if the number of events in } (0, t] \text{ is even} \\ -1 & \text{if the number of events in } (0, t] \text{ is odd} \end{cases} \quad (2.3)$$



## 2.2 Analysis of Stochastic Processes

- **Signal Theory:** Normally interested in the **spectral analysis** of the signal. The basic tool is the **Fourier transform** of the **auto-correlation function** of the process (**energy spectral density**). We will not do this analysis.

$$R(t) = E[X(\tau)X(\tau - t)] \quad \text{autocorrelation}$$

$$F(f) = \mathcal{F}[R(t)] = \int_{-\infty}^{\infty} R(t) e^{-j2\pi f t} dt \quad \text{Fourier transform}$$

$$(2.4) \quad \text{(energy spectral density)}$$

- **Computer Networks:** Normally interested in probabilistic models using **Markov Chains** and **Queueing Theory**.

## Part II

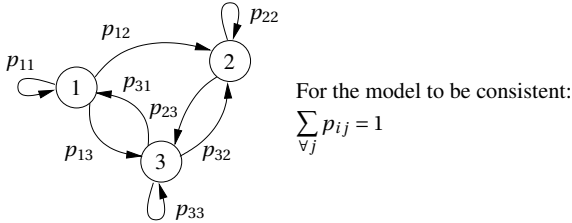
## Discrete Time Markov Chains

### Chapter 3

### Definition of a DTMC

#### 3.1 State Transition Diagram

- We are interested in a **process that evolve in stages**.
- For the model to be tractable, it is convenient to represent the SP by giving all **possible states** (there may be  $\infty$ ), and the **possible transitions** between them:



- Mathematically:

$$p_{ij} = P(X(n) = j | X(n-1) = i) \quad (3.1)$$

#### 3.2 Properties of a DTMC

- The event  $X(n) = i$  (at step  $n$  the system is in state  $i$ ) must satisfy (**memoryless property**):

$$P(X(n) = j | X(n-1) = i, X(n-2) = k, \dots) = P(X(n) = j | X(n-1) = i) \quad (3.2)$$

- If  $P(X(n) = j | X(n-1) = i) = P(X(1) = j | X(0) = i)$  for any  $n$  we have an **homogeneous** DTMC. We shall only consider homogeneous DTMC.
- We call **one-step transition probabilities** to:

$$p_{ij} = P(X(n) = j | X(n-1) = i) \quad (3.3)$$

- The SP is called a Markov Process (MP) or Markov Chain (MC) depending on the state being continuous or discrete.

#### 3.3 Transition Matrix

- Transition probabilities:

$$p_{ij} = P(X(n) = j | X(n-1) = i) \quad (3.4)$$

- In matrix form:

$$\mathbf{P} = \begin{bmatrix} p_{11} & p_{12} & \dots \\ p_{21} & p_{22} & \dots \\ \dots & \dots & \dots \end{bmatrix} \quad (3.5)$$

- For the model to be consistent, the probability to move from  $i$  to any state must be 1. Mathematically:

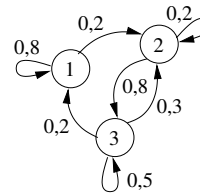
$$\sum_{\forall j} p_{ij} = \sum_{\forall j} P(X(n) = j | X(n-1) = i) = \sum_{\forall j} \frac{P(X(n-1) = i | X(n) = j) P(X(n) = j)}{P(X(n-1) = i)} = \frac{P(X(n-1) = i)}{P(X(n-1) = i)} = 1 \quad (3.6)$$

- $\mathbf{P}$  is a **stochastic matrix**, i.e. a matrix which rows sum 1.

#### Example

- Assume a terminal can be in **3 states**:

- State 1: Idle.
- State 2: Active without sending data.
- State 3: Active and sending data at a rate  $\nu$  bps.



		to state			
		1	2	3	
$\mathbf{P} =$	1	0,8	0,2	0	from state
	2	0	0,2	0,8	
	3	0,2	0,3	0,5	

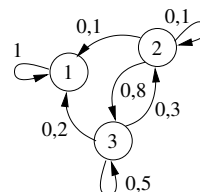
(3.7)

- The **average transmission rate** (throughput),  $\nu_a$ , is:

$$\nu_a = P(\text{the terminal is in state 3}) \times \nu \quad (3.8)$$

#### 3.4 Absorbing Chains

- It is possible to have chains with **absorbing states**.
- A state  $i$  is absorbing if  $p_{ii} = 1$ .
- Example: State 1 is absorbing.

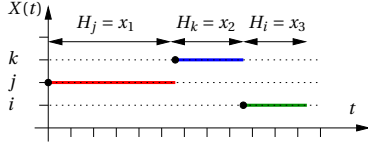


		to state			
		1	2	3	
$\mathbf{P} =$	1	1	0	0	from state
	2	0,1	0,1	0,8	
	3	0,2	0,3	0,5	

(3.9)

### 3.5 Sojourn or Holding Time

- **Sojourn or holding time** in state  $k$ : Is the RV  $H_k$  equal to the number of steps that the chain remains in state  $k$  before leaving to a different state:



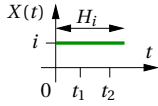
- The Markov property implies:

$$H_i(n) = P(H_i = n) = p_{ii}^{n-1} (1 - p_{ii}), n \geq 1 \quad (3.10)$$

- Which is a **geometric** distribution with mean:

$$E[H_i] = \sum_{n=1}^{\infty} n P(H_i = n) = \frac{1}{1 - p_{ii}}. \quad (3.11)$$

The geometric distribution satisfies the Markov property



*Proof.*

- Markov property:  
 $P(X(n_2) = i | X(n_1) = i, X(n_0) = i) = P(X(n_2) = i | X(n_1) = i)$
- Thus, the Markov property **in terms of the sojourn time** can be written as:

$$P(H_i > n_2 - n_0 | H_i > n_1 - n_0) = P(H_i > n_2 - n_1) \quad (3.12)$$

- Since

$$P(H_i > k) = 1 - P(H_i \leq k) = 1 - \sum_{n=1}^k p^{n-1} (1 - p) = 1 - (1 - p) \frac{1 - p^k}{1 - p} = p^k \quad (3.13)$$

- We have:

$$\begin{aligned} P(H_i > n_2 - n_0 | H_i > n_1 - n_0) &= \frac{P(H_i > n_2 - n_0, H_i > n_1 - n_0)}{P(H_i > n_1 - n_0)} = \\ &= \frac{P(H_i > n_2 - n_0)}{P(H_i > n_1 - n_0)} = \frac{p^{n_2 - n_0}}{p^{n_1 - n_0}} = p^{n_2 - n_1} = P(H_i > n_2 - n_1) \end{aligned} \quad (3.14)$$

□

### 3.6 n-step transition probabilities

- Transition probabilities:  $p_{ij} = P(X(n) = j | X(n-1) = i)$
- In matrix form:

$$\mathbf{P} = \begin{bmatrix} p_{11} & p_{12} & \cdots \\ p_{21} & p_{22} & \cdots \\ \cdots & \cdots & \cdots \end{bmatrix} \quad (3.15)$$

- We define the **n-step** transition probabilities:

$$p_{ij}(n) = P(X(n) = j | X(0) = i) \quad (3.16)$$

$$\mathbf{P}(n) = \begin{bmatrix} p_{11}(n) & p_{12}(n) & \cdots \\ p_{21}(n) & p_{22}(n) & \cdots \\ \cdots & \cdots & \cdots \end{bmatrix} \quad (3.17)$$

- $\mathbf{P}$  and  $\mathbf{P}(n)$  are **stochastic matrices**: Their rows sum 1.

### 3.7 State Probabilities

- Define the probability of being in state  $i$  at step  $n$ :

$$\pi_i(n) = P(X(n) = i) \quad (3.18)$$

- In vector form (row vector)

$$\boldsymbol{\pi}(n) = (\pi_1(n), \pi_2(n), \dots) = (P(X(n) = 1), P(X(n) = 2), \dots). \quad (3.19)$$

- Thus, the vector  $\boldsymbol{\pi}(n)$  is the distribution of the random variable  $X(n)$ , and it is called the **state probability at step  $n$** .

- Law of total prob.  $P(A) = \sum_n P(A \cap B_n) = \sum_n P(A | B_n) P(B_n)$ :

$$\begin{aligned} \pi_i(n) &= \sum_k P(X(n-1) = k) P(X(n) = i | X(n-1) = k) = \sum_k \pi_k(n-1) p_{ki} \\ \pi_i(n) &= \sum_k P(X(0) = k) P(X(n) = i | X(0) = k) = \sum_k \pi_k(0) p_{ki}(n) \end{aligned} \quad (3.20)$$

- In matrix form:

$$\begin{aligned} \boldsymbol{\pi}(n) &= \boldsymbol{\pi}(n-1) \mathbf{P} \\ \boldsymbol{\pi}(n) &= \boldsymbol{\pi}(0) \mathbf{P}(n) \end{aligned} \quad (3.21)$$

where  $\boldsymbol{\pi}(0)$  is the **initial distribution**.

- Iterating

$$\boldsymbol{\pi}(n) = \boldsymbol{\pi}(n-1) \mathbf{P} = \boldsymbol{\pi}(n-2) \mathbf{P} \mathbf{P} = \boldsymbol{\pi}(n-3) \mathbf{P} \mathbf{P} \mathbf{P} = \dots = \boldsymbol{\pi}(0) \mathbf{P}^n \quad (3.22)$$

- Thus:

$$\boldsymbol{\pi}(n) = \boldsymbol{\pi}(0) \mathbf{P}(n) = \boldsymbol{\pi}(0) \mathbf{P}^n \quad (3.23)$$

### 3.8 Chapman-Kolmogorov Equations

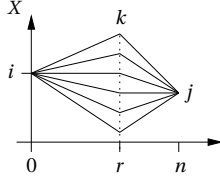
$$p_{ij}(n) = \sum_k p_{ik}(r) p_{kj}(n-r) \quad (3.24)$$

*Proof.*

$$\begin{aligned} p_{ij}(n) &= P(X(n) = j | X(0) = i) = \\ &= \sum_k P(X(n) = j, X(r) = k | X(0) = i) \\ &= \sum_k \frac{P(X(n) = j, X(r) = k, X(0) = i)}{P(X(0) = i)} \times \\ &= \sum_k \frac{P(X(r) = k, X(0) = i)}{P(X(r) = k, X(0) = i)} \times \\ &= \sum_k P(X(n) = j | X(r) = k, X(0) = i) \times \\ &= \sum_k P(X(r) = k | X(0) = i) \times \\ &= \sum_k P(X(n) = j | X(r) = k) P(X(r) = k | X(0) = i) \\ &= \sum_k p_{ik}(r) p_{kj}(n-r) \quad \square \end{aligned}$$



- Graphical interpretation:



- In matrix form:

$$\mathbf{P}(n) = \mathbf{P}(r) \mathbf{P}(n-r) \quad (3.25)$$

- Particularly:

$$\mathbf{P}(n) = \mathbf{P}(1) \mathbf{P}(n-1) = \mathbf{P} \mathbf{P}(n-1) = \mathbf{P}(n-1) \mathbf{P} \quad (3.26)$$

- Iterating:

$$\mathbf{P}(n) = \mathbf{P}^n \quad (3.27)$$

- Thus:

$$\boldsymbol{\pi}(n) = \boldsymbol{\pi}(0) \mathbf{P}(n) = \boldsymbol{\pi}(0) \mathbf{P}^n \quad (3.28)$$

## Chapter 4

### Transient Solution

#### 4.1 Close Form Solution

- If we are interested in the **transient evolution** we shall study  $\boldsymbol{\pi}(n) = \boldsymbol{\pi}(0) \mathbf{P}^n$ .
- If we can **diagonalize**  $\mathbf{P}$ , we can obtain the transient evolution in **close form**.
- $\mathbf{P}$  can be **diagonalized** if  $\mathbf{P}$  can be decomposed as:

$$\mathbf{P} = \mathbf{L}^{-1} \boldsymbol{\Lambda} \mathbf{L} \quad (4.1)$$

where  $\mathbf{L}$  is some invertible matrix and  $\boldsymbol{\Lambda}$  is the diagonal matrix

$$\boldsymbol{\Lambda} = \text{diag}(\lambda_1, \dots, \lambda_N) = \begin{bmatrix} \lambda_1 & 0 & \dots & 0 \\ 0 & \lambda_2 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & \lambda_N \end{bmatrix} \quad (4.2)$$

with  $\lambda_l$ ,  $l = 1, \dots, N$  the **eigenvalues** of  $\mathbf{P}$ .

- Assume a **finite DTMC** with  $N$  states. Then  $\mathbf{P} = \mathbf{P}^{N \times N}$ .
- Assume that  $\mathbf{P}$  can be **diagonalized**:  $\mathbf{P} = \mathbf{L}^{-1} \boldsymbol{\Lambda} \mathbf{L}$ , where  $\boldsymbol{\Lambda}$  is the diagonal matrix  $\boldsymbol{\Lambda} = \text{diag}(\lambda_1, \dots, \lambda_N)$ , with  $\lambda_l$ ,  $l = 1, \dots, N$  the eigenvalues of  $\mathbf{P}$ .
- But  $\mathbf{L}^{-1} \text{diag}(\lambda_1^n, \dots, \lambda_N^n) \mathbf{L}$  are linear combinations of  $\lambda_1^n, \dots, \lambda_N^n$ . Thus, the probability of being in state  $i$  is given by:

$$\pi_i(n) = (\boldsymbol{\pi}(n))_i = \sum_{l=1}^N a_i^{(l)} \lambda_l^n \quad (4.3)$$

where the **unknown coefficients**  $a_i^{(l)}$  can be obtained solving the system of equations:

$$\sum_{l=1}^N a_i^{(l)} \lambda_l^n = (\boldsymbol{\pi}(n))_i = (\boldsymbol{\pi}(0) \mathbf{P}^n)_i, \quad n = 0, \dots, N-1 \quad (4.4)$$

#### 4.2 Eigenvalues

- The **eigenvalues**  $\lambda_l$  of a matrix  $\mathbf{A}$  are scalars that satisfy:  $\mathbf{A} \mathbf{l} = \lambda_l \mathbf{l}$  (or  $\mathbf{A} \mathbf{r} = \lambda_l \mathbf{r}$ ) for some row vectors  $\mathbf{l}$  (column vectors  $\mathbf{r}$ ), referred to as **left and right eigenvectors**, respectively.

$$\begin{aligned} \mathbf{A} \mathbf{l} = \lambda_l \mathbf{l} &\Rightarrow \mathbf{l}(\mathbf{A} - \lambda_l \mathbf{I}) = 0 \Rightarrow \det(\lambda_l \mathbf{I} - \mathbf{A}) = 0 \\ \mathbf{A} \mathbf{r} = \lambda_l \mathbf{r} &\Rightarrow (\mathbf{A} - \lambda_l \mathbf{I}) \mathbf{r} = 0 \Rightarrow \det(\lambda_l \mathbf{I} - \mathbf{A}) = 0 \end{aligned} \quad (4.5)$$

- Thus,  $\lambda_l$  solve the **characteristic polynomial**  $\det(\lambda \mathbf{I} - \mathbf{A}) = 0$ .
- Note that, in general, **left and right eigenvectors** are different, but eigenvalues are the same (they solve the same **characteristic polynomial**).
- A matrix can be **diagonalized** if all eigenvalues are single (multiplicity = 1). If a matrix cannot be diagonalized it is called **defective**.

#### 4.3 Determinant

$$\det \begin{bmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{bmatrix} = a_{11} a_{22} - a_{12} a_{21} \quad (4.6)$$

$$\det \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{bmatrix} = \begin{aligned} &+a_{11} a_{22} a_{33} + a_{12} a_{23} a_{31} + a_{21} a_{32} a_{13} \\ &-a_{31} a_{22} a_{13} - a_{12} a_{21} a_{33} - a_{23} a_{32} a_{11} \end{aligned} \quad (4.7)$$

- Cofactor Formula**: expanding along a row  $i$ :

$$\det \mathbf{A} = \sum_{j=1}^N a_{ij} (-1)^{i+j} \det M_{ij}, \quad (4.8)$$

where the **minor matrices**  $M_{ij}$  are obtained removing the row  $i$  and column  $j$  from  $\mathbf{A}$ .  $(-1)^{i+j} \det M_{ij}$  is called the **cofactor** of  $a_{ij}$ .

$$\det \mathbf{A} = \prod \text{eigenvalues of } \mathbf{A} \quad (4.9)$$

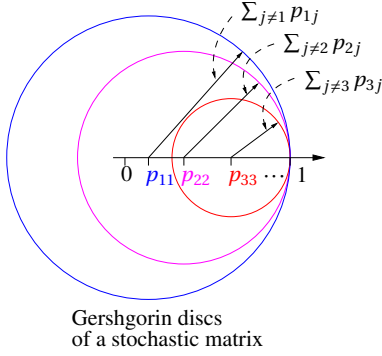
$$\text{trace } \mathbf{A} = \sum \text{eigenvalues of } \mathbf{A} \quad (4.10)$$

where  $\text{trace } \mathbf{A} = \sum$  elements of the diagonal of  $\mathbf{A}$ .

#### 4.4 Eigenvalues of a Stochastic Matrix

- $\mathbf{P}$  has an **eigenvalue equal to 1** ( $\mathbf{P} \mathbf{x} = \lambda \mathbf{x}$ , for  $\lambda = 1$ ).
- Proof**:  $\mathbf{P} \mathbf{e} = \mathbf{e}$ , where  $\mathbf{e} = \begin{bmatrix} 1 & 1 & \dots \end{bmatrix}^T$  is a column vector of 1 (all rows of  $\mathbf{P}$  add to 1).
- All eigenvalues of  $\mathbf{P}$  are  $|\lambda_l| \leq 1$ .

*Proof.* Using Gerschgorin's theorem *The eigenvalues of a matrix  $\mathbf{P}_{n \times n}$  lie within the union of the  $n$  circular disks with center  $p_{ii}$  and radius  $\sum_{j \neq i} |p_{ij}|$  in  $\mathbb{C}$ .* Since  $\sum_j p_{ij} = 1$ , the property is proved.



*Proof.* of Gerschgorin's theorem

From  $\mathbf{P}\mathbf{x} = \lambda\mathbf{x}$  we have

$$\sum_j p_{ij} x_j = \lambda x_i \quad \forall i \in \{1, \dots, n\}. \quad (4.11)$$

We choose  $i$  such that  $|x_i| = \max_j |x_j|$ . Thus,  $\sum_{j \neq i} p_{ij} x_j = \lambda x_i - p_{ii} x_i$ , and

$$|\lambda - p_{ii}| = \left| \sum_{j \neq i} p_{ij} \frac{x_j}{x_i} \right| \leq \sum_{j \neq i} \left| p_{ij} \frac{x_j}{x_i} \right| \leq \sum_{j \neq i} |p_{ij}| \quad (4.12)$$

and the equation  $|\mathbf{x} - \mathbf{c}| \leq \mathbf{r}$ ,  $\mathbf{x}, \mathbf{c} \in \mathbb{C}, \mathbf{r} \in \mathbb{R}$  is a disk of center  $\mathbf{c}$  and radius  $\mathbf{r}$  in  $\mathbb{C}$ .

#### Example

- Assume a DTMC with

$$\mathbf{P} = \begin{bmatrix} 4/5 & 1/5 \\ 2/5 & 3/5 \end{bmatrix}$$

- We want the probability of being in state 2 in  $n$  steps starting from state 1:  $\pi_2(n)$  with  $\boldsymbol{\pi}(0) = \begin{bmatrix} 1 & 0 \end{bmatrix}$ .

#### Solution

- It can be easily found that the **eigenvalues** of  $\mathbf{P}$  are  $\lambda_1 = 1$  and  $\lambda_2 = 2/5$ .

$$\pi_2(n) = \lambda_1^n a + b \lambda_2^n = a + b(2/5)^n \quad (4.13)$$

- Imposing the **boundary conditions**  $\pi_i(n) = (\boldsymbol{\pi}(0) \mathbf{P}^n)_i$ :

$$\begin{aligned} \pi_2(0) &= a + b = \begin{bmatrix} 1 & 0 \end{bmatrix} \mathbf{P}^0 = (\mathbf{P}^0)_{12} = 0 \\ \pi_2(1) &= a + b(2/5) = \begin{bmatrix} 1 & 0 \end{bmatrix} \mathbf{P}^1 = (\mathbf{P})_{12} = 1/5 \end{aligned} \quad (4.14)$$

we have that  $a = 1/3$ ,  $b = -1/3$ , thus:

$$\begin{aligned} \pi_2(n) &= 1/3 - 1/3(2/5)^n, \quad n \geq 0 \\ \pi_1(n) &= 1 - \pi_2(n) = 2/3 + 1/3(2/5)^n, \quad n \geq 0 \end{aligned} \quad (4.15)$$

## 4.5 Chain with a Defective Matrix

- What if  $\mathbf{P}$  cannot be diagonalized? (**defective** matrix).
- Let  $\lambda_l$ ,  $l = 1, \dots, L$  be the eigenvalues of  $\mathbf{P}^{N \times N}$ , each with multiplicity  $k_l$  ( $k_l \geq 1$ ,  $\sum_l k_l = N$ ), and a possible eigenvalue  $\lambda_1 = 0$  with multiplicity  $k_1$ . Then [1]:

$$\pi_j(n) = \sum_{m=0}^{k_1-1} a_j^{(1,m)} I(n=m) + \sum_{l=2}^L \lambda_l^n \sum_{m=0}^{k_l-1} a_j^{(l,m)} n^m, \quad 1 \leq j \leq N, n \geq 0 \quad (4.16)$$

$I(n=m)$  is the indicator func.:  $I(n) = 1$  if  $n = m$ ,  $I(n) = 0$  if  $n \neq m$ .

#### Example

- Assume a DTMC with

$$\mathbf{P} = \begin{bmatrix} 3/4 & 1/4 & 0 \\ 0 & 3/4 & 1/4 \\ 1 & 0 & 0 \end{bmatrix} \quad (4.17)$$

- We want the probability of being in state 1 in  $n$  steps starting from state 1:  $\pi_1(n)$  with  $\pi_1(0) = 1$ .
- It can be easily found that the **eigenvalues** of  $\mathbf{P}$  are  $\lambda_1 = 1$  and  $\lambda_2 = 1/4$  with multiplicity 2. We guess:

$$\pi_1(n) = a + 1/4^n(b + cn) \quad (4.18)$$

- Imposing  $\pi_1(0) = 1$ ,  $\pi_1(1) = 3/4$ ,  $\pi_1(2) = (3/4)^2$ , we have:

$$\pi_1(n) = \frac{4}{9} + \frac{1}{4^n} \left( \frac{5}{9} + \frac{2}{3} n \right) \quad (4.19)$$

## Chapter 5

### Classification of States

#### Objective

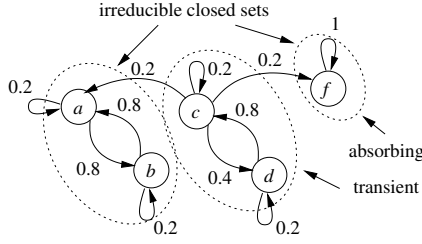
- Identify the different **types of behavior** that the chain can have.
- Introduce the concepts of **first passage probability** and **mean recurrence time**.

#### 5.1 Irreducibility

- A state  $j$  is said to **communicate** with  $i$ ,  $i \leftrightarrow j$ , if  $p_{ij}(m_1) > 0$ ,  $p_{ji}(m_2) > 0$  for some  $m_1, m_2 \geq 0$ .
- We define an **irreducible closed set, ICS**  $C_k$  as a set where all states communicate with each other, and have no transitions to other states out of the set:  
 $i \leftrightarrow j, \forall i, j \in C_k$  and  $p_{ij} = 0, \forall i \in C_k, j \notin C_k$   
 (note that for  $i \in C_k, j \notin C_k$  we have:  $p_{ij}(2) = \sum_k p_{ik} p_{kj} = 0$ , since  $p_{ik} = 0$  if  $k \notin C_k$ , and  $p_{kj} = 0$  if  $k \in C_k$ . Thus,  $p_{ij}(n) = 0, \forall n$ .)
- An **absorbing state** form an ICS of only one element. This state,  $i$ , must have  $p_{ii} = 1$ ,  $p_{ij} = 0 \forall j \neq i$ .
- Transient states** do not belong to any ICS.
- A MC is **irreducible** if all the states form a unique ICS.
- Assume a MC has  $M$  ICSs: By properly numbering the states, we can write  $\mathbf{P}$  as an  $M$  block diagonal matrix with the probabilities of the transient states in the last rows.
- Example**, if  $M = 3$ :

$$\mathbf{P} = \begin{bmatrix} \mathbf{P}_1 & \mathbf{0} \\ \mathbf{0} & \mathbf{P}_2 \\ \text{at least one } > 0 & \mathbf{T} \end{bmatrix} \Rightarrow \boldsymbol{\pi}(n) = \boldsymbol{\pi}(0) \mathbf{P}^n = \boldsymbol{\pi}(0) \begin{bmatrix} \mathbf{P}_1^n & \mathbf{0} \\ \mathbf{0} & \mathbf{P}_2^n \\ \text{at least one } > 0 & \mathbf{T}^n \end{bmatrix}$$

- Note that the  $M$  sub-matrices are **stochastic** (their rows sum 1).

**Example**

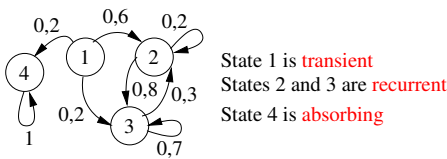
$$\mathbf{P} = \begin{array}{c|cc|cc|c} & a & b & f & c & d \\ \hline a & 0,2 & 0,8 & 0 & 0 & 0 \\ b & 0,8 & 0,2 & 0 & 0 & 0 \\ f & 0 & 0 & 1,0 & 0 & 0 \\ \hline c & 0,2 & 0 & 0,2 & 0,2 & 0,4 \\ d & 0 & 0 & 0 & 0,8 & 0,2 \end{array}$$

$$\mathbf{P}^\infty = \begin{array}{c|cc|cc|c} & a & b & f & c & d \\ \hline a & 0,5 & 0,5 & 0 & 0 & 0 \\ b & 0,5 & 0,5 & 0 & 0 & 0 \\ f & 0 & 0 & 1,0 & 0 & 0 \\ \hline c & 0,25 & 0,25 & 0,5 & 0 & 0 \\ d & 0,25 & 0,25 & 0,5 & 0 & 0 \end{array}$$

- What is the meaning of the probabilities in  $\mathbf{P}^\infty$ ? (recall that  $(\mathbf{P}^n)_{ij} = p_{ij}(n) = P(X(n) = j | X(0) = i)$ ).

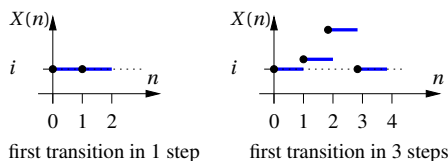
**5.2 Transient and Recurrent**

- Recurrent:** States that, being visited, they are visited again with probability 1. They are visited an infinite number of times when  $n \rightarrow \infty$ .
- Transient:** States that, being visited, have a probability  $> 0$  of never being visited again. They are visited a finite number of times when  $n \rightarrow \infty$ .
- Absorbing:** A single (recurrent) state where the chain remains with probability = 1.

**5.3 First Passage (Transition) Probabilities**

- To derive a classification criteria, we shall study the distribution of the number of steps to **go for the first time from a state  $i$  another state  $j$** . Definition:

$$f_{ji}(n) = P \left( \begin{array}{l} \text{first transition into state } i \\ \text{in } n \text{ steps starting from } j \end{array} \right) \quad (5.1)$$



- Do **not confuse** with the  $n$ -step transition probability  $p_{ii}(n)$ , where the state  $i$  can be visited in the intermediate states.

**Relation between  $f_{ii}(n)$  and  $p_{ii}(n)$** 

- $f_{ii}(n)$  and  $p_{ii}(n)$  satisfy:

$$\begin{aligned} f_{ii}(1) &= p_{ii}(1) \\ p_{ii}(n) &= \sum_{l=1}^n f_{ii}(l) p_{ii}(n-l), \quad n \geq 1 \end{aligned} \quad (5.2)$$

- The probability that the MC **eventually enters state  $i$  starting from  $i$**  is given by:

$$f_{ii} = \sum_{n=1}^{\infty} f_{ii}(n) \quad (5.3)$$

- If  $f_{ii} = 1$  we say  $i$  is a **recurrent state**.
- If  $f_{ii} < 1$  we say  $i$  is a **transient state**.

**5.4 Mean Recurrence Time**

- When  $f_{ii} = 1$ , we define the **mean recurrence time**  $m_{ii}$  as:

$$m_{ii} = \sum_{n=1}^{\infty} n f_{ii}(n) \quad (5.4)$$

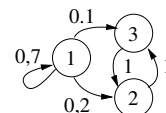
- $m_{ii}$  is the **average number of steps to eventually reach  $i$  starting from  $i$** . If  $f_{ii} < 1$  (**transient state**) then we define  $m_{ii} = \infty$ .
- Classification of **recurrent states** ( $f_{ii} = 1$ ):
  - If  $m_{ii} = \infty$  the state is **null recurrent**: it takes an  $\infty$  time to reach the state after leave it. Can only happen in chains with an infinite number of states.
  - If  $m_{ii} < \infty$  the state is **positive recurrent**: the state is reached in a finite time after leave it.

**5.5 Property of States**

In **finite MC**:

- States can be only of type positive recurrent or transient.
- At least one state must be positive recurrent.
- There are not null recurrent states.

- Example:**



- State 1 is transient. States 2 and 3 are positive recurrent.

**Generalization to Any State Pair**

- Analogously to  $f_{ii}(n)$ , we define the probability of the **first passage to state  $j$  starting from any state  $i$**  in  $n$  steps:  $f_{ij}(n)$ .
- $f_{ij}(n)$  and  $p_{ij}(n)$  satisfy:

$$p_{ij}(n) = \sum_{l=1}^n f_{ij}(l) p_{ij}(n-l), n \geq 1 \quad (5.5)$$

- When  $f_{ij} = 1$** , the average number of steps to eventually reach  $j$  starting from  $i$ ,  $m_{ij}$  is given by:

$$m_{ij} = \sum_{n=1}^{\infty} n f_{ij}(n) \quad (5.6)$$

- If state  $j$  can not be reached starting from state  $i$  with probability one (if  $f_{ij} < 1$ ), then we define  $m_{ij} = \infty$ .

**5.6 Recursive Equation for the First Passage Probabilities**

- Recall that the The probability that the MC **eventually enters state  $j$  starting from  $i$**  is given by:  $f_{ij} = \sum_{n=1}^{\infty} f_{ij}(n)$
- $f_{ij}$  can be computed as follows: Assume we are in  $i$ . With probability  $p_{ij}$  we will go to  $j$  in one step. Otherwise, we will go to  $k$ ,  $k \neq j$ , and then we will reach  $j$  with probability  $f_{kj}$ . Thus:

$$f_{ij} = p_{ij} + \sum_{k \neq j} p_{ik} f_{kj} \quad (5.7)$$

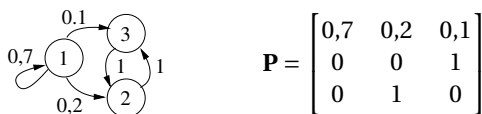
- If there are more than 1 **absorbing states**, we can compute the probability to reach them using this method (if there is only 1, say  $j$ , then  $f_{ij} = 1, \forall i$ ).

**5.7 Recursive Equation for the Mean Recurrence Time**

- Recall that the **mean recurrence time**  $m_{ij} = \sum_{n=1}^{\infty} n f_{ij}(n)$  is the average number of steps to eventually reach  $j$  starting from  $i$ , i.e. it is the mean first passage time from state  $i$  to  $j$ .
- When  $f_{ij} = 1$** ,  $m_{ij}$  can be computed as follows: Assume we are in  $i$ . With probability  $p_{ij}$  we will go to  $j$  in one step. Otherwise, we will go to  $k$ ,  $k \neq j$ , and then it will take  $m_{kj}$  steps to reach  $j$ . Thus:

$$m_{ij} = p_{ij} + \sum_{k \neq j} p_{ik} (1 + m_{kj}) = 1 + \sum_{k \neq j} p_{ik} m_{kj} \quad (5.8)$$

since  $\sum_j p_{ij} = 1$ .

**Example: Recurrence Times Using the Definition**

$$\begin{aligned} f_{21}(n) &= f_{31}(n) = 0 \\ f_{11}(n) &= 0,7 I(n=1) \\ f_{22}(n) &= f_{33}(n) = I(n=2) \\ f_{23}(n) &= f_{32}(n) = I(n=1) \end{aligned}$$

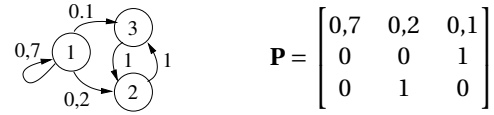
$$f_{12}(n) = \begin{cases} 0,2, & n=1 \\ 0,7^{n-1} 0,2 + 0,7^{n-2} 0,1, & n>1 \end{cases}$$

$$f_{13}(n) = \begin{cases} 0,1, & n=1 \\ 0,7^{n-1} 0,1 + 0,7^{n-2} 0,2, & n>1 \end{cases}$$

$$\begin{aligned} f_{11} &= 0,7 \\ f_{12} &= f_{13} = 1 \\ f_{32} &= f_{33} = 1 \end{aligned} \quad \begin{aligned} f_{22} &= f_{23} = 1 \\ f_{21} &= f_{31} = 0 \end{aligned}$$

$$\mathbf{M} = (m_{ij}) = \begin{bmatrix} \infty & 11/3 & 12/3 \\ \infty & 2 & 1 \\ \infty & 1 & 2 \end{bmatrix}$$

- State 1 is **transient**. States 2 and 3 are **recurrent**.

**Example: First Passage Probability Using Recursion**

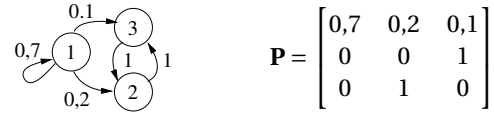
- We have:

$$f_{12} = p_{11} f_{12} + p_{12} + p_{13} f_{32} \quad (5.9)$$

- Clearly  $f_{32} = 1$ , thus:

$$f_{12} = 0,7 f_{12} + 0,2 + 0,1 \times 1 \Rightarrow f_{12} = 1 \quad (5.10)$$

as before.

**Example: Mean Recurrence Time Using Recursion**

- We have:

$$m_{12} = p_{12} + p_{11} (1 + m_{12}) + p_{13} (1 + m_{32}) = 1 + p_{11} m_{12} + p_{13} m_{32} \quad (5.11)$$

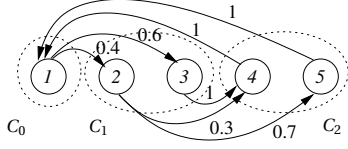
- Clearly  $m_{32} = 1$ , thus:

$$m_{12} = 1 + 0,7 m_{12} + 0,1 \times 1 \Rightarrow m_{12} = 11/3. \quad (5.12)$$

**5.8 Periodic states**

- A recurrent state  $j$  is **periodic** with period  $d > 1$  if  $j$  can only be reached after leaving it with a multiple of  $d$  steps.
- If  $d = 1$  the state is aperiodic.
- Any **periodic irreducible chain** can be partitioned in  $d$  **cyclic classes**  $C_0, \dots, C_{d-1}$  such that at each step a transition occur from class  $C_i$  to  $C_{(i+1) \bmod d}$ .
- By properly numerating the states, the transition matrix can be written as (the sub-matrices  $\mathbf{A}_i$  may not be square):

$$\mathbf{P} = \begin{matrix} & C_0 & C_1 & C_2 & \dots & C_{d-1} \\ \begin{matrix} C_0 \\ C_1 \\ \dots \\ C_{d-1} \end{matrix} & \begin{bmatrix} 0 & \mathbf{A}_1 & 0 & \dots & 0 \\ 0 & 0 & \mathbf{A}_2 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots \\ \mathbf{A}_{d-1} & 0 & 0 & \dots & 0 \end{bmatrix} \end{matrix} \quad (5.13)$$

**Example**

$$\mathbf{P} = \begin{bmatrix} 0 & 0.4 & 0.6 & 0 & 0 \\ 0 & 0 & 0 & 0.3 & 0.7 \\ 0 & 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \end{bmatrix}$$

$$\mathbf{P}^2 = \begin{bmatrix} 0 & 0 & 0 & 0.72 & 0.28 \\ 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 0.4 & 0.6 & 0 & 0 \\ 0 & 0.4 & 0.6 & 0 & 0 \end{bmatrix},$$

$$\mathbf{P}^3 = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 0.4 & 0.6 & 0 & 0 \\ 0 & 0.4 & 0.6 & 0 & 0 \\ 0 & 0 & 0 & 0.72 & 0.28 \\ 0 & 0 & 0 & 0.72 & 0.28 \end{bmatrix},$$

$$\mathbf{P}^4 = \begin{bmatrix} 0 & 0.4 & 0.6 & 0 & 0 \\ 0 & 0 & 0 & 0.72 & 0.28 \\ 0 & 0 & 0 & 0.72 & 0.28 \\ 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \end{bmatrix}, \dots$$

- In periodic chains  $\mathbf{P}^n$  does not converge.

## Chapter 6

### Steady State

#### 6.1 Limiting Distribution

- Probability of being in state  $i$  at step  $n$ :

$$\pi_i(n) = P(X(n) = i). \quad (6.1)$$

In vector form (row vector)

$$\boldsymbol{\pi}(n) = (\pi_1(n), \pi_2(n), \dots). \quad (6.2)$$

- The evolution of the chain depends on the initial distribution  $\boldsymbol{\pi}(0)$ .
- If we are interested in the **transient** evolution we shall study

$$\boldsymbol{\pi}(n) = \boldsymbol{\pi}(0) \mathbf{P}^n. \quad (6.3)$$

- If we are interested the **steady state** we shall be interested in the **limiting distribution** (if the limit exists):

$$\boldsymbol{\pi}(\infty) = (\pi_1(\infty), \pi_2(\infty), \dots) \quad (6.4)$$

Assume an **irreducible** chain with **positive recurrent** states.

- With infinite steps, we look for a probability converging to a value that depends only on the final state:

$$\pi_j(\infty) = \sum_i \pi_i(0) \lim_{n \rightarrow \infty} p_{ij}(n), \forall j \text{ and for any } \boldsymbol{\pi}(0), \quad (6.5)$$

which implies:

$$\pi_j(\infty) = \lim_{n \rightarrow \infty} p_{ij}(n) \sum_i \pi_i(0) = p_{ij}(\infty), \forall j \Rightarrow$$

$$\mathbf{P}(\infty) = \lim_{n \rightarrow \infty} \mathbf{P}^n = \begin{bmatrix} \boldsymbol{\pi}(\infty) \\ \dots \\ \boldsymbol{\pi}(\infty) \end{bmatrix} \quad (6.6)$$

- If this limit exists, we call  $\mathbf{P}(\infty)$  the **limiting matrix**, and  $\boldsymbol{\pi}(\infty)$  the **limiting distribution**.

**Example**

$$\mathbf{P} = \begin{bmatrix} 0.8 & 0.15 & 0.05 \\ 0.7 & 0.2 & 0.1 \\ 0.5 & 0.3 & 0.2 \end{bmatrix}$$

$$\mathbf{P}^2 = \begin{bmatrix} 0.764 & 0.168 & 0.068 \\ 0.760 & 0.170 & 0.070 \\ 0.752 & 0.174 & 0.074 \end{bmatrix}$$

$$\mathbf{P}^4 = \begin{bmatrix} 0.7628 & 0.1686 & 0.0686 \\ 0.7620 & 0.1690 & 0.0690 \\ 0.7604 & 0.1698 & 0.0698 \end{bmatrix}$$

$$\mathbf{P}^8 = \begin{bmatrix} 0.762500 & 0.168750 & 0.068750 \\ 0.762499 & 0.168750 & 0.068750 \\ 0.762497 & 0.168752 & 0.068752 \end{bmatrix}$$

$$\dots$$

$$\Rightarrow \boldsymbol{\pi}(\infty) = (0.76250, 0.16875, 0.06875)$$

#### 6.2 Stationary distribution

- We have:

$$\pi_i(n) = P(X(n) = i) = \sum_k P(X(n-1) = k) P(X(n) = i | X(n-1) = k) = \sum_k \pi_k(n-1) p_{ki} \quad (6.7)$$

- In matrix form:  $\boldsymbol{\pi}(n) = \boldsymbol{\pi}(n-1) \mathbf{P}$
- If  $\pi_i(n) = \pi_i(n-1) = \pi_i \forall i$ , we call  $\pi_i$  the **stationary probability of state  $i$** , and  $\boldsymbol{\pi} = (\pi_1, \pi_2, \dots)$ , the **stationary distribution** of the chain.
- In matrix form (**Global balance equations**):

$$\boldsymbol{\pi} = \boldsymbol{\pi} \mathbf{P}$$

$$\boldsymbol{\pi} \mathbf{e} = 1, \mathbf{e} = \begin{bmatrix} 1 & 1 & \dots \end{bmatrix}^T \quad (6.8)$$

- Thus, the stationary distribution is the left-hand eigenvector corresponding to the unit eigenvalue of  $\mathbf{P}$ .
- $\boldsymbol{\pi}(n) = \boldsymbol{\pi} \Rightarrow \boldsymbol{\pi}(n+1) = \boldsymbol{\pi}(n) \mathbf{P} = \boldsymbol{\pi} \mathbf{P} = \boldsymbol{\pi} \Rightarrow \boxed{\boldsymbol{\pi}(k) = \boldsymbol{\pi}, k \geq n}$
- Do not confuse the **limiting distribution**  $\boldsymbol{\pi}(\infty)$  and the **stationary distribution**  $\boldsymbol{\pi} = \boldsymbol{\pi} \mathbf{P}$ .
- $\boldsymbol{\pi}(\infty)$  and  $\boldsymbol{\pi}$  may not be the same, e.g. in **periodic chains**  $\boldsymbol{\pi}(\infty)$  does not exist ( $\mathbf{P}$  does not converge), but we can compute the stationary distribution.

- **Example:** the periodic chain

$$\mathbf{P} = \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{bmatrix} \quad (6.9)$$

has the stationary distribution

$$\boldsymbol{\pi} = \left[ 1/3 \quad 1/3 \quad 1/3 \right]. \quad (6.10)$$

### 6.3 Numerical Solution

#### Replace one equation method

$$\boldsymbol{\pi} = \boldsymbol{\pi} \mathbf{P}$$

$$\boldsymbol{\pi} \mathbf{e} = 1, \mathbf{e} = \begin{bmatrix} 1 & 1 & \dots \end{bmatrix}^T$$

We solve the equation  $\boldsymbol{\pi}(\mathbf{I} - \mathbf{P}) = 0$  replacing the last equation by  $\boldsymbol{\pi} \mathbf{e} = 1$ :

$$\boldsymbol{\pi} \begin{bmatrix} p_{11}-1 & p_{12} & \dots & p_{1n-1} & 1 \\ p_{21} & p_{22}-1 & \dots & p_{2n-1} & 1 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ p_{n1} & p_{n2} & \dots & p_{nn-1} & 1 \end{bmatrix} = \begin{bmatrix} 0 & 0 & \dots & 0 & 1 \end{bmatrix} \quad (6.11)$$

#### Examples

- **Replace one equation method:**  $\mathbf{P} = \begin{bmatrix} 0.8 & 0.15 & 0.05 \\ 0.7 & 0.2 & 0.1 \\ 0.5 & 0.3 & 0.2 \end{bmatrix}$

- With **octave** (matlab clone):

```
octave:1> P
      =[0.8,0.15,0.05;0.7,0.2,0.1;0.5,0.3,0.2];
octave:2> s=size(P,1); # number of rows.
octave:3> [zeros(1,s-1),1] / ...
> [eye(s,s-1)-P(1:s,1:s-1), ones(s,1)]
ans =
0.762500 0.168750 0.068750
```

- With **R**

```
> P <- matrix(nc=3, byr=T,
+ c(0.8,0.15,0.05,0.7,0.2,0.1,0.5,0.3,0.2))
> s <- nrow(P)
> solve(t(cbind(P[,1:(s-1)]-diag(nr=s,nc=s-1),
+ rep(1,s))), c(rep(0,s-1),1))
[1] 0.76250 0.16875 0.06875
```

NOTE:  $\boldsymbol{\pi} = \boldsymbol{\pi} \mathbf{P} \Rightarrow \boldsymbol{\pi}^T = \mathbf{P}^T \boldsymbol{\pi}^T$ . The transpose operator in R is `t()`.

### 6.4 Global balance equations

- Why are they called Global balance equations?

$$\left. \begin{aligned} \boldsymbol{\pi} = \boldsymbol{\pi} \mathbf{P} &\Rightarrow \pi_j = \sum_{i=0}^{\infty} \pi_i p_{ij} \\ \sum_{i=0}^{\infty} p_{ji} = 1 &\Rightarrow \pi_j \sum_{i=0}^{\infty} p_{ji} = \pi_j \end{aligned} \right\} \Rightarrow \sum_{i=0}^{\infty} \pi_i p_{ij} = \pi_j \sum_{i=0}^{\infty} p_{ji} \quad (6.12)$$

$$\sum_{i=0}^{\infty} \pi_i p_{ij} \Rightarrow \text{Frequency of transitions entering state } j$$

$$\pi_j \sum_{i=0}^{\infty} p_{ji} \Rightarrow \text{Frequency of transitions leaving state } j \quad (6.13)$$

- In **stationary regime**, the frequency of transitions leaving state  $j$  is equal to the frequency of transitions entering state  $j$ .

#### Flux Balancing

- Define the **flux**  $F_{uv}$  from state  $u$  to  $v$ :

$$F_{uv} = \pi_u p_{uv} \quad (6.14)$$

- and the flux from set of states  $U$  to  $V$ :

$$F(U,V) = \sum_{u \in U} \sum_{v \in V} F_{uv} \quad (6.15)$$

- From the Global balance equations we have:

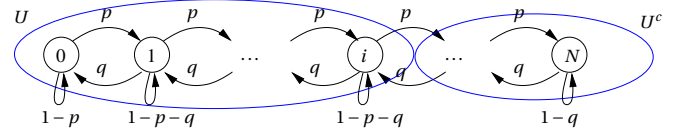
$$\sum_{i=0}^{\infty} \pi_i p_{ij} = \pi_j \sum_{i=0}^{\infty} p_{ji} \Rightarrow \sum_{i \in U} F_{ij} + \sum_{i \notin U} F_{ij} = \sum_{i \in U} F_{ji} + \sum_{i \notin U} F_{ji} \quad (6.16)$$

- Adding for  $j \in U$ :

$$\sum_{j \in U} \sum_{i \in U} F_{ij} + \sum_{j \in U} \sum_{i \notin U} F_{ij} = \sum_{j \in U} \sum_{i \in U} F_{ji} + \sum_{j \in U} \sum_{i \notin U} F_{ji} \Rightarrow$$

$$\sum_{j \in U} \sum_{i \notin U} F_{ij} = \sum_{j \in U} \sum_{i \notin U} F_{ji} \Rightarrow F(U, U^c) = F(U^c, U)$$

#### Solution Using Flux Balancing



- Flux balancing  $\Rightarrow p \pi_i = q \pi_{i+1}$
- Iterating:  $\pi_1 = \rho \pi_0, \pi_2 = \rho \pi_1 = \rho^2 \pi_0, \dots, \Rightarrow$

$$\pi_i = \rho^i \pi_0, i = 1, \dots, N \quad \text{where: } \rho = \frac{p}{q},$$

- Normalizing:  $\sum_{i=0}^N \pi_i = 1$

$$\pi_0 = \frac{1-\rho}{1-\rho^{N+1}}, \quad p \neq q$$

$$\pi_0 = \frac{1}{N+1}, \quad p = q$$

### 6.5 Ergodic Chains

**Ergodic state** positive recurrent and aperiodic state.

**Ergodic chain** if all states are ergodic.

**Theorem:** All states of an irreducible Markov chain are of the same type: Transient or positive/null recurrent, and aperiodic/periodic with the same period [2, chapter XV].

**Consequences:**

- **Finite aperiodic and irreducible** chains are **ergodic** (since all states are positive recurrent).
- **Infinite aperiodic and irreducible** chains can be:
  - **Ergodic:** all the states are positive recurrent (stable chains).
  - **Non ergodic:** all states are null recurrent or transient (unstable chains).

### 6.5.1 Theorems for ergodic chains

- $\pi = \pi(\infty)$

*Proof.* For an **aperiodic irreducible** chain with **positive recurrent** states:

$$\begin{aligned} \begin{cases} \pi(\infty) &= \pi(0) \mathbf{P}(\infty) \\ \mathbf{P}(\infty) &= \lim_{n \rightarrow \infty} \mathbf{P}^n = \begin{bmatrix} \pi(\infty) \\ \dots \\ \pi(\infty) \end{bmatrix} \end{cases} \Rightarrow \\ \pi(\infty) \mathbf{P} = (\pi(0) \lim_{n \rightarrow \infty} \mathbf{P}^n) \mathbf{P} = \pi(0) \mathbf{P}(\infty) = \pi(\infty) \\ \Rightarrow \begin{cases} \pi(\infty) \mathbf{P} = \pi(\infty) \\ \pi(\infty) \mathbf{e} = 1 \end{cases} \quad \pi(\infty) \text{ satisfies the GBE} \Rightarrow \pi = \pi(\infty) \end{aligned} \quad (6.17)$$

□

- In stationary regime (when  $\pi(n) \mathbf{P} = \pi(n)$ ), the **mean number of steps the system remains in state  $j$**  during  $k$  steps is given by

$$k \pi_j. \quad (6.18)$$

*Proof.* Assume the chain in stationary regime at time  $t = 0$  ( $\pi(0) \mathbf{P} = \pi(0)$ ), and let  $j(k)$  be the number of visits to  $j$  in  $k$  steps:  $j(k) = \sum_{i=0}^{k-1} I(X(i) = j)$  ( $I(A)$  is the indicator function:  $I(A) = 1$  if  $A$  occurs,  $I(A) = 0$  otherwise):

$$E[j(k)] = \sum_{i=0}^{k-1} E[I(X(i) = j)] = \sum_{i=0}^{k-1} P(X(i) = j) = k \pi_j \quad (6.19)$$

□

- In stationary regime the **mean recurrence time** (mean number of steps between two consecutive visits to state  $j$ ) is given by

$$m_{jj} = 1/\pi_j \quad (6.20)$$

*Proof.* Let  $j(k)$  be the number of visits to  $j$  in  $k$  steps:

$$\pi_j = \lim_{k \rightarrow \infty} \frac{j(k)}{k} = \lim_{k \rightarrow \infty} \frac{1}{k/j(k)} = 1/m_{jj} \quad (6.21)$$

□

## Chapter 7

### Reversed Chain

#### Definition

- Let  $X(n)$  be an **ergodic** MC. The chain  $X^r(n) = X(-n)$  is referred to as the **time reversal chain** of  $X(n)$ .
- **Example**, consider a possible sample path of  $X(n)$ :

$$\dots (i_0, n_0), (i_1, n_1), (i_2, n_2), \dots \quad (7.1)$$

The same path in the time reversal chain  $X^r(n)$  would be:

$$\dots (i_2, -n_2), (i_1, -n_1), (i_0, -n_0), \dots \quad (7.2)$$

#### Properties

- Let  $p_{ij}, p_{ij}^r$  be the transition probabilities of  $X(n)$  respectively  $X^r(n)$ , and  $\pi_i, \pi_i^r$  the stationary distributions of  $X(n)$  respectively  $X^r(n)$ , then:

$$\pi_i = \pi_i^r \quad (7.3)$$

- **Proof:** the mean time in each state is the same for both chains.
- However, **in general**  $p_{ij} \neq p_{ij}^r$ . For example,  $X(n)$  may be able to jump from state  $i$  to  $j$ , but not from  $j$  to  $i \Rightarrow X^r(n)$  can jump from  $j$  to  $i$ , but not from  $i$  to  $j$ .
- But it must be  $p_{ii} = p_{ii}^r$ , since self-state transitions are the same in the direct and reversed chains.

### 7.1 Computation of $p_{ij}^r$

The transition probabilities in the time reversal chain ( $p_{ij}^r$ ) satisfy:

$$\pi_i p_{ij} = \pi_j p_{ji}^r \quad (7.4)$$

*Proof.* Assume the chain in **steady state**. We have:

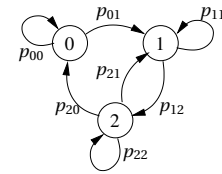
$$\begin{aligned} P\{X(n+1) = j, X(n) = i\} &= \\ P\{X^r(-n) = i, X^r(-n-1) = j\} &= \\ P\{X^r(n+1) = i, X^r(n) = j\} &\Rightarrow \\ P\{X(n) = i\} P\{X(n+1) = j \mid X(n) = i\} &= \\ P\{X(n) = j\} P\{X^r(n+1) = i \mid X^r(n) = j\} &\Rightarrow \\ \pi_i p_{ij} &= \pi_j p_{ji}^r. \quad \square \end{aligned}$$

We can **compute**  $p_{ji}^r$  using the **reversed balance equations**:

$$\pi_i p_{ij} = \pi_j p_{ji}^r \Rightarrow \sum_{i \in U} \sum_{j \in V} \pi_i p_{ij} = \sum_{i \in U} \sum_{j \in V} \pi_j p_{ji}^r \Rightarrow$$

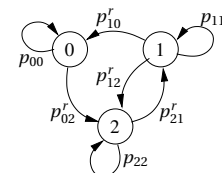
$$F(U, V) = F^r(V, U) \quad (7.5)$$

#### Example



$$\Rightarrow \begin{cases} \pi_0 = \frac{p_{12} p_{20}}{p_{12} p_{20} + p_{01} (p_{20} + p_{21}) + p_{01} p_{12}} \\ \pi_1 = \frac{p_{01} (p_{20} + p_{21})}{p_{12} p_{20} + p_{01} (p_{20} + p_{21}) + p_{01} p_{12}} \\ \pi_2 = \frac{p_{01} p_{12}}{p_{12} p_{20} + p_{01} (p_{20} + p_{21}) + p_{01} p_{12}} \end{cases}$$

Time reversal chain:





$$\Rightarrow \begin{cases} \pi_0 p_{01} = \pi_1 p_{10}^r \\ \pi_1 p_{12} = \pi_2 p_{21}^r \\ \pi_2 p_{21} = \pi_1 p_{12}^r \\ \pi_2 p_{20} = \pi_0 p_{02}^r \end{cases} \Rightarrow \begin{cases} p_{10}^r = \frac{p_{12} p_{20}}{p_{20} + p_{21}} \\ p_{21}^r = \frac{p_{20} + p_{21}}{p_{12} p_{21}} \\ p_{12}^r = \frac{p_{12} p_{21}}{p_{20} + p_{21}} \\ p_{02}^r = p_{01} \end{cases}$$

## Chapter 8

### Reversible Chains

#### Definition

- A chain is reversible if:

$$p_{ij} = p_{ij}^r \quad (8.1)$$

- This equality implies the **reversibility balance equations**:

$$\pi_i p_{ij} = \pi_j^r p_{ji}^r \Rightarrow F(U, V) = F^r(U, V) \quad (8.2)$$

- Using both reversed ( $F^r(U, V) = F(V, U)$ ) and reversibility balance equations, the following relation holds for a reversible chain (**detailed balance equations**):

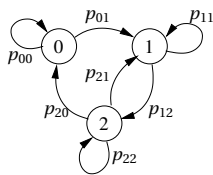
$$F(U, V) = F(V, U) \quad (8.3)$$

- NOTE: Compare with the **global balance equations**:  $F(U, U^C) = F(U^C, U)$ .

### 8.1 Kolmogorov Criteria

#### Definition of path

- Define a **path** as a possible sequence of transitions of the chain. For example, in the figure it could be  $0 \rightsquigarrow 0 \rightsquigarrow 1 \rightsquigarrow 2$ .



- We denote the **sequence of states** of one path  $l$  as:

$$(l, 1) \rightsquigarrow (l, 2) \rightsquigarrow \dots \rightsquigarrow (l, m) \quad (8.4)$$

- For instance, if  $l$  is  $0 \rightsquigarrow 0 \rightsquigarrow 1 \rightsquigarrow 2$ , then  $(l, 1) = 0$ ,  $(l, 2) = 0$ ,  $(l, 3) = 1$ ,  $(l, 4) = 2$ .
- We say that a path is **closed** if it starts and ends in the same state. For instance, a path starting and ending in state  $(l, 1)$ :

$$(l, 1) \rightsquigarrow (l, 2) \rightsquigarrow \dots \rightsquigarrow (l, m) \rightsquigarrow (l, 1) \quad (8.5)$$

#### Kolmogorov Criteria

- Take a **closed path**  $l$  with  $m \geq 0$  transitions, i.e.:

$$(l, 1) \rightsquigarrow (l, 2) \rightsquigarrow \dots \rightsquigarrow (l, m) \rightsquigarrow (l, 1), m \geq 0 \quad (8.6)$$

- The chain is **reversible iff for all**  $l$ :

$$p_{(l,1)(l,2)} p_{(l,2)(l,3)} \dots p_{(l,m)(l,1)} = p_{(l,1)(l,m)} p_{(l,m)(l,m-1)} \dots p_{(l,2)(l,1)} \quad (8.7)$$

#### • Proof:

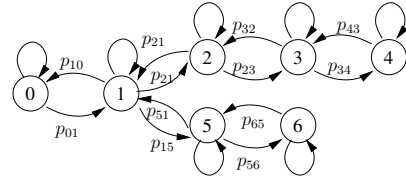
- If the chain is reversible  $\pi_i p_{ij} = \pi_j p_{ji}$  (detailed balance equations):  $\Rightarrow \pi_{(l,k)} p_{(l,k)(l,k+1)} = \pi_{(l,k+1)} p_{(l,k+1)(l,k)}$
- Multiplying for  $k = 1, 2, \dots, m$  and simplifying we obtain the previous relation.

#### Corollary

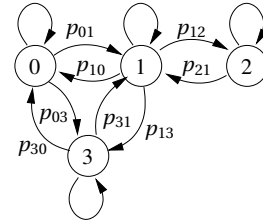
- A reversible chain must satisfy:

$$\begin{aligned} p_{ij} > 0 &\Rightarrow p_{ji} > 0 \\ p_{ij} = 0 &\Rightarrow p_{ji} = 0 \end{aligned} \quad (8.8)$$

- An **ergodic tree chain is always reversible**. We define a tree chain as chain having no cycles, i.e. the only possible transitions are bidirectional arcs between states, and self transitions.



#### Example



- The chain is **reversible iff**:

$$p_{01} p_{13} p_{30} = p_{10} p_{03} p_{31}$$

### 8.2 Product Form Solution

- Let  $X(n)$  be a reversible MC with space state  $S \Rightarrow$  the **stationary probabilities** of  $X(n)$  can be computed as follows:
- Choose a state  $\mathbf{s} \in S$ ,
- For every other state  $\mathbf{i} \in S$ ,  $\mathbf{i} \neq \mathbf{s}$  look for a possible path  $l_i$  from state  $\mathbf{s}$  to state  $\mathbf{i}$ :

$$\mathbf{s} = (l_i, 1) \rightsquigarrow (l_i, 2) \rightsquigarrow \dots \rightsquigarrow (l_i, m_{l_i}) = \mathbf{i}, m_{l_i} \geq 1 \quad (8.9)$$

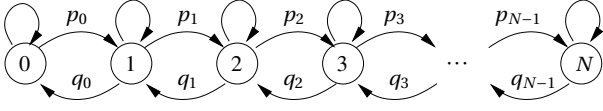
- The stationary probabilities are given by:

$$\pi_i = \frac{\psi_i}{\sum_{j \in S} \psi_j}, i \in S \quad \text{where } \psi_i = \begin{cases} 1, & i = \mathbf{s} \\ \prod_{k=1}^{m_{l_i}-1} \frac{p_{(l_i,k)(l_i,k+1)}}{p_{(l_i,k+1)(l_i,k)}}, & i \neq \mathbf{s} \end{cases} \quad (8.10)$$

- Proof** Use the detailed balance equations  $\pi_i p_{ij} = \pi_j p_{ji}$ .



### 8.3 Birth and Death Chains



- Birth and death chains are reversible.
- Applying the product form solution for reversible chains we obtain the general solution of birth death chains. Choosing  $s = 0$ :

$$\pi_i = \frac{\psi_i}{\sum_{j=0}^N \psi_j}, i \geq 0 \quad \text{where } \psi_i = \begin{cases} 1, & i = 0 \\ \prod_{k=0}^{i-1} \frac{p_k}{q_k}, & i = 1, \dots, N \end{cases} \quad (8.11)$$

#### Truncated Reversible Chain

- Consider a reversible MC  $X$  with a stationary distribution  $\pi_i$ .
- Suppose that we **truncate the chain**  $X$  and we obtain another irreducible chain  $X'$ .
- Then,  $X'$  is also reversible with stationary distribution:

$$\pi'_i = \frac{\pi_i}{G}, \quad \sum_k \pi'_k = 1 \quad (8.12)$$

## Chapter 9

### Research Example: Aloha

**Access Protocol** (see the paper of Kleinrock and Lam [5]).

#### • Pure Aloha:

- Broadcast radio system.
- **Single hop** system (all stations are in coverage).
- Whenever a station has a frame ready, it is transmitted.
- If two or more frames Tx overlap in time there is a **collision**, otherwise the frame is received correctly.
- Colliding frames are reTx after a **random time (backoff)**.

#### • Slotted Aloha:

- Time is slotted.
- Tx can only occur at the beginning of a slot.
- Collisions occur when 2 or more stations Tx in the same slot.

### 9.1 Analysis with finite population

Assumptions

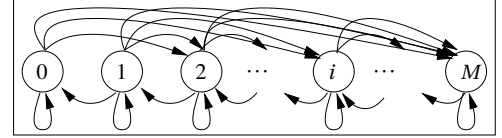
- **Slotted Aloha**.
- **Acks** are sent immediately after the reception of a frame, and are never lost.
- $M$  nodes with a **buffer** of 1 frame.
- The **nodes** can be in 2 states:
  - **Thinking**: when the buffer is empty
  - **Backlogged**: when there is a frame in the buffer.

- A thinking node generate one frame in each slot with probability  $\sigma$ . When a frame collides, the frame is stored and the node becomes backlogged.
- A backlogged node ReTx the frame in each slot with probability  $\nu$ .

#### Markov Chain

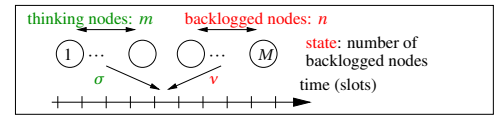
- The system **state**,  $X(n)$ , is the **number of backlogged nodes**:

$$p_{ij} = P(X(n) = j \text{ baklogged} | X(n-1) = i \text{ baklogged}) \quad (9.1)$$



#### Transition probabilities

- 0 for  $j < i - 1$ .
- for  $j = i - 1$ : no thinking Tx and only 1 backlogged Tx.
- for  $j = i$ :
  1. no thinking Tx and none or more than 1 backlogged Tx,
  2. only 1 thinking Tx and no backlogged Tx.
- for  $j = i + 1$ : 1 thinking and 1 or more backlogged Tx.
- for  $j > i + 1$ :  $j - i$  thinking Tx, regardless of backlogged Tx.



In order to compute the previous events, define the probabilities:

- **Arrivals:**  $Q_a(m, n)$ , Probability of  $m$  thinking nodes Tx in a slot given that  $n$  nodes are backlogged:

$$Q_a(m, n) = P\left(\begin{matrix} m \text{ think.} \\ \text{nodes Tx} \end{matrix} \middle| \begin{matrix} n \text{ nodes are} \\ \text{backlogged} \end{matrix}\right) = \binom{M-n}{m} \sigma^m (1-\sigma)^{M-n-m} \quad (9.2)$$

- **Retransmissions:**  $Q_r(m, n)$ , Probability of  $m$  backlogged nodes Tx in a slot given that  $n$  nodes are backlogged:

$$Q_r(m, n) = P\left(\begin{matrix} m \text{ backl.} \\ \text{nodes Tx} \end{matrix} \middle| \begin{matrix} n \text{ nodes are} \\ \text{backlogged} \end{matrix}\right) = \binom{n}{m} \nu^m (1-\nu)^{n-m} \quad (9.3)$$

- and we have:

$$p_{ij} = \begin{cases} 0, & j < i - 1 \\ Q_a(0, i) Q_r(1, i), & j = i - 1 \\ Q_a(0, i) (1 - Q_r(1, i)) + Q_a(1, i) Q_r(0, i), & j = i \\ Q_a(1, i) (1 - Q_r(0, i)), & j = i + 1 \\ Q_a(j - i, i), & j > i + 1 \end{cases} \quad (9.4)$$

### 9.1.1 Stationary distribution

- Solving the global balance equations:

$$\begin{aligned}\pi &= \pi \mathbf{P} \\ \pi \mathbf{e} &= 1\end{aligned}\quad (9.5)$$

- We obtain the probability of having  $i$  backlogged nodes:

$$\pi_i = P(i \text{ backlogged nodes}) \quad (9.6)$$

NOTE: there is **no closed form solution** of the chain. The matrix  $\mathbf{P}$  must be constructed using the expression of  $p_{ij}$ , and solved numerically.

## 9.2 Throughput

- Define the probabilities:

$$P_{succ}(i) = P(\text{successful Tx} \mid i \text{ backlogged}) \quad (9.7)$$

- The **normalized throughput**, i.e. proportion of steps with a successful transmission) is:

$$S = \sum_{i=0}^M P(\text{successful Tx} \mid i \text{ backlogged}) \times P(i \text{ backlogged}) = \sum_{i=0}^M P_{succ}(i) \pi_i \quad (9.8)$$

- For a slot to be successful: (i) 1 thinking Tx and no backlogged Tx, or (ii) no thinking Tx and 1 backlogged Tx:

$$P_{succ}(i) = Q_a(1, i) Q_r(0, i) + Q_a(0, i) Q_r(1, i) \quad (9.9)$$

### Notes on the throughput

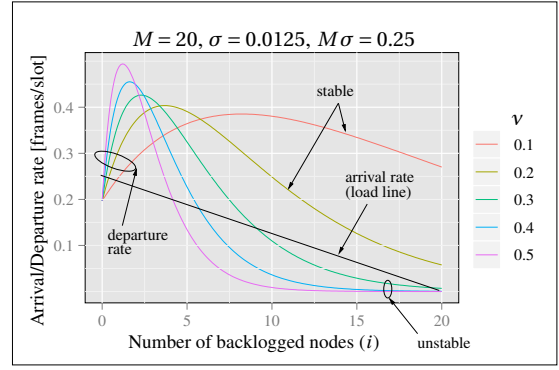
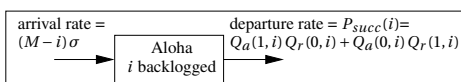
$$S = \sum_{i=0}^M P_{succ}(i) \pi_i \quad (9.10)$$

- For the **special case**  $\sigma = \nu$  (thinking Tx with the same probability as backlogged):  $P_{succ}(i) = M\sigma(1-\sigma)^{M-1}$ , which does not depend on  $i$ , thus:  $S = M\sigma(1-\sigma)^{M-1}$ .
- The **offered load** (i.e. proportion of arrivals per slot)  $G$  is now:  $G = M\sigma$ , thus:

$$S = G \left(1 - \frac{G}{M}\right)^{M-1} \Rightarrow \lim_{M \rightarrow \infty} S = G e^{-G} \quad (9.11)$$

- We conclude that the **infinite population model** is the limit of the finite population if backlogged Tx with the same probability as thinking, and  $M \rightarrow \infty$ .

### 9.2.1 Dynamics



Note on the arrival rate (expected value of a binomial distribution):

$$\sum_{k=0}^{M-i} k \binom{M-i}{k} \sigma^k (1-\sigma)^{M-i-k} = (M-i) \sigma$$

Solving the chain:  $S = \sum_{i=0}^M P_{succ}(i) \pi_i$

$\nu$	$S$
0.1	2.38e-01
0.2	2.42e-01
<b>0.3</b>	<b>1.30e-02</b>
<b>0.4</b>	<b>4.98e-04</b>
<b>0.5</b>	<b>1.90e-05</b>

### 9.2.2 Stabilizing Aloha

- The **retransmission probabilities** must adapt in accordance with the state of the system.
- Example: **binary exponential backoff** (ethernet). The retransmission rate at retransmission  $i$  is adapted as  $\nu = 2^{-i}$ . Thus, the higher are the number of retransmission trials  $i$ , the lower (exponentially) is the retransmission rate.

## Chapter 10

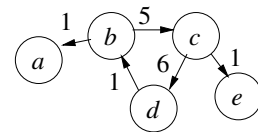
### Finite Absorbing Chains

#### Canonical Form

- Let  $\mathbf{P}^{rxr}$  be the transition probability matrix of a chain with  $r$  states:  $s$  **transient** states and  $r-s$  **absorbing** states. We can write  $\mathbf{P}^{rxr}$  in the **canonical form**:

$$\mathbf{P}^{rxr} = \begin{bmatrix} \mathbf{Q}^{s \times s} & \mathbf{R}^{s \times r-s} \\ \mathbf{0}^{r-s \times s} & \mathbf{I}^{r-s \times r-s} \end{bmatrix} \quad (10.1)$$

#### Example



$$\mathbf{P} = \begin{array}{c|cc|cc} & b & c & d & a & e \\ \hline b & 0 & 0.1 & 0 & 0.9 & 0 \\ c & 0 & 0 & 0.8 & 0 & 0.2 \\ d & 0.3 & 0 & 0.7 & 0 & 0 \\ \hline a & 0 & 0 & 0 & 1 & 0 \\ e & 0 & 0 & 0 & 0 & 1 \end{array}$$

## 10.1 Results

- Define:

$$\begin{aligned} \mathbf{n}_{ij} &= \begin{cases} \text{number of steps in state } j \text{ before} \\ \text{absorption, starting from state } i \end{cases}, \\ \mathbf{t}_i &= \begin{cases} \text{number of steps in transient states before} \\ \text{absorption, starting from state } i \end{cases}, \\ \mathbf{b}_{ij} &= P(\text{probability to be absorbed } j \text{ starting } i) \end{aligned} \quad (10.2)$$

- Then:

$$\begin{aligned} \{E[n_{ij}]\} &= \mathbf{N} = (\mathbf{I} - \mathbf{Q})^{-1} \\ \{\text{Var}[n_{ij}]\} &= \mathbf{N}(2\mathbf{N}_{\text{diag}} - \mathbf{I}) - \mathbf{N}\mathbf{sqr} \\ \{E[t_i]\} &= \boldsymbol{\tau} = \mathbf{N}\mathbf{e} \\ \{\text{Var}[t_i]\} &= (2\mathbf{N} - \mathbf{I})\boldsymbol{\tau} - \boldsymbol{\tau}_{\text{sqr}} \\ \{b_{ij}\} &= \mathbf{B} = \mathbf{N}\mathbf{R}. \end{aligned} \quad (10.3)$$

where  $\{a_{ij}\}$  is a matrix with  $a_{ij}$  as element  $ij$  and  $\mathbf{e}$  is a column vector of 1s.  $\mathbf{N}$  is called the **fundamental matrix**.

### Proofs

- $\{E[n_{ij}]\} = \mathbf{N} = (\mathbf{I} - \mathbf{Q})^{-1}$

*Proof.*

$$\begin{aligned} E[n_{ij}] &= \sum_{k \in A} p_{ik} \delta_{ij} + \sum_{k \in T} p_{ik} E[n_{kj} + \delta_{ij}] = \\ &\quad \delta_{ij} + \sum_{k \in T} p_{ik} E[n_{kj}] \\ \Rightarrow \{E[n_{ij}]\} &= \mathbf{N} = \mathbf{I} + \mathbf{Q}\mathbf{N} \Rightarrow \mathbf{N} = (\mathbf{I} - \mathbf{Q})^{-1} \end{aligned}$$

where  $A$  is the set of absorbing states and  $T$  is the set of transient states.

$$\text{Notation: } \delta_{ij} = I(i = j) = \begin{cases} 1, & i = j, \\ 0, & \text{otherwise.} \end{cases} \quad \square$$

- $\{\text{Var}[n_{ij}]\} = \mathbf{N}(2\mathbf{N}_{\text{diag}} - \mathbf{I}) - \mathbf{N}\mathbf{sqr}$

*Proof.*

$$\begin{aligned} \text{Var}[n_{ij}] &= E[n_{ij}^2] - E[n_{ij}]^2 \Rightarrow \\ &\quad \{\text{Var}[\mathbf{n}_{ij}]\} = \{E[\mathbf{n}_{ij}^2]\} - \mathbf{N}\mathbf{sqr} \\ E[\mathbf{n}_{ij}^2] &= \sum_{k \in A} p_{ik} \delta_{ij}^2 + \sum_{k \in T} p_{ik} E[(n_{kj} + \delta_{ij})^2] = \\ &\quad \sum_{k \in A} p_{ik} \delta_{ij} + \sum_{k \in T} p_{ik} (E[n_{kj}^2] + 2E[n_{kj}] \delta_{ij} + \delta_{ij}) = \\ &\quad \delta_{ij} + \sum_{k \in T} (p_{ik} E[n_{kj}^2] + 2p_{ik} E[n_{kj}] \delta_{ij}) \Rightarrow \\ \{E[\mathbf{n}_{ij}^2]\} &= \mathbf{I} + \mathbf{Q}\{E[\mathbf{n}_{ij}^2]\} + 2(\mathbf{Q}\mathbf{N})_{\text{diag}} = \\ &\quad (\mathbf{I} - \mathbf{Q})^{-1} (\mathbf{I} + 2(\mathbf{Q}\mathbf{N})_{\text{diag}}) = \\ &\quad \mathbf{N}(\mathbf{I} + 2(\mathbf{N} - \mathbf{I})_{\text{diag}}) = \mathbf{N}(2\mathbf{N}_{\text{diag}} - \mathbf{I}) \quad \square \end{aligned}$$

- $\{E[t_i]\} = \boldsymbol{\tau} = \mathbf{N}\mathbf{e}$

*Proof.*

$$E[t_i] = \sum_{k \in T} E[n_{ik}] \Rightarrow \{E[t_i]\} = \boldsymbol{\tau} = \mathbf{N}\mathbf{e} \quad \square$$

- $\{\text{Var}[t_i]\} = (2\mathbf{N} - \mathbf{I})\boldsymbol{\tau} - \boldsymbol{\tau}_{\text{sqr}}$

*Proof.*

$$\begin{aligned} \text{Var}[t_i] &= E[t_i^2] - E[t_i]^2 \Rightarrow \{\text{Var}[\mathbf{t}_i]\} = \{E[\mathbf{t}_i^2]\} - \boldsymbol{\tau}_{\text{sqr}} \\ E[\mathbf{t}_i^2] &= \sum_{k \in A} p_{ik} + \sum_{k \in T} p_{ik} E[(t_k + 1)^2] = \\ &\quad \sum_{k \in A} p_{ik} + \sum_{k \in T} p_{ik} (E[t_k^2] + 2E[t_k] + 1) = \\ &\quad 1 + \sum_{k \in T} (p_{ik} E[t_k^2] + 2p_{ik} E[t_k]) \Rightarrow \\ \{E[\mathbf{t}_i^2]\} &= \mathbf{e} + \mathbf{Q}\{E[\mathbf{t}_i^2]\} + 2\mathbf{Q}\boldsymbol{\tau} = \\ &\quad (\mathbf{I} - \mathbf{Q})^{-1} (\mathbf{e} + 2\mathbf{Q}\boldsymbol{\tau}) = \\ &\quad \mathbf{N}(\mathbf{e} + 2\mathbf{Q}\boldsymbol{\tau}) = \boldsymbol{\tau} + 2\mathbf{N}\mathbf{Q}\boldsymbol{\tau} = \\ &\quad \boldsymbol{\tau} + 2(\mathbf{N} - \mathbf{I})\boldsymbol{\tau} = (2\mathbf{N} - \mathbf{I})\boldsymbol{\tau} \quad \square \end{aligned}$$

- $\{b_{ij}\} = \mathbf{B} = \mathbf{N}\mathbf{R}$

*Proof.*

$$\begin{aligned} b_{ij} &= p_{ij} + \sum_{k \in T} p_{ik} b_{kj}, \quad j \in A \Rightarrow \\ \{b_{ij}\} &= \mathbf{B} = \mathbf{R} + \mathbf{Q}\mathbf{B} \\ \Rightarrow \mathbf{B} &= (\mathbf{I} - \mathbf{Q})^{-1} \mathbf{R} = \mathbf{N}\mathbf{R}. \quad \square \end{aligned}$$

## 10.2 Extension of the Results

- The previous results can be generalized to any group of states of  $\mathbf{P}$ :
- A set  $S$  is referred to as **open** if the chain can reach some state of  $S^c$  starting from any state of  $S$ . Let

$$\mathbf{Q} = \{p_{ij}, i \in S, j \in S\} \quad (10.4)$$

$$\mathbf{R} = \{p_{ij}, i \in S, j \in S^c\} \quad (10.5)$$

Let assume that the process starts from  $i \in S$ . Define:

$$\begin{aligned} \mathbf{n}_{ij} &= \begin{cases} \text{number of steps in state } j \text{ before} \\ \text{leaving } S, \text{ starting from state } i \end{cases}, \\ \Rightarrow \{E[n_{ij}]\} &= \mathbf{N} = (\mathbf{I} - \mathbf{Q})^{-1}. \end{aligned}$$

- Similarly for the other results, e.g.  $\boldsymbol{\tau} = \{E[t_i]\} = \mathbf{N}\mathbf{e}$  and  $\mathbf{B} = \{b_{ij}\} = \mathbf{N}\mathbf{R}$ .

## 10.3 Inverse of a matrix

### Cofactors

$$\mathbf{A}^{-1} = \frac{1}{\det \mathbf{A}} \mathbf{C}^T \quad (10.6)$$

where  $\mathbf{C}^T$  is the transposed cofactor matrix:  $c_{ij} = (-1)^{i+j} \det \mathbf{M}_{ij}$ , and  $\mathbf{M}_{ij}$  are the minor matrices obtained removing the row  $i$  and column  $j$  from  $\mathbf{A}$ .

### Gaussian Elimination

Do the transformation:

$$[\mathbf{A} | \mathbf{I}] \rightarrow [\mathbf{I} | \mathbf{A}^{-1}] \quad (10.7)$$

using the elementary row operations:

- Swapping two rows.
- Multiplying a row by a nonzero number.
- Adding a multiple of one row to another row.

## Part III

# Continuous Time Markov Chains

## Chapter 11

### Definition of a CTMC

#### 11.1 Properties of a continuous time MC

- The states must be a numerable set.
- Let  $X(t)$  be the event {at time  $t$  the system is in state  $i$ }, then it must hold the **memoryless property**:

$$P(X(t_n) = i | X(t_1) = j, X(t_2) = k, \dots) = P(X(t_n) = i | X(t_1) = j) \text{ for any } t_n > t_1 > t_2 > t_3 \dots \quad (11.1)$$

#### 11.2 Transition Matrix

- **Transition probabilities:**

$$p_{ij}(t_1, t_2) = P(X(t_2) = j | X(t_1) = i) \quad (11.2)$$

- For an **homogeneous chain**:

$$p_{ij}(t) = P(X(t_1 + t) = j | X(t_1) = i) = P(X(t) = j | X(0) = i), \forall t_1 \quad (11.3)$$

- In matrix form (**transition probability matrix**):

$$\mathbf{P}(t) = \begin{bmatrix} p_{11}(t) & p_{12}(t) & \dots \\ p_{21}(t) & p_{22}(t) & \dots \\ \dots & \dots & \dots \end{bmatrix}, t \geq 0 \quad (11.4)$$

- Notes:

- Compare with the n-step prob. matrix of a DTMC:  $\mathbf{P}(n)$ .
- $\mathbf{P}(t)$  must be a **stochastic matrix** (all rows add to 1).

- We look for an equivalent 1-step prob. matrix  $\mathbf{P}$  of DTMCs.

- For consistency:  $\lim_{t \rightarrow 0} p_{ij}(t) = \delta_{ij}$ . In matrix form:

$$\lim_{t \rightarrow 0} \mathbf{P}(t) = \mathbf{I}. \quad (11.5)$$

- And assume that the following **transition rates** exist:

$$q_{ij} = \lim_{t \rightarrow 0} \frac{p_{ij}(t)}{t}, i \neq j; \quad q_{ii} = \lim_{t \rightarrow 0} \frac{p_{ii}(t) - 1}{t} \quad (11.6)$$

- In matrix form:  $\mathbf{Q} = \lim_{t \rightarrow 0} \frac{\mathbf{P}(t) - \mathbf{I}}{t}$

- Note that  $\sum_j p_{ij}(t) = 1 \Rightarrow p_{ii}(t) = 1 - \sum_{j \neq i} p_{ij}(t)$ , thus:

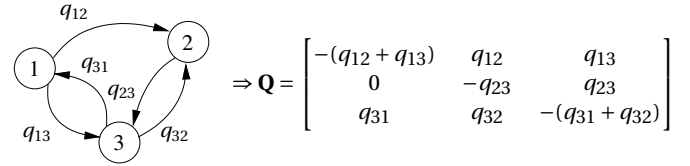
$$q_{ii} = \lim_{t \rightarrow 0} \frac{p_{ii}(t) - 1}{t} = \lim_{t \rightarrow 0} \frac{-\sum_{j \neq i} p_{ij}(t)}{t} = -\sum_{j \neq i} q_{ij} \quad (11.7)$$

- The matrix  $\mathbf{Q}$  is called the **transition rate or infinitesimal generator** of the chain.

- Since  $q_{ii} = -\sum_{j \neq i} q_{ij}$ , **all the rows of  $\mathbf{Q}$  add to 0**.
- The rate  $q_{ij}$ ,  $i \neq j$  measures “how fast” the chain moves from state  $i$  to  $j$ : the higher is  $q_{ij}$ , the faster it moves from  $i$  to  $j$ .
- For  $q_{ii} = -\sum_{j \neq i} q_{ij}$ , the higher  $-q_{ii}$  is, the faster the chain leaves state  $i$ .
- If  $q_{ij} = 0, \forall j \Rightarrow q_{ii} = 0$ , then  $i$  is an **absorbing state**: the chain “moves with rate 0 from  $i$  to other states”, i.e. never leaves  $i$ .

#### 11.3 State Transition Diagram

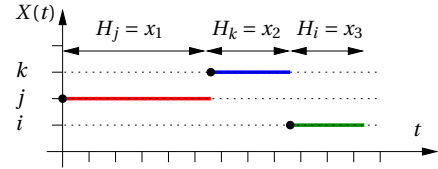
- A continuous MC is characterized by the transition rate or infinitesimal generator: the  $\mathbf{Q}$ -matrix.
- The state transition diagram is now represented as:



- Note that now we have **transition rates** ( $0 \leq q_{ij} < \infty, i \neq j$ ) and **not probabilities**.
- The rates  $q_{ii}$  are **not written** in the diagram, **no self transitions**.

#### 11.4 Sojourn Time

- Sojourn or holding time: Is the RV  $H_k$  equal to the sojourn time in state  $k$ :



- The Markov property implies that **the sojourn time is exponentially distributed with parameter  $q_{ii}$** :

$$P(H_i \leq x) = 1 - e^{-q_{ii}x} \Rightarrow P(H_i > x) = e^{-q_{ii}x}, q_{ii} = -\sum_{j \neq i} q_{ij}, x \geq 0 \quad (11.8)$$

#### The exponential distribution satisfies the Markov property

- Markov property (**memoryless**):

$$P(X(t_2) = i | X(t_1) = i, X(0) = i) = P(X(t_2) = i | X(t_1) = i), t_2 > t_1 > 0 \quad (11.9)$$

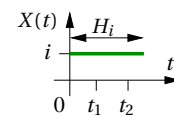


Figure 11.1: Sojourn time.

- In terms of the sojourn time:

$$P(H_i > t_2 | H_i > t_1) = P(H_i > t_2 - t_1)$$

- But:

$$P(H_i > t_2 | H_i > t_1) = \frac{P(H_i > t_2, H_i > t_1)}{P(H_i > t_1)} = \frac{P(H_i > t_2)}{P(H_i > t_1)} = \frac{e^{q_{ii} t_2}}{e^{q_{ii} t_1}} = e^{q_{ii} (t_2 - t_1)} = P(H_i > t_2 - t_1)$$

- The **exponential distribution** is the only one satisfying the **memoryless property**.

## 11.5 Exponential Jumps Description of a CTMC

Assume a process built as follows:

- Upon reaching a state  $i$ 
  1. the process can jump to a state  $j \in \{1, 2, \dots, l\}$ .
  2. A set of **independent exponential RVs**,  $\{H_{i1}, H_{i2}, \dots, H_{il}\}$ , with parameters  $\{q_{i1}, q_{i2}, \dots, q_{il}\}$  are triggered. That is,  $P(H_{ij} \leq t) = 1 - e^{-q_{ij} t}$ ,  $t \geq 0$ .
- If  $\min\{H_{i1}, H_{i2}, \dots, H_{il}\} = H_{ij} \Rightarrow$  the process jumps to the state  $j$ . In other words, a transition occurs to state  $j$  if the RV  $H_{ij}$  is the minimum of  $\{H_{i1}, H_{i2}, \dots, H_{il}\}$  (see figure 11.2).

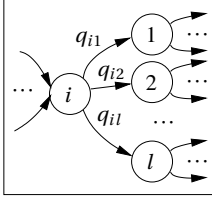


Figure 11.2: Jumps in a CTMC.

**Theorem:** This process is a CTMC with transition rates  $q_{ij}$ .

$$P(H_{ij} \leq t) = 1 - e^{-q_{ij} t}. \quad (11.10)$$

*Proof.* • The RV  $H_i = \min\{H_{i1}, H_{i2}, \dots, H_{il}\}$  (sojourn time in state  $i$ ) is **exponentially distributed** with parameter  $q_i = \sum_j q_{ij}$ :  $P(H_i \leq t) = 1 - e^{-q_i t}$ .

- $P(\min\{H_{i1}, H_{i2}, \dots, H_{il}\} = H_{ij}) = q_{ij} / \sum_j q_{ij}$ . Thus, the **transition rate to state  $j$**  is:

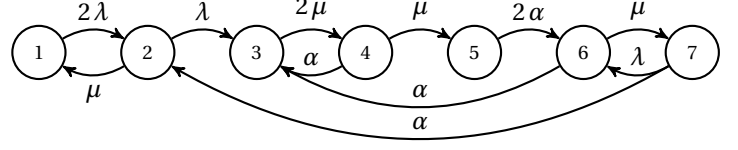
$$\lim_{t \rightarrow 0} \frac{p_{ij}(t)}{t} = \lim_{t \rightarrow 0} \frac{P(\min\{H_{i1}, H_{i2}, \dots, H_{il}\} = H_{ij}) \times P(H_i \leq t)}{t} = \frac{q_{ij}}{\sum_j q_{ij}} \frac{\partial P(H_i \leq t)}{\partial t} \bigg|_{t=0} = \frac{q_{ij}}{\sum_j q_{ij}} \sum_j q_{ij} = q_{ij} \quad (11.11)$$

□

### Example: Pure Aloha System

- Consider a **Pure Aloha System** with **2 nodes**:
  - Nodes in **thinking state** Tx a packet in a time exponentially distributed with rate  $\lambda$ .
  - **Transmission time** is exponentially distributed with rate  $\mu$ .

- If two transmissions overlap, the packet is lost and stations become backlogged (after the packet transmission) until the packet is successfully transmitted.
- Nodes in **backlogged state** Tx a packet in a time exponentially distributed with rate  $\alpha$ .



State	Condition	Legend
1	$T, T$	$T$ Thinking
2	$X, T$	$X$ Transmitting
3	$C, C$	$C$ Collided transmission
4	$B, C$	$B$ Backlogged
5	$B, B$	
6	$X, B$	
7	$T, B$	

## Chapter 12

### Transient Solution

#### 12.1 Chapman-Kolmogorov Equations

- **Chapman-Kolmogorov:**  $p_{ij}(t) = \sum_k p_{ik}(t - \alpha) p_{kj}(\alpha)$ ,  $0 \leq \alpha \leq t$

- Thus:

$$\frac{p_{ij}(t + \Delta t) - p_{ij}(t)}{\Delta t} = \sum_k \left\{ \frac{p_{ik}(t + \Delta t - \alpha) - p_{ik}(t - \alpha)}{\Delta t} p_{kj}(\alpha) \right\}$$

- Taking the limit

$$\alpha \rightarrow t, \Delta t \rightarrow 0 \Rightarrow \begin{cases} p_{ik}(t - \alpha) \rightarrow 0, & i \neq k \\ p_{ik}(t - \alpha) \rightarrow 1, & i = k \end{cases}$$

and using:

$$\begin{cases} q_{ij} = \lim_{t \rightarrow 0} \frac{p_{ij}(t)}{t}, & i \neq j \\ q_{ij} = \lim_{t \rightarrow 0} \frac{p_{ij}(t) - 1}{t}, & i = j \end{cases}$$

we have:

$$\frac{\partial p_{ij}(t)}{\partial t} = \sum_k q_{ik} p_{kj}(t), \quad t \geq 0, \forall i, j \quad (12.1)$$

- In matrix form:

$$\frac{\partial \mathbf{P}(t)}{\partial t} = \mathbf{Q} \mathbf{P}(t), \quad t \geq 0 \quad (12.2)$$

known as the master equations of a CTMC.

- The solution of the previous matrix differential equation is the **exponential matrix**:

$$\mathbf{P}(t) = e^{\mathbf{Q}t} = \sum_{i=0}^{\infty} \frac{(\mathbf{Q}t)^i}{i!} = \mathbf{I} + \mathbf{Q}t + \frac{\mathbf{Q}^2 t^2}{2!} + \frac{\mathbf{Q}^3 t^3}{3!} + \dots, \quad t \geq 0 \quad (12.3)$$

- Due to rounding errors, the previous series is difficult to compute numerically (the powers of  $\mathbf{Q}$  have positive and negative entries).

## 12.2 State Probabilities

- Define the probability of being in state  $i$  at time  $t$ :

$$\pi_i(t) = P(X(t) = i) \quad (12.4)$$

- In vector form (row vector)

$$\boldsymbol{\pi}(t) = (\pi_1(t), \pi_2(t), \dots). \quad (12.5)$$

- Clearly:

$$\pi_i(t) = \sum_k P(X(0) = k) P(X(t) = i | X(0) = k) = \sum_k \pi_k(0) p_{ki}(t) \quad (12.6)$$

- In matrix form:

$$\boldsymbol{\pi}(t) = \boldsymbol{\pi}(0) \mathbf{P}(t) = \boldsymbol{\pi}(0) e^{\mathbf{Q}t}, t \geq 0 \quad (12.7)$$

where  $\boldsymbol{\pi}(0)$  is the **initial distribution**.

- **NOTE:** Compare with DTMC

$$\boldsymbol{\pi}(n) = \boldsymbol{\pi}(0) \mathbf{P}^n, n \geq 0 \quad (12.8)$$

## 12.3 Transient Solution

- If we are interested in the **transient evolution** we shall study  $\boldsymbol{\pi}(t) = \boldsymbol{\pi}(0) \mathbf{P}(t) = \boldsymbol{\pi}(0) e^{\mathbf{Q}t}, t \geq 0$ .
- Assume a **finite CTMC** with  $N$  states (infinitesimal generator  $\mathbf{Q}^{N \times N}$ ).
- Assume that  $\mathbf{Q}$  can be **diagonalized**:  $\mathbf{Q} = \mathbf{L}^{-1} \boldsymbol{\Lambda} \mathbf{L}$ , where  $\boldsymbol{\Lambda}$  is the diagonal matrix  $\boldsymbol{\Lambda} = \text{diag}(\lambda_1, \dots, \lambda_N)$ , with  $\lambda_l, l = 1, \dots, N$  the eigenvalues of  $\mathbf{Q}$ .
- **NOTE:** the **eigenvalues**  $\lambda_l$  of a matrix  $\mathbf{A}$  are scalars that satisfy:  $\mathbf{I}\mathbf{A} = \lambda_l \mathbf{I}$  (or  $\mathbf{A}\mathbf{r} = \lambda_l \mathbf{r}$ ) for some row vectors  $\mathbf{I}$  (column vectors  $\mathbf{r}$ ), referred to as **left and right eigenvectors**, respectively. Thus, solve the **characteristic polynomial**  $\det(\lambda \mathbf{I} - \mathbf{A}) = 0$ .
- Since

$$\mathbf{P}(t) = e^{\mathbf{Q}t} = \sum_{i=0}^{\infty} \frac{(\mathbf{Q}t)^i}{i!} = \sum_{i=0}^{\infty} \frac{(\mathbf{L}^{-1} \boldsymbol{\Lambda} \mathbf{L} t)^i}{i!} = \mathbf{L}^{-1} \text{diag}\left(\sum_{i=0}^{\infty} \frac{(\lambda_1 t)^i}{i!}, \dots\right) \mathbf{L} = \mathbf{L}^{-1} \text{diag}(e^{\lambda_1 t}, \dots) \mathbf{L} \quad (12.9)$$

we have that

$$\boldsymbol{\pi}(t) = \boldsymbol{\pi}(0) e^{\mathbf{Q}t} = \boldsymbol{\pi}(0) \mathbf{L}^{-1} \text{diag}(e^{\lambda_1 t}, e^{\lambda_2 t}, \dots, e^{\lambda_L t}) \mathbf{L} \quad (12.10)$$

- Thus, the probability of being in state  $i$  is given by:

$$\pi_i(t) = (\boldsymbol{\pi}(t))_i = \sum_{l=1}^N a_i^{(l)} e^{\lambda_l t}, t \geq 0 \quad (12.11)$$

where the **unknown coefficients**  $a_i^{(l)}$  can be obtained solving the system of equations:

$$\left. \frac{\partial^n \pi_i(t)}{\partial t^n} \right|_{t=0} = (\boldsymbol{\pi}(0) \mathbf{Q}^n)_i, n = 0, \dots, N-1 \quad (12.12)$$

**NOTE:** Compare with  $(\boldsymbol{\pi}(n))_i = (\boldsymbol{\pi}(0) \mathbf{P}^n)_i, n = 0, \dots, N-1$

### 12.3.1 Eigenvalues of an Infinitesimal Generator

- $\mathbf{Q}$  has an **eigenvalue equal to 0** ( $\mathbf{Q}\mathbf{x} = \lambda\mathbf{x}$ , for  $\lambda = 0, \mathbf{x} \neq \mathbf{0}$ ).  
**Proof:**  $\mathbf{Q}\mathbf{e} = \mathbf{0}$ , where  $\mathbf{e} = (1, 1, \dots)^T$  is a column vector of 1 (all rows of  $\mathbf{Q}$  add to 0).
- The eigenvalue  $\lambda = 0$  is **single** if  $\mathbf{Q}$  is **irreducible** (Perron-Frobenius theorem).  $\mathbf{Q}$  is irreducible if all states communicate: for  $t > 0, p_{ij}(t) > 0, \forall i, j$ .
- All eigenvalues of  $\mathbf{Q}$  are  $\lambda_l \leq 0$ .  
**Proof:** Using Gerschgorin's theorem (see figure 12.1) and the fact that the rows of  $\mathbf{Q}$  add to 0.

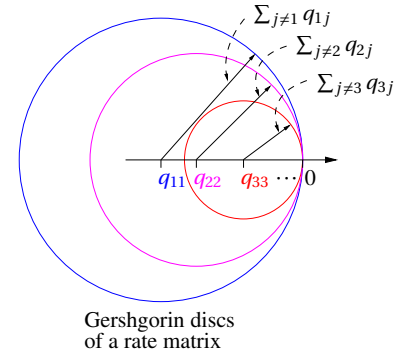


Figure 12.1: Gerschgorin discs of a CTMC

#### Example

- Assume a CTMC with

$$\mathbf{Q} = \begin{bmatrix} -1 & 1 \\ 1/2 & -1/2 \end{bmatrix}$$

- We want the probability of being in state 2 at time  $t$  starting from state 1:  $\pi_2(t)$  with  $\boldsymbol{\pi}(0) = \begin{bmatrix} 1 & 0 \end{bmatrix}$ .

#### Solution

- It can be easily found that the **eigenvalues** of  $\mathbf{Q}$  are  $\lambda_1 = 0$  and  $\lambda_2 = -3/2$ .

$$\pi_2(t) = a e^{\lambda_1 t} + b e^{\lambda_2 t} = a + b e^{-(3/2)t} \quad (12.13)$$

- Imposing the **boundary conditions**:

$$\pi_2(0) = a + b = (\boldsymbol{\pi}(0) \mathbf{Q}^0)_2 = (\boldsymbol{\pi}(0) \mathbf{I})_2 = (\boldsymbol{\pi}(0))_2 = 0$$

$$\left. \frac{\partial \pi_2(t)}{\partial t} \right|_{t=0} = b(-3/2) = (\boldsymbol{\pi}(0) \mathbf{Q})_2 = \mathbf{Q}_{12} = 1 \quad (12.14)$$

we have that  $a = 2/3, b = -2/3$ , thus:

$$\pi_2(t) = 2/3(1 - e^{-(3/2)t}), t \geq 0 \quad (12.15)$$



### 12.3.2 Chain with a Defective Matrix

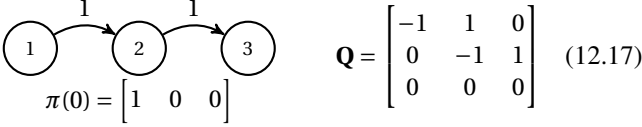
- What if  $\mathbf{Q}$  cannot be diagonalized? (**defective** matrix).
- Let  $\lambda_l$ ,  $l = 1, \dots, L$  be the eigenvalues of  $\mathbf{Q}^{N \times N}$ , each with multiplicity  $k_l$  ( $k_l \geq 1$ ,  $\sum_l k_l = N$ ). Then [1]:

$$\pi_j(t) = \sum_{l=1}^L e^{\lambda_l t} \sum_{m=0}^{k_l-1} a_j^{(l,m)} t^m \quad (12.16)$$

where  $a_j^{(l,m)}$  are constants. So, exponentials associated with eigenvalues  $\lambda_l$  of multiplicity  $k_l > 1$  are multiplied by polynomials in  $t$  of degree  $k_l - 1$ .

#### Example

- Assume the CTMC:



- We have  $\lambda_1 = 0$  and  $\lambda_2 = -1$  with multiplicity 2. Thus:

$$\pi_3(t) = a + e^{-t}(b + c t) \quad (12.18)$$

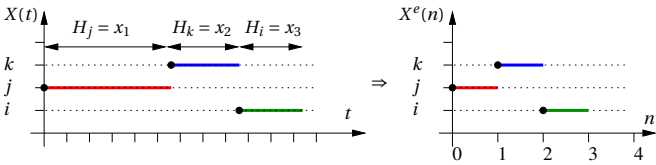
- We have that  $a = 1$ , because state 3 is absorbing. Imposing  $\pi_3(0) = 0$  and  $\pi'_3(0) = 0$ , we have  $b = c = -1$ , and

$$\pi_3(t) = 1 - e^{-t}(1 + t), t \geq 0 \quad (12.19)$$

## Chapter 13

### Embedded MC of a CTMC

#### 13.1 Definition



- We form a discrete time process  $X^e(n)$ , called the **Embedded MC (EMC)**, by looking at a CTMC at the transition time instants.

**Theorem:** This process is a **DTMC** with transition probabilities:

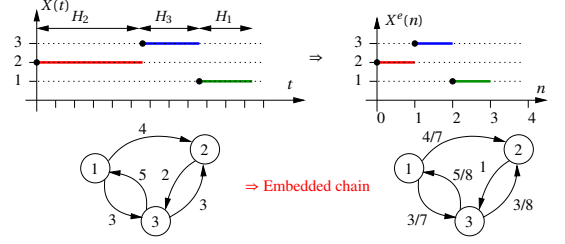
$$p_{ij}^e = \begin{cases} 0, & i = j \\ \frac{q_{ij}}{\sum_{k \neq i} q_{ik}}, & i \neq j \end{cases} \quad (13.1)$$

- NOTE:** If  $i$  is **absorbing** ( $q_{ii} = 0$ ), we define  $p_{ii}^e = 1$ .

*Proof.*

- The EMC satisfies the **memoryless** property.
- Since we look the system only upon transition to a different state:  $p_{ii}^e = 0$ . NOTE: it might be  $p_{ii}^e \neq 0$  if we look at transitions that end up in the same state.
- The probability that there is a transition from state  $i$  to  $j$  in the CTMC is the probability that the exponentially distributed RV with parameter  $q_{ij}$  is the **minimum from the independent exponentially distributed RVs** with parameters  $\{q_{ik}\}_{k \neq i}$ . This probability is  $q_{ij} / \sum_{k \neq i} q_{ik}$ .

#### Example



$$\mathbf{Q} = \begin{bmatrix} -7 & 4 & 3 \\ 0 & -2 & 2 \\ 5 & 3 & -8 \end{bmatrix} \Rightarrow \mathbf{P}_e = \begin{bmatrix} 0 & 4/7 & 3/7 \\ 0 & 0 & 1 \\ 5/8 & 3/8 & 0 \end{bmatrix} \quad (13.2)$$

- Each **transition** in the CTMC is a transition in the EMC.
- One step in  $i$  in the EMC is a **sojourn time**  $H_i$  in the CTMC.

## Chapter 14

### Classification of States

#### 14.1 Irreducibly

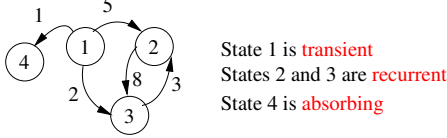
- A state  $j$  is said to **communicate** with  $i$ ,  $i \leftrightarrow j$ , if  $p_{ij}(t_1) > 0$ ,  $p_{ji}(t_2) > 0$  for some  $t_1 \geq 0$ ,  $t_2 \geq 0$ .
- We define an **irreducible closed set, ICS**  $C_k$  as a set where all states communicate with each other, and have no transitions to other states out of the set:  
 $i \leftrightarrow j, \forall i, j \in C_k$  and  $q_{ij} = 0, \forall i \in C_k, j \notin C_k$
- An **absorbing state** form an ICS of only one element. This state,  $i$ , must have  $q_{ij} = 0 \forall i, j$ .
- Transient states** do not belong to any ICS.
- A MC is **irreducible** if all the states form a unique ICS.
- Assume a MC has  $M$  **ICSs**: By properly numbering the states, we can write  $\mathbf{P}$  as an  $M$  block diagonal matrix with the probabilities of the transient states in the last rows.
- Example**, if  $M = 3$ :

$$\mathbf{Q} = \begin{bmatrix} \mathbf{Q}_1 & & \mathbf{0} \\ \mathbf{0} & \mathbf{Q}_2 & \\ & & \mathbf{Q}_3 \\ \text{at least one } > 0 & & \mathbf{T} \end{bmatrix} \Rightarrow \pi(t) = \pi(0) e^{\mathbf{Q}t} = \pi(0) \begin{bmatrix} e^{\mathbf{Q}_1 t} & & \mathbf{0} \\ \mathbf{0} & e^{\mathbf{Q}_2 t} & \\ & & e^{\mathbf{Q}_3 t} \\ \text{at least one } > 0 & & e^{\mathbf{T}t} \end{bmatrix}$$

- Note that the  $M$  **sub-matrices are infinitesimal generators** (their rows add to 0).

## 14.2 Transient and Recurrent

- **Recurrent:** States that, being visited, have a probability  $> 0$  of being visited again. They are visited an infinite number of times when  $t \rightarrow \infty$ .
- **Transient:** States that, being visited, have a probability  $> 0$  of never being visited again. They are visited a finite number of times when  $t \rightarrow \infty$ .
- **Absorbing:** A single (recurrent) state where the chain remains with probability  $= 1$ .



- To derive a classification criteria, we shall study the **embedded MC (EMC)**, and proceed as in DTMC: Let  $f_{ij}^e(n)$  the first passage prob. of the EMC, and  $f_{ij}^e = \sum_{n=1}^{\infty} f_{ij}^e(n)$ .
- If  $f_{ii}^e = 1$  we say  $i$  is a **recurrent state**.
- If  $f_{ii}^e < 1$  we say  $i$  is a **transient state**.
- When  $f_{ij}^e = 1$ , we define the **mean recurrence time of the EMC**  $m_{ij}^e = \sum_{n=1}^{\infty} n f_{ij}^e(n)$ . **NOTE:** in **steps**, not time units.
- If  $m_{ii}^e = \infty$  the state is **null recurrent**.
- If  $m_{ii}^e < \infty$  the state is **positive recurrent**.
- **NOTES:** (i) Even if the EMC is periodic, **there are not periodic CTMC** (it has no sense). (ii)  $f_{ij}^e$  and  $m_{ij}^e$  can be computed using the **recursive equations** for DTMC.

## 14.3 Mean recurrence time of the CTMC

- If the chain is in  $i$  at a time  $t$ , it takes an **expected time to leave  $i$**  equal to  $1/(-q_{ii}) = 1/\sum_{j \neq i} q_{ij}$  (**sojourn time exponentially distributed** with rate  $q_i = -q_{ii} = \sum_{j \neq i} q_{ij}$ ).
- Thus, if the chain is in **state  $i$** , it takes a **mean time to enter state  $j$**  (**mean first passage time**):

$$m_{ij} = \frac{1}{q_i} + \sum_{k \neq i} p_{ik}^e m_{kj} \quad (14.1)$$

- Since:  $p_{ij}^e = \begin{cases} 0, & i = j \\ \frac{q_{ij}}{\sum_{j \neq i} q_{ij}} = \frac{q_{ij}}{q_i}, & i \neq j \end{cases}$  we have:

$$m_{ij} = \frac{1}{q_i} + \sum_{k \neq i} p_{ik}^e m_{kj} = \frac{1}{q_i} + \sum_{k \neq i} \frac{q_{ik}}{q_i} m_{kj} \text{ [time units]} \quad (14.2)$$

## Chapter 15

### Steady State

#### 15.1 Limiting Distribution

- The probability to be in state  $i$  at time  $t$  is:

$$\pi_i(t) = P(X(t) = i) = \sum_k \pi_k(0) p_{ki}(t), t \geq 0 \quad (15.1)$$

- In matrix form:

$$\boldsymbol{\pi}(t) = \boldsymbol{\pi}(0) \mathbf{P}(t) = \boldsymbol{\pi}(0) e^{\mathbf{Q}t}, t \geq 0 \quad (15.2)$$

- Assume that the following limit exists:

$$\boldsymbol{\pi}(\infty) = \lim_{t \rightarrow \infty} \boldsymbol{\pi}(t) = \lim_{t \rightarrow \infty} \boldsymbol{\pi}(0) \mathbf{P}(t) = \boldsymbol{\pi}(0) \lim_{t \rightarrow \infty} e^{\mathbf{Q}t} \quad (15.3)$$

- for any  $\boldsymbol{\pi}(0)$ , which implies

$$\lim_{t \rightarrow \infty} e^{\mathbf{Q}t} = \mathbf{P}(\infty) = \begin{bmatrix} \boldsymbol{\pi}(\infty) & \cdots & \boldsymbol{\pi}(\infty) \end{bmatrix}^T \quad (15.4)$$

- If this limit exists, we call  $\mathbf{P}(\infty)$  the **limiting matrix**, and  $\boldsymbol{\pi}(\infty)$  the **limiting distribution**.
- $\mathbf{P}(\infty) = \begin{bmatrix} \boldsymbol{\pi}(\infty) & \cdots & \boldsymbol{\pi}(\infty) \end{bmatrix}^T$  does not exist if the CTMC has more than one irreducible closed set (each ICS will converge to a diagonal block, and  $\boldsymbol{\pi}(\infty)$  will depend on  $\boldsymbol{\pi}(0)$ ).

## 15.2 Stationary Distribution

- We have:  $\boldsymbol{\pi}(t) = \boldsymbol{\pi}(0) e^{\mathbf{Q}t}, t \geq 0$ .
- In steady state the probabilities do not change. We look for a probability vector  $\boldsymbol{\pi} = \boldsymbol{\pi}(t_1)$  satisfying:  $\boldsymbol{\pi}(t_1) e^{\mathbf{Q}t} = \boldsymbol{\pi}(t_1)$ . In other words, for  $t \geq t_1$  the probability vector reach the steady state  $\boldsymbol{\pi}$ , and do not change anymore. Thus:

$$\boldsymbol{\pi} \frac{\partial e^{\mathbf{Q}t}}{\partial t} = \boldsymbol{\pi} \mathbf{Q} e^{\mathbf{Q}t} = \mathbf{0} \quad (15.5)$$

- and we obtain that the **stationary distribution  $\boldsymbol{\pi}$**  can be **computed with the Global balance equations**:

$$\begin{aligned} \boldsymbol{\pi} \mathbf{Q} &= \mathbf{0} \\ \boldsymbol{\pi} \mathbf{e} &= 1, \mathbf{e}^T = (1, 1, \dots) \end{aligned} \quad (15.6)$$

- **NOTE:** Compare with DTMC  $\boldsymbol{\pi} = \boldsymbol{\pi} \mathbf{P}, \boldsymbol{\pi} \mathbf{e} = 1$ .

## 15.3 Numerical Solution

- **Replace one equation method:**

$$\begin{aligned} \boldsymbol{\pi} \mathbf{Q} &= \mathbf{0} \\ \boldsymbol{\pi} \mathbf{e} &= 1, \mathbf{e}^T = (1, 1, \dots) \end{aligned} \quad (15.7)$$

- We solve the equation  $\boldsymbol{\pi} \mathbf{Q} = \mathbf{0}$  replacing the last equation by  $\boldsymbol{\pi} \mathbf{e} = 1$ :

$$\boldsymbol{\pi} \begin{bmatrix} q_{11} & q_{12} & \cdots & q_{1n-1} & 1 \\ q_{21} & q_{22} & \cdots & q_{2n-1} & 1 \\ \vdots & \vdots & & \vdots & \vdots \\ q_{n1} & q_{n2} & \cdots & q_{nn-1} & 1 \end{bmatrix} = \begin{bmatrix} 0 & 0 & \cdots & 0 & 1 \end{bmatrix} \quad (15.8)$$

- **Replace one equation method:**

$$\mathbf{Q} = \begin{bmatrix} -2 & 1 & 1 \\ 1 & -2 & 1 \\ 1 & 1 & -2 \end{bmatrix} \quad \begin{aligned} \boldsymbol{\pi} \mathbf{Q} &= \mathbf{0} \\ \boldsymbol{\pi} \mathbf{e} &= 1 \end{aligned}$$



- Solving with **octave** (matlab clone):

```
octave:1> Q=[ -2,1,1;1,-2,1;1,1,-2];
octave:2> s=size(Q,1); # number of rows.
octave:3> [zeros(1,s-1),1] / ...
> [Q(1:s,1:s-1), ones(s,1)]
ans =
0.33333 0.33333 0.33333
```

- With **R**

```
> Q <- matrix(nc=3, byr=T, c
(-2,1,1,1,-2,1,1,1,-2))
> s <- nrow(Q)
> solve(t(cbind(Q[,1:(s-1)], rep(1,s))), c(rep
(0,s-1),1))
[1] 0.3333333 0.3333333 0.3333333
```

## 15.4 Global balance equations

- Why are they called Global balance equations?

$$\left. \begin{aligned} \pi Q = 0 &\Rightarrow \sum_{i=0}^{\infty} \pi_i q_{ij} = 0 \\ \sum_{i=0}^{\infty} q_{ji} = 0 &\Rightarrow \pi_j \sum_{i=0}^{\infty} q_{ji} = 0 \end{aligned} \right\} \Rightarrow \pi_j \sum_{i=0}^{\infty} q_{ji} = \sum_{i=0}^{\infty} \pi_i q_{ij} \quad (15.9)$$

$$\sum_{i=0}^{\infty} \pi_i q_{ij} \Rightarrow \text{Frequency of transitions entering state } j$$

$$\pi_j \sum_{i=0}^{\infty} q_{ji} \Rightarrow \text{Frequency of transitions leaving state } j \quad (15.10)$$

- In **stationary regime**, the frequency of transitions leaving state  $j$  is equal to the frequency of transitions entering state  $j$ .

### 15.4.1 Solving using flux balancing

- Define the **flux**  $F_{uv}$  from state  $u$  to  $v$ :

$$F_{uv} = \pi_u q_{uv} \quad (15.11)$$

- and the flux from set of states  $U$  to  $V$ :

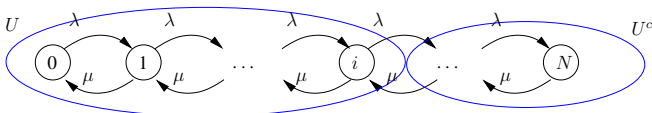
$$F(U,V) = \sum_{u \in U} \sum_{v \in V} F_{uv} \quad (15.12)$$

- From the Global balance equations, and reasoning exactly as in DTMC:

$$F(U,U^c) = F(U^c,U) \quad (15.13)$$

- **NOTE:** Same equation as in DTMC, changing  $p_{ij}$  by  $q_{ij}$ .

#### Example: Birth-dead Process



- Flux balancing  $\Rightarrow \lambda \pi_i = \mu \pi_{i+1}$
- Iterating:

$$\pi_i = \pi_0 \rho^i, i = 0, 1, \dots, N-1, \rho = \frac{\lambda}{\mu} \quad (15.14)$$

- Normalizing:

$$\pi_0 = \frac{1-\rho}{1-\rho^N} \quad (15.15)$$

## 15.5 Ergodic Chains

- **Ergodic state:** positive recurrent ( $f_{ii}^e = 1, m_{ii}^e < \infty$ ).
- **Ergodic chain** if all states are ergodic.
- **Theorem:** All states of an irreducible Markov chain are of the same type: Transient or positive/null recurrent (see [2, chapter XV]).
- **Consequences:**

- **Finite irreducible chains are ergodic** (since all states are positive recurrent).
- **Infinite irreducible chains can be:**
  - \* **Ergodic:** all the states are positive recurrent (stable chains).
  - \* **Non ergodic:** all states are null recurrent or transient (unstable chains).

### Theorems for ergodic chains

- $\pi = \pi(\infty)$ . Proof:  $\pi(\infty)$  satisfies the GBE.
- In stationary regime (when  $\pi = \pi e^{Qt}$ ), the **mean number of time the system remains in state  $j$**  during  $T$  time units is given by

$$T \pi_j \quad (15.16)$$

thus,  $\pi_j$  is the fraction of time the chain remains in state  $j$ . The proof is analogous to DTMC.

- **NOTE:** The relation of DTMC between **mean recurrence time** and stationary probabilities does not hold for CTMC. I.e., the mean number of time units between two consecutive visits to state  $j$ ,  $m_{jj}$ , **cannot be computed as  $1/\pi_j$** . It must be computed with the **recursive equations** (slide 21).

## 15.6 Reversible Chains

- Let  $X(t)$  be an **ergodic** MC. The chain  $X^r(t) = X(-t)$  is referred to as the **time reversal chain** of  $X(t)$ .
- The same results obtained for DTMC reversed chains apply to CTMC, changing  $p_{ij}$  by  $q_{ij}$ :
  - The reversed chain transition rates  $q_{ij}^r$ , given by:

$$\pi_i q_{ij} = \pi_j q_{ji}^r \quad (15.17)$$

satisfy the **reversed balance equations**:  $F(U,V) = F^r(V,U)$

- A chain is **reversible** if:

$$q_{ij} = q_{ji}^r \quad (15.18)$$

- Reversible chains satisfy the **detailed balance equations**:

$$F(U, V) = F(V, U), \forall (V, U), V \cap U = \emptyset \quad (15.19)$$

- The same results obtained for DTMC Reversible Chains apply to CTMC: **Kolmogorov Criteria** and **Product Form Solution** for the stationary distribution (changing  $p_{ij}$  by  $q_{ij}$ ).

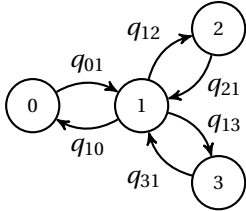
- E.g. the **stationary probabilities** are given by:

- Choose a state  $\mathbf{s} \in S$ ,
- For every other state  $\mathbf{i} \in S$ ,  $\mathbf{i} \neq \mathbf{s}$  look for a possible path  $l_i$  from state  $\mathbf{s}$  to state  $\mathbf{i}$ :

$$\mathbf{s} = (l_i, 1) \rightsquigarrow (l_i, 2) \rightsquigarrow \dots \rightsquigarrow (l_i, m_{l_i}) = \mathbf{i}, m_{l_i} \geq 1$$

$$\pi_i = \frac{\psi_i}{\sum_{j \in S} \psi_j}, i \in S \quad \text{where } \psi_i = \begin{cases} 1, & i = s \\ \prod_{k=1}^{m_{l_i}-1} \frac{q_{(l_i, k)(l_i, k+1)}}{q_{(l_i, k+1)(l_i, k)}}, & i \neq s \end{cases} \quad (15.20)$$

### 15.6.1 Example



- An ergodic tree is always reversible, thus

$$\pi_0 = \frac{1}{G}, \pi_1 = \frac{1}{G} \frac{q_{01}}{q_{10}}, \pi_2 = \frac{1}{G} \frac{q_{01}}{q_{10}} \frac{q_{12}}{q_{21}}, \pi_3 = \frac{1}{G} \frac{q_{01}}{q_{10}} \frac{q_{13}}{q_{31}}. \quad (15.21)$$

- Normalizing:

$$G = 1 + \frac{q_{01}}{q_{10}} + \frac{q_{01}}{q_{10}} \frac{q_{12}}{q_{21}} + \frac{q_{01}}{q_{10}} \frac{q_{13}}{q_{31}} \quad (15.22)$$

## Chapter 16

### Semi-Markov Process

- Define the continuous RV  $H_i$  equal to the **sojourn time** in state  $i$ .
- In a **semi-Markov process** we leave the  $H_i$  distribution to be **arbitrary**. If  $H_i$  is exponentially distributed, we have a CTMC.
- **NOTE: If  $H_i$  is not exponentially distributed**, considering only the current state does not satisfy the Markov property (memoryless) since **the evolution of the process depends on the current state and the sojourn time in the state**:  $(i, t_i)$ .
- If we consider  $(i, t_i)$  as the state, the state would satisfy the Markov property, but we would have a **Markov process** (since  $t_i$  is not a discrete RV).

### 16.1 Embedded MC (EMC) of a semi-Markov process

- **Embedded MC (EMC) of the process: We only look at the state transition instants.**

- **The EMC is a DTMC** with transition probabilities  $p_{ij}^e$ .
- The **time step is variable**.
- There are not self transitions ( $p_{ii}^e = 0$ ), unless we look at some memoryless event that produce a self transition.
- **Theorem:** let  $\pi_i^e$  and  $\pi_i$  be the stationary distribution of the EMC and the semi-Markov process respectively. Let  $E[H_i]$  be the mean sojourn time in state  $i$ , then:

$$\pi_i = \frac{\pi_i^e E[H_i]}{\sum_j \pi_j^e E[H_j]} \quad (16.1)$$

**NOTE:** By *stationary distribution* for the semi-Markov process we mean to the long-run proportion of time that the process is in each state.

*Proof.*

- **For  $n$  steps of the EMC**, define:

- $f_i(n)$ : **proportion of time** the process is in state  $i$ .
- $N_i(n)$ : **number of visits** to state  $i$ .
- $H_i(l)$ : **sojourn time** in state  $i$  in the **visit number**  $l$ .

$$f_i(n) = \frac{\sum_{l=1}^{N_i(n)} H_i(l)}{\sum_j \sum_{l=1}^{N_j(n)} H_j(l)} = \frac{\frac{N_i(n)}{n} \sum_{l=1}^{N_i(n)} \frac{H_i(l)}{N_i(n)}}{\sum_j \frac{N_j(n)}{n} \sum_{l=1}^{N_j(n)} \frac{H_j(l)}{N_j(n)}} \Rightarrow \pi_i = \frac{\pi_i^e E[H_i]}{\sum_j \pi_j^e E[H_j]}$$

- since:

$$\lim_{n \rightarrow \infty} f_i(n) = \pi_i$$

$$\lim_{n \rightarrow \infty} \sum_{l=1}^{N_i(n)} \frac{H_j(l)}{N_i(n)} = E[H_i]$$

$$\lim_{n \rightarrow \infty} \frac{N_i(n)}{n} = \pi_i^e. \quad \square$$

### 16.1.1 Example

Suppose the system:



- Packets arrive deterministically every  $T$  time units.
- Upon a packet arrival it goes immediately into service if the server is empty, and it is lost if the server is busy.
- Services are exponentially distributed with rate  $\mu$ .

Define a semi-Markov process with states  $\begin{cases} 0 & \text{server empty,} \\ 1 & \text{server busy.} \end{cases} \quad (16.2)$

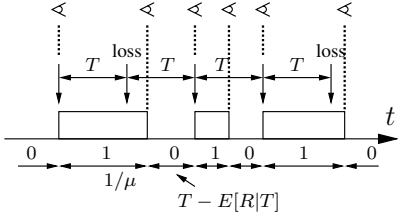
1. Derive the EMC and stationary distribution of the EMC and continuous time process.
2. Compute the throughput and loss probability.

**Hint:** The distribution of an event  $R$  exponentially distributed with rate  $\mu$ , given that occurs in an interval  $t \in [0, T]$ , is

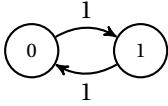
$$F_R(t|T) = \frac{P(R \leq t, R \leq T)}{P(R \leq T)} = \frac{P(R \leq t)_{t \in [0, T]}}{P(R \leq T)} = \frac{1 - e^{-\mu t}}{1 - e^{-\mu T}}, \quad t \in [0, T], \text{ and}$$

$$E[R|T] = \int_0^T (1 - F_R(t|T)) dt = \frac{1 - \alpha - \alpha \mu T}{\mu(1 - \alpha)}, \text{ where } \alpha = e^{-\mu T}.$$

### Solution



EMC:



- $\pi_0^e = \pi_1^e = 1/2$ ,
- $E[H_0] = T - E[R|T] = \frac{T\mu - (1 - \alpha)}{\mu(1 - \alpha)}$ ,  $E[H_1] = \frac{1}{\mu}$ .

And the continuous time process:

$$G = \pi_0^e E[H_0] + \pi_1^e E[H_1] = \frac{T}{2(1 - \alpha)} \left[ \frac{\text{time units}}{\text{step}} \right] \quad (16.3)$$

$$\pi_0 = \frac{\pi_0^e E[H_0]}{G} = \frac{\mu T - (1 - \alpha)}{\mu T}, \quad \pi_1 = \frac{\pi_1^e E[H_1]}{G} = \frac{1 - \alpha}{\mu T}$$

$$\text{Throughput: } S = \mu \pi_1 = \frac{1 - \alpha}{T} \quad (\text{check } S = \frac{1}{E[H_0] + E[H_1]})$$

$$\text{Loss probability: } S = \frac{1}{T} (1 - p_L), \quad p_L = 1 - S T = \alpha = e^{-\mu T}. \quad (16.4)$$

### 16.2 Embedded MC of a CTMC

- Assume that the semi-Markov process is a **CTMC** (sojourn times are exponentially distributed).
- The transition probabilities  $p_{ij}^e$  of the EMC, and the stationary distribution  $\pi_i$  of the CTMC are given by:

$$\begin{cases} p_{ij}^e = q_{ij}/q_i, & i \neq j \\ p_{ij}^e = 0, & i = j \end{cases} \quad \pi_i = \frac{\pi_i^e/q_i}{\sum_k \pi_k^e/q_k} \quad (16.5)$$

$$\text{where: } q_i = \sum_{k \neq i} q_{ik} = -q_{ii}.$$

*Proof.*

- The equations for  $p_{ij}^e$  was proven in a previous section.

- The distribution of the **sojourn time in state  $i$  in the CTMC** is the distribution of the **minimum of independent exponentially distributed RV** with parameters  $\{q_{ik}\}_{k \neq i}$ . This distribution is **exponentially distributed with parameter  $\sum_{k \neq i} q_{ik}$** . Thus  $E[H_i] = 1/\sum_{k \neq i} q_{ik} = 1/q_i$ ,  $q_i = \sum_{k \neq i} q_{ik}$ .
- Substituting:

$$\pi_i = \frac{\pi_i^e E[H_i]}{\sum_j \pi_j^e E[H_j]} = \frac{\pi_i^e/q_i}{\sum_j \pi_j^e/q_j} \quad (16.6)$$

$$\text{where: } q_i = \sum_{k \neq i} q_{ik} = -q_{ii}. \quad \square$$

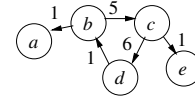
## Chapter 17

### Finite Absorbing Chains

#### 17.1 Canonical Form

- Let  $\mathbf{Q}^{r \times r}$  be the transition probability matrix of a chain with a set  $S$  of  $s$  **transient** states and a set  $A$  of  $r - s$  **absorbing** states. We can write  $\mathbf{Q}^{r \times r}$  in the **canonical** form:

$$\mathbf{Q}^{r \times r} = \begin{bmatrix} \mathbf{T}^{s \times s} & \mathbf{R}^{s \times r-s} \\ \mathbf{0}^{r-s \times s} & \mathbf{0}^{r-s \times r-s} \end{bmatrix} \quad (17.1)$$



$$\mathbf{P} = \begin{array}{c} \begin{matrix} b & c & d & a & e \end{matrix} \\ \begin{matrix} b \\ c \\ d \\ a \\ e \end{matrix} \left[ \begin{array}{ccccc} -6 & 5 & 0 & 1 & 0 \\ 0 & -7 & 6 & 0 & 1 \\ 1 & 0 & -1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{array} \right] \end{array} \quad (17.2)$$

#### 17.2 Results

- Define:

$$\mathbf{n}_{ij} = \left\{ \begin{array}{l} \text{time in state } j \text{ before absorption,} \\ \text{starting from state } i \end{array} \right\}, \quad (17.3)$$

$$\mathbf{t}_i = \left\{ \begin{array}{l} \text{time in transient states before} \\ \text{absorption, starting from state } i \end{array} \right\},$$

$$\mathbf{b}_{ij} = P(\text{probability to be absorbed } j \text{ starting } i).$$

- Then:

$$\begin{aligned} \{E[n_{ij}]\} &= \mathbf{N} = -\mathbf{T}^{-1} \\ \{E[t_i]\} &= \mathbf{\tau} = \mathbf{N}\mathbf{e} \\ \{b_{ij}\} &= \mathbf{B} = \mathbf{N}\mathbf{R} \end{aligned} \quad (17.4)$$

where  $\{a_{ij}\}$  is a matrix with  $a_{ij}$  as element  $ij$  and  $\mathbf{e}$  is a column vector of 1s.  $\mathbf{N}$  is called the **fundamental matrix**.

*Proof.*

Let  $S$  be the set of **transient** states and  $A$  the set of **absorbing** states.

- $\{\mathbf{E}[n_{ij}]\} = \mathbf{N} = -\mathbf{T}^{-1}$

$$\begin{aligned} \mathbf{E}[n_{ij}] &= \frac{\delta_{ij}}{-q_{ii}} + \sum_{\substack{k \neq i \\ k \in S}} p_{ik}^e \mathbf{E}[n_{kj}] = \\ &= \frac{\delta_{ij}}{-q_{ii}} + \sum_{\substack{k \neq i \\ k \in S}} \frac{q_{ik}}{-q_{ii} - q_{ii}} \mathbf{E}[n_{kj}] \Rightarrow \\ &= - \sum_{k \in S} q_{ik} \mathbf{E}[n_{kj}] = \delta_{ij} \Rightarrow \\ &= -\mathbf{T} \{\mathbf{E}[n_{ij}]\} = \mathbf{I} \Rightarrow \{\mathbf{E}[\mathbf{n}_{ij}]\} = -\mathbf{T}^{-1} = \mathbf{N} \end{aligned}$$

- $\{\mathbf{E}[t_i]\} = \boldsymbol{\tau} = \mathbf{N}\mathbf{e}$

$$\mathbf{E}[t_i] = \sum_{k \in S} \mathbf{E}[n_{ik}] \Rightarrow \{\mathbf{E}[\mathbf{t}_i]\} = \boldsymbol{\tau} = \mathbf{N}\mathbf{e} \quad (17.5)$$

- $\{b_{ij}\} = \mathbf{B} = \mathbf{N}\mathbf{R}$

$$\begin{aligned} b_{ij} &= p_{ij}^e + \sum_{\substack{k \neq i \\ k \in S}} p_{ik}^e b_{kj} = \\ &= \frac{q_{ij}}{-q_{ii}} + \sum_{\substack{k \neq i \\ k \in S}} \frac{q_{ik}}{-q_{ii} - q_{ii}} b_{kj}, i \in S, j \in \mathcal{A} \Rightarrow \\ &= - \sum_{k \in S} q_{ik} b_{kj} = q_{ij} \Rightarrow -\mathbf{T} \{b_{ij}\} = \mathbf{R} \Rightarrow \\ &= \{\mathbf{b}_{ij}\} = \mathbf{B} = -\mathbf{T}^{-1} \mathbf{R} = \mathbf{N}\mathbf{R}. \quad \square \end{aligned}$$

### 17.3 Extension of the Results

- The previous results can be generalized to any group of states of  $\mathbf{Q}$ :
- A set  $S$  is referred to as **open** if the chain can reach some state of  $S^c$  starting from any state of  $S$ . Let

$$\begin{aligned} \mathbf{T} &= \{q_{ij}, i \in S, j \in S\} \\ \mathbf{R} &= \{q_{ij}, i \in S, j \in S^c\} \end{aligned} \quad (17.6)$$

Let assume that the process starts from  $i \in S$ . Define:

$$\begin{aligned} \mathbf{n}_{ij} &= \left\{ \begin{array}{l} \text{number of steps in state } j \text{ before} \\ \text{leaving } S, \text{ starting from state } i \end{array} \right\}, \\ &\Rightarrow \boxed{\{\mathbf{E}[n_{ij}]\} = \mathbf{N} = -\mathbf{T}^{-1}}. \end{aligned} \quad (17.7)$$

- Similarly for the other results, e.g.  $\boldsymbol{\tau} = \{\mathbf{E}[t_i]\} = \mathbf{N}\mathbf{e}$  and  $\mathbf{B} = \{b_{ij}\} = \mathbf{N}\mathbf{R}$ .

## Part IV

## Queuing Theory

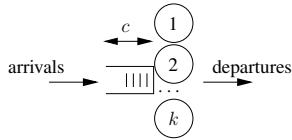
### Chapter 18

#### Introduction

Kendal Notation:  $A/S/k/[c/p]$

$$A/S/k/[c/p] \quad (18.1)$$

- **A**: arrival process,
- **S**: service process,
- **k**: number of servers,
- **c**: maximum number in the system (number of servers + queue size). Note: some authors use the queue size.
- **p**: population.  
If “c” or “p” are missing, they are assumed to be **infinite**.



#### Common arrivals/departures processes

- **G**: general (non specific process is assumed),
- **M**: Markovian (exponentially or geometrically distributed),
- **D**: deterministic,
- **P**: Poisson (discrete RV,  $N$ , equal to the number of arrivals exponentially dist. in a time  $t$ ):

$$P_p(N = n, t) = \frac{(\lambda t)^n e^{-\lambda t}}{n!}, \quad n \geq 0, t \geq 0. \quad (18.2)$$

- **Er**: Erlang (continuous RV equal to the time  $t$  that last  $n$  arrivals exponentially dist.):

$$f_e(t) = \lambda P_p(N = n - 1, t) = \frac{\lambda^n t^{n-1} e^{-\lambda t}}{(n-1)!}, \quad t \geq 0, n \geq 1 \quad (18.3)$$

#### Examples

- **M/M/1**: M. arr. / M. serv. / 1 server,  $\infty$  queue and population.
- **M/G/1**: M. arr. / Gen. serv. / 1 server,  $\infty$  queue and population.

### Chapter 19

#### Fundamental Theorems

##### 19.1 Little Theorem

- Define the stochastic processes:
  - $A(t)$ : number of arrivals  $[0, t]$ .
  - $T_n$ : time in the system (response time) for customer  $n$ .
  - $N(t)$ : number in the system at time  $t$ .

- And the mean values:

- Mean number of customers in the system:

$$N = \lim_{t \rightarrow \infty} \frac{1}{t} \int_0^t N(s) ds \quad (19.1)$$

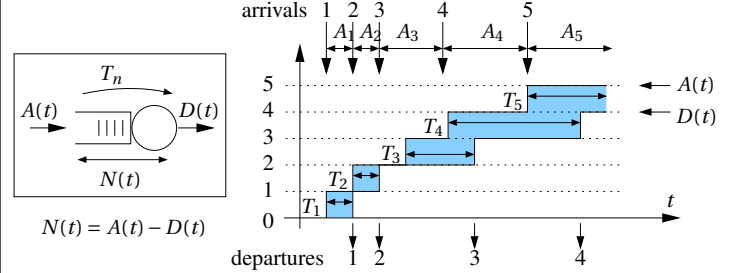
- Arrival rate:  $\lambda = \lim_{t \rightarrow \infty} A(t)/t$
- Mean time in the system:  $T = \lim_{t \rightarrow \infty} (\sum_n T_n) / A(t)$

- The following relation follows:

$$N = \lambda T \quad (19.2)$$

**Mnemonic: NAT** (Number = Arrivals x Time).

*Proof.* (Graphical proof)



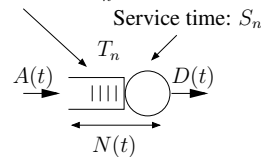
- From the graph we have:

$$\frac{1}{t} \int_0^t N(s) ds = \frac{1}{t} \sum_{i=1}^{A(t)} T_i = \frac{A(t)}{t} \frac{\sum_{i=1}^{A(t)} T_i}{A(t)} \quad (19.3)$$

- Taking the limit  $t \rightarrow \infty$ :  $N = \lambda T$  □

#### Application to the waiting line and the server

Waiting time in the queue  
of customer  $n$ :  $W_n$



Time in the system:

$$T_n = W_n + S_n$$

Expected value:

$$T = W + S$$

where

$$T = E[T_n], W = E[W_n],$$

$$S = E[S_n]$$

- We can apply the Little theorem to the **waiting line** and the **server**:

- **Mean number of customers in the queue**:  $N_Q = \lambda W$ .

- **Mean number of customers in the server**:  $N_S = \rho = \lambda S$ .

#### Mean number in the Server

- In a **single server queue** (even if not Markovian):

$$\rho = N_S = E[N_S(t)] = \lambda E[S] \quad (19.4)$$

$$E[N_S(t)] = 0 \times \pi_0 + 1 \times (1 - \pi_0) = 1 - \pi_0 \Rightarrow \pi_0 = 1 - \rho$$

- $\rho = N_S = \lambda E[S] = 1 - \pi_0$  is the proportion of time the system is busy, in other words, is the **server utilization or load**.

## 19.2 PASTA Theorem

PASTA: Poisson Arrivals See Time Averages

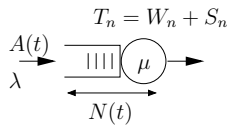
- The mean time the chain is in state  $i$  is  $\pi_i \Rightarrow$  using **PASTA**, the **probability that a Markovian arrival see the system in state  $i$  is  $\pi_i$**  (proof: see [9]).
- The equivalent theorem in **discrete time** is the **arrival theorem**, **RASTA**: Random Arrivals See Time Averages: the **probability that a random arrival see the system in state  $i$  is  $\pi_i$** .

### Example of PASTA

- Assume that a system can have, at most,  $N$  **customers** (e.g  $N-1$  in the queue and 1 in service).
- Assume that an arrival is **lost** when the system is full.
- By **PASTA** the proportion of Poisson arrivals that see the system full, and are lost, is equal to the proportion of time the system has  $N$  in the system,  $\pi_N$ .
- Thus, the **loss probability** is  $\pi_N$ .

## Chapter 20

### The M/M/1 Queue



- Markovian **arrivals** with rate  $\lambda \Rightarrow$  the **interarrival time** is exponentially distributed with mean  $1/\lambda$ :

$$P\{A_n \leq x\} = 1 - e^{-\lambda x}, x \geq 0 \quad (20.1)$$

$\Rightarrow A(t)$  is a **Poisson process**:  $P(A(t) = i) = \frac{(\lambda t)^i}{i!} e^{-\lambda t}, i \geq 0, t \geq 0$

- Markovian Services** with rate  $\mu \Rightarrow$  **service time** exponentially distributed with mean  $1/\mu$ :

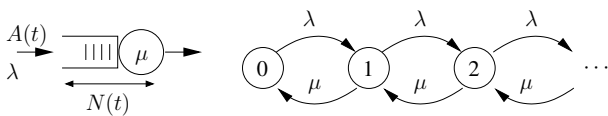
$$P\{S_n \leq x\} = 1 - e^{-\mu x}, x \geq 0 \quad (20.2)$$

y

### Q-matrix

- The process  $N(t) = \{\text{number in the system at time } t \geq 0\}$  is a **CTMC**.

OBSERVATION: for a non Markovian service, the process  $N(t)$  would not be a MC! State transition diagram:



- Q-matrix:

$$Q = \begin{bmatrix} -\lambda & \lambda & 0 & 0 & \dots \\ \mu & -(\mu + \lambda) & \lambda & 0 & \dots \\ 0 & \mu & -(\mu + \lambda) & \lambda & \dots \\ \dots & \dots & \dots & \dots & \dots \end{bmatrix} \quad (20.3)$$

### Stationary Distribution

- Solving the M/M/1 queue using flux balancing (or the general solution of a reversible chain):

$$\pi_i = (1 - \rho) \rho^i, i = 0, \dots, \infty \quad (20.4)$$

where  $\rho = \frac{\lambda}{\mu}$

### Properties

- Mean **customers in the system**:

$$N = \lim_{t \rightarrow \infty} \frac{1}{t} \int_0^t N(s) ds = \sum_{i=0}^{\infty} i \pi_i = \sum_{i=0}^{\infty} i (1 - \rho) \rho^i = \frac{\rho}{1 - \rho} \quad (20.5)$$

- Mean **time in the system** (response time):

$$\text{Little: } N = \lambda T \Rightarrow T = \frac{N}{\lambda} = \frac{\rho}{\lambda(1 - \rho)} = \frac{1}{\mu - \lambda}$$

- Mean **time in the queue**:  $W = T - \frac{1}{\mu} = \frac{\rho}{\mu - \lambda}$

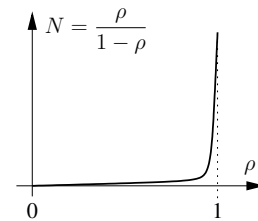
- Mean **Number in the queue**:  $N_Q = \lambda W = \frac{\rho^2}{1 - \rho}$

- Mean **number in the server**:  $N_s = N - N_Q = \rho$

NOTE:  $\pi_0 = 1 - \rho$

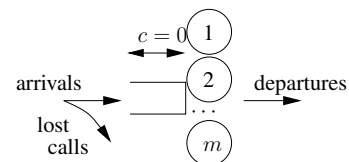
### Stability

- $N$  and  $T$  are proportional to  $1/(1 - \rho) \Rightarrow$  when  $\rho \rightarrow 1 \Rightarrow N, T \rightarrow \infty$ .
- The process  $N(t)$  is **positive recurrent**, **null recurrent** or **transient** according to whether  $\rho = \lambda/\mu$  is below, equal or greater than 1, respectively.

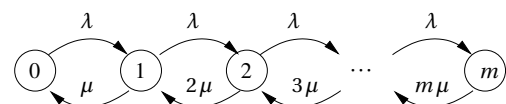


### Example: Loss probability in a telephone switching center

- Hypothesis: Switching center with  $m$  **circuits** and “lost call”, infinite population, Markovian arrivals with rate  $\lambda$  and exponentially distributed call duration with mean  $1/\mu \Rightarrow$  **M/M/m/m** queue.



- Since the minimum of  $i$  independent and identically exponentially distributed RV with parameter **service time** is exponentially distributed with parameter  $i\mu$ :



- Stationary Distribution of the queue M/M/m/m:

- Solving using the **general solution of a reversible chain**:

$$\text{Define } \rho_k = \frac{\lambda}{(k+1)\mu}, k = 0, \dots, m-1$$

$$\pi_0 = \frac{1}{G}, \pi_i = \frac{1}{G} \prod_{k=0}^{i-1} \rho_k = \frac{1}{G} \left(\frac{\lambda}{\mu}\right)^i \frac{1}{i!}, 0 < i \leq m \Rightarrow \quad (20.6)$$

$$\pi_i = \frac{1}{G} \left(\frac{\lambda}{\mu}\right)^i \frac{1}{i!}, 0 \leq i \leq m. G = \sum_{k=0}^m \left(\frac{\lambda}{\mu}\right)^k \frac{1}{k!}. \quad (20.7)$$

- Using **PASTA** Theorem (Poisson Arrivals See Time Average): the **loss call probability** is the probability that the queue is in state  $m$ :  $\pi_m$ , “**Erlang B Formula**”.

## Chapter 21

### M/G/1 Queue

- The process  $N(t) = \{\text{number in the system at time } t \geq 0\}$  in general it is not a MC (it is so only if G is Markovian).
- We can build a **semi-Markov process** observing the system at **departure times**  $t_n$  (note that  $t_n$  are also the service completion times). Define the discrete time process:

$$X(n) = \{\text{number in the system at time } t_n \geq 0, n = 0, 1, \dots\} \quad (21.1)$$

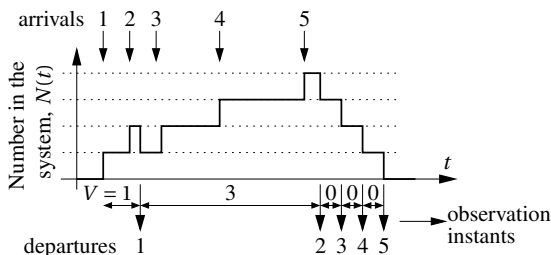
- Theorem:** The process  $X(n)$  is a **DTMC**.
- Proof:**  $X(n)$  only depends on the number of **arrivals in non overlapping intervals**. Since arrivals are Markovian, this is a **memoryless** process.
- NOTE:** Looking at **departure times** the chain may have **self transitions** (in contrast to observing at transition times): we can have the same number in the system after a departure.

#### Transition Probability Matrix

- Let  $f_S(x)$ ,  $x \geq 0$  be the **service time** density function.
- Define the RV  $V = \{\text{number of arrivals during a service time}\}$ , and the probabilities:  $v_i = P\{V = i\}$ .
- Conditioning on the service duration:

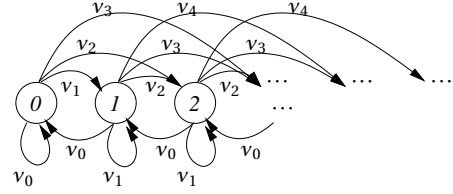
$$v_i = \int_{x=0}^{\infty} P\{i \text{ arrivals in time } x \mid S = x\} f_S(x) dx \Rightarrow \quad (21.2)$$

$$v_i = \int_{x=0}^{\infty} \frac{(\lambda x)^i}{i!} e^{-\lambda x} f_S(x) dx \quad (21.3)$$



- $v_i = P\{\text{number of arrivals during a service time} = i\} \Rightarrow$

$$p_{ij} = \begin{cases} 0, & j < i-1 \quad (N(t) \text{ can only be decreased by 1}) \\ v_j, & i = 0, j \geq 0 \quad (i = 0 \rightarrow \text{the queue was empty}) \\ v_{j-i+1}, & i > 0, j \geq i-1 \quad (i > 0 \rightarrow \text{the queue was busy}) \end{cases} \quad (21.4)$$



$$p_{ij} = \begin{cases} 0, & j < i-1 \\ v_j, & i = 0, j \geq 0 \\ v_{j-i+1}, & i > 0, j \geq i-1 \end{cases} \Rightarrow \mathbf{P} = \begin{bmatrix} v_0 & v_1 & v_2 & v_3 & \dots \\ v_0 & v_1 & v_2 & v_3 & \dots \\ 0 & v_0 & v_1 & v_2 & \dots \\ 0 & 0 & v_0 & v_1 & \dots \\ \dots & \dots & \dots & \dots & \dots \end{bmatrix} \quad (21.5)$$

- Stationary distribution:  $\pi = \pi \mathbf{P}, \pi \mathbf{e} = 1$ .

#### Properties of the stationary distribution ( $\pi = \pi \mathbf{P}, \pi \mathbf{e} = 1$ )

- Using the “**Level Crossing Law**” theorem: a queue with **unitary arrivals and departures** satisfies:

$$\begin{aligned} P\{\text{an arriving customer finds } i \text{ in the system}\} &= \\ P\{\text{a departing customer leaves } i \text{ in the system}\} &\Rightarrow \end{aligned} \quad (21.6)$$

$$\pi_i = P\{\text{an arriving customer find } i \text{ in the system}\} \quad (21.7)$$

- Using **PASTA**:

$$\pi_i = P\{\text{there are } i \text{ customers in the system at an arbitrary time}\} \quad (21.8)$$

So, in an M/G/1 the stationary distribution of the EMC obtained observing the departures, is the stationary distribution of the continuous time process.

*Proof.* Level Crossing Law Theorem

- Define:

- $\mathbf{A}_i(t) = \{\text{number of arrivals finding } i \text{ in the system at } t \geq 0\}$
- $\mathbf{D}_i(t) = \{\text{number of departures leaving } i \text{ in the system at } t \geq 0\}$
- $\mathbf{P}\{\text{a customer finds } i \text{ in the system}\} = \lim_{t \rightarrow \infty} \mathbf{A}_i(t) / A(t)$
- $\mathbf{P}\{\text{a customer leave } i \text{ in the system}\} = \lim_{t \rightarrow \infty} \mathbf{D}_i(t) / D(t)$

- An arriving customer that finds  $i$  in the system produce a transition  $i \rightarrow i+1$ . A customer leaving  $i$  in the system produce a transition  $i+1 \rightarrow i$ .

- Since arrivals and departures are unitary, the number of transitions  $i \rightarrow i+1$  and  $i+1 \rightarrow i$  can only differ in 1:  $|A_i(t) - D_i(t)| \leq 1$ . Note that  $N(t) = A(t) - D(t)$ .

- For a **stable queue**:  $A(t) - D(t) < \infty$

- We have:

- $\mathbf{A}_i(t) = \{\text{number of arrivals finding } i \text{ customer in the system}\}$



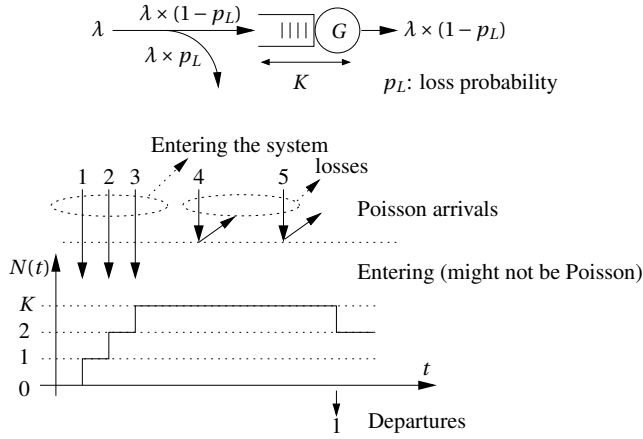
- $D_i(t) = \{\text{number of departures leaving } i \text{ customers in the system}\}$
- $P\{\text{a customer finds } i \text{ in the system}\} = \lim_{t \rightarrow \infty} A_i(t) / A(t)$
- $P\{\text{a customer leave } i \text{ in the system}\} = \lim_{t \rightarrow \infty} D_i(t) / D(t)$
- $A_i(t) - D_i(t) \in \{0, 1\}$ ,  $N(t) = A(t) - D(t) < \infty$ .
- $\lim_{t \rightarrow \infty} A(t) = \infty$ ,  $\lim_{t \rightarrow \infty} D(t) = \infty$ .

• Thus:

$$\begin{aligned} \lim_{t \rightarrow \infty} \left\{ \frac{A_i(t)}{A(t)} - \frac{D_i(t)}{D(t)} \right\} &= \\ \lim_{t \rightarrow \infty} \left\{ \frac{A_i(t)}{A(t)} - \frac{D_i(t)}{A(t)} - \left( \frac{D_i(t)}{D(t)} - \frac{D_i(t)}{A(t)} \right) \right\} &= \\ \lim_{t \rightarrow \infty} \left\{ \frac{A_i(t) - D_i(t)}{A(t)} - \frac{D_i(t)}{D(t)} \frac{A(t) - D(t)}{A(t)} \right\} &= 0 \quad \square \end{aligned}$$

## 21.1 M/G/1/K Queue

### Problem Formulation



### Stationary Distribution

- Using the **general solution of an M/G/1/K** we obtain the stationary distribution of the number in the system left by a **departing** customer:  $\pi_i^d$ .
- By the **Level Crossing Law** this is the stationary distribution of the number in the system found by the **successful arrivals**:

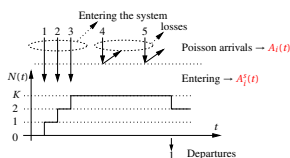
$$\pi_i^s = \pi_i^d, i = 0, 1, \dots, K-1. \quad (21.9)$$

and

$$\pi_i^s = P(\text{a customer entering the system finds } i) \quad (21.10)$$

- **NOTE:** a departing customer cannot leave the system full (nor an arrival can enter the system when it is full).

### Loss Probability



Define:

- $A_i^a(t)$ : Number of **arrivals** (lost or not) finding  $i$  in the system.
- $A_i^s(t)$ : Number of **successful arrivals** finding  $i$  in the system.
- $\pi_i^a, \pi_i^s$  the stationary distribution of the embedded Markov chains  $A_i^a(t), A_i^s(t)$ . By **PASTA**  $\pi_i^a$  is also the stationary distribution of the continuous time process. Thus,

$$\begin{aligned} \pi_i^s &= P(\text{a customer entering the system finds } i), i = 0, 1, \dots, K-1 \Rightarrow \\ \pi_i^s &= \lim_{t \rightarrow \infty} \frac{A_i^s(t)}{\sum_{k=0}^{K-1} A_k^s(t)} \frac{\sum_{k=0}^K A_k^a(t)}{\sum_{k=0}^K A_k^a(t)} = \frac{\pi_i^a}{\sum_{k=0}^{K-1} \pi_k^a} = \frac{\pi_i^a}{1 - \pi_K^a} = \frac{\pi_i^a}{1 - p_L} \Rightarrow \end{aligned}$$

$$\pi_i^a = \pi_i^s (1 - p_L) = \pi_i^d (1 - p_L), i = 0, 1, \dots, K-1 \quad (21.11)$$

- Applying **Little**:  $\rho_s = \mathbf{E}[N_s] = 1 - \pi_0 = \lambda (1 - p_L) \mathbf{E}[S] = \rho (1 - p_L)$ . Where  $\rho = \lambda \mathbf{E}[S]$  and  $\pi_0$  is the proportion of time the server is empty.
- Using **PASTA**:  $\pi_0 = \pi_0^a$  (Poisson arrivals). Using  $\pi_i^a = \pi_i^d (1 - p_L)$ :

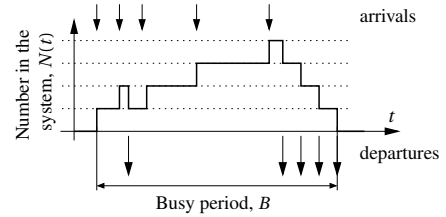
$$\left. \begin{aligned} 1 - \pi_0 &= 1 - \pi_0^a = 1 - \pi_0^d (1 - p_L) \\ 1 - \pi_0 &= \rho (1 - p_L) \end{aligned} \right\} \Rightarrow$$

$$p_L = \frac{\rho + \pi_0^d - 1}{\rho + \pi_0^d}, \rho = \lambda \mathbf{E}[S] \quad (21.12)$$

- Where  $\pi_0^d$  is computed using the general solution of an M/G/1/K.

## 21.2 M/G/1 Busy Period

### Expected Length of a Busy Period



- Define the RV:

- **Busy period**,  $B$ .
- **Idle period**,  $I$ . Poisson arrivals with rate  $\lambda \Rightarrow \mathbf{E}[I] = 1/\lambda$

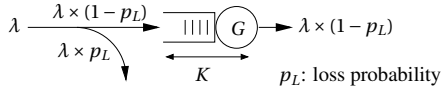
- Clearly:

$$\text{System load } \rho = \lambda \mathbf{E}[S] = \frac{\mathbf{E}[B]}{\mathbf{E}[I] + \mathbf{E}[B]} \Rightarrow \quad (21.13)$$

$$\mathbf{E}[B] = \frac{1}{\lambda} \frac{\rho}{1 - \rho} \quad (21.14)$$



### 21.3 M/G/1/K Busy Period



- **Busy period,  $B$ .**
- **Idle period,  $I$ .** Poisson arrivals with rate  $\lambda \Rightarrow E[I] = 1/\lambda$
- Clearly:

$$\text{System load } \rho_s = \lambda(1-p_L)E[S] = \frac{E[B]}{E[I] + E[B]} \Rightarrow \quad (21.15)$$

$$E[B] = \frac{1}{\lambda} \frac{\rho(1-p_L)}{1-\rho(1-p_L)}, \rho = \lambda E[S]$$

- Or, in terms of  $\pi_0 = \pi_0^d(1-p_L)$ :

$$\text{System load } \rho_s = 1 - \pi_0 = \frac{E[B]}{E[I] + E[B]} \Rightarrow$$

$$E[B] = \frac{1}{\lambda} \frac{1-\pi_0}{\pi_0} \quad (21.16)$$

### 21.4 M/G/1 Delays

#### M/G/1 Mean Time in the Queue

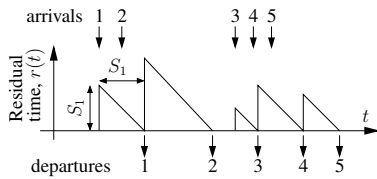
- **Method of the moments:** Using **PASTA**, the **mean time in the queue** ( $W$ ) for an arriving customer, is the mean time to finish the current service (**mean residual time,  $R$** ) plus the **mean time to service the customers in the queue** ( $E[S]N_Q$ ):

$$W = R + E[S]N_Q \quad (21.17)$$

- Using **Little for the queue length**:

$$N_Q = \lambda W \Rightarrow W = R + E[S]\lambda W \Rightarrow W = \frac{R}{1-\rho}, \rho = \lambda E[S]. \quad (21.18)$$

#### M/G/1 Residual Time



- From the figure (note the **right triangles with two equal cathetus**), we have:

$$R = \frac{1}{t} \int_0^t r(\tau) d\tau = \frac{1}{t} \sum_{i=1}^{A(t)} \frac{S_i^2}{2} = \frac{1}{2} \frac{A(t)}{t} \sum_{i=1}^{A(t)} \frac{S_i^2}{A(t)} \xrightarrow{t \rightarrow \infty} \frac{1}{2} \lambda E[S^2] \quad (21.19)$$

- For instance, for an M/M/1

$$E[S^2] = \text{Var}(S) + E[S]^2 = \frac{1}{\mu^2} + \left(\frac{1}{\mu}\right)^2 = \frac{2}{\mu^2}, \quad (21.20)$$

thus, the residual time is:

$$R = \frac{1}{2} \lambda E[S^2] = \frac{\lambda}{\mu^2} = \frac{\rho}{\mu}, \rho = \frac{\lambda}{\mu}. \quad (21.21)$$

- We can check that  $E[R|S \text{ idle}] = 0$  and  $E[R|S \text{ busy}] = 1/\mu$ , thus

$$R = E[R|S \text{ idle}] \pi_0 + E[R|S \text{ busy}] (1-\pi_0) = \frac{\rho}{\mu}, \rho = 1 - \pi_0, \quad (21.22)$$

as expected.

#### Pollaczek-Khinchin, P-K formula

- We have:

$$W = \frac{R}{1-\rho}, \rho = \lambda E[S] \quad (21.23)$$

$$R = \frac{1}{2} \lambda E[S^2]$$

- Substituting we get the **Pollaczek-Khinchin, P-K formula**:

$$W = \frac{\lambda E[S^2]}{2(1-\rho)}, \rho = \lambda E[S] \quad (21.24)$$

- **Mean time in the system** (response time):

$$T = E[S] + W = E[S] + \frac{\lambda E[S^2]}{2(1-\rho)} \quad (21.25)$$

- For an **M/M/1** queue:  $E[S^2] = \frac{2}{\mu^2} \Rightarrow W = \frac{\rho}{\mu(1-\rho)}$

- For an **M/D/1** queue:  $E[S^2] = \frac{1}{\mu^2} \Rightarrow W = \frac{\rho}{2\mu(1-\rho)}$

- **Observation:** The M/D/1 has the minimum value of  $E[S^2] \Rightarrow$  is a lower bound of  $W$ ,  $T$ ,  $N_Q$  and  $N$  for an M/G/1.

#### P-K Formula Does Not Apply to an M/G/1/K Queue

- **P-K formula is not applicable** to an **M/G/1/K** queue because the **customers entering the system** might not be Poisson. Thus, they **does not observe the mean residual time**.

- **Example:** Customers entering an **M/G/1/1** queue (0 queue size) observe the system always empty. Thus, in an M/G/1/1 queue the expected time in the queue is  $W = 0$  (P-K formula does not apply), and the expected time in the system is  $T = E[S]$  (mean service time).

- **With an M/G/1/K** we can compute  $N = \sum_{n=1}^K n \pi_n^a$ , and use Little:  $N = \lambda(1-p_L)T$ . For instance, for an M/G/1/1 we have  $\pi_0^d = 1$ , and  $N = 0\pi_0^a + 1\pi_1^a = \pi_1^a = p_L$ . Thus,  $p_L = \frac{\rho + \pi_0^d - 1}{\rho + \pi_0^d} = \frac{\rho}{\rho + 1}$ , and  $T = \frac{N}{\lambda(1-p_L)} = \frac{p_L}{\lambda(1-p_L)} = \frac{\rho}{\lambda} = E[S]$ , as expected.

## Chapter 22

### Queues in Tandem

#### 22.1 Burke theorem

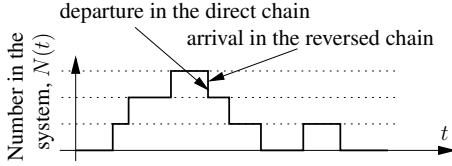
- The **departure process in an M/M/m** queue,  $1 \leq m \leq \infty$ , is a **Poisson** process with the same parameter than the arrival process.
- At each time  $t$ , the **number of customers in the system** is independent of the sequence of departures previous to  $t$ .

- Relation between the arrival and departure process:

The **departure process** in a reversible queue has the same joint distribution than the **arrival process**.

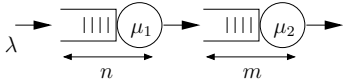
*Proof.*

- If the queue is reversible:  $q_{ij} = q_{ij}^r \Rightarrow$  the arrival process in the reversed chain has the same distribution than the arrival process in the direct chain,
- but: □



- The queue **M/M/m is reversible**  $\Rightarrow$  The **departures are Poisson** with the same parameter than the arrivals.
- The arrivals in the reversed chain previous to  $t$  are Markovian, thus, independent of the number of customers in the system after  $t$ . This implies that the **departures** in the direct chain are **independent of the number in the system before  $t$** .

## 22.2 Tandem M/M/m Queues



- Define the chain:

$$\mathbf{X}(\mathbf{n}, \mathbf{m}) = \{n \text{ in the system 1, } m \text{ in the system 2}\} \quad (22.1)$$

- The **stationary distribution** is the **product of the stationary distributions of the isolated queues**:

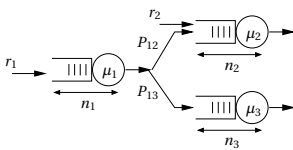
$$\pi_{nm} = (1 - \rho_1) \rho_1^n (1 - \rho_2) \rho_2^m, \rho_1 = \lambda / \mu_1, \rho_2 = \lambda / \mu_2 \quad (22.2)$$

- *Proof.* Using Burke, the departures of system 1 are Poisson and the number in the system 1 is independent of the previous departures (arrivals to system 2), thus, independent from the number of customers in system 2. □

## Chapter 23

### Networks of Queues

#### 23.1 Feed Forward Queues



- Suppose **M/M/1 queues** with outside arrivals with rate  $r_i$  randomly forwarded with probabilities  $P_{ij}$  (see figure).
- The network has the following **product form solution**:

$$\pi(n_1, n_2, \dots, n_K) = (1 - \rho_1) \rho_1^{n_1} (1 - \rho_2) \rho_2^{n_2} \dots (1 - \rho_K) \rho_K^{n_K},$$

$$\rho_i = \lambda_i / \mu_i. \quad (23.1)$$

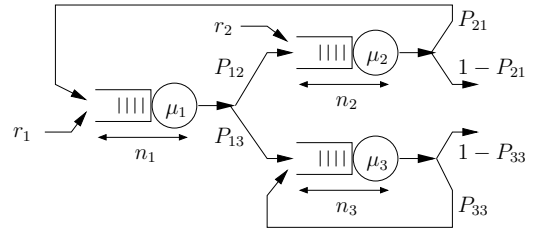
- The rates  $\lambda_i$  are computed solving:  $\lambda_i = r_i + \sum_j \lambda_j P_{ji}$ .
- Stability condition:  $\rho_i < 1$ .

*Proof.* (draft)

- Burke theorem.
- **Superposition of Poisson** processes with rates  $\lambda_i$  is Poisson with rate  $\sum_i \lambda_i$ .
- A **Poisson** process with rate  $\lambda$  **randomly split** with probabilities  $p_i$ ,  $\sum_i p_i = 1$ , produce Poisson processes with rates  $p_i \lambda$ . □

#### 23.2 Jackson Theorem

- Suppose **M/M/m queues**. In queue  $i$  the customers **arrive** from outside with rate  $r_i$  and **depart** to queue  $j$  with probability  $P_{ij}$ , or leave the system with probability  $1 - \sum_j P_{ij}$ :



- The network has the following **product form solution**:

$$\pi(n_1, n_2, \dots, n_K) = \pi_{n_1} \pi_{n_2} \dots \pi_{n_K} \quad (23.2)$$

where  $\pi_{n_i}$  is the solution of the queue  $i$  with arrival rates  $\lambda_i$  obtained solving:

$$\lambda_i = r_i + \sum_j \lambda_j P_{ji} \quad (23.3)$$

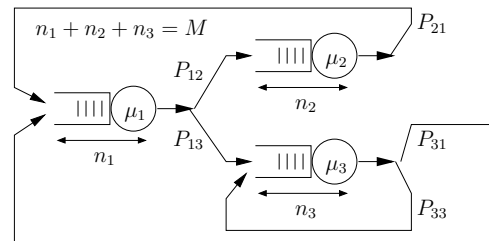
- Stability condition:  $\rho_i = \lambda_i / \mu_i < 1$ .
- For example, for M/M/1 queues:

$$\pi(n_1, n_2, \dots, n_K) = (1 - \rho_1) \rho_1^{n_1} (1 - \rho_2) \rho_2^{n_2} \dots (1 - \rho_K) \rho_K^{n_K} \quad (23.4)$$

- **Proof:** The solution satisfies the global balance equations.
- **NOTE:** The proof is different from feed forward queues, since routing loops make arrivals not necessarily Poisson.

#### 23.3 Closed Networks of Queues

- **M/M/m networks** without arrivals and departures to outside of the system:



### Jackson Theorem for Closed Networks of Queues

- The network has the following **product form solution**:

$$\pi(n_1, n_2, \dots, n_K) = \frac{1}{G} \rho_1^{n_1} \rho_2^{n_2} \dots \rho_K^{n_K}, \rho_i = \lambda_i / \mu_i. \quad (23.5)$$

- Where the rates  $\lambda_i$  are any solution to the equations:

$$\lambda_i = \sum_j \lambda_j P_{ji} \quad (\text{in matrix form: } \boldsymbol{\lambda} = \boldsymbol{\lambda} \mathbf{P}) \quad (23.6)$$

- And the normalization factor is given by:

$$G = \sum_{n_1 + n_2 + \dots + n_K = M} \rho_1^{n_1} \rho_2^{n_2} \dots \rho_K^{n_K} \quad (23.7)$$

- Proof:** The solution satisfies the global balance equations.
- NOTE:** the equation  $n_1 + n_2 + \dots + n_K = M$  has  $\binom{M+K-1}{M} = \binom{M+K-1}{K-1}$  solutions (ways to allocate  $M$  items in  $K$  boxes).

## Chapter 24

### Matrix Geometric Method

#### 24.1 Squared coefficient of variation

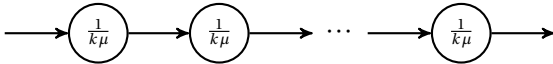
- Idea:** almost all distributions can be approximated by a mixture of exponentials.
- Squared coefficient of variation**, characterization of a distribution variability (for distributions with  $E[X] > 0$ ):

$$C_X^2 = \frac{\text{Var}(X)}{E[X]^2} = \frac{E[X^2] - E[X]^2}{E[X]^2} = \frac{E[X^2]}{E[X]^2} - 1 \quad (24.1)$$

- Deterministic** distribution:  $C_D^2 = 0$ .
- Exponential** distribution:  $E[X] = 1/\mu$ ,  $\text{Var}(X) = 1/\mu^2$ . Thus  $C_{\text{exp}}^2 = 1$ .
- What if we want a distribution more *deterministic* than an exponential,  $C_X^2 < 1$ ? or with larger variability,  $C_X^2 > 1$ ?

#### 24.2 $C_X^2 < 1$ : Erlang-k

- $k$  stages exponentially distributed with parameter  $k\mu$ :



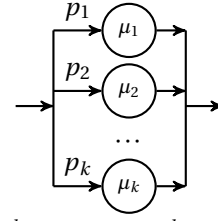
$$f_E(t) = \frac{(k\mu)^k t^{k-1} e^{-k\mu t}}{(k-1)!}, t \geq 0, k \geq 1 \quad (24.2)$$

$$E[X] = k \frac{1}{k\mu} = \frac{1}{\mu}$$

$$\text{Var}(X) = k \times \text{Var}(\text{exp}(k\mu)) = k \frac{1}{(k\mu)^2} = \frac{1}{k\mu^2} \quad (24.3)$$

$$C_X^2 = \frac{\text{Var}[X]}{E[X]^2} = \frac{1}{k} < 1$$

#### 24.3 $C_X^2 > 1$ : Hyper-exponential



$$f_H(t) = \sum_{i=1}^k p_i \mu_i e^{-\mu_i t}, \sum_{i=1}^k p_i = 1, t \geq 0$$

$$E[X] = \sum_{i=1}^k p_i \frac{1}{\mu_i}, E[X^2] = \sum_{i=1}^k p_i \frac{2}{\mu_i^2}$$

$$\text{Var}(X) = E[X^2] - E[X]^2 =$$

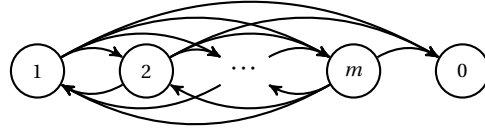
$$\sum_{i=1}^k p_i \frac{2}{\mu_i^2} - \left( \sum_{i=1}^k p_i \frac{1}{\mu_i} \right)^2 =$$

$$\left( \sum_{i=1}^k p_i \frac{1}{\mu_i} \right)^2 + \sum_{i=1}^k \sum_{j \neq i} p_i p_j \left( \frac{1}{\mu_i} - \frac{1}{\mu_j} \right)^2$$

$$C_X^2 = \frac{\text{Var}[X]}{E[X]^2} > 1$$

#### 24.4 Phase type distribution

- General mixture of exponentials.
- The service ends upon reaching the absorbing state.
- Can approximate arbitrary distributions.
- Representation:  $PH(\mathbf{a}, \mathbf{T})$ .



$$\mathbf{Q}^{m+1 \times m+1} = \begin{bmatrix} \mathbf{T}^{m \times m} & \mathbf{c}^{m \times 1} \\ \mathbf{0}^{1 \times m} & 0 \end{bmatrix} \quad (24.4)$$

$$\text{Initial prob. } \begin{bmatrix} \mathbf{a}^{1 \times m} & a_0 \end{bmatrix}. \quad (24.5)$$

$$\mathbf{f}_{PH}(t) = \mathbf{a} e^{\mathbf{T}t} \mathbf{c}, t \geq 0 \quad (24.6)$$

$$E[\mathbf{X}^k] = k! \mathbf{a} (-\mathbf{T}^{-1})^k \mathbf{e}$$

where  $\mathbf{e}$  is a column vector of 1s.

#### 24.5 Quasi Birth Death Processes

- Assume a two dimensional MC with **states**  $(n, i)$  (e.g. an **M/PH/1** queue). We call  $n$  the **level** and  $i$  the **phase**. We group the states of the **stationary distribution**:

$$\boldsymbol{\pi} = [\pi_0 \quad \pi_1 \quad \pi_2 \quad \dots]$$

$$\begin{cases} \pi_0 = [(0,0) \quad (0,1) \quad \dots (0,m')] & \text{initial part (level 0)} \\ \pi_i = [(i,1) \quad \dots (i,m)], i \geq 1 & \text{repetitive part (level } i \geq 1) \end{cases} \quad (24.7)$$

$$Q = \begin{bmatrix} L_0 & F_0 & & & \\ B_0 & L & F & & \\ & B & L & F & \\ & & B & L & \dots \\ & & & \dots & \dots \end{bmatrix} \quad (24.8)$$

$$\begin{cases} B & \text{governs the transitions to previous level} \\ L & \text{governs the change of phase inside a level} \\ F & \text{governs the transitions to next level} \end{cases} \quad (24.9)$$

## 24.6 Matrix Geometric Solution

- Due to similarity with an M/M/1 ( $\pi_i = \pi_0 \rho^i$ ) we guess for the repetitive part:

$$\pi_{i+1} = \pi_1 R^i, i \geq 0 \quad (24.10)$$

which gives:

$$\begin{aligned} \pi_1 F + \pi_2 L + \pi_3 B &= 0 \Rightarrow \\ \pi_1 F + \pi_1 R L + \pi_1 R^2 B &= 0 \Rightarrow \\ F + R L + R^2 B &= 0 \end{aligned} \quad (24.11)$$

- Isolating  $R$  we have that it can be found iterating

$$R_{n+1} = -(F + R_n^2 B) L^{-1}, \quad (24.12)$$

starting e.g. with  $R_0 = I$ .

- Better iterative algorithms can be found in [6].

```
##
## Basic iterative algorithm to compute the
## matrix R
##
## B, L, F: repetitive part matrices
##
invL <- -solve(L) # -1/L
C1 <- F %*% invL # -F/L
C2 <- B %*% invL # -B/L
R <- diag(nrow(B))
epsilon <- 1e-15
MaxIter <- 500
IterB <- 1
while (IterB < MaxIter) {
  prev <- R
  R <- C1 + R %*% C2 # -(F + R^2 B)/L
  if (max(abs(prev - R)) < epsilon) { break }
  IterB = IterB + 1
}
```

- Solving  $\pi_0$  and  $\pi_1$ :

$$Q = \begin{bmatrix} L_0 & F_0 & & & \\ B_0 & L & F & & \\ & B & L & F & \\ & & B & L & \dots \\ & & & \dots & \dots \end{bmatrix} \quad (24.13)$$

$$\begin{bmatrix} \pi_0 & \pi_1 & \pi_2 & \pi_3 & \dots \end{bmatrix} Q = 0$$

$$\pi_{i+1} = \pi_1 R^i, i \geq 0.$$

- Thus:

$$\begin{aligned} \pi_0 L_0 + \pi_1 B_0 &= 0 \\ \pi_0 F_0 + \pi_1 L + \pi_1 R B &= 0 \Rightarrow \\ \begin{bmatrix} \pi_0 & \pi_1 \end{bmatrix} \begin{bmatrix} L_0 & F_0 \\ B_0 & L + R B \end{bmatrix} &= 0 \end{aligned} \quad (24.14)$$

- and the normalization condition:

$$\pi_0 e_0 + \sum_{i=0}^{\infty} \pi_1 R^i e_1 = 1 \Rightarrow \pi_0 e_0 + \pi_1 (I - R)^{-1} e_1 = 1 \Rightarrow$$

$$\begin{bmatrix} \pi_0 & \pi_1 \end{bmatrix} \begin{bmatrix} e_0 \\ (I - R)^{-1} e_1 \end{bmatrix} = 1 \quad (24.15)$$

where  $e_i$  are column vectors of 1s of appropriate size.

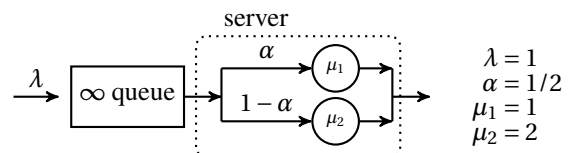
```
##
## Basic algorithm to compute the initial
## probabilities
##
## B0, L0, F0: initial part matrices
## B, L, F: repetitive part matrices
##
IMRinv <- solve(diag(nrow(B)) - R) # 1/(I-R)
M0 <- rbind(cbind(L0, F0), cbind(B0, L + R %*% B))

## Normalization column
NE <- c(rep(1, nrow(L0)), IMRinv %*% rep(1, nrow(
  IMRinv)))

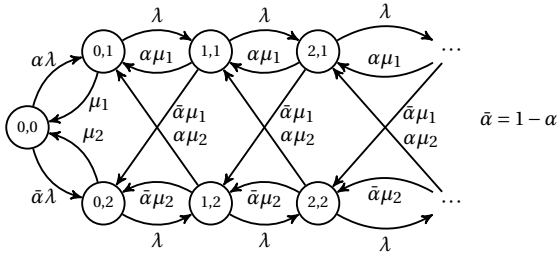
## solve using the replace 1 equation method
M0 <- cbind(NE, M0[, 2:ncol(M0)]) # replace first
  column of M0 by NE
stat <- solve(t(M0), c(1, rep(0, nrow(M0)-1)))
```

### 24.6.1 Example

- Consider an M/G/1 queue where service time is hyper-exponentially distributed:



- Derive the rate matrix,  $Q$ , ordering the states lexicographically. Identify the states that form the initial and repetitive part. Identify the submatrices that would be used for a matrix geometric solution:  $B_0, L_0, F_0, B, L, F$ .
- Solve the Chain using the matrix geometric method. Compute the number in the system. Check it with the PK formula.

**Solution**

$$\bar{\alpha} = 1 - \alpha$$

$$\mathbf{P} = \begin{array}{c|cccccc|c} & \begin{matrix} 0,0 & 0,1 & 0,2 & 1,1 & 1,2 & 2,1 & 2,2 & \dots \end{matrix} \\ \hline \begin{matrix} -\lambda \\ \mu_1 \\ \mu_2 \end{matrix} & \begin{matrix} -\lambda \\ \alpha\lambda \\ 0 \end{matrix} & \begin{matrix} \alpha\lambda \\ 0 \\ -(\lambda + \mu_2) \end{matrix} & \begin{matrix} 0 \\ \lambda \\ 0 \end{matrix} & \begin{matrix} 0 \\ 0 \\ \lambda \end{matrix} & \begin{matrix} 0 \\ 0 \\ 0 \end{matrix} & \begin{matrix} 0 \\ 0 \\ 0 \end{matrix} & \begin{matrix} \dots \\ \dots \\ \dots \end{matrix} \\ \hline \begin{matrix} 0 \\ 0 \end{matrix} & \begin{matrix} \alpha\mu_1 \\ \alpha\mu_2 \end{matrix} & \begin{matrix} \bar{\alpha}\mu_1 \\ \bar{\alpha}\mu_2 \end{matrix} & \begin{matrix} -(\lambda + \mu_1) \\ 0 \end{matrix} & \begin{matrix} 0 \\ -(\lambda + \mu_2) \end{matrix} & \begin{matrix} \lambda \\ 0 \end{matrix} & \begin{matrix} 0 \\ \lambda \end{matrix} & \begin{matrix} \dots \\ \dots \end{matrix} \\ \hline \begin{matrix} 0 \\ 0 \end{matrix} & \begin{matrix} 0 \\ 0 \end{matrix} & \begin{matrix} 0 \\ 0 \end{matrix} & \begin{matrix} \alpha\mu_1 \\ \alpha\mu_2 \end{matrix} & \begin{matrix} \bar{\alpha}\mu_1 \\ \bar{\alpha}\mu_2 \end{matrix} & \begin{matrix} -(\lambda + \mu_1) \\ 0 \end{matrix} & \begin{matrix} 0 \\ -(\lambda + \mu_2) \end{matrix} & \begin{matrix} \dots \\ \dots \end{matrix} \\ \hline \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \ddots \end{array}$$

$$\Rightarrow \mathbf{B}_0 = \begin{bmatrix} \mu_1 \\ \mu_2 \end{bmatrix}, \mathbf{L}_0 = \begin{bmatrix} -\lambda \end{bmatrix}, \mathbf{F}_0 = \begin{bmatrix} \alpha\lambda & \bar{\alpha}\lambda \end{bmatrix}, \mathbf{B} = \begin{bmatrix} \alpha\mu_1 & \bar{\alpha}\mu_1 \\ \alpha\mu_2 & \bar{\alpha}\mu_2 \end{bmatrix},$$

$$\mathbf{L} = \begin{bmatrix} -(\lambda + \mu_1) & 0 \\ 0 & -(\lambda + \mu_2) \end{bmatrix}, \mathbf{F} = \begin{bmatrix} \lambda & 0 \\ 0 & \lambda \end{bmatrix}.$$

- Iterating  $\mathbf{R}_{n+1} = -(\mathbf{F} + \mathbf{R}_n^2 \mathbf{B}) \mathbf{L}^{-1}$  we get:

$$\mathbf{R} = \begin{bmatrix} 5/7 & 1/7 \\ 1/7 & 3/7 \end{bmatrix}$$

- Using  $\begin{bmatrix} \pi_0 & \pi_1 \end{bmatrix} \begin{bmatrix} \mathbf{L}_0 & \mathbf{F}_0 \\ \mathbf{B}_0 & \mathbf{L} + \mathbf{R}\mathbf{B} \end{bmatrix} = \mathbf{0}$ ,  $\begin{bmatrix} \pi_0 & \pi_1 \end{bmatrix} \begin{bmatrix} \mathbf{e}_0 \\ (\mathbf{I} - \mathbf{R})^{-1} \mathbf{e}_1 \end{bmatrix} = 1$  we get:

$$\begin{bmatrix} \pi_0 & \pi_1 \end{bmatrix} = \begin{bmatrix} 1/4 & 3/28 & 1/14 \end{bmatrix}$$

- Number in the system:

$$N = \sum_{n=1}^{\infty} n \pi_1 \mathbf{R}^n \mathbf{e}_1 = \pi_1 (\mathbf{I} - \mathbf{R})^{-2} \mathbf{e}_1 = \frac{13}{4}$$

- Using the PK Formula:

$$E[S] = \frac{\alpha}{\mu_1} + \frac{1-\alpha}{\mu_2} = \frac{1}{4}, \rho = \lambda E[S] = \frac{3}{4}, E[S^2] = \frac{2\alpha}{\mu_1^2} + \frac{2(1-\alpha)}{\mu_2^2} = \frac{5}{4}$$

thus,

$$T = E[S] + \frac{\lambda E[S^2]}{2(1-\rho)} = \frac{13}{4}, N = \lambda T = \frac{13}{4}, \text{ as expected.}$$

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**Table with some distributions**

Distribution	Parametres	Density	Mean	Variance	Characteristic Function
Bernoulli	$0 \leq p \leq 1$ $q = 1 - p$	$p^k (1 - p)^{1-k}$ $k = 0, 1$	$p$	$p(1 - p)$	$q + p e^{it}$
Binomial	$0 \leq p \leq 1$ $q = 1 - p$	$\binom{n}{k} p^k (1 - p)^{n-k}$ $k = 0, 1, \dots, n$	$np$	$np(1 - p)$	$(q + p e^{it})^n$
Geometric	$0 \leq p \leq 1$ $q = 1 - p$	$p(1 - p)^k$ $k \geq 0$	$\frac{1 - p}{p}$	$\frac{1 - p}{p^2}$	$\frac{p}{1 - q e^{it}}$
Negative binomial	$r > 0$ $0 \leq p \leq 1$ $q = 1 - p$	$\binom{k+r-1}{k} p^r q^k$ $k \geq 0$	$r \frac{1 - p}{p}$	$r \frac{1 - p}{p^2}$	$\left( \frac{p}{1 - q e^{it}} \right)^r$
Poisson	$\lambda > 0$	$\frac{\lambda^k}{k!} e^{-\lambda}, \quad k \geq 0$	$\lambda$	$\lambda$	$\exp \{ \lambda (e^{it} - 1) \}$
Normal $N(\mu, \sigma)$	$\mu \in \mathbb{R}$ $\sigma > 0$	$\frac{1}{\sqrt{2\pi}\sigma^2} \exp \left\{ \frac{-(x - \mu)^2}{2\sigma^2} \right\}$ $x \in \mathbb{R}$	$\mu$	$\sigma^2$	$\exp \left\{ \mu i t - \frac{t^2 \sigma^2}{2} \right\}$
Uniform	$a < b$	$\frac{1}{b - a}, \quad a \leq x \leq b$	$\frac{a + b}{2}$	$\frac{(b - a)^2}{12}$	$\frac{e^{itb} - e^{ita}}{it(b - a)}$
Exponential	$\alpha$	$\alpha e^{-\alpha x}, \quad x \geq 0$	$\frac{1}{\alpha}$	$\frac{1}{\alpha^2}$	$\left( 1 - \frac{it}{\alpha} \right)^{-1}$
Gamma $\gamma(n, \alpha)$	$\alpha > 0,$ $n > 0$	$\frac{\alpha^n x^{n-1} e^{-\alpha x}}{\Gamma(n)}, \quad x \geq 0$	$\frac{n}{\alpha}$	$\frac{n}{\alpha^2}$	$\left( 1 - \frac{it}{\alpha} \right)^{-n}$
Beta $\beta(p, q)$	$p > 0,$ $q > 0$	$\frac{x^{p-1} (1 - x)^{q-1}}{B(p, q)},$ $0 \leq x \leq 1$	$\frac{p}{p + q}$	$\frac{pq}{(p + q)^2(p + q + 1)}$	

$$\Gamma(x) = \int_0^\infty e^{-t} t^{x-1} dt, \Gamma(n) = (n-1)! \quad B(p, q) = \int_0^1 t^{p-1} (1-t)^{q-1} dt, B(p, q) = \frac{\Gamma(p)\Gamma(q)}{\Gamma(p+q)}$$