Computability via Recursive Functions

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1 Effective Calcubility and Computability

2 Primitive Recursive Functions

2.1 Functions

For this paper, $\mathbb{N}$ refers to the set $\{0, 1, 2, 3, ...\}$

3 The Ackermann Function

Definition (The Ackermann Function)**.** *Let $n, m \in \mathbb{N}$. Then define $A(n, m)$ as follows:*

$$A(m,n) = \begin{cases} n+1 & m = 0 \\ A(m-1, 1) & m > 0 \land n = 0 \\ A(m-1, A(m, n-1)) & m > 0 \land n > 0 \end{cases}$$

Lemma 1. *For any $m, n \in \mathbb{N}, A(m, n) \in \mathbb{N}$*

Proof. Proof Here □

Theorem 1. *A(m, n) is a total function*

Proof. We will proceed inductively to show that $A(m, n)$ is defined for all $m, n \in \mathbb{N}$.

Clearly $A(0, n)$ is defined for all $n \in \mathbb{N}$. Assume $A(k, n)$ is defined for some

$k \in \mathbb{N}$ and every $n \in \mathbb{N}$. Since $k+1>0$, $A(k+1,0) = A(k,1)$, which is defined.

Now we assume $A(k+1,j)$ is defined for some $j \in \mathbb{N}$. By Lemma 1, $A(k+1,j) = a$ for some $a \in \mathbb{N}$. Then since $j+1>0$, $A(k+1,j+1) = A(k, A(k+1,j)) = A(k,a)$. Since $A(k,n)$ is defined for every $n \in \mathbb{N}$ by our inductive hypothesis, $A(k,a) = A(k+1,j+1)$ is defined. □

Definition. *Let P be the set of all primitive recursive functions so that if $f(x_1, x_2, ..., x_n) \in P$ and $m = max\{x_1, x_2, ..., x_n\}$, then there exists $t \in \mathbb{N}$ so that $f(x_1, x_2, ..., x_n) < A(t,m)$*

Theorem 2. *$c(x)$, $s(x)$, $p_i(x) \in P$*

Proof. □

Theorem 3. *P is closed under composition*

Proof. □

Theorem 4. *P is closed under primitive recursion*

Proof. □

Theorem 5. *P is precisely the primitive recursive functions*

Proof. □

Theorem 6. *$A(m,n)$ is not a primitive recursive function*

Proof. Proof Here □

4 General Recursive Functions

4.1 Partial Functions

4.2 Definition of General Recursive Functions