

```

1  ┌────────────────────────── MODULE op_counter ───────────────────┐
2  EXTENDS Integers, Sequences, TLC
3  CONSTANTS N
4
5  Procs  $\triangleq$  1 .. N
6
7
8  --algorithm op_counter
9  variables
10   ops = [m ∈ Procs ↦ ⟨⟩];  utilized to broadcast the operations
11
12  macro Broadcast(o) begin
13   ops := [m ∈ Procs ↦ Append(ops[m], o)] ;  send the operation to all
14  end macro ;
15
16  macro Update(v) begin  receive the operation
17   if Len(ops[self]) > 0 then
18    if Head(ops[self]) = "I" then
19     v := v + 1 ;
20    elseif Head(ops[self]) = "D" then
21     v := v + 1 ;
22    end if ;
23    ops[self] := Tail(ops[self]) ;  clean received msg
24   end if ;
25  end macro ;
26
27  process Counter ∈ Procs
28  variables
29   count = 0 ;  local counter
30  begin Main:
31   while TRUE do
32    Update(count) ;
33    either Increment:
34     Broadcast("I") ;
35    or Decrement:
36     Broadcast("D") ;
37   end either ;
38  end while ;
39  end process ;
40
41  end algorithm ;
42
43  BEGIN TRANSLATION
44  VARIABLES ops, pc, count
45
46  vars  $\triangleq$  ⟨ops, pc, count⟩
47
48  ProcSet  $\triangleq$  (Procs)
49
50  Init  $\triangleq$  Global variables

```

```

51       $\wedge ops = [m \in Procs \mapsto \langle \rangle]$ 
52      Process Counter
53       $\wedge count = [self \in Procs \mapsto 0]$ 
54       $\wedge pc = [self \in ProcSet \mapsto \text{"Main"}]$ 
56   $Main(self) \triangleq \wedge pc[self] = \text{"Main"}$ 
57       $\wedge \text{IF } Len(ops[self]) > 0$ 
58           $\text{THEN } \wedge \text{IF } Head(ops[self]) = \text{"I"}$ 
59               $\text{THEN } \wedge count' = [count \text{ EXCEPT } ![self] = count[self] + 1]$ 
60               $\text{ELSE } \wedge \text{IF } Head(ops[self]) = \text{"D"}$ 
61                   $\text{THEN } \wedge count' = [count \text{ EXCEPT } ![self] = count[self] + 1]$ 
62                   $\text{ELSE } \wedge \text{TRUE}$ 
63                       $\wedge count' = count$ 
64                       $\wedge ops' = [ops \text{ EXCEPT } ![self] = Tail(ops[self])]$ 
65                   $\text{ELSE } \wedge \text{TRUE}$ 
66                       $\wedge \text{UNCHANGED } \langle ops, count \rangle$ 
67           $\wedge \vee \wedge pc' = [pc \text{ EXCEPT } ![self] = \text{"Increment"}]$ 
68           $\vee \wedge pc' = [pc \text{ EXCEPT } ![self] = \text{"Decrement"}]$ 
70   $Increment(self) \triangleq \wedge pc[self] = \text{"Increment"}$ 
71       $\wedge ops' = [m \in Procs \mapsto Append(ops[m], \text{"I"})]$ 
72       $\wedge pc' = [pc \text{ EXCEPT } ![self] = \text{"Main"}]$ 
73       $\wedge count' = count$ 
75   $Decrement(self) \triangleq \wedge pc[self] = \text{"Decrement"}$ 
76       $\wedge ops' = [m \in Procs \mapsto Append(ops[m], \text{"D"})]$ 
77       $\wedge pc' = [pc \text{ EXCEPT } ![self] = \text{"Main"}]$ 
78       $\wedge count' = count$ 
80   $Counter(self) \triangleq Main(self) \vee Increment(self) \vee Decrement(self)$ 
82   $Next \triangleq (\exists self \in Procs : Counter(self))$ 
84   $Spec \triangleq Init \wedge \Box [Next]_{vars}$ 
86  END TRANSLATION
88  Eventual Convergence:
89  Safety:  $i, j : C(xi) = C(xj)$  implies that the abstract states of  $i$  and  $j$  are equivalent.
90   $Safety \triangleq (\forall i, j \in Procs : count[i] = count[j])$ 
91  Liveness:  $i, j : \text{f } C(xi)$  implies that, eventually,  $\text{f } C(xj)$ .
92   $Liveness \triangleq \Diamond (\forall i, j \in Procs : count[i] = count[j])$ 
94  |
    \ * Modification History
    \ * Last modified Wed Dec 12 15:22:32 PST 2018 by ocosta
    \ * Created Sat Dec 01 16:58:15 PST 2018 by ocosta

```