# On Frame Synchronization in Aeronautical Telemetry

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Frame synchronizers, suitable for use with continuous phase modulation (CPM) in the presence of phase and frequency offsets, are derived for periodically inserted preamble bits in continuous-mode transmission. The resulting frame synchronizers involve single or double correlations of the received samples with a local copy of the CPM samples corresponding to the preamble bits. Of particular interest in aeronautical telemetry is shaped offset QPSK, telemetry group version (SOQPSK-TG), with the preamble defined in the integrated network enhanced telemetry (iNET) standard. Two low-complexity versions that leverage the special properties of the iNET preamble and SOQPSK-TG operating at 2 samples/bit are developed. Simulation results show that most frame synchronizers are capable of achieving acceptable performance for  $E_b/N_0 \geq 0$  dB and the maximum frequency offset that can be tolerated increases with  $E_b/N_0$  to about 3% of the bit rate.

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#### I. INTRODUCTION

The traditional aeronautical telemetry link is a continuous-mode, one-way downlink from an airborne test article to a ground-based receiver. The recently introduced standard for a networked version of aeronautical telemetry, called iNET for integrated network enhanced telemetry [1], defines packetized transmission in which the preamble and attached sync marker (ASM) fields precede the data field. As of this writing, there is interest in periodically inserting the iNET preamble and ASM fields into continuous-mode links to enable data-aided estimators for synchronization, equalization, and channel quality monitoring. A prerequisite to enabling the data-aided estimators is identification of the start of the preamble in the received signal. Following the terminology most often used in the open literature, we refer to finding the start of the preamble as frame synchronization. Because the received signal corresponding to the preamble and ASM fields is to be used to estimate such things as timing, frequency, and phase offsets, frame synchronization must be performed in the presence of these offsets. Furthermore, frame synchronization must function with continuous phase modulation (CPM) as described in Section II.

Most previously published work on frame synchronization has been devoted to linear modulation. The usual approach is to start with symbol-synchronous matched filter outputs of the form

$$r_k = e^{j\theta} e^{j\omega k} a_k + z_k \tag{1}$$

where  $a_k$  is the kth symbol drawn from a constellation A;  $\theta$  and  $\omega$  are the phase and frequency offsets, respectively; and  $z_k$  is a sample of the additive noise. The basic idea is to find the starting point of a known word (whose elements are members of A) based on observing  $r_k$ . For  $\theta = \omega = 0$ in (1), the intuitive approach is to correlate the sequence  $r_k$ with a locally stored copy of the known word. Barker [2] was the first to describe the use of a correlator for binary phase shift keying (BPSK), i.e.,  $a_k \in \{-1, +1\}$ . Massev [3] derived the maximum likelihood (ML) frame synchronizer for the same case. The ML frame synchronizer comprised Barker's correlator and a nonlinear correction term to account for the random data preceding and following the known word. Simulation results presented by Nielsen [4] showed that the ML frame synchronizer achieves significant improvement over the simple correlator. Chiani [5] derived the ML synchronizer for BPSK for  $\omega = 0$ ,  $\theta \neq 0$ . The problem was formulated as a likelihood ratio test based on the marginalized probability density function of  $r_k$ , assuming  $\theta$  to be a uniform random variable on the interval  $[-\pi, \pi)$  and assuming uniformly distributed data symbols. The result involved a noncoherent correlation followed by a correction term. Generalizations to M-ary phase shift keying (MPSK) for M > 2 were also discussed. Lui and Tan [6] generalized Massey's ML result [3] to nonbinary quadrature amplitude modulation (QAM). Both the coherent case,  $\theta = \omega = 0$  in (1), and the noncoherent case,

 $\omega=0, \theta\neq0$  in (1), were examined. For the coherent case, the ML frame synchronizer is a correlator with a data correction term. For the noncoherent case,  $\theta$  was assumed to be a uniform random variable on the interval  $[-\pi,\pi)$ , and the corresponding ML frame synchronizer was derived from the marginalized probability density function of  $r_k$ . The resulting frame synchronizer comprises a form of noncoherent correlation and a data correction term. Chiani and Martini [7] derived analytical expressions bounding the false alarm and miss probabilities of the Lui and Tan approach [6], and Bastaki et al. [8] extended the work of Lui and Tan [6] to the observation of multiple frames.

For the case in (1), where  $\omega \neq 0$ , frame synchronization must be performed in the presence of a frequency offset. Three basic approaches have been used to deal with the frequency offset. The first, published by Choi and Lee [9] and devoted to MPSK, assumed  $\theta$  and  $\omega$ are independent uniform random variables on  $[-\pi, \pi)$  and derived the ML frame synchronizer from the marginalized probability density function of  $r_k$ . The result was a double correlator replacing the traditional correlator and a data correction term. The second approach, from Koo and Lee [10], incorporated the unknown phase and frequency in a joint estimation problem. A three-step process was described involving estimates of  $\theta$  and  $\omega$  for each possible starting location, followed by maximization over the starting locations. The third approach, described by Pedone et al. [11], is based on a novel stepwise (or stepwise) approximation of the phase ramp because of the frequency offset. The result is a partition of the required correlations into coherent and noncoherent segments.

Notable variations on the problem include "one shot" estimators for burst-mode communications [12], the closely related problem of variable length frames [13, 14], and the case in which the known word is distributed through the data [15, 16].

Because aeronautical telemetry uses CPM, the model for the observed random variables is different from (1). Whereas considerable attention has been devoted to the frame synchronization problem for linear modulations, only a few publications have tackled the problem for CPM. Stantchev and Fettweis [17] developed the ML frame synchronizer for noncoherent M-ary frequency shift keying (FSK) based on the magnitudes of the M-matched filter outputs. For the special case of M = 4, Bobula et al. [18] devised an ad hoc frame synchronization technique based on the sign of limiter-discriminator outputs. Both approaches leveraged the special properties of orthogonal FSK, and neither easily generalizes to arbitrary CPM. Huh and Krogmeier [19] modeled the observed symbols as noisy Markov chain outputs and derived frame synchronizers with excellent performance. Their approach was general, with application to both linear modulations with memory (e.g., coded modulation) or CPM, but considered only additive white Gaussian noise (AWGN) with  $\theta = \omega = 0$ . Hosseini and Perrins [20] discussed the problem of one-shot frame synchronization for CPM based on a clever preamble design described in [21]. The

result was a double correlator similar to that of Choi and Lee [9] and including a term that accounted for the noise-only measurements preceding the burst.

In this paper, we return to basic principles and revisit the problem of frame synchronization for CPM for continuous-mode transmission with a periodically inserted known word and in the presence of phase and frequency offsets and AWGN. After identifying suitable approximations, we follow the approaches of both Choi and Lee [9] and Pedone et al. [11] in dealing with the frequency offset. We show that correlation-type functions both with and without correction terms operating on samples of the received waveform and samples of the waveform corresponding to the known word follow from the analysis. Finally, by leveraging the particular properties of shaped offset QPSK, telemetry group version (SOQPSK-TG) and the iNET preamble, we develop low-complexity frame synchronizers whose performance is comparable to that of their full-complexity counterparts. Some initial simulation results using a subset of these frame synchronizers over mulitpath channels are described by the authors in [22].

This paper is organized as follows. A detailed problem formulation is described in Section II. In Section III, eight frame synchronizers that follow from the problem formulation are described, together with two low-complexity counterparts that leverage the special properties of the iNET preamble and SOQPSK-TG operating at 2 samples/bit. The performance of these frame synchronizers is outlined in Section IV, followed by the conclusions and a short discussion in Section V. The frame synchronizers described in Section III are derived in Appendices A and B. In these sections, boldface variables denote column vectors and  $(\cdot)^{\top}$  denotes the transpose operation.

## II. PROBLEM FORMULATION

The bit sequence used to create the CPM waveform is illustrated in Fig. 1a. The motivation for this particular structure is the iNET standard from aeronautical telemetry [1], but the generalization to other word fields is straightforward. The bits are used to modulate a CPM carrier, whose continuous-time complex-valued low-pass equivalent representation is [23–25]

$$s_c(t) = \exp\{j\phi(t)\}\tag{2}$$

where

$$\phi(t) = 2\pi \int_{-\infty}^{t} \sum_{k} h_k a_k g(x - kT_s) dx.$$
 (3)

Here,  $a_k$  is the kth symbol,  $h_k$  is the modulation index applied during the kth symbol interval,  $T_s$  is the symbol interval, and g(t) is the frequency pulse, spanning  $L_g$  symbol times and normalized to have area 1/2. Three forms of CPM are defined in the aeronautical telemetry standard IRIG 106 [26]:

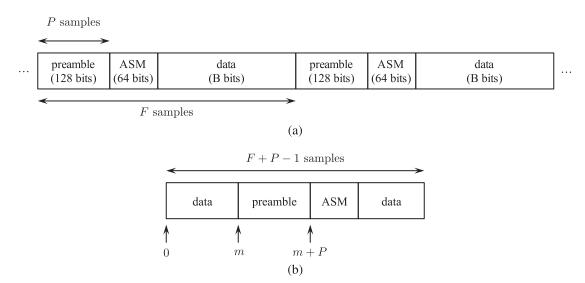


Fig. 1. Frame structure: (a) diagram illustrating periodic insertion of  $L_p$  preamble bits every  $L_f$  bits. Note relationship between bits and waveform samples. (b) Possible arrangement of P waveform samples corresponding to preamble bits in block of F samples.

- PCM/FM (for pulse code modulation with frequency modulation) is defined as follows:
  - $-a_k \in \{-1, +1\}$
  - $T_s = T_b$ , where  $T_b$  is the bit interval
  - $-h_k = 0.7$  for all k
  - -g(t) is the length-2 raised cosine pulse [23–25]
  - SOQPSK-TG is defined as follows:
- $-a_k$  is drawn from a constrained ternary alphabet, as described in [27, 28]
  - $-T_s = T_b$
  - $-h_k = 1/2$  for all k
- -g(t) is a pulse spanning 8 bit times, as described in [27, 28]
  - ARTM CPM (for advanced range telemetry CPM)
  - $-a_k \in \{-3, -1, +1, +3\}$
  - $-T_s=2T_b$
  - $h_k = \frac{4}{16}$  for k even and  $h_k = \frac{5}{16}$  for k odd
  - -g(t) is the length-3 raised cosine pulse [23–25]

The received signal is given by

$$r_c(t) = e^{j\theta} e^{j2\pi vt} s_c(t) + z_c(t)$$
 (4)

where  $\theta$  is the phase offset,  $\nu$  is the frequency offset in cycles per second, and  $z_c(t)$  is the thermal noise, modeled as a complex-valued circularly symmetric wide-sense stationary Gaussian random process with zero mean and autocorrelation

$$E\left\{z_{c}\left(t+\tau\right)z_{c}^{*}\left(t\right)\right\}=2N_{0}\delta\left(\tau\right).\tag{5}$$

In preparation for sampling at a rate of 1/T samples/s (equivalent to  $N = T_b/T$  samples/bit), the received signal is filtered by an antialiasing filter. We assume the antialiasing filter is an ideal low-pass filter with bandwidth  $^{1}/_{2T}$  and that the sample rate is sufficiently high so as not to produce noticeable aliasing in the spectrally shifted version of the CPM waveform  $s_c(t)e^{j2\pi\nu t}$ . The nth sample

of the resulting sample sequence is given by

$$r(nT) = e^{j\theta} e^{j\omega n} e^{j\phi(nT)} + z(nT)$$
 (6)

where  $\omega = 2\pi vT$  rads/sample and z(nT) is a sequence of uncorrelated complex-valued Gaussian random variables with zero mean and variance  $2\sigma^2 = 2N_0/T$ .

We assume continuous-mode transmission with an  $L_p$ -bit preamble sequence periodically inserted every  $L_f$  –  $L_p$  bits to produce a frame with a length of  $L_f$  bits. The situation is illustrated in Fig. 1a for framed transmissions using the iNET frame structure [1]. Assuming a sample rate of N samples/bit, the problem may be restated in terms of samples:  $P = NL_p$  samples, corresponding to the CPM waveform modulated by the preamble bits, are periodically inserted every  $N(L_f - L_p)$  samples to produce a frame comprising  $F = NL_f$  samples. Because the frame length is fixed and known, continuous transmission means that the preamble samples are guaranteed to be present in any contiguous block of F samples of the received signal. If the start of the preamble samples is assumed to be m and  $m \in \{0, 1, ..., F-1\}$ , then a block of F + P - 1consecutive samples of the received signal must be available [19]. The concept is illustrated in Fig. 1b.

For notational convenience, we stack the F+P-1 observed samples in the  $(F+P-1)\times 1$  vector

$$\mathbf{r} = \begin{bmatrix} r(0) & r(T) & \cdots & r((F+P-2)T) \end{bmatrix}^{\top}$$
 (7)

and let

$$\mathbf{p} = \begin{bmatrix} p(0) & p(T) & \cdots & p((P-1)T) \end{bmatrix}^{\top}$$
 (8)

be the  $P \times 1$  vector of samples of the CPM waveform corresponding to the preamble bits. Let  $\phi_m(nT)$  for n = 0,  $1, \ldots, F + P - 2$  be the sequence of phase samples corresponding to the case in which the preamble samples start at index m. Assuming the preamble samples occur only once in the block of F + P - 1 samples, the

elements of  $\phi_m(nT)$  are arranged as follows:

$$\underbrace{\phi_{m}(0)\cdots\phi_{m}((m-1)T)}_{\text{random data samples}}$$

$$\times \underbrace{\phi_{m}(mT)\cdots\phi_{m}((m+P-1)T)}_{\text{preamble samples}}$$

$$\times \underbrace{\phi_{m}((m+P)T)\cdots\phi_{m}((F+P-2)T)}_{\text{random data samples}}.$$
(9)

With this notation, we have  $p(kT) = e^{j\phi_m((m+k)T)}$  for k = 0, ..., P-1. Consequently, under the hypothesis the preamble sequence begins at index m, the model for the received samples is

$$\begin{split} r\left(nT\right) &= \begin{cases} e^{j\theta}e^{j\omega n}e^{j\phi_{m}(nT)} + z\left(nT\right) & 0 \leq n < m \\ e^{j\theta}e^{j\omega n}p\left((n-m)T\right) + z\left(nT\right) & m \leq n < m + P \\ e^{j\theta}e^{j\omega n}e^{j\phi_{m}(nT)} + z\left(nT\right) & m + P \leq n < F + P - 1. \end{cases} \end{split}$$

Let  $f(\mathbf{r}|m)$  be the conditional probability density function of the vector  $\mathbf{r}$  given that the preamble samples start at index m. The ML frame synchronizer produces the decision

$$\hat{m} = \underset{0 \le m \le F-1}{\operatorname{argmax}} \left\{ f\left(\mathbf{r} \mid m\right) \right\}. \tag{11}$$

The difficulty lies in computing a usable form for  $f(\mathbf{r}|m)$  from the conditional probability density function

$$f\left(\mathbf{r}|m,\theta,\omega,\phi_{d}\right) = \prod_{n=0}^{F+P-2} \frac{1}{2\pi\sigma^{2}} \exp\left\{-\frac{1}{2\sigma^{2}} \left|r\left(nT\right) - e^{j\theta}e^{j\omega n}e^{j\phi(nT)}\right|^{2}\right\}$$

$$(12)$$

where the vector

$$\phi_d = \left[ \phi(0) \cdots \phi((m-1)T) \ \phi((m+P)T) \ \cdots \phi((F+P-2)T) \right]^{\top}$$

$$(13)$$

represents the phase samples corresponding to random data bits. Here, the phase, frequency, and random data are nuisance parameters.

We follow two approaches in dealing with the nuisance parameters and compare their performance. The first approach follows that of Choi and Lee [9] and is detailed in Appendix A. The second approach follows Pedone et al. [11] and is detailed in Appendix B. Each approach is based on a number of assumptions about the nuisance parameters. These assumptions are spelled out at the beginning of each appendix.

The derivations in both appendices rely on the following high-level assumptions:

1) The samples with the ASM bits are treated as part of the random data. In other words, the CPM samples with the ASM bits are not used for frame synchronization. This omission is motivated by the complexity reduction

techniques associated with the repetitive structure of the preamble bits, as described in Section III.B. Frame synchronization performance improves if the ASM bits are used for frame synchronization. But this performance improvement is achieved at the expense of computational complexity.

- 2) In the block of F + P 1 received samples, the preamble samples are assumed to occur only once, even though the block may contain one entire occurrence of the preamble samples and part of another occurrence.
- 3) The first phase sample of the preamble sample sequence,  $\phi_m(mT)$ , is due only to the first preamble bit. This is clearly a coarse approximation, because for  $L_g > 1$  (the case for all three IRIG 106 modulations), each phase sample is influenced by the phase state, defined by previous input symbols, and the correlative state, defined by the previous  $L_g 1$  symbols and g(t). See [23–25] for more details. The assumption is made for mathematical convenience. To do otherwise leads to an intractable result or a torturously difficult form.

The simulation results presented in Section IV help assess the impact of these approximations.

## III. ML FRAME SYNCHRONIZERS

The frame synchronizers considered in this section are of the form

$$\hat{m} = \underset{0 \le m \le F-1}{\operatorname{argmax}} \{ L_i(m) \}$$
 (14)

where  $L_i(m)$  derives from ML considerations described in the sections that follow. It is understood that frame synchronization is based on  $L_i(\cdot)$  as described in [14].

## A. Full-Complexity Synchronizers

The first synchronizer considered is the simple correlator operating on the samples as discussed in [5]:

$$L_0(m) = \left| \sum_{k=0}^{P-1} r((k+m)T) p^*(kT) \right|.$$
 (15)

The absolute value is used to account for unknown phase rotations because of phase and frequency offsets. A closely related frame synchronizer is the heuristic function [5]

$$L_{0h}(m) = \left| \sum_{k=0}^{P-1} r((k+m)T) p^*(kT) \right| - \sum_{n=m}^{m+P-1} |r(nT)|.$$
(16)

The heuristic function comprises the correlator of  $L_0$  and a correction term. It is well known that these functions do not perform well in the presence of a frequency offset. They are included here to demonstrate the improvement realized by the more complicated functions that follow.

The next four frame synchronizers are based on the development outlined in Appendix A. The function that follows from ML principles based on the marginalized

probability density function of the vector  $\mathbf{r}$  is

$$L_{1}(m) = \sum_{i=1}^{P-1} \left[ \left| \sum_{k=i}^{P-1} r^{*} ((k+m)T) p(kT) r((k+m-i)T) \right| \times p^{*} ((k-i)T) \left| -\sum_{n=i+m}^{m+P-1} |r(nT)| |r((n-i)T)| \right].$$
(17)

The first term on the right-hand side of (17) was dubbed a "double correlation" by Choi and Lee [9]. The second term plays the role of Massey's data correction term [3]. Whereas the Choi-Lee function [9] computed the double correlation on the matched filter outputs and a local copy of the preamble symbols [cf. (1)], the double correlation here operates on samples of the received waveform and a local copy of the samples corresponding to the preamble bits. In this regard, the result is almost the same as that from Hosseini and Perrins [20], except the correction term is different.

Following Choi and Lee [9], three simplifications are investigated. First, the data correction term in (17) may be eliminated:

$$L_{2}(m) = \sum_{i=1}^{P-1} \left| \sum_{k=i}^{P-1} r^{*} ((k+m)T) p (kT) \right| \times r ((k+m-i)T) p^{*} ((k-i)T) \right|.$$
(18)

Second, only the i = 1 term in (17) is used:

$$L_{3}(m) = \left| \sum_{k=1}^{P-1} r^{*} ((k+m)T) p(kT) r((k+m-1)T) \right| \times p^{*} ((k-1)T) \left| -\sum_{n=1+m}^{m+P-1} |r(nT)| |r((n-1)T)| \right|.$$
(19)

Finally, only the i = 1 term in (18) is used:

$$L_{4}(m) = \left| \sum_{k=1}^{P-1} r^{*} ((k+m)T) p(kT) r((k+m-1)T) \right| \times p^{*} ((k-1)T) \right|.$$
 (20)

Each of these reduces the computational complexity of frame synchronization but does so with a loss in performance. The complexity–performance tradeoff is examined in Section IV.

The next two frame synchronizers follow from the development in Appendix B. The function that follows from a high signal-to-noise ratio (SNR) approximation

and ML principles is

$$L_{5}(m) = \sqrt{\Lambda_{0}(m) + \sum_{n=1}^{L_{PDI}-1} \bar{\Lambda}_{n}(m)} - \sum_{n=m}^{m+P-1} |r(nT)|$$
(21)

where

$$\Lambda_0(m) = \sum_{i=0}^{L_{\text{PDI}}-1} \left| \sum_{k=iL_{\text{coh}}}^{(i+1)L_{\text{coh}}-1} r((k+m)T) p^*(kT) \right|^2$$
 (22)

and

$$\bar{\Lambda}_{n}(m) = 2 \left| \sum_{i=0}^{L_{PDI}-1-n} \sum_{k=(i+n)L_{coh}-1}^{(i+n+1)L_{coh}-1} r((k+m)T) p^{*}(kT) \right| \times \sum_{k'=iL_{coh}-1}^{(i+1)L_{coh}-1} r^{*} ((k'+m)T) p(k'T) \right|. \quad (23)$$

Here,  $L_{PDI}$  and  $L_{coh}$  are design parameters such that P = $L_{\rm PDI}$   $L_{\rm coh}$  and are described later in this paragraph. In addition,  $\lambda_0(m)$  is the sample-based version of the noncoherent postdetection integration (NCPDI) concept, a standard result in code division multiple access acquisition [29, 30], and  $\bar{\Lambda}_n(m)$  is the sample-based version of the *n*-span differential postdetection integration (*n*-span DPDI) principle [11]. The NCPDI term partitions the correlation involving P consecutive samples of the received signal and the pilot samples into segments of length-L<sub>coh</sub> coherent correlation and noncoherently combines the  $L_{PDI}$  coherent correlations (the subscript PDI refers to postdetection integration). This would seem to solve the problem for the phase and frequency offsets. Therefore, the last frame synchronizer we consider is based just on the NCPDI term  $\Lambda_0(m)$ :

$$L_6(m) = \sum_{i=0}^{L_{\text{PDI}}-1} \left| \sum_{k=iL_{\text{coh}}}^{(i+1)L_{\text{coh}}-1} r((k+m)T) p^*(kT) \right|^2. \tag{24}$$

## B. Reduced-Complexity Synchronizers

The eight frame synchronizers described in the previous section involve correlations between the samples of the received signal and a local copy of preamble samples. For each m, the correlation requires Pmultiply-accumulate operations (and quite a number more for the double correlation functions). For SOQPSK-TG and the iNET preamble, an interesting complexityreducing technique is available. The 128-bit iNET preamble is the 16-bit sequence CD98hex repeated eight times [1, pp. 48]. A plot of the complex-valued samples of SOQPSK-TG for N = 2 samples/bit corresponding to the iNET preamble is shown in Fig. 2. The samples appear to be clustered about eight equally spaced points on the unit circle. This suggests a quantized version of the local copy of the preamble samples, where the quantized version comprises values only at  $\pm 1$ ,  $\pm j$ , and  $(\pm 1 \pm j)/\sqrt{2}$ . Let  $\tilde{p}(nT)$  for n = 0, ..., 255 be the quantized version of p(nT).

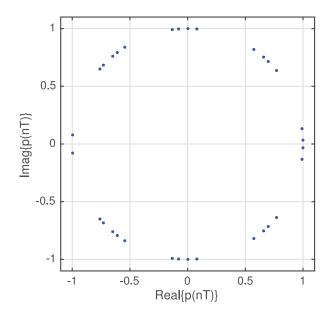


Fig. 2. Complex-valued samples of SOQPSK-TG, at sample rate of 2 samples/bit, corresponding to iNET preamble.

The reduction in computational complexity is best illustrated by considering the correlation

$$R(k) = \sum_{i=0}^{255} r((i+k)T) p^*(iT).$$
 (25)

This correlation, based on the unquantized preamble samples, requires 256 complex-valued multiply-accumulate operations, which translates to 1024 real-valued multiplications and 1024 real-valued additions. Using  $\tilde{p}(iT)$  (the quantized versions) in place of p(iT) (the unquantized versions) all but eliminates the multiplications. If  $r(nT) = r_R(nT) + jr_I(nT)$ , the real and imaginary parts of R(k) may be expressed as

$$\operatorname{Re} \left\{ R(k) \right\} \approx \sum_{i \in \mathcal{I}_{1}} r_{R} ((i+k)T) - \sum_{i \in \mathcal{I}_{2}} r_{R} ((i+k)T) + \sum_{i \in \mathcal{I}_{3}} r_{I} ((i+k)T) - \sum_{i \in \mathcal{I}_{4}} r_{I} ((i+k)T) + \frac{1}{\sqrt{2}} \left[ \sum_{i \in \mathcal{I}_{5}} r_{R} ((i+k)T) - \sum_{i \in \mathcal{I}_{6}} r_{R} ((i+k)T) + \sum_{i \in \mathcal{I}_{7}} r_{I} ((i+k)T) - \sum_{i \in \mathcal{I}_{8}} r_{I} ((i+k)T) \right]$$

$$+ \sum_{i \in \mathcal{I}_{7}} r_{I} ((i+k)T) - \sum_{i \in \mathcal{I}_{8}} r_{I} ((i+k)T)$$

$$(26)$$

and

$$\operatorname{Im}\left\{R\left(k\right)\right\} \approx \sum_{i \in \mathcal{I}_{1}} r_{I}\left(\left(i+k\right)T\right) - \sum_{i \in \mathcal{I}_{2}} r_{I}\left(\left(i+k\right)T\right)$$
$$- \sum_{i \in \mathcal{I}_{3}} r_{R}\left(\left(i+k\right)T\right) + \sum_{i \in \mathcal{I}_{4}} r_{R}\left(\left(i+k\right)T\right)$$

$$+\frac{1}{\sqrt{2}}\left[\sum_{i\in\mathcal{I}_{5}}r_{I}\left(\left(i+k\right)T\right)-\sum_{i\in\mathcal{I}_{6}}r_{I}\left(\left(i+k\right)T\right)\right]$$
$$-\sum_{i\in\mathcal{I}_{7}}r_{R}\left(\left(i+k\right)T\right)+\sum_{i\in\mathcal{I}_{8}}r_{R}\left(\left(i+k\right)T\right)\right]$$
(27)

respectively, where

$$\mathcal{I}_{1} = \{i, 0 \leq i < 256 | \tilde{p}(iT) = +1\} 
\mathcal{I}_{2} = \{i, 0 \leq i < 256 | \tilde{p}(iT) = -1\} 
\mathcal{I}_{3} = \{i, 0 \leq i < 256 | \tilde{p}(iT) = +j\} 
\mathcal{I}_{4} = \{i, 0 \leq i < 256 | \tilde{p}(iT) = -j\} 
\mathcal{I}_{5} = \left\{i, 0 \leq i < 256 | \operatorname{Re}\left[\tilde{p}(iT)\right] = +1 / \sqrt{2}\right\} 
\mathcal{I}_{6} = \left\{i, 0 \leq i < 256 | \operatorname{Re}\left[\tilde{p}(iT)\right] = -1 / \sqrt{2}\right\} 
\mathcal{I}_{7} = \left\{i, 0 \leq i < 256 | \operatorname{Im}\left[\tilde{p}(iT)\right] = +1 / \sqrt{2}\right\} 
\mathcal{I}_{8} = \left\{i, 0 \leq i < 256 | \operatorname{Im}\left[\tilde{p}(iT)\right] = -1 / \sqrt{2}\right\}.$$

This approach requires 512 real-valued additions and only 2 real-valued multiplications. The principle may be extended to the double correlation terms of (17)–(20) in a straightforward way.

Further complexity reductions are available for  $L_6$ . Because the preamble bit sequence comprises eight repetitions of the 16-bit pattern CD98<sub>hex</sub>, the corresponding sample sequence is, to a first-order approximation, eight repetitions of the Q=16N sample sequence corresponding to the 16-bit sequence CD98<sub>hex</sub>. Let q(nT) for  $n=0,\ldots,16N-1$  be this sequence. The relationship between the short sample sequence q(nT) and the entire preamble sample sequence is illustrated in Fig. 3. In the figure, the start of the preamble sequence corresponds to the mth sample of the received signal. Thus, r(mT) is aligned with both p(0) and the first occurrence of q(0). Q samples later, r((m+Q)T) is aligned with p(QT) and with the second occurrence of q(0). And so on.

A remarkable reduction in computational complexity occurs with  $L_6$  for N=2 samples/bit. Here, the assignments  $L_{\rm coh}=Q=32$  and  $L_{\rm PDI}=8$  may be used. Recognizing that  $\tilde{p}((\ell+32i)T)=\tilde{q}(\ell T)$ ,  $L_6$  becomes

$$L_{6q}(m) = \sum_{i=0}^{7} \left| \sum_{\ell=0}^{31} r((\ell+32i+m)T)\tilde{q}^*(\ell T) \right|^2. \quad (29)$$

Using the quantized values for the short preamble sample sequences and collecting similar terms gives

$$L_{6q}(m) = \sum_{i=0}^{7} \left[ I^{2}(m,i) + Q^{2}(m,i) \right]$$
 (30)

<sup>&</sup>lt;sup>1</sup> This is not quite true for the first *Q*-sample sequence because of the memory in SOQPSK-TG.

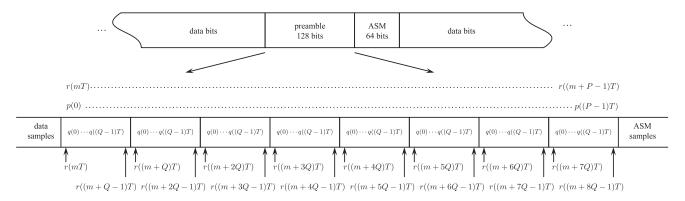


Fig. 3. Illustration of relationship between samples of received signal, samples of SOQPSK-TG signal corresponding to preamble, and samples of SOQPSK-TG signal corresponding to constituent bit sequence CD98<sub>hex</sub>.

where

$$I(m, i) \approx \sum_{i \in \mathcal{L}_{1}} r_{R} (\ell + 32i + m) - \sum_{\ell \in \mathcal{L}_{2}} r_{R} (\ell + 32i + m) + \sum_{\ell \in \mathcal{L}_{3}} r_{I} (\ell + 32i + m) - \sum_{\ell \in \mathcal{L}_{4}} r_{I} (\ell + 32i + m) + \frac{1}{\sqrt{2}} \left[ \sum_{\ell \in \mathcal{L}_{5}} r_{R} (\ell + 32i + m) - \sum_{\ell \in \mathcal{L}_{6}} r_{R} (\ell + 32i + m) + \sum_{\ell \in \mathcal{L}_{7}} r_{I} (\ell + 32i + m) - \sum_{\ell \in \mathcal{L}_{8}} r_{I} (\ell + 32i + m) - \sum_{\ell \in \mathcal{L}_{8}} r_{I} (\ell + 32i + m) \right]$$

$$(31)$$

and

$$Q(m, i) \approx \sum_{\ell \in \mathcal{L}_{1}} r_{I} (\ell + 32i + m) - \sum_{\ell \in \mathcal{L}_{2}} r_{I} (\ell + 32i + m) - \sum_{\ell \in \mathcal{L}_{3}} r_{R} (\ell + 32i + m) + \sum_{\ell \in \mathcal{L}_{4}} r_{R} (\ell + 32i + m) + \frac{1}{\sqrt{2}} \left[ \sum_{\ell \in \mathcal{L}_{5}} r_{I} (\ell + 32i + m) - \sum_{\ell \in \mathcal{L}_{6}} r_{I} (\ell + 32i + m) - \sum_{\ell \in \mathcal{L}_{7}} r_{R} (\ell + 32i + m) + \sum_{\ell \in \mathcal{L}_{8}} r_{R} (\ell + 32i + m) \right]$$

$$(32)$$

with

$$\mathcal{L}_{1} = \{0, 8, 16, 24\}$$

$$\mathcal{L}_{2} = \{4, 20\}$$

$$\mathcal{L}_{3} = \{2, 10, 14, 22\}$$

$$\mathcal{L}_{4} = \{6, 18, 26, 30\}$$

$$\mathcal{L}_{5} = \{1, 7, 9, 15, 17, 23, 25, 31\}$$

$$\mathcal{L}_{6} = \{3, 5, 11, 12, 13, 19, 21, 27, 28, 29\}$$

$$\mathcal{L}_{7} = \{1, 3, 9, 11, 12, 13, 15, 21, 23\}$$

$$\mathcal{L}_{8} = \{5, 7, 17, 19, 25, 27, 28, 29, 31\}.$$
(33)

To further simplify the computation, it is straightforward to show that

$$I(m+32, i) = I(m, i+1)$$

$$Q(m+32, i) = Q(m, i+1)$$
(34)

for i = 0, 1, ..., 6. The consequence of this relationship is that (30) may be computed by

$$L_{6q}(m) = L_{6q}(m-32) - [I^{2}(m-32,0) + Q^{2}(m-32,0)] + I^{2}(m,7) + Q^{2}(m,7)$$
(35)

for  $m \ge 32$ . Thus,  $L_{6q}$  may be computed recursively, where each recursion requires four real-valued multiplications and 67 real-valued additions.

The performance penalty associated with equating  $L_{\rm coh}$  to Q can be measured using the energy loss due to the phase rotation over the coherent correlation interval  $L_{\rm coh}$ . The energy loss is proportional to  ${\rm sinc^2}(L_{\rm coh}T\nu)$  [31], where  $\nu$  is the continuous-time frequency offset; see (4). A plot of this loss versus the normalized frequency offset for SOQPSK-TG operating at N=2 samples/bit is shown in Fig. 4. An energy loss of 1 dB occurs when the frequency offset is approximately 1.5% of the bit rate. The energy loss of 3 dB occurs at a frequency offset of 2.75% of the bit rate. This shows that as long as the frequency offset is on the order of 1% to 2% of the bit rate, the assignment  $L_{\rm coh}=Q$  has the potential to be a robust frame synchronizer.

A summary of the computational complexity, measured by the number of real-valued multiplication and square-root operations, is given in Table I for a normalized sample rate of N=2 samples/bit. In the table,  $L_{1q}(m)$  denotes (17), with  $\tilde{p}(kT)$  replacing p(kT). The quantized version of  $L_6$  has clear advantages from a computational complexity point of view. The performance of these frame synchronizers is summarized in the next section.

## IV. PERFORMANCE RESULTS

We evaluated the performance of the frame synchronizers  $L_0$  through  $L_6$  and the quantized version  $L_{6q}$  using computer simulation. The computer simulations followed the usual procedure: samples of an SOQPSK-TG

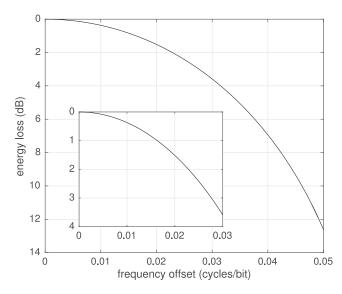


Fig. 4. Energy loss because of coherent correlation over interval  $L_{\rm coh} = Q = 32$  samples as function of normalized frequency. These results correspond to N = 2 samples/bit.

TABLE I Computational Complexity for Examined Frame Synchronizers for SOQPSK-TG with iNET Preamble and Operating at N=2 Samples/Bit

Synchronizer	Multiplication	Square roo
$L_0(m)$	1024	1
$L_{0h}(m)$	1024	2
$L_1(m)$	424 320	511
$L_{1q}(m)$	163 200	511
$L_2(m)$	391 680	255
$L_3(m)$	3315	257
$L_4(m)$	3060	1
$L_5(m)$	1168	293
$L_6(m)$	1024	0
$L_{6q}(m)$	4	0

signal operating at N=2 samples/bit were created using a sequence of 1000 random data bits, the iNET preamble, the iNET ASM, and another sequence of 1000 random data bits. A frequency offset, equivalent to 0%, 1%, or 5% of the bit rate, was applied to the SOQPSK-TG samples. Finally, AWGN samples, whose variance was defined by  $E_b/N_0$ , were added. The frame synchronizer (14) was applied to the received samples. For each value of the frequency offset and  $E_b/N_0$ , the process was repeated 10 000 times.

We now move to the question of performance measure. For linear modulations, most frame synchronizers operate on matched filter outputs sampled at 1 sample/symbol. Correctly identifying the matched filter output corresponding to the first preamble symbol is a well-defined detection event. Otherwise, a false-alarm event occurs. In contrast, the frame synchronizers in this paper operate at more than 1 sample/bit. Correspondingly, the notions of detection and false alarm are less clear. Consequently, we measure the difference between the true starting index m and its estimate  $\hat{m}$ . The well-established

convention for assessing the performance of estimation problems such as this one is to examine the mean and variance of the estimate error. Plots of the simulated mean error are omitted here in the interest of space. The frame synchronizers with low error variance also have close-to-zero mean error. The frame synchronizers with high error variance have large mean error. For this particular problem, the mean error does not provide information (regarding the suitability of the approach to the problem) that cannot be obtained from the error variance.

The simulated estimation error variance for  $\nu T_b = 0$  is plotted in Fig. 5a. The relationship between  $L_1$  and  $L_2$  is noteworthy. At low values of  $E_b/N_0$ ,  $L_2$  has a much lower variance than  $L_1$ , although  $L_2$  is an approximation of  $L_1$ . As  $E_b/N_0$  increases, the situation reverses.  $L_3$  and  $L_4$  have poor performance. This is because the respective approximations of  $L_1$  and  $L_2$  stretch the assumptions used to derive  $L_1$  too far.  $L_5$ ,  $L_6$ , and  $L_{6q}$  all have similar performance except that the variance of  $L_5$  is about one order of magnitude lower than the variances of  $L_6$  and  $L_{6q}$ . The variance of  $L_{6q}$  is only about two times that of  $L_6$ . Curiously,  $L_0$  and  $L_{0h}$  achieve competitive estimator error variances; the variance of  $L_0$  is near the lowest and the variance of  $L_{0h}$  is, for all practical purposes, the lowest. This is to be expected, given that there is no frequency offset. The fact that  $L_{0h}$  outperforms  $L_0$  is why the heuristic function is popular.

Most index errors for  $L_0$ ,  $L_0$ ,  $L_1$ ,  $L_2$ ,  $L_5$ ,  $L_6$ , and  $L_{6q}$  are small integer multiples of Q (the length, in samples, of the short sequences that comprise the preamble). Why this is so is illustrated by the plot in Fig. 6, which shows a plot of  $L_0$  for the case of zero frequency offset and no noise. The iNET preamble is periodic (with period Q samples), which creates mini correlation peaks every Q samples when the occurrence of the preamble in the received samples is involved in the correlation. The highest peak corresponds to the case in which the received sample data and the preamble are aligned. The lower peaks, either side of the highest peak, correspond to cases in which fewer than eight of the preamble short sequences are aligned with the locally stored preamble template (cf. Fig. 3).

The frame synchronizer in (14) can be adjusted to look for index errors that are small integer multiples of Q. This adjustment removes these index errors. As an example, the adjustment based on looking for index errors  $\pm Q$  and  $\pm 2Q$  produces the estimators whose variance is plotted in Fig. 5b. Except for  $L_3$  and  $L_4$ , this adjustment all but eliminates index errors for  $E_b/N_0 > 0$  dB. There are no circle or star markers in this plot; this is because all index errors for  $L_0$  and  $L_{0h}$  in Fig. 5a were either  $\pm Q$  or  $\pm 2Q$ .

The next set of plots, shown in Fig. 7, are for  $\nu T_b = 0.01$  cycles/bit. The variances of the unadjusted synchronizers are plotted in Fig. 7a and the variances of the adjusted synchronizers are plotted in Fig. 7b. The general trends are the same as before:  $L_3$  and  $L_4$  display unusably high index error variance, and the variance of  $L_1$  is higher than that of  $L_2$ ,  $L_5$ ,  $L_6$ , and  $L_{6g}$  for low values of

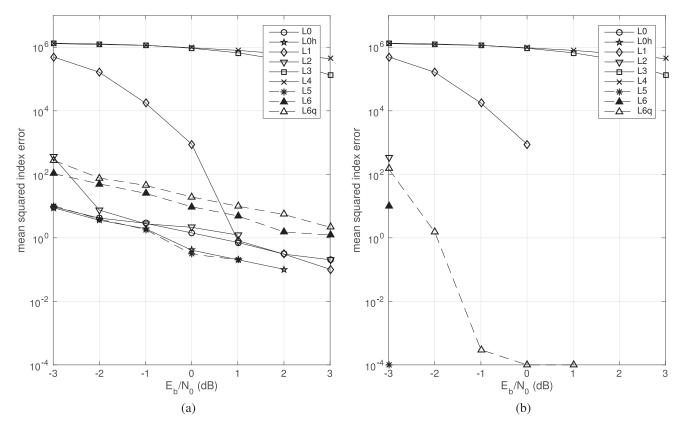


Fig. 5. Simulated mean-squared error performance for frame synchronizers described in this paper for frequency offset 0: (a) performance without adjustments, (b) performance with adjustments for errors of  $\pm Q$  and  $\pm 2Q$ .

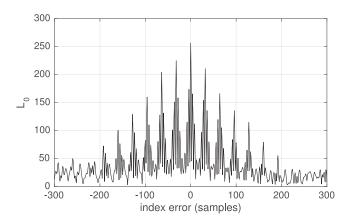


Fig. 6. Plot of  $L_0(m)$  for SOQPSK-TG operating at N = 2 samples/bit and using iNET preamble without frequency offset and without noise.

 $E_b/N_0$ ; however, the situation reverses at a higher SNR. The major difference here is that  $L_0$  and  $L_{0h}$  are quite poor. This is to be expected, because  $L_0$  and  $L_{0h}$  are not designed to account for a frequency offset.

Finally, the plots shown in Fig. 8 display the variances of the frame synchronizers for  $\nu T_b = 0.05$  cycles/bit. The same trends observed in Fig. 7 are present. In this plot, the value of  $E_b/N_0$  for which  $L_1$  is better than  $L_5$ ,  $L_6$ , and  $L_{6q}$  is lower. This is because the energy loss in  $L_5$  and  $L_6$ 

because of frequency offset (see Fig. 4) limits performance. In contrast,  $L_1$  and  $L_2$  were derived assuming the possibility of a larger frequency offset

Because the simulated system operates at N=2 samples/bit, an estimation error of  $\pm 1$  sample is not catastrophic; data-aided estimators see a shift by  $^1/_2$  bit time. Consequently, any frame synchronizer that achieves an RMS index error less than 1 sample is a viable solution. The plots in Figs. 5, 7, and 8 show that for  $vT_b \leq 0.01$  cycles/bit, the frame synchronizers  $L_1, L_2, L_5, L_6$ , and  $L_{6q}$  are capable of meeting this requirement for  $E_b/N_0 > 0$  dB. Of these, Table I shows that the quantized NCPDI synchronizer  $L_{6q}$  is the best choice by a considerable margin.

The computational advantage of  $L_{6q}$  motivates further investigation. The simulated index error variance for the adjusted version of (14) based on  $L_{6q}$  as a function of the normalized frequency offset for fixed values of  $E_b/N_0$  is plotted in Fig. 9. The performance of the unquantized version in (14) based on  $L_6$  and the standard correlator  $L_0$  are also plotted for reference. Using our rule of thumb that the RMS index error not exceeding 1 sample is acceptable, these plots show that as  $E_b/N_0$  increases, the maximum frequency offset, which  $L_{6q}$  is able to tolerate also increases. For example, at  $E_b/N_0 = 0$  dB, the simulation

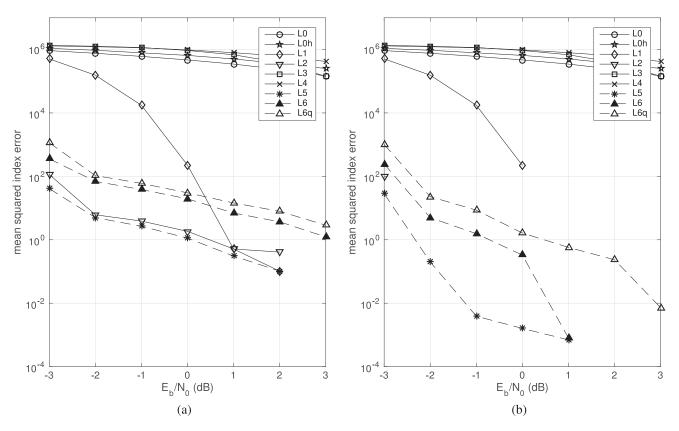


Fig. 7. Simulated mean-squared error performance for frame synchronizers described in this paper for frequency offset 1% of bit rate: (a) performance without adjustments, (b) performance with adjustments for errors of  $\pm Q$  and  $\pm 2Q$ .

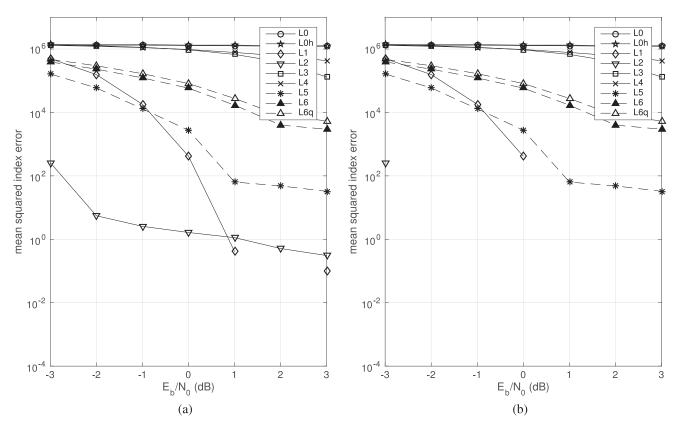
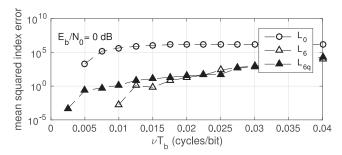
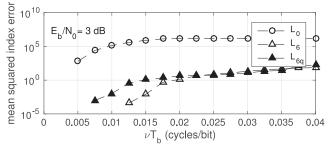


Fig. 8. Simulated mean-squared error performance for frame synchronizers described in this paper for frequency offset 5% of bit rate: (a) performance without adjustments, (b) performance with adjustments for errors of  $\pm Q$  and  $\pm 2Q$ .





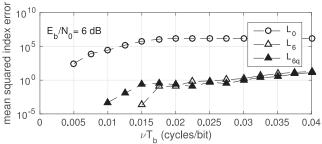


Fig. 9. Simulated index error variance for adjusted version of (14) using  $L_0$ ,  $L_6$ , and  $L_{6q}$  as function of normalized frequency offset  $\nu T_b$  cycles/bit for three values of  $E_b/N_0$ .

results show acceptable performance for  $L_{6q}$  for  $\nu T_b \le 0.01$  cycle/bit, at  $E_b/N_0 = 3$  dB for  $\nu T_b \le 0.015$  cycle/bit, and at  $E_b/N_0 = 6$  dB for  $\nu T_b \le 0.03$  cycle/bit.

## V. DISCUSSION AND CONCLUSIONS

By returning to basic principles, we derived ML frame synchronizers for CPM operating on samples of the received waveform, where the sampling rate satisfies the Nyquist sampling theorem. The derivation followed two approaches for dealing with the nuisance parameters: the first by Choi and Lee [9] and the second by Pedone et al. [11]. Simulation results showed that both approaches, and a number of reduced-complexity approximations, are capable of performing frame synchronization in the aeronautical telemetry environment. By leveraging the structural properties of the iNET preamble and for SOQPSK-TG operating at 2 samples/bit, low-complexity frame synchronizers with excellent performance were obtained.

The techniques outlined in this paper could be applied to any of the CPMs described in IRIG 106. As of the writing of this paper, SOQPSK-TG is the only one of the three being considered for use with the iNET preamble. Therefore, only results using SOQPSK-TG were simulated.

## APPENDIX A. DERIVATION OF THE ML FRAME SYNCHRONIZER FOLLOWING CHOI AND LEE [9]

The outline and corresponding assumptions for this approach are as follows:

- 1) Average the likelihood function in (12) with respect to the phase  $\theta$ . This step assumes the phase is uniformly distributed over the interval  $[-\pi, \pi)$ .
- 2) Average the result of the previous step with respect to the frequency offset  $\omega$ . This step assumes the frequency offset is uniformly distributed over the interval  $[-\pi, \pi)$ .
- 3) Average the result of the previous step with respect to the random data. This step assumes the phase samples due to random data are independent and identically distributed random variables, each with a uniform distribution over the interval  $[-\pi, \pi)$ .

Assumption 2 is an approximation, because  $\omega=\pi$  is equivalent to a frequency offset equal to half the sample rate. A frequency offset this high causes irreparable distortion. To see that this is so, assume the bandwidth of the continuous-time complex-valued low-pass equivalent signal is B Hz. Then, the relationship between the continuous-time frequency offset  $\nu$  and the sample rate must be

$$|\nu| + B < \frac{1}{2T}. (36)$$

Using  $\omega = 2\pi vT$ , the corresponding relationship for  $\omega$  is

$$|\omega| < \pi \left(1 - 2BT\right). \tag{37}$$

Thus, any signal with B>0 cannot accommodate  $\omega=\pi$ . Assumption 2 is made purely for mathematical convenience. The independence component of Assumption 3 is clearly not true for CPM. Independence is assumed purely in the interest of tractability. However, the uniform distribution is a good assumption for the marginal distribution of each phase sample.

The starting point is to rewrite the conditional probability density function in (12) as

$$f(\mathbf{r}|m,\theta,\omega,\phi_d) = \frac{1}{(2\pi\sigma^2)^{F+P-1}} \exp\left\{-\frac{1}{2\sigma^2} \sum_{n=0}^{F+P-2} \left| r(nT) - e^{j\theta} e^{j\omega n} e^{j\phi_m(nT)} \right|^2\right\}$$
(38)

where the hypothesis is that the samples corresponding to the iNET preamble coincide with the received samples starting at index m. Averaging over the phase angle  $\theta$  gives

$$f(\mathbf{r}|m,\omega,\phi_d) = \frac{1}{2\pi} \int_{-\pi}^{\pi} f(\mathbf{r}|m,\theta,\omega,\phi_d) d\theta \quad (39)$$

$$=C(\mathbf{r})I_0\left(\frac{1}{\sigma^2}\left|\sum_{n=0}^{F+P-2}r^*(nT)e^{j\omega n}e^{j\phi_m(nT)}\right|\right)$$
(40)

where

$$C(\mathbf{r}) = \frac{1}{(2\pi\sigma^2)^{F+P-1}} \exp\left\{-\frac{1}{2\sigma^2} \sum_{n=0}^{F+P-2} \left[ |r(nT)|^2 + 1 \right] \right\}.$$
(41)

Next, the dependency on the frequency offset  $\omega$  is removed by averaging:

$$f(\mathbf{r}|m,\phi_d) = \frac{1}{2\pi} \int_{-\pi}^{\pi} f(\mathbf{r}|m,\omega,\phi_d) d\omega$$
 (42)

$$= \frac{C(\mathbf{r})}{2\pi} \int_{-\pi}^{\pi} I_0 \left( \frac{1}{\sigma^2} \left| \sum_{n=0}^{F+P-2} r^*(nT) e^{j\omega n} e^{j\phi_m(nT)} \right| \right) d\omega.$$
(43)

The integral is evaluated using the approximation

$$I_0(z) \approx 1 + \frac{z^2}{4} + \frac{z^4}{64}.$$
 (44)

Thus,

$$f(\mathbf{r}|m,\phi_d) \approx \frac{C(\mathbf{r})}{2\pi} \int_{-\pi}^{\pi} d\omega + \frac{C(\mathbf{r})}{16\pi^2} \int_{-\pi}^{\pi} \left| \sum_{n=0}^{F+P-2} r^*(nT) e^{j\omega n} e^{j\phi_m(nT)} \right|^2 d\omega + \frac{C(\mathbf{r})}{1024\pi^4} \int_{-\pi}^{\pi} \left| \sum_{n=0}^{F+P-2} r^*(nT) e^{j\omega n} e^{j\phi_m(nT)} \right|^4 d\omega. \tag{45}$$

The first integral is simply  $C(\mathbf{r})$ . The second and third integrals are reduced using the property

$$\int_{-\pi}^{\pi} e^{jk\omega} d\omega = \begin{cases} 2\pi & k = 0\\ 0 & k \neq 0. \end{cases}$$
 (46)

The second integral is

$$\int_{-\pi}^{\pi} \left| \sum_{n=0}^{F+P-2} r^*(nT) e^{j\theta} e^{j\omega n} e^{j\phi_m(nT)} \right|^2 d\omega$$

$$= \sum_{n_1=0}^{F+P-2} \sum_{n_2=0}^{F+P-2} r^*(n_1T) r(n_2T) e^{j\phi_m(n_1T)} e^{-j\phi_m(n_2T)}$$

$$\times \int_{-\pi}^{\pi} e^{j\omega n_1} e^{-j\omega n_2} d\omega$$

$$= 2\pi \sum_{n=0}^{F+P-2} |r(nT)|^2. \tag{47}$$

The third integral is

$$\int_{-\pi}^{\pi} \left| \sum_{n=0}^{F+P-2} r^*(nT) e^{j\omega n} e^{j\phi(nT)} \right|^4 d\omega$$

$$= \sum_{n_1=0}^{F+P-2} \sum_{n_2=0}^{F+P-2} \sum_{n_3=0}^{F+P-2} \sum_{n_4=0}^{F+P-2} G(m, n_1, n_2, n_3, n_4)$$

$$\times \int_{-\pi}^{\pi} e^{j\omega n_1} e^{-j\omega n_2} e^{j\omega n_3} e^{-j\omega n_4} d\omega$$
(48)

where

$$G(m, n_1, n_2, n_3, n_4)$$

$$= r^* (n_1 T) r (n_2 T) r^* (n_3 T) r (n_4 T) e^{j\phi_m(n_1 T)} e^{-j\phi_m(n_2 T)}$$

$$\times e^{j\phi_m(n_3 T)} e^{-j\phi_m(n_4 T)}.$$
(49)

The integral is zero except where  $n_1 - n_2 + n_3 - n_4 = 0$ , in which case the integral evaluates to  $2\pi$ . Parameterizing the nonzero condition using  $i = n_1 - n_4 = n_2 - n_3$  and substituting for the integral gives

$$\int_{-\pi}^{\pi} \left| \sum_{n=0}^{F+P-2} r^*(nT) e^{j\omega n} e^{j\phi_m(nT)} \right|^4 d\omega$$

$$= 2\pi \sum_{n_1=0}^{F+P-2} \sum_{n_2=0}^{F+P-2} G(m, n_1, n_2, n_2, n_1)$$

$$+ 4\pi \sum_{i=1}^{F+P-2} \sum_{n_1=i}^{F+P-2} \sum_{n_2=i}^{F+P-2} G(m, n_1, n_2, n_2 - i, n_1 - i)$$

$$= 2\pi \sum_{n_1=0}^{F+P-2} \sum_{n_2=0}^{F+P-2} |r(n_1T)|^2 |r(n_2T)|^2$$

$$+ 4\pi \sum_{i=1}^{F+P-2} \sum_{n_1=i}^{F+P-2} \sum_{n_2=i}^{F+P-2} G(m, n_1, n_2, n_2 - i, n_1 - i).$$
(50)

Assembling the results gives

$$f(\mathbf{r}|m,\phi_d) \approx C(\mathbf{r}) + \frac{C(\mathbf{r})}{8\pi} \sum_{n=0}^{F+P-2} |r(nT)|^2 + \frac{C(\mathbf{r})}{512\pi^3} \sum_{n_1=0}^{F+P-2} \sum_{n_2=0}^{F+P-2} |r(n_1T)|^2 |r(n_2T)|^2 + \frac{C(\mathbf{r})}{256\pi^3} \sum_{i=1}^{F+P-2} \sum_{n_1=i}^{F+P-2} \sum_{n_2=i}^{F+P-2} \times G(m, n_1, n_2, n_2 - i, n_1 - i).$$
 (51)

Finally, we average over the data. Here, we have

$$f(\mathbf{r}|m) = \int \cdots \int f(\mathbf{r}|m, \phi_d) f(\phi_d) d\phi_d.$$
 (52)

The only term in (51) involving data is the fourth term, and each summand in this term contains only four signal samples. We are interested in

$$\int \cdots \int G(m, n_1, n_2, n_2 - i, n_1 - i) f(\widetilde{\phi}) d\widetilde{\phi}$$

$$= r^*(n_1 T) r(n_2 T) r^*((n_2 - i) T) r((n_1 - i) T)$$

$$\times I(m, i, n_1, n_2)$$
(53)

where

$$I(m, i, n_1, n_2) = \int \cdots \int e^{-j\phi_m(n_1T)} e^{j\phi_m(n_2T)} e^{j\phi_m((n_2-i)T)}$$
$$\times e^{-j\phi_m(n_1-i)} f(\widetilde{\phi}) d\widetilde{\phi}$$
(54)

and

$$\widetilde{\phi} = F + \left[ \phi_m(n_1 T) \ \phi_m(n_2 T) \ \phi_m((n_2 - i) T) \ \phi_m((n_1 - i) T) \right]^{\top}.$$
(55) 
$$1 \le i \le P - 1$$

When  $m \le n_1 < P + m$ ,  $e^{j\phi_m(n_1T)}$  is a sample corresponding to the known preamble and is treated as a constant, not a random variable. An identical situation holds for  $n_2$ ,  $n_2 - i$ ,  $n_1 - i$ . The integral is zero any time i,  $n_1$ , and  $n_2$  are such that at least one sample corresponding to a random data symbol is included the integrand or except when the integrand is a constant (this occurs when  $n_1 = n_2$ ).

It is convenient to consider the cases  $i \ge P$  and i < P separately. To this end, we write

$$\sum_{i=1}^{F+P-2} \sum_{n_1=i}^{F+P-2} \sum_{n_2=i}^{F+P-2} r^* (n_1 T) r (n_2 T) r^* ((n_2 - i) T)$$

$$\times r ((n_1 - i) T) I (m, i, n_1, n_2)$$

$$= \sum_{i=1}^{P-1} \sum_{n_1=i}^{F+P-2} \sum_{n_2=i}^{F+P-2} r^* (n_1 T) r (n_2 T) r^* ((n_2 - i) T)$$

$$\times r ((n_1 - i) T) I (m, i, n_1, n_2)$$

$$+ \sum_{i=P}^{F+P-2} \sum_{n_1=i}^{F+P-2} \sum_{n_2=i}^{F+P-2} r^* (n_1 T) r (n_2 T)$$

$$\times r^* ((n_2 - i) T) r ((n_1 - i) T) I (m, i, n_1, n_2). (56)$$

The first term on the right-hand side of (56) represents the case in which i < P. For i < P, the integral is zero except when the following conditions are simultaneously achieved:

$$m \leq n_1 < m + P$$
  
 $m \leq n_1 - i < m + P$   
 $m \leq n_2 < m + P$   
 $m \leq n_2 - i < m + P$ . (57)

Because these conditions are separable, the relationship between i and  $n_1$  and the relationship between i and  $n_2$  may be considered separately. The relationship between i and  $n_1$  is illustrated in Fig. 10. The two medium gray areas in the figure illustrate the constraints of the first two conditions of (57). The intersection of these constraints, shown by the dark gray region in the figure, illustrates the values of i and  $n_1$  that produce a nonzero result. An identical relationship exists for i and  $n_2$ . Consequently, the triple summation of the first term on the right-hand side of (56) reduces to

$$\begin{split} &\sum_{i=1}^{P-1} \sum_{n_{1}=i}^{F+P-2} \sum_{n_{2}=i}^{F+P-2} r^{*} \left( n_{1}T \right) r \left( n_{2}T \right) \\ &\times r^{*} \left( \left( n_{2}-i \right) T \right) r \left( \left( n_{1}-i \right) T \right) I \left( m,i,n_{1},n_{2} \right) \\ &= \sum_{i=1}^{P-1} \left[ \sum_{n_{1}=i+m}^{m+P-1} r^{*} \left( n_{1}T \right) e^{j\phi_{m}(n_{1}T)} r \left( \left( n_{1}-i \right) T \right) e^{-j\phi_{m}((n_{1}-i)T)} \right] \end{split}$$

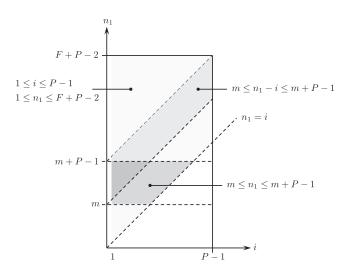


Fig. 10. Relationship of indexes involved in summations in (56) [cf. (57)].

$$\times \sum_{n_{2}=i+m}^{m+P-1} r(n_{2}T) e^{-j\phi_{m}(n_{2}T)} r^{*} ((n_{2}-i)T) e^{j\phi_{m}((n_{2}-i)T)}$$

$$= \sum_{i=1}^{P-1} \left| \sum_{n=i+m}^{m+P-1} r^{*} (nT) e^{j\phi_{m}(nT)} r((n-i)T) e^{-j\phi_{m}((n-i)T)} \right|^{2}.$$
(58)

The second term on the right-hand side of (56) represents the case in which  $i \ge P$ . Here, one of the elements of  $\widetilde{\phi}$  is guaranteed to be a sample corresponding to a random data symbol. Consequently, the integral is zero except when  $n_1 = n_2$ :

$$I(m, i, n_1, n_2) = \begin{cases} 0 & n_1 \neq n_2 \\ 1 & n_1 = n_2. \end{cases}$$
 (59)

The triple summation of the second term on the right-hand side of (56) reduces to

$$\sum_{i=P}^{F+P-2} \sum_{n_1=i}^{F+P-2} \sum_{n_2=i}^{F+P-2} r^*(n_1T) r(n_2T)$$

$$\times r^*((n_2-i)T) r((n_1-i)T) I(m,i,n_1,n_2)$$

$$= \sum_{i=P}^{F+P-2} \sum_{n=i}^{F+P-2} |r(nT)|^2 |r((n-i)T)|^2.$$
 (60)

The double summation represents all correlations of samples corresponding to random data symbols. Adding and subtracting the double summation for all correlations of samples involving only samples corresponding to preamble symbols gives

$$\sum_{i=P}^{F+P-2} \sum_{n=i}^{F+P-2} |r(nT)|^2 |r((n-i)T)|^2$$

$$= \sum_{i=1}^{F+P-2} \sum_{n=i}^{F+P-2} |r(nT)|^2 |r((n-i)T)|^2$$

$$-\sum_{i=1}^{P-1} \sum_{n=i+m}^{m+P-1} |r(nT)|^2 |r((n-i)T)|^2.$$
 (61)

Assembling these results gives

$$f(\mathbf{r}|m) \approx \frac{C(\mathbf{r})}{8\pi} + 2\pi \sum_{n=0}^{F+P-2} |r(nT)|^{2} + \frac{C(\mathbf{r})}{512\pi^{3}} \sum_{n_{1}=0}^{F+P-2} \sum_{n_{2}=0}^{F+P-2} |r(n_{1}T)|^{2} |r(n_{2}T)|^{2} + \frac{C(\mathbf{r})}{256\pi^{3}} \times \sum_{i=1}^{P-1} \left| \sum_{n=i+m}^{m+P-1} r^{*}(nT)e^{j\phi_{m}(nT)}r((n-i)T)e^{-j\phi_{m}((n-i)T)} \right|^{2} + \frac{C(\mathbf{r})}{256\pi^{3}} \sum_{i=1}^{F+P-2} \sum_{n=i}^{F+P-2} |r(nT)|^{2} |r((n-i)T)|^{2} - \frac{C(\mathbf{r})}{256\pi^{3}} \sum_{i=1}^{P-1} \sum_{n=i+m}^{m+P-1} |r(nT)|^{2} |r((n-i)T)|^{2}.$$
 (62)

Eliminating terms that do not depend on m and using the change of variables k = n - m gives

$$L'_{1}(m) = \sum_{i=1}^{P-1} \left| \sum_{k=i}^{P-1} r^{*} ((k+m)T) e^{j\phi_{m}((k+m)T)} \right|^{2}$$

$$\times r ((k+m-i)T) e^{-j\phi_{m}((k+m-i)T)} \Big|^{2}$$

$$- \sum_{i=1}^{P-1} \sum_{n=i+m}^{m+P-1} |r(nT)|^{2} |r((n-i)T)|^{2}. \quad (63)$$

The function  $L'_1(m)$  is unbalanced [9]; that is,  $L'_1(m) \neq 0$  in the noise-free case, where m is the true starting index of the preamble sequence. As explained in [9], a balanced function is obtained from  $L'_1(m)$  by removing the squares applied to the summands in (63) to produce

$$L_{1}(m) = \sum_{i=1}^{P-1} \left| \sum_{k=i}^{P-1} r^{*} \left( (k+m) T \right) e^{j\phi_{m}((k+m)T)} \right|$$

$$\times r \left( (k+m-i) T \right) e^{-j\phi_{m}((k+m-i)T)}$$

$$- \sum_{i=1}^{P-1} \sum_{n=i+m}^{m+P-1} |r(nT)| |r((n-i) T)|.$$
 (64)

Finally, using the relationship  $p(k) = e^{j\phi_m((k+m)T)}$  gives the form for  $L_1(m)$  given by (17).

APPENDIX B. DERIVATION OF THE ML FRAME SYNCHRONIZER FOLLOWING PEDONE ET AL. [11]

The outline and corresponding assumptions for this approach are as follows:

1) Average the likelihood function in (12) with respect to the random data. This step assumes the phase samples

due to random data are independent and identically distributed random variables, each with a uniform distribution over the interval  $[-\pi, \pi)$ .

- 2) Average the result of the previous step with respect to the phase  $\theta$ . This step assumes the phase is uniformly distributed over the interval  $[-\pi, \pi)$ .
- 3) Compute the logarithm of the previous step. This produces a log-likelihood function for m and  $\omega$ .
- 4) Estimate  $\omega$  using some simplifying assumptions, and substitute the estimate into the previous step.

The independence component of Assumption 1 is clearly not true for CPM. Independence is assumed purely in the interest of tractability. However, the uniform distribution is a good assumption for the marginal distribution of each phase sample.

Expanding the squares in the exponents of the conditional probability density function in (12) and partitioning the product into three products corresponding to samples produced by random data and samples produced by the pilot symbols gives

$$f(\mathbf{r}|m,\theta,\omega,\phi_{d})$$

$$= C(\mathbf{r}) \prod_{n=0}^{m-1} \exp\left\{\frac{1}{\sigma^{2}} \operatorname{Re}\left[r(nT)e^{-j\theta}e^{-j\omega n}e^{-j\phi_{m}(nT)}\right]\right\}$$

$$\times \prod_{n=m}^{m+P-1} \exp\left\{\frac{1}{\sigma^{2}} \operatorname{Re}\left[r(nT)e^{-j\theta}e^{-j\omega n}e^{-j\phi_{m}(nT)}\right]\right\}$$

$$\times \prod_{n=m+P}^{F+P-2} \exp\left\{\frac{1}{\sigma^{2}} \operatorname{Re}\left[r(nT)e^{-j\theta}e^{-j\omega n}e^{-j\phi_{m}(nT)}\right]\right\}$$
(65)

where  $C(\mathbf{r})$  is given by (41).

The first step is to average with respect to the data:

$$f(\mathbf{r}|m,\theta,\omega) = \int \cdots \int f(\mathbf{r}|m,\theta,\omega,\phi_d) f(\phi_d) d\phi_d$$
(66)

Because the phase samples due to random data are assumed to be independent uniformly distributed random variables,

$$f(\phi_d) = \prod_{n=0}^{m-1} f(\phi(nT)) \prod_{n=m+P}^{F+P-2} f(\phi(nT))$$
 (67)

where

$$f(\phi(nT)) = \begin{cases} \frac{1}{2\pi} & -\pi \le \phi(nT) < \pi \\ 0 & \text{otherwise.} \end{cases}$$
 (68)

Consequently, we have

$$\begin{split} f\left(\mathbf{r}|\,m,\theta,\omega\right) \\ &= C\left(\mathbf{r}\right) \prod_{n=0}^{m-1} \frac{1}{2\pi} \int_{-\pi}^{\pi} \exp\left\{\frac{1}{\sigma^{2}} \operatorname{Re}\left[r\left(nT\right) e^{-j\theta} e^{-j\omega n} e^{-j\alpha}\right]\right\} d\alpha \\ &\times \prod_{n=m+P}^{F+P-2} \frac{1}{2\pi} \int_{-\pi}^{\pi} \exp\left\{\frac{1}{\sigma^{2}} \operatorname{Re}\left[r\left(nT\right) e^{-j\theta} e^{-j\omega n} e^{-j\alpha}\right]\right\} d\alpha \end{split}$$

$$\times \prod_{n=-\infty}^{m+P-1} \exp \left\{ \frac{1}{\sigma^2} \operatorname{Re} \left[ r \left( nT \right) e^{-j\theta} e^{-j\omega n} e^{-j\phi_m(nT)} \right] \right\}$$
 (69)

which leads to

$$f(\mathbf{r}|m,\theta,\omega)$$

$$= C(\mathbf{r}) \prod_{n=0}^{m-1} I_0 \left( \frac{1}{\sigma^2} |r(nT)| \right) \prod_{n=m+P}^{F+P-2} I_0 \left( \frac{1}{\sigma^2} |r(nT)| \right)$$

$$\times \prod_{n=m}^{m+P-1} \exp \left\{ \frac{1}{\sigma^2} \operatorname{Re} \left[ r(nT) e^{-j\theta} e^{-j\omega n} e^{-j\phi_m(nT)} \right] \right\}.$$
(70)

The next step is to average over the phase. This gives

$$f(\mathbf{r}|m,\omega) = \frac{1}{2\pi} \int_{-\pi}^{\pi} f(\mathbf{r}|m,\theta,\omega) d\theta$$

$$= C(\mathbf{r}) \prod_{n=0}^{m-1} I_0 \left( \frac{1}{\sigma^2} |r(nT)| \right) \prod_{n=m+P}^{F+P-2} I_0 \left( \frac{1}{\sigma^2} |r(nT)| \right)$$

$$\times \frac{1}{2\pi} \int_{-\pi}^{\pi} \exp \left\{ e^{j\theta} \frac{1}{\sigma^2} \sum_{n=m}^{m+P-1} \operatorname{Re} \left[ r(nT) e^{-j\omega n} e^{-j\phi_m(nT)} \right] \right\} d\theta$$
(71)

which evaluates to

$$f(\mathbf{r}|m,\omega)$$

$$= C(\mathbf{r}) \prod_{n=0}^{m-1} I_0 \left( \frac{1}{\sigma^2} |r(nT)| \right) \prod_{n=m+P}^{F+P-2} I_0 \left( \frac{1}{\sigma^2} |r(nT)| \right)$$

$$\times I_0 \left( \frac{1}{\sigma^2} \left| \sum_{n=m}^{m+P-1} r(nT) e^{-j\omega n} e^{-j\phi_m(nT)} \right| \right). \tag{72}$$

The log-likelihood function is obtained by computing the logarithm of the conditional probability density function in (72). Dropping the scale constant that is not a function of m gives

$$\Lambda''(m,\omega) = \sum_{n \in \mathcal{D}_m} \log \left\{ I_0 \left( \frac{1}{\sigma^2} |r(nT)| \right) \right\} + \log \left\{ I_0 \left( \frac{1}{\sigma^2} \left| \sum_{n=m}^{m+P-1} r(nT) e^{-j\omega n} e^{-j\phi_m(nT)} \right| \right) \right\}$$
(73)

where  $\mathcal{D}_m$  is the set of indexes of samples corresponding to random data:

$$\mathcal{D}_m = \{0, \dots, m-1\} \cup \{m+P, \dots, F+P-2\}.$$
 (74)

Rewrite the first term as

$$\sum_{n \in \mathcal{D}_{m}} \log \left\{ I_{0} \left( \frac{1}{\sigma^{2}} |r(nT)| \right) \right\}$$

$$= \sum_{n \in \mathcal{D}} \log \left\{ I_{0} \left( \frac{1}{\sigma^{2}} |r(nT)| \right) \right\}$$

$$+\sum_{n=m}^{m+P-1} \log \left\{ I_0 \left( \frac{1}{\sigma^2} |r(nT)| \right) \right\}$$

$$-\sum_{n=m}^{m+P-1} \log \left\{ I_0 \left( \frac{1}{\sigma^2} |r(nT)| \right) \right\}$$

$$=\sum_{n=0}^{F+P-2} \log \left\{ I_0 \left( \frac{1}{\sigma^2} |r(nT)| \right) \right\}$$

$$-\sum_{n=m}^{m+P-1} \log \left\{ I_0 \left( \frac{1}{\sigma^2} |r(nT)| \right) \right\}. \tag{75}$$

Recognizing the first term is not a function of *m* and may be dropped, the log-likelihood function becomes

$$\Lambda'(m,\omega) = \log \left\{ I_0 \left( \frac{1}{\sigma^2} \left| \sum_{n=m}^{m+P-1} r(nT) e^{-j\omega n} e^{-j\phi_m(nT)} \right| \right) \right\} - \sum_{n=m}^{m+P-1} \log \left\{ I_0 \left( \frac{1}{\sigma^2} |r(nT)| \right) \right\}.$$
 (76)

Following Pedone et al. [11], the frequency offset is treated as a nuisance parameter whose estimate is used to produce a log-likelihood function for m. Because the first term on the right-hand side of (76) is the only one containing  $\omega$ , the starting point is the correlation in the argument of the Bessel function:

$$\varrho = \sum_{n=m}^{m+P-1} r(nT)e^{-j\phi_m(nT)}e^{-j\omega n}$$
 (77)

$$= e^{j\omega m} \sum_{k=0}^{P-1} r((k+m)T) e^{-j\phi_m((k+m)T)} e^{-j\omega k}.$$
 (78)

Now, partition the length-P summation over n using  $P = L_{PDI} L_{coh}$ . This partitions the correlation sum into  $L_{PDI} = P/L_{coh}$  nonoverlapping segments, each of length  $L_{coh}$ :

$$\varrho = e^{j\omega m} \sum_{k'=0}^{L_{PDI}} \sum_{k=k'L_{coh}}^{(k'+1)L_{coh}-1} r((k+m)T) e^{-j\phi_m((k+m)T)} e^{-j\omega k}.$$
(79)

If the  $\omega$  and  $L_{\rm coh}$  are such that  $e^{-j\omega k}\approx e^{-j\omega k'L_{\rm coh}}$  for  $k=k'L_{\rm coh},\ldots,(k'+1)L_{\rm coh}-1$ , then  $\varrho$  may be rewritten

$$\varrho = e^{j\omega m} \sum_{k'=0}^{L_{PDI}} e^{-j\omega k' L_{coh}}$$

$$\times \sum_{k=k'L_{coh}} r\left((k+m)T\right) e^{-j\phi_m((k+m)T)}$$

$$(80)$$

from which we have

$$|\varrho|^{2} = \sum_{k_{1}=0}^{L_{\text{PDI}}-1} \sum_{k_{2}=0}^{L_{\text{PDI}}-1} e^{-j\omega k_{1}L_{\text{coh}}} e^{j\omega k_{2}L_{\text{coh}}} \rho\left(m, k_{1}\right) \rho^{*}\left(m, k_{2}\right). \tag{81}$$

Rearranging the double summation along the diagonal lines  $k_2 = k_1 + n$  organizes the correlations in terms of delay n. The result is

$$|\varrho|^{2} = \sum_{n=-L_{\text{PDI}}+1}^{-1} e^{-j\omega n L_{\text{coh}}} \sum_{i=0}^{L_{\text{PDI}}-1-n} \rho(m, i+n) \rho^{*}(m, i)$$

$$+ \sum_{i=0}^{L_{\text{PDI}}-1} |\rho(m, i)|^{2}$$

$$+ \sum_{n=1}^{L_{\text{PDI}}-1} e^{j\omega n L_{\text{coh}}} \sum_{i=0}^{L_{\text{PDI}}-1-n} \rho(m, i) \rho^{*}(m, i+n)$$
(82)

$$= \sum_{i=0}^{L_{\text{PDI}}-1} |\rho(m, i)|^{2} + \sum_{n=1}^{L_{\text{PDI}}-1} 2\text{Re} \left\{ e^{-j\omega n L_{\text{coh}}} \sum_{i=0}^{L_{\text{PDI}}-1-n} \rho^{*}(m, i) \rho(m, i+n) \right\}$$
(83)

$$= \Lambda_0(m) + \sum_{n=1}^{L_{\text{PDI}}-1} \Lambda_n(m, \omega).$$
 (84)

Here,  $\Lambda_n(m, \omega)$  is interpreted as an *n*-delay correlation of  $\rho(m, i)$ , followed by a phase rotation proportional to  $\omega$ . When n = 0, the phase rotation term is not present and  $\Lambda_0(m)$  corresponds to NCPDI [11].

Maximizing the second term with respect to  $\omega$  is intractably difficult. Instead, Pedone et al. [11] develop the notion of maximizing each summand with respect to  $\omega$ . This produces a sequence of frequency offset estimates:

$$\hat{\omega}_{n} = \frac{\arg\left\{\sum_{i=0}^{L_{\text{PDI}}-1-n} \rho\left(m, i+n\right) \rho^{*}\left(m, i\right)\right\}}{nL_{\text{coh}}}.$$
 (85)

These estimates are used to remove the dependency on  $\omega$  from (76). Using the relationship

$$e^{j\hat{\omega}_{n}nL_{\text{coh}}} = \frac{\sum_{i=0}^{L_{\text{PDI}}-1-n} \rho(m, i+n) \rho^{*}(m, i)}{\left|\sum_{i=0}^{L_{\text{PDI}}-1-n} \rho(m, i+n) \rho^{*}(m, i)\right|}$$
(86)

we have

$$\bar{\Lambda}_n(m) = \Lambda_n(m, \hat{\omega}_n) \tag{87}$$

$$= 2\operatorname{Re}\left\{e^{-j\hat{\omega}nL_{\operatorname{coh}}} \sum_{i=0}^{L_{\operatorname{PDI}}-1-n} \rho(m,i+n) \rho^{*}(m,i)\right\}$$
(88)

$$= 2\operatorname{Re} \left\{ \frac{\left(\sum_{i=0}^{L_{\text{PDI}}-1-n} \rho(m, i+n) \rho^{*}(m, i)\right)^{*}}{\left|\sum_{i=0}^{L_{\text{PDI}}-1-n} \rho(m, i+n) \rho^{*}(m, i)\right|} \right.$$

$$\times \sum_{i=0}^{L_{PDI}-1-n} \rho(m, i+n) \rho^{*}(m, i)$$
 (89)

$$= 2 \left| \sum_{i=0}^{L_{PDI}-1-n} \rho(m, i+n) \rho^*(m, i) \right|.$$
 (90)

 $\bar{\Lambda}_n(m)$  is called *n*-span DPDI [11].

The correlation  $\varrho$  may be expressed as

$$\varrho = \sqrt{\Lambda_0(m) + \sum_{n=1}^{L_{\text{PDI}} - 1} \bar{\Lambda}_n(m)}.$$
 (91)

Inserting these results into the log-likelihood function in (76) gives

$$\Lambda(m) = \log \left\{ I_0 \left( \frac{1}{\sigma^2} \sqrt{\Lambda_0(m) + \sum_{n=1}^{L_{PDI}-1} \bar{\Lambda}_n(m)} \right) \right\} \\
- \sum_{n=m}^{m+P-1} \log \left\{ I_0 \left( \frac{1}{\sigma^2} |r(nT)| \right) \right\}.$$
(92)

The high SNR approximation based on  $I_0(z) \approx e^z$  is

$$L_{5}(m) \approx \sqrt{\Lambda_{0}(m) + \sum_{n=1}^{L_{PDI}-1} \bar{\Lambda}_{n}(m) - \sum_{n=m}^{m+P-1} |r(nT)|}$$
 (93)

given by (21) in the main body.

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