

An overview of Self-Justifying Axiom Systems, and related questions

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What and Why

Goedel's Second Incompleteness Theorem (G2):

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What and Why

Self-Justifying Axiom System (SJAS):

- i. A theory S includes a statement of the consistency of S .
- ii. S is consistent relative to a known theory T .

Parametrizing Expressivity

Generalized Arithmetic Configuration:

- i. Axiom Basis (language, consistency statement, et al.)
- ii. Deductive Apparatus
 - Logical Axioms
 - Rules of Inference

Parametrizing Consistency

“An axiom system α owns a Level-1 appreciation of its own self- consistency (under a deductive apparatus D) iff it can verify that D produces no two simultaneous proofs for a Π_1 sentence and its negation.”

Axiom Basis

First Order Language without Induction:

Integer Subtraction(x, y)

Integer Division(x, y)

Maximum(x, y)

Length(x) = $|x|$

Root(x, y) = $\lceil x^{1/y} \rceil$

Count(x, j) := the number of “1”s in x ’s rightmost j bits

Mult(x, y, z):=

$[(x = 0 \vee y = 0) \Rightarrow z = 0] \wedge [(x \neq 0 \wedge y \neq 0) \Rightarrow (z/x = y \wedge ((z - 1)/x) < y)]$

Construct $\Delta_k, \Pi_k, \Sigma_k$ sentences as per normal.

Axiom Basis

Group 0: Constants 0, 1, Addition(x, y), Double(x)

Group 1: a finite set of Π_1 sentences, proving any Δ_0 sentence that holds true under the standard model

Group 2: For each Π_1 sentence Φ , $\forall p \{ \text{HilbPrf} (<\Phi>, p) \Rightarrow \Phi \}$

Group 3: Kleene Fixed Point encoding of consistency:

$$\forall x \forall y \forall p \forall q \neg [\text{Pair}(x, y) \wedge \text{Prf}(x, p) \wedge \text{Prf}(y, q)]$$

Deductive Apparatus

HilbPrf – Hilbert style proofs, using axiom schemae and modus ponens

Tab – standard semantic tableaux

X-Tab – semantic tableaux with an axiom scheme for LEM

Tab-1 – allows application of modus ponens for previously derived sentences.

Consistency Preservation

“An operation $I(\bullet)$ that maps an initial axiom system A onto an alternate system $I(A)$ will be called Consistency Preserving iff $I(A)$ is consistent whenever all of A 's axioms hold true under the standard model of the natural numbers.

Suppose the symbol D denotes either semantic tableaux deduction or its Tab-1 generalization. Then the $I_D(\bullet)$ mapping operation is consistency preserving (e.g. $I_D(A)$ will be consistent whenever all of A 's axioms hold true under the standard model of the natural numbers).”

Trade-Offs

(1) $\forall x \exists z \text{ Add}(x, 1, z)$	NS: None
(2) $\forall x \forall y \exists z \text{ Add}(x, y, z)$	S: (1)
(3) $\forall x \forall y \exists z \text{ Mult}(x, y, z)$	A: (1), (2)
	M: (1), (2), (3)

$$\text{Tab-1} + \text{A} = \text{Hilb} + \theta = \text{SJAS}$$

The Local Cluster

Artemov – proves each of an infinite tower of consistency statements for subsets of PA

Beklemishev – unpublished simplification of Willard's SJAS

Ganea – SJAS does not define an intermediate r.e. degree

Niebergall – non-arithmetically definable, consistency proving theories

Pakhomov – “A weak set theory that proves its own consistency”

Further Questions

Is it correct?

Other parameter values

- Deduction systems: cirquent calculus, resolution

Executable realization

- Relational logic programming

Autarkic Formal Systems

To what extent are the aspects of a formal system determined by its own powers?

- Consistency
- Definability
- Decidability
- Interpretation
 - Brown and Palsberg, self-interpretation in System F
- Replication
 - Von Neumann automata

Autarkic Formal Systems

To what extent are the aspects of a formal system determined by its own powers?

- Consistency
 - Definability
 - Decidability
- \leq Lawvere's Fixed Point Theorem
(per Yanofsky)

Further motivation for intensional exploration.