### MAE 5930 - Optimization Fall 2019

### Homework 7 Jared Hansen

Due: 11:59 PM, Tuesday November 26, 2019

A-number: A01439768

 $e\hbox{-}mail\hbox{:} \verb"jrdhansen@gmail.com"$ 

Purpose: the problems assigned help develop your ability to

- solve optimization problems using optimality conditions.
- formulate and solve larger problems in MATLAB.

NOTE: please write or type your formulations clearly so that a reader can understand what you are doing. You are welcome to use the equivalent functions in Python.

```
minimize x_1^2 + x_2^2
subject to x_1 + x_2 - 2 \le 0
x_1^2 - x_2 - 4 \le 0
```

```
%== PROBLEM 1
% Initial guess (we know x^* = [0,0], so we pick a close point.)
x0 = [0.5; 0.3];
% Solve the problem using fmincon.
[x,f] = fmincon(@obj,x0,[],[],[],[],[],[],@con)
% Problem 1 Objective function
function J = obj(x)
x1 = x(1);
x2 = x(2);
J = (x1)^2 + (x2)^2;
% Problem 1 Constraint functions
[cin, ceq] = con(x)
x1 = x(1);
x2 = x(2);
cin = x1 + x2 - 2;
cin = (x1)^2 - x2 - 4;
ceq = [];
end
```

Figure 1: Here is my MATLAB code that uses fmincon to solve the problem.

```
<stopping criteria details>
x =
    1.0e-06 *
    0.0247
    0.2348

f =
    5.5761e-14
```

Figure 2: Output for the code gets very very close to the correct answer of  $x^* = [x_1 = 0, x_2 = 0]^T$ . As the initial guess  $x_0$  was made closer and closer to (0,0) the fmincon function got closer and closer to (0,0). I believe that the output is giving the answer  $(x_1 = 0.0247 \times 10.0^{-6}, x_2 = 0.2348 \times 10.0^{-6}) \approx (x_1 = 0, x_2 = 0)$ 

- First, let's do some thinking that will help us see if an opten hort, #1 (1) answer exists & what it might be: we know  $\chi_1^2 \ge 0$  and  $\chi_2^2 \ge 0 \implies$  the lowest  $\chi_1^2 + \chi_2^2$  can be (unconstrained) is 0+0=0, achieved by  $(\chi_1=0,\chi_2=0.)$
- Is  $(x_1=0, x_2=0)$  feasible? 0+0-2=0 and  $0^2-0-4=0$ , so yes, it is.
- · Therefore, our process should lead us to (X\*=0 and X2 =0).
- Also, note that this is a convex program.  $\nabla_x f = \begin{bmatrix} 2x_1 \\ 2x_2 \end{bmatrix} \rightarrow \nabla_x^2 f = \begin{bmatrix} 2 & 0 \\ 0 & 2 \end{bmatrix}$  which is PD  $\Rightarrow$  objective is convex. We can see  $g_1(x)$  is linear, and  $\nabla_x^2 g_2 = \begin{bmatrix} 2 & 0 \\ 0 & 0 \end{bmatrix}$  which is PSD  $\Rightarrow$   $g_2(x)$  is convex. h(x) = 0 is trivially affine.
- · Since this is a convex program, if we find a solution condidate such that the optimality conditions system is sockisfied, these are also sufficient (for CP's) so it is guaranteed to be a global minimum.
- $L(x) = \lambda_0 x_1^2 + \lambda_0 x_2^2 + \lambda_1 (x_1 + x_2 2) + \lambda_2 (x_1^2 x_2 4)$
- $\begin{array}{c}
  \left( \sqrt{\chi} \sum_{i=0}^{\text{set}} 0 \right) \rightarrow \frac{\partial L}{\partial x_{1}} = 2 \lambda_{0} \chi_{1} + \lambda_{1} + 2 \lambda_{2} \chi_{1} \stackrel{\text{set}}{=} 0 \qquad \left[ eq. 1 \right] \\
  \frac{\partial L}{\partial x_{2}} = 2 \lambda_{0} \chi_{2} + \lambda_{1} \lambda_{2} \stackrel{\text{set}}{=} 0 \qquad \left[ eq. 2 \right]
  \end{array}$
- · Complimentarity:  $\lambda_1(x_1+x_2-2)=0$  and  $\lambda_2(x_1^2-x_2-4)=0$
- CASE 1:  $\lambda_0 = 0$ 
  - [eq1] becomes  $0 + \lambda_1 + 2\lambda_2 \times_1 = 0 \implies \lambda_1 = -2\lambda_2 \times_1$
  - [eq2] becomes  $0 + \lambda_1 \lambda_2 = 0 \implies \lambda_1 = \lambda_2$
  - · Thus Feg1 now becomes  $\lambda_1 = -2\lambda_1 \times_1 \longrightarrow (x_1 = -1/2)$
  - \* Now solve for  $\chi_2$ :  $\lambda_1 \left( -\frac{1}{2} + \chi_2 2 \right) = 0$  and  $\lambda_1 \left( \left( -\frac{1}{2} \right)^2 \chi_2 4 \right) = 0$   $\Rightarrow -\frac{5}{2} + \chi_2 = -\chi_2 \frac{15}{4} \Rightarrow \chi_2 = -\frac{5}{8}$
  - · To soutisfy compl. cond's, since (x1+x2-2) +0 = (1=0) and (x12-x2-4) +0 = (1=0)
  - This can't work since it requires  $\lambda_0 = \lambda_1 = \lambda_2 = 0$  which violates the non-triviality condition. Therefore, if  $\exists$  a solution,  $\lambda_0 = 1$ .

• CASE 2:  $\lambda_0 = 1$  with  $\lambda_1 = 0$ 

- opt hw7, #1 (2)
- [eq1] becomes  $2x_1 + 2\lambda_2 x_1 = 0 \rightarrow 2x_1(1+\lambda_2)=0 \Rightarrow x_1=0$  or  $\lambda_2=-1$ , but since  $\lambda \ge 0 \Rightarrow \lambda_2 \ne -1 \Rightarrow x_1=0$  here.
- $eq^2$  becomes  $2x_1-\lambda_2=0 \Rightarrow \lambda_2=2x_2$  or  $x_2=\frac{\lambda_2}{2}$
- . To satisfy complim. cords.  $\lambda_2(\chi_1^2-\chi_2-4)=0 \rightarrow \text{sub in } \lambda_2=2\chi_2 \text{ and } \chi_1=0$ , giving  $2\chi_2(-\chi_2-4)=0 \Rightarrow \text{either } (\chi_2=0) \text{ or } \chi_2=-4$ . If  $\chi_2=0=\lambda_2=0$ , if  $\chi_2=-4 \Rightarrow \lambda_2=-8$  which  $13n^4$  allowed due to  $\lambda=0$  constraint.
- . Therefore we have a candidate minimizer of:

 $\chi_1 = 0$ ,  $\chi_2 = 0$ ,  $\lambda_0 = 1$ ,  $\lambda_1 = 0$ ,  $\lambda_2 = 0$ 

Since we showed above that this is a convex program, we know that a cardidate that satisfies the optimality condition system is a global minimizer. Also, we showed that our intuition is correct, verifying that  $[x_i=0, x_s=0]$  with  $\lambda_0=1$  and  $\lambda=[\lambda_i=0, \lambda_2=0]$  satisfies the opt. conditions and a chieves the unconstrained optimal objective value of 0.

Since (0,0) is a unique point & uniquely achieves f(x)=0 & satisfies the conditions, we can say that  $\begin{bmatrix} x_1^* = 0 \end{bmatrix}$  globally infinitizes

 $f(x) = \chi_1^2 + \chi_2^2$  for the constraints  $g_1(x) \leq 0$  and  $g_2(x) \leq 0$ .

```
\begin{array}{ll} \text{minimize} & x_1^2 + x_2^2 \\ \text{subject to} & x_1 - 10 \leq 0 \\ & x_1 - x_2^2 - 4 \geq 0 \end{array}
```

```
%== PROBLEM 2
% Initial guess (we know x^* = [4,0], so we pick a close point.)
x0 = [4.2; 0.3];
% Solve the problem using fmincon.
[x,f] = fmincon(@obj,x0,[],[],[],[],[],[],@con)
% Problem 2 Objective function
function J = obj(x)
x1 = x(1);
x2 = x(2);
J = (x1)^2 + (x2)^2;
% Problem 2 Constraint functions
function [cin, ceq] = con(x)
x1 = x(1);
x2 = x(2);
cin = x1 - 10;
cin = -x1 + (x2)^2 +4;
ceq = [];
end
```

Figure 3: Here is my MATLAB code that uses fmincon to solve the problem.

```
<stopping criteria details>
x =
     4.0000
     0.0000

f =
     16.0000
```

Figure 4: Output for the code achieves exactly the correct answer of  $x^* = [x_1 = 4, x_2 = 0]^T$ . As with problem 1, if  $x_0$  isn't close enough to  $x^*$  fmincon won't quite converge to (4,0), but a guess sufficiently close to (4,0) will converge almost exactly.

- optzn hw7,#2 (1) · First let's do some analysis of the problem. Is it convex? Yes, it is. = We know from problem #1 that f(x) = X1+ x2 is convex.
  - =  $g_1(x) = x_1 10$  is linear and thus convex,  $g_2(x) = -x_1 + x_2^2 + H$  has  $\sqrt{g_2(x)} = \begin{bmatrix} 0 & 0 \\ 0 & 2 \end{bmatrix}$  which is PSD making 92(X) convex also. So the Thequality constraints are convex.
  - = hrx = 0 is trivilly affine.
  - = Thus, if I a cardidate that satisfies optimality conditions it is a guaranteed global min.
- $L(x) = \lambda_0 x_1^2 + \lambda_0 x_2^2 + \lambda_1 (x_1 10) + \lambda_2 (-x_1 + x_2^2 + 4)$
- · 2L/2x, = 2 / 0 × 1 + / / 2 = 0 [ eq 1]  $\partial L/\partial x_2 = 2 \lambda_0 x_2 + 2 \lambda_2 x_2 \stackrel{\text{set}}{=} 0$  [eq2]
- · Complimentarity:  $\lambda_1(x_1-10)=0$  and  $\lambda_2(-x_1+x_2^2+4)=0$
- · CASE1: \ > = 0
- eq1 becomes  $\lambda_1 = \lambda_2$ , eq2 becomes  $2\lambda_2 \times_2 = 0 \implies$  either  $\lambda_2 = 0$  or  $\times_2 = 0$ . We can't let  $\lambda_2 = 0$  since that makes  $\lambda_1 = 0$  which gives  $\lambda_0 = \lambda_1 = \lambda_2 = 0$  which is trivial. Thus let  $x_2 = 0$ .
- With  $x_2=0$  let's book at complimentarity:  $\lambda_1(x_1-\iota_0)=0$  and  $\lambda_2(-x_1+4)=0$   $\Longrightarrow$  $\lambda_1(x_1-10)=0$  and  $\lambda_2(x_1+4)=0$ . Since  $\lambda_1$  and  $\lambda_2$  cannot =0 this means that  $x_1-10=0 \Rightarrow x_1=10$  and  $-x_1+4=0 \Rightarrow x_1=4$ . Since  $x_1$  cannot = 4 and 10 we know that  $\lambda_0 = 0$  won't work, and any feasible candidates must be normal ( $\lambda_0 = 1$ ).
- · CASE2: ho=1
- eq1 becomes:  $2x_1 + \lambda_1 \lambda_2 = 0$ , eq2 becomes:  $2x_2 + 2\lambda_2 x_2 = 0 \rightarrow 2x_2(1 + \lambda_2) = 0$  means either  $x_2=0$  or  $\lambda_2=-1$ , but  $\lambda=0$  is required, so we must let  $x_2=0$ .
- . Now the complimentarity conditions:  $\lambda_2(-x_1+0^2+4)=0$  and  $\lambda_1(x_1-10)=0$ These conditions will give us 4 sub-cases to check when  $\lambda_0 = 1$ : (A)  $x_1 = 4$ , (B)  $x_1 = 10$ , (C)  $\lambda_1 = 0$ , (D)  $\lambda_2 = 0$

- CASE 2A:  $(\lambda_0=1)$ ,  $X_1=4$
- · If x=4 > x + 10 > (from complimentanty)
- · Using eg1: 2x,+ \(\lambda\_1 \lambda\_2 = 0 \rightarrow 2(4) + 0 = \lambda\_2 8)

• This candidate,  $(x_1 = 4, x_2 = 0, \lambda_0 = 1, \lambda_1 = 0, \lambda_2 = 8)$  satisfies all conditions. Since the problem is convex, we know [x; = 4] is a global optimizer of f in this case. We'll check the remaining 3 cases to see if I any other optima that achieve f([4,0]) = 16 optimum given the constraints.

- · CASE 2B: (\(\lambda\_0=1\), \(\text{X}\_1=10\)
- If  $x_1=10 \Rightarrow x_1 \neq 4 \Rightarrow (\lambda_2=0)$  (from complomentarity). Using eq1:  $2x_1 + \lambda_1 = 0 \Rightarrow 20 = -\lambda_1$  $\Rightarrow \lambda_1 = -20$  which violates  $\lambda \ge 0 \Rightarrow X_1 \ne 10$ .
- · CASE 2C: (\lambda\_0=1), \lambda\_1=0
- If  $\lambda_1=0$ , eq1 says  $2x_1=\lambda_2 \rightarrow \text{plugged Tho compl. conds gives } 2x_1(-x_1+4) \Rightarrow x_1 is either$ 0 or 4. We've already Cooked at X1=4 (see CASE 2A), so let X1=0.
- If  $x_1=0$ , in the other compl. cond we have  $\lambda_2$  (0+4)=0  $\Rightarrow$   $\lambda_2=0$ . But is  $x_1=x_2=0$ feasible? No. Consider  $g_2(x) = x_1 - x_2^2 - 4 \ge 0 \Rightarrow g_2([0,0]) = -4 \not\ge 0 \Rightarrow \lambda_1 = 0$  isn't feasible here.
- CASE 2D:  $(\lambda_0=1)$ ,  $\lambda_2=0$
- · If  $\lambda_2 = 0$ , eq1 gives  $2x + \lambda_1 = 0$  ⇒  $\lambda_1 = -2x_1$ . Plugging This compl. and 1 gives:  $-2x_1(X_1 10) = 0$ . From CASE 2c we know XI=X2=0 13 n't feasible, so let XI=10 to sootisty this.
- As above in CASE 2B, if  $x_i=10$  we have  $\lambda_i=-2(10)=-20$ , and  $\lambda_i=-20$  violates  $\lambda \geq 0$ .
  - · FINAL: based on all of these cases, and the fact that our problem is convex, we can guarantee that the point  $[(x_1^* = 4, x_2^* = 0)]$  is a global optimizer of  $f(x) = x_1^2 + x_2^2$  for given constraints. This  $x^*$  gives  $f(x^*) = 16$  as the global aptimum, and soutisfies optimality conditions with  $\lambda_0 = 1, \quad \lambda^T = [\lambda_1 = 0, \lambda_2 = 8].$

```
\begin{array}{ll} \text{minimize} & x_1^2 + x_2^2 \\ \text{subject to} & 4 - x_1 - x_2^2 \leq 0 \\ & 3x_2 - x_1 \leq 0 \\ & -3x_2 - x_1 \leq 0 \end{array}
```

```
%== PROBLEM 3
% Initial guess (we know x^* = [3,1], so we pick a close point.)
% NOTE: not even giving the correct answer could get fmincon to give the
       right answer.
x0 = [3.0; 1.0];
% Solve the problem using fmincon.
[x,f] = fmincon(@obj,x0,[],[],[],[],[],[],@con)
% Problem 3 Objective function
function J = obj(x)
x1 = x(1);
x2 = x(2);
J = (x1)^2 + (x2)^2;
end
% Problem 3 Constraint functions
function [cin, ceq] = con(x)
x1 = x(1);
x2 = x(2);
cin = 4 -x1 - (x2)^2 ;
cin = 3*x2 -x1 ;
cin = -3*x2 -x1;
ceq = [];
end
```

Figure 5: Here is my MATLAB code that uses fmincon to solve the problem.

## <stopping criteria details> x = 1.0e-03 \* 0.1573 0.4719 f = 2.4748e-07

Figure 6: We can see from the output that fmincon falls flat on this problem. Even when we make  $x_0$  the analytical solution the function can't converge to  $x^* = (3,1)$ , but instead returns a point that is actually infeasible. The same held true for initial guesses  $x_0$  further from  $x^*$ . Here fmincon is returning  $(x_1 = 0.1573 \times 10.0^{-3}, x_2 = 0.4719 \times 10.0^{-3}) \approx (x_1 = 0, x_2 = 0)$  which is obviously not close to  $x^* = (3,1)$ .

- · First, is this problem a convex program like #1 and #2 were? No. Look at  $g_1(x) \leq 0 \rightarrow 4-x_1-x_2^2 \leq 0$ . Is  $g_1(x)$  convex? No.  $\nabla_x g_1 = \begin{bmatrix} -1 \\ -2x_2 \end{bmatrix}$  and  $\nabla_{x}^{2}g_{1} = \begin{bmatrix} 0 & 0 \\ 0 & -2 \end{bmatrix} \Rightarrow \nabla_{x}^{2}g_{1}$  is not PSD (since ends  $\{0, -2\} \neq 0$ ).
- · Since not all Thequality constraints are convex, we won't have the strong quarantee of any optimality-condition-soutistying points being global minimizers.
- $L(x) = \lambda_0 x_1^2 + \lambda_0 x_2^2 + \lambda_1 (4 x_1 x_2^2) + \lambda_2 (-x_1 + 3x_2) + \lambda_3 (-x_1 3x_2)$

$$\partial L/\partial x_1 = 2\lambda_0 x_1 - \lambda_1 - \lambda_2 - \lambda_3 \stackrel{\text{set}}{=} 0$$
 [eq1]

$$\partial L/\partial x_2 = 2\lambda_0 x_2 - 2\lambda_1 x_2 + 3\lambda_2 - 3\lambda_3 \stackrel{\text{set}}{=} 0$$
 [eq2]

- Complimentarity:  $\lambda_1 (4-x_1-x_2^2) = 0$  [comp<sup>1</sup>],  $\lambda_2 (-x_1+3x_2)$  [comp<sup>2</sup>],  $\lambda_3 (-x_1-3x_2)$  [comp<sup>3</sup>]
- CASE 1:  $\lambda_0 = 0$ :
- · This yields eqt: -\1-\2-\3=0 → \1+\2+\3=0. Since \20 \Xi >> \1+\2+\3=0 only holds if  $\lambda_1 = \lambda_2 = \lambda_3 = 0$ . But we can't have this, since it gives  $\lambda_0 = 0$ ,  $\lambda = 0$  which violates the non-trivality condition.
- · Therefore, if  $\exists$  a solution it must be "normal", i.e.  $\lambda_0 = 1$ .
- When  $\lambda_0=1$ : we have eqt:  $\chi_1=\frac{\lambda_1+\lambda_2+\lambda_3}{2}$ , eq2:  $\chi_2=\frac{3}{2}\left(\frac{-\lambda_2+\lambda_3}{1-\lambda_1}\right)$ 
  - · We will have numerous cases to check, generated by compil conditions. We can Start by checking  $\lambda_1=0$ ,  $\lambda_2=0$ ,  $\lambda_3=0$ , also checking when the other terms in the comp i) conditions are = 0, allowing  $\lambda_1$ ,  $\lambda_2$ ,  $\lambda_3$  to be some value > 0.
- · Since ho must = 1, we now must check all combos of  $\lambda_1, \lambda_2, \lambda_3$  being = 0.
- · CASE 2: \(\lambda\_0 = 1, \text{ while } \lambda\_1 = \lambda\_2 = \lambda\_3 = 0
  - · Learnes eg1:  $x_1 = (0+0+0)/2 \longrightarrow x_1 = 0$ , eq2:  $x_2 = \frac{3}{2} \left( \frac{0}{1} \right) = 0$ , but is  $(x_1 = x_2 = 0)$  feasible? No: see  $g_1(x) \neq 0 \rightarrow g_1([0,0]) = 4-0-0=4 \neq 0 \Rightarrow 2 \neq 2$

```
• CASE 3: \lambda_0 = 1, while \lambda_2 = \lambda_3 = 0 and (4 - x_1 - x_2^2) = 0
```

• Eqt is now: 
$$2x_1 - \lambda_1 - 0 - 0 = 0 \Rightarrow x_1 = \frac{\lambda_2}{2}$$

· Eq2 TS now: 
$$2x_2 - 2\lambda_1 x_2 + 0 - 0 = 0 \Rightarrow x_2(2 - 2\lambda_1) = 0 \Rightarrow x_2 = 0$$
 or  $\lambda_1 = 1$ 

• If 
$$x_2=0$$
, [comp1] is  $(4-x_1-0^2)=0 \Rightarrow x_1=4 \Rightarrow \lambda_1=8$ . Is  $(x_1=4, x_2=0)$  feasible?

(i) 
$$g_1(x) = 4 - 4 - 0 = 0 \le 0$$

(ii) 
$$g_2(x) = 3(0) - 4 = -4 = 0$$

(iii) 
$$g_3(x) = -3(0) - 4 = -4 = 0$$

So 
$$(x_1 = 4, x_2 = 0)$$
 with  $\lambda_0 = 1, \lambda_1 = [\lambda_1 = 8, \lambda_2 = \lambda_3 = 0]$ 

is feasible & satisfies optimality conditions.

(Not necessarily a local or global min though.)

. What if 
$$\lambda_1=1$$
? Eq.1 is:  $2x_1-1=D \rightarrow x_1=\frac{1}{2}$ . Then, to satisfy comp. I we need  $(4-x_1-x_2^2)=0 \rightarrow (\frac{8}{2}-\frac{1}{2}=x_2^2) \Rightarrow x_2=\pm \sqrt{\frac{7}{2}}$ . Is  $\left(x_1=\frac{1}{2}, x_2=\pm \sqrt{\frac{7}{2}}\right)$  feasible?

No for both points. It obviously satisfies  $g_1(x) \in O$  by construction, but we can't satisfy  $g_2(x) = O$  and  $g_3(x) \leq O$ .

$$g_2\left(\left[x_1=\frac{1}{2},x_2=\sqrt{\frac{1}{2}}\right]\right)=3\left(\sqrt{\frac{1}{2}}\right)-\frac{1}{2}=5.11\neq0 \Rightarrow \left(x_1=\frac{1}{2},x_2=\sqrt{\frac{1}{2}}\right)\text{ is n't feasible.}$$

$$93([x_1=\frac{1}{2}, x_2=-\sqrt{\frac{1}{2}})=++3([\frac{1}{2})-\frac{1}{2}=5.11 \pm 0 \Rightarrow (x_1=\frac{1}{2}, x_2=-\sqrt{\frac{1}{2}})$$
 is r'+ feasible.

· CASE 4: 
$$\lambda_0 = 1$$
, while  $\lambda_1 = \lambda_2 = 0$  and  $(-x_1 - 3x_2 = 0)$ 

• Eq1 is now: 
$$2x_1 - \lambda_3 = 0 \rightarrow x_1 = \frac{\lambda_3}{2}$$
, Eq2 is now:  $2x_2 - 3\lambda_3 \rightarrow x_2 = \frac{3}{2}\lambda_3$ 

. Sub in to 
$$\overline{\text{comp3}}$$
 cord:  $\left(-\frac{\lambda_3}{2} - 3\left(\frac{3}{2}\lambda_3\right)\right) \rightarrow -\lambda_3 = 9\lambda_3 \rightarrow \text{iff } \lambda_3 = 0 \Rightarrow \chi_1 = \chi_2 = 0$ 

· But 
$$x_1 = x_2 = 0$$
 violates  $g_1(x) = 0 \rightarrow 4 - 0 - 0^2 = 4 \neq 0 \Rightarrow$  this is infeasible.

• CASE 5: 
$$\lambda_0 = 1$$
, while  $\lambda_1 = \lambda_3 = 0$  and  $(-x_1 + 3x_2 = 0)$ 

. Eg1 is now: 
$$2x_1-\lambda_2=0$$
 and Eg2 is now:  $2x_2+3\lambda_2=0 \Rightarrow x_1=\frac{\lambda_2}{2}$ ,  $x_2=\frac{-3\lambda_2}{2}$ 

Plugging in to comp2 condition gives: 
$$\frac{\lambda_2}{2} = 3\left(-\frac{3\lambda_2}{2}\right) \rightarrow \lambda_2 = -9\lambda_2 \rightarrow \text{iff } \lambda_2 = 0$$

$$\Rightarrow \chi_1 = \chi_2 = 0 \text{, which again violates } g_1(\chi) \leq 0 \rightarrow 4-0-0^2 = 4 \neq 0.$$
Thus, this route is infeasible.

- CASE 6:  $\lambda_0 = 1$ ,  $\lambda_1 = 0$ ,  $(-x_1 + 3x_2 = 0)$ ,  $(-x_1 3x_2 = 0)$
- We can see that  $(-x_1+3x_2 \text{ must} = -x_1-3x_2) \rightarrow -x_1+x_1 = 3x_2++3x_2$   $\rightarrow (6x_2=0) \rightarrow \text{iff } x_2=0 \Rightarrow x_1=0, \text{ but we know } x_1=x_2=0 \text{ violates } g_1(x)=0,$   $5me \ 4-0-0^2=4 \neq 0. \text{ Thus, this solution can't work.}$
- CASE 7:  $\lambda_0 = 1$ ,  $\lambda_2 = 0$ ,  $(4 x_1 x_2^2 = 0)$ ,  $(-x_1 3x_2) = 0$ 
  - · Using the last equation in the CASE,  $\Rightarrow x_1 = -3x_2$ . Plugging in to the other CASE equation girts:  $(4+t 3x_2-x_2^2=0) \equiv x_2^2-3x_2-4=0 \rightarrow (x_2+1)(x_2+1)=0$  This gives  $x_2=4$  and  $x_2=-1 \rightarrow (x_1=3,x_2=1)$ ,  $(x_1=-12,x_2=4)$ . Feasible?  $(x_1=3,x_2=-1)$ : this is a feasible point. Looking at eq1:  $-2-\lambda_1-\lambda_3=0$  and eq2:  $-2+2\lambda_1-3\lambda_3=0$ . Gives  $\lambda_1=-2-\lambda_3 \rightarrow \text{into eq2}$ :  $-2-4-2\lambda_3-3\lambda_3=0 \rightarrow -5\lambda_3=6 \rightarrow \lambda_3=-6/5$  which violates  $\lambda_1 \geq 0$   $\forall i \in \{1,7,3\}$ , so this solution won't work. In short,  $(x_1=3,x_2=1)$  is feasible, but doesn't satisfy FJ conditions here.
- $(X_1 = -12, X_2 = 4)$ : this is a feasible point. However, it has a much greater objective value than the candidate from CASE 3, so it's not worth finding  $\lambda_1$  &  $\lambda_3$ .
- CASE 8:  $\lambda_0 = 1$ ,  $\lambda_3 = 0$ ,  $(4 x_1 x_2^2 = 0)$ ,  $(-x_1 + 3x_2 = 0)$
- Using the equations in the case:  $X_1 = 3X_2 \rightarrow [4-3x_2-x_2^2=0] = [X_2^2+3x_2-4=0]$   $= [(X_2+4)(X_2-1)=0] \implies \text{ either } (X_2=1, X_1=3) \text{ or } (X_2=-4, X_1=-12).$
- · (X1=-12, X2=-4): doesn't matter, much larger objective value than case 3 candidate.
- $(X_1 = 3, X_2 = 1)$ : feasible point. Our new eq1:  $(a \lambda_1 \lambda_2 = 0, \text{ and our new})$ eq2:  $2 - 2\lambda_1 + 3\lambda_2 = 0$ .  $\lambda_1 = 6 - \lambda_2 \rightarrow 2 - 2(6 - \lambda_2) + 3\lambda_2 = 0 \rightarrow 5\lambda_2 = 10 \rightarrow 2 = 2$  $\lambda_1 = 4$  Thus,  $(X_1 = 3, X_2 = 1)$  with  $\lambda_0 = 1$ ,  $\lambda_1 = [\lambda_1 = 4, \lambda_2 = 2, \lambda_3 = 0]$  is

our best candidate with  $f(x) = 3^2 + 1^2 = 10$ . We can't guarantee it to be a global max, but it is the best solution using the FJ conditions.  $(x_1=3, x_2=-1)$  also achieves f(x)=10, but diesn't satisfy FJ conditions (although it Is a feasible point).

```
\begin{array}{ll} \text{minimize} & x_1x_2 \\ \text{subject to} & x_1+x_2 \geq 2 \\ & x_2 \geq x_1 \end{array}
```

```
%== PROBLEM 4
% Initial quess...answer DNE, so doesn't really matter what we put in.
x0 = [-10; 14];
% Solve the problem using fmincon.
[x,f] = fmincon(@obj,x0,[],[],[],[],[],[],@con)
% Problem 4 Objective function
function J = obj(x)
x1 = x(1);
x2 = x(2);
J = x1 * x2 ;
end
% Problem 4 Constraint functions
function [cin, ceq] = con(x)
x1 = x(1);
x2 = x(2);
cin = -x1 - x2 + 2;
cin = x1 - x2;
ceq = [];
end
```

Figure 7: Here is my MATLAB code that uses fmincon to solve the problem.

## <stopping criteria details> x = 1.0e+12 \* -1.5362 1.4928 f = -2.2931e+24

Figure 8: As we would expect, fmincon can't find the correct answer since this problem does not have an answer (see analytical examination below; min DNE since it continually approaches  $-\infty$ ). But, we can see that the function behaves as we've predicted, making  $x_2$  very large and positive and making  $x_1$  very large and negative in order to achieve a very large, negative value of f(x). Here the function achieves  $(x_1 = -1.5362 \times 10^{12}, x_2 = 1.4928 \times 10^{12})$  and  $f(x) = -2.2931 \times 10^{24}$ . I'd bet that as we continue to make guesses of very negative  $x_1$  and very positive  $x_2$  for  $x_0$  then  $f(x) \longrightarrow -\infty$ .

- First, is the problem convex? No.  $\nabla_x f = \begin{bmatrix} x_2 \\ x_1 \end{bmatrix} \rightarrow \nabla_x^2 f = \begin{bmatrix} 0 & 1 \end{bmatrix}$  which has eigrals = -1, 1 which aren't both  $\geq 0 \Rightarrow \nabla_x^2 f$  not PSD  $\Rightarrow f$  not convex  $\Rightarrow$  this is not a convex program.
- Let's build some additional intuition before jumping in to the math. In general, we'll institute  $X, X_2$  if one of  $\{X_1, X_2\} > 0$  and the other of  $\{X_1, X_2\} < 0$ , with  $\{X_1\}$  or  $\{X_2\}$  a very large value, e.g. a positive times a regative is a regative, and the larger in magnitude the values  $X_1$  d  $X_2$  the larger in (absolute) magnitude their product, which will be regative. We know  $X_2 \ge X_1$   $\Rightarrow$   $X_2$  will be the positive and  $X_1$  will be the regative. Then we just need  $X_1 + X_2 \ge 2$   $\Rightarrow$   $|X_2| \ge |X_1| + 2$ .
  - For example: let  $x_2=12$  and  $x_1=-10$ . This satisfies the constraints and gives f(x)=(-10)(12)=-120, a prefly low (minimal) objective!
  - · But what if  $X_2 = 102$  and  $X_1 = -100$ . This also satisfies the constraints but gives f(x) = (-100)(102) = -10,200 an even lower (more min.) objective.
- We can see this trend can continue interminably. Therefore we know a min of  $f(x) = X_1 \times 2$  DNE (it continually approaches  $-\infty$  as  $X_1$  and  $X_2$  grow in magnitude.)
- Now let's use more formal math to substantiate our intuition and line of reasoning that min f(x)=X,X2 DNE.
- $L(x) = \lambda_0 \times_1 \times_2 + \lambda_1 (-x_1 x_2 + 2) + \lambda_2 (x_1 x_2)$   $\partial L/\partial x_1 = \lambda_0 \times_2 - \lambda_1 + \lambda_2 \stackrel{\text{set}}{=} 0 \stackrel{\text{eq} 1}{=}$  $\partial L/\partial x_2 = \lambda_0 \times_1 - \lambda_1 - \lambda_2 \stackrel{\text{set}}{=} 0 \stackrel{\text{eq} 2}{=}$

- · Complimentanty:  $\lambda_1(-x_1-x_2+2)=0$  (comp1), [comp2]  $\rightarrow \lambda_2(x_1-x_2)=0$
- · Case 1: \ \ > = 0
  - Eq1 becomes:  $-\lambda_1 + \lambda_2 = 0 \implies \lambda_1 = \lambda_2$ , Eq2 is now:  $-\lambda_1 = \lambda_2$ . For both to hold, we must have  $\lambda_1 = \lambda_2 = 0$ , But  $\lambda_0 = 0$  and  $\lambda^T = [\lambda_1 = 0, \lambda_2 = 0]$  violates the non-trivality condition. Thus,  $\lambda_0$  must = 1 for any potential solutions.
- Case 2:  $\lambda_0 = 1$ ,  $\lambda_1 = 0$ ,  $\lambda_2 = 0$ 
  - \* Eq.1 becomes:  $x_2 = 0$ , Eq.2 becomes  $x_1 = 0$ . Is (0,0) feasible? No, it's not. Consider  $x_1 + x_2 \ge 2 \implies 0 + 0 = 0 \not\equiv 2 \implies \text{this case is } n \not\equiv \text{ feasible}.$
- Case 3:  $\lambda_0=1$ ,  $\lambda_1=0$ ,  $(x_1-x_2=0)$ 
  - If  $\chi_1 \chi_2 = 0 \Rightarrow \chi_1 = \chi_2$ . Eq.1 becomes:  $\chi_2 + \lambda_2 = 0$ , Eq.2 becomes:  $\chi_1 \lambda_2 = 0$ .
  - Now we have  $\chi_2 = -\lambda_2$  and  $\chi_2 = \lambda_2$  which only works if  $\lambda_1 = 0$  which then makes  $\chi_2 = 0 \Rightarrow \chi_1 = 0$  and we just showed (0,0) to be infeasible In Case 2. Thus this case, case 3, is also infeasible.
- · Case 4:  $\lambda_0 = 1$ ,  $\lambda_2 = 0$ ,  $(-x_1 x_2 + 2 = 0)$ 
  - Eq1 becomes:  $\chi_2 \lambda_1 = 0$ , Eq2 becomes:  $\chi_1 \lambda_1 = 0 \implies \chi_1 = \chi_2 = \lambda_1$
  - · Remitting  $(x_1+x_2=2) \rightarrow x_1+x_1=2 \rightarrow x_1=1 \Rightarrow x_2=1 \Rightarrow \lambda_1=1$ .
  - The candidate  $x = [x_1 = 1, x_2 = 1]$  is feasible and satisfies optimality conditions with  $\lambda_0 = 1$ ,  $\lambda_1 = [\lambda_1 = 1, \lambda_2 = 0]$ . But we've already shown the pt (-100, 102) has a much more minimal objective value, thus (1,1) is not a minimizer.
- Case 5:  $\lambda_0 = 1$ ,  $(x_1 x_2 = 0)$ ,  $(-x_1 x_2 + 2 = 0)$ 
  - $X_1-X_2=0 \Rightarrow X_1=X_2 \rightarrow \text{into other egn: } -X_1-X_1=-2 \Rightarrow X_1=1$ ,  $X_2=1$ , same as case 4.
- thus,  $\not\equiv$  a minimizer of  $f(x)=x_1x_2$  by using FJ optimality conditions, and our exploration at the beginning is upheld:  $\not\equiv$  a minimizer for this problem.

```
minimize -x_1

subject to -x_1 \le 0

-x_2 \le 0

x_2 + (x_1 - 1)^3 \le 0
```

```
%== PROBLEM 5
% Initial guess. We know x^* = [1,0] so we'll pick x^0 close to that.
% NOTE: even giving an initial quess that is the answer can't get
        fmincon to converge to the correct answer.
        However, I accidentally made q3(x) equality instead of ineq
       and it did give the correct answer.
x0 = [1.0; 0.0];
% Solve the problem using fmincon.
[x,f] = fmincon(@obj,x0,[],[],[],[],[],[],@con)
% Problem 5 Objective function
function J = obj(x)
x1 = x(1);
x2 = x(2);
J = -x1 + 0*x2;
end
% Problem 5 Constraint functions
function [cin, ceq] = con(x)
x1 = x(1);
x2 = x(2);
cin = -x1;
cin = -x2;
cin = x2 + (x1 - 1)^3;
ceq = [] ;
end
```

Figure 9: Here is my MATLAB code that uses fmincon to solve the problem.

```
<stopping criteria details>
x =
    1.0e+19 *
    0.0000
    -4.4751

f =
    -3.5503e+06
```

Figure 10: We can see that fmincon gets this problem very wrong even though we gave it the exactly-correct answer as  $x_0$ . It is giving an infeasible solution since  $-x_2 \nleq 0$ . Also, it is very far from the correct answer of  $x^* = (1,0)$  instead getting  $x = (x_1 = 0.0, x_2 = -4.48 \times 10^{19})$ .

- First, is the problem convex? No.
- $\nabla_x g_3 = \begin{bmatrix} 3(x_1-1)^2 \end{bmatrix} \rightarrow \nabla_x^2 g_3 = \begin{bmatrix} 6(x_1-1) & 0 \end{bmatrix}$  Is this a PSD matrix? No. Let  $x_1=0$ ,  $y_1=0$  of  $y_2=0$  of  $y_3=0$  if  $y_2=0$  of  $y_3=0$  is not convex  $y_3=0$  of  $y_3=0$  of
- Since this isn't a CP we can't make guarantees on candidate points being global optima or not.
- · Let's solve the problem without using the FJ conditions to make sure we get the right answer when we do so:
- · In general, to minimize -x, we must make x, a positive value as large as possible, relative to the given constraints.
- \* g,(x): -x,=0 This fits into the desire to make -x, as small (x, as large) as possible. But may leave x, unbounded? (Taken care of in constraint g3(x).)
- $g_2(x)$ :  $-x_2 = 0$  We know that  $x_2 \ge 0$ .
- $g_3(x)$ :  $\chi_2 + (\chi_1-1)^3 \le 0 \rightarrow (\chi_1-1)^3 \le -\chi_2$  and  $g_2(x)$  stipulates  $-\chi_2 \le 0$   $\Rightarrow (\chi_1-1)^3 \le 0$  How do we satisfy this? With some  $\chi_1 \le 1$ . But, our means of minimizing  $-\chi_1$  is to make  $\chi_1$  as large as possible. Therefore, if it must be that  $\chi_1 \le 1$  we'd choose  $\chi_1 = 1$  to minimize  $-\chi_1$ .
- So using a rough approach (that won't work in many cases, but does here) we see  $(X_1^*=1, X_2^*=0)$  which we'll verify using FJ-conditions. (Technically, we may only be showing that  $(X_1=1, X_2=0)$  is the best conditions from applying the FJ conditions. Since the program is non-convex, we can't mathematically guarantee it's the global optimum. Rother, our deductive method helps us with this.)

•  $L(x) = -\lambda_0 x_1 - \lambda_1 x_1 - \lambda_2 x_2 + \lambda_3 (x_2 + (x_{1}-1)^3)$   $\partial L/\partial x_1 = -\lambda_0 - \lambda_1 + 3\lambda_3 (x_1-1)^2 \stackrel{\text{set}}{=} 0 \stackrel{\text{[eq1]}}{=}$  $\partial L/\partial x_2 = -\lambda_2 + \lambda_3 = 0 \stackrel{\text{[eq2]}}{=} \rightarrow \lambda_3 = \lambda_2$ 

### · Complimentanty conditions:

$$[e1]: -\lambda_1 \times_{1} = 0 \longrightarrow either \lambda_1 = 0$$
 or  $\times_{1} = 0$ 

$$\frac{1}{(2)} = -\lambda_2 \times 2 = 0 \implies \text{ either } \lambda_2 = 0 \text{ or } \times_2 = 0$$

$$\begin{array}{c} (X_1 + (X_1 - 1)^3) = 0 \\ \end{array} \rightarrow \text{either } \lambda_3 = 0 \text{ or } (X_1 - 1)^3 = -X_2 \\ \end{array}$$

- Since [eg 2] requires that  $\lambda_3 = \lambda_2$  we need only check cases where  $\lambda_3 = \lambda_2$ .
- CASE 1:  $\lambda_0 = 0$ ,  $\lambda_1 \neq 0$ ,  $\lambda_2 = \lambda_3 = 0$
- \*Eq1:  $-\lambda_1 + 0 + 0 (\chi_1-1)^2 = 0 \Rightarrow \lambda_1 = 0 \Rightarrow \lambda_0 = 0$  and  $\lambda_1 = 0$  which violates the non-triviality condition. Therefore, this case is not feasible:
- CASE 2:  $\lambda_0=0$ ,  $\lambda_1=0$ ,  $\lambda_2\neq 0$ ,  $\lambda_3\neq 0$
- · Eq1: 3)3 (x1-1)2=0, 23 must be +0 => (x,-1)=0 => X1=0
- · [c2]: since  $\lambda_2 \neq 0$  here  $\Rightarrow \lambda_2 \times_2 = 0$  iff  $x_2 = 0$
- [3]: Since  $\lambda_3 \neq 0$  we have  $\lambda_3 (0 + (1-1)^3) = 0 \rightarrow \lambda_3 (0) = 0 \Rightarrow \lambda_3 \in \mathbb{R}$ and since  $\lambda_3 \in \mathbb{R}$  and  $\lambda_3 = \lambda_2 \Rightarrow \lambda_2 \in \mathbb{R}$

Thus,  $(x_1=1, x_2=0)$  with  $X=[\lambda_1=0, \lambda_2\geq 0, \lambda_3\geq 0]$  with  $\lambda_2=\lambda_3$  is a condidate according to the FJ conditions. Strictly speaking, we can't declare this to be a global ininimizer, but from our procedural reasoning above (at problems outset) we know that this is a best solution (there may exist others, we'll continue checking other cases).

- CASE 3:  $\lambda_0 = 1$ ,  $\lambda_1 = 0$  e.g.  $\lambda_1 = \lambda_2 = \lambda_3 = 0$
- · Eq1: -1-0+0=0 → -1=0 which isn't true. Therefore, having these

  > values fails to satisfy FJ coorditions ("FJ infeasible").
- CASE 4: λ₀=1, λ₁ ≠0, λ₂=λ₃=0
- Eq1:  $-1 \lambda_1 + 0 = 0 \implies \lambda_1 = -1$  which can't be, since all elements of  $\lambda_1 = -1 + 0 = 0 \implies \lambda_2 = -1$  which can't be, since all elements of  $\lambda_1 = -1 + 0 = 0 \implies \lambda_2 = -1 + 0 = 0 \implies \lambda_3 = -1 + 0 = 0 \implies \lambda_4 = -1 + 0 = 0 \implies \lambda_4 = -1 + 0 = 0 \implies \lambda_5 = -1 + 0 = 0 \implies \lambda$
- · CASE 5:  $\lambda_0 = 1$ ,  $\lambda_1 = 0$ ,  $\lambda_2 \neq 0$ ,  $\lambda_3 \neq 0$ ,  $(\lambda_2 = \lambda_3 \text{ still required by } [eq 2])$
- Eq1:  $-1 0 + 3\lambda_3 (x_1 1)^2 = 0 \implies 3\lambda_3 (x_1 1)^2 = 1$
- · c2: Since we're said 12 +0, for 12 x2=0 => X2=0
- · c3: Since we've said \( \lambda\_3 \neq 0, \lambda\_3 (\chi\_1)^2 = 0 \Rightarrow \chi\_1 1 = 0 \Rightarrow \chi\_1 = 1
- · Back to Eq1: if  $3\lambda_3(X_1-1)^2=1 \rightarrow 3\lambda_3(1-1)^2=1 \Rightarrow 3\lambda_3(0)=0=1$  which isn't possible  $\Rightarrow$  no feasible solution under FJ conditions for  $\lambda_0=1$ ,  $\lambda_1=0$ ,  $[\lambda_2=\lambda_3]\neq 0$ .

CONCLUSION: we have demonstrated the existence of one viable candidate,  $(X_1=1, X_2=0)$  with  $\lambda_0=0$  and  $\lambda^T=[\lambda_1=0, \lambda_2=0, \lambda_3=0]$  and  $\lambda_2=\lambda_3$ . Since the program 13 nt convex, we can't declare this to be a global minimizer based on the FJ conditions being satisfied. However, based on the procedural explanation at the problem's outset, we concluded the best solution 13  $(X_1=1, X_2=0)$ . (Done via some logic and combination of Constraints  $g_2(x)$  and  $g_3(x)$ .) This is substantiated by our firdings in CASE2, showing this conditate has a satisfactory formulation for the FJ conditions. Thus, based on some not-typical reasoning, we can conclude  $\chi^*=(\chi_1^*=1, \chi_2^*=0)$  achieving  $f(\chi^*)=-1$ .

```
%== PROBLEM 6
% Initial guess. We know x^* = [0,0] so we'll pick x^0 close to that.
x0 = [0.005; 0.005];
% Solve the problem using fmincon.
[x,f] = fmincon(@obj,x0,[],[],[],[],[],[],@con)
% Problem 6 Objective function
function J = obj(x)
x1 = x(1);
x2 = x(2);
J = 2*(x1^2) - (x2^2);
% Problem 6 Constraint functions
function [cin, ceq] = con(x)
x1 = x(1);
x2 = x(2);
cin = (x1^2)*x2 - (x2^3);
ceq = [];
end
```

Figure 11: Here is my MATLAB code that uses fmincon to solve the problem.

```
<stopping criteria details>
x =
    1.0e+10 *
    0.2810
    2.3552

f =
    -5.3892e+20
```

Figure 12: Just as above in problem 5, fmincon gets this problem very wrong, even when we give it almost the exactly-correct answer for  $x_0$ . It is giving an infeasible solution since  $x_1^2x_2 - x_2^3 \neq 0$  but instead for the outputted  $x_1$  and  $x_2$  we get  $x_1^2x_2 - x_2^3 = 1.8597 \times 10^{29}$ . Also, it is very far from the correct answer of  $x^* = (0,0)$  instead getting  $x = (x_1 = 0.2810 \times 10^{10}, x_2 = 2.3552 \times 10^{10})$  but getting close to the correct optimal objective value of f(x) = 0.

- First, is the problem convex? No.  $\nabla_x^2 f(x) = \begin{bmatrix} 2 & 0 \\ 0 & -1 \end{bmatrix} \Rightarrow \text{eignals} = -1, 2$ , since  $-1 \not\equiv 0 \Rightarrow \text{not}$  all positive eignals  $\Rightarrow \nabla_x^2 f(x)$  not PSD  $\Rightarrow f(x)$  objective is not convex  $\Rightarrow$  this is not a convex program.
- Let's see if we can solve the problem without using the FJ conditions.

  This will help us know if candidates are good, Correct, lend intuition.
- Intuition, to minimize  $2x_1^2 x_2^2$  we should make  $x_1^2$  as small as possible (occurs when  $x_1=0$ ) and make  $|x_2|$  as large as possible. Thus we'll have 0- (some big number).
- If  $x_1=0 \Rightarrow h(x)=x_1^2 \times_2 x_2^3=0 \rightarrow 0 \times_2 x_2^3=0 \rightarrow x_2^3=0 \rightarrow x_2=0$ . Is  $(x_1=x_2=0)$  the best we can do? (obj value = 0)
- From h(x) we see  $(\chi_1^2)(\chi_2) = (\chi_2)(\chi_2^2)$ .  $\Rightarrow \pm \chi_1 = \pm \chi_2$ . But since both  $\chi_1$  and  $\chi_2$  are squared in fix), let's just pick  $[\chi_1 = \chi_2] > 0$  for ease. Say  $\chi_1 = \chi_2 = 2$   $\Rightarrow$  obj:  $f(\chi) = 2(2^2) (2^2) = 8 4 = 4 \pm 0$  (best obj. so far).
- Intuitively, since  $\chi_1^2$  must =  $\chi_2^2$ , in the objective function the first term will always be double the second term. Since the terms are always non-negative, the best objective value is when the  $\chi_1^2 = \chi_2^2 = 0$ , which occurs when  $\chi_1 = \chi_2 = 0$ .
- Thus, our best objective value is f(x)=0 at  $[x_1=0, x_2=0]$ . We didn't arrive here through rigorous proofs, but rather through intuition, logical deduction, and some simple mouth. (Thus, can't declare global optimality.) Let's see if use of the FJ conditions substantiates this finding.

• 
$$L(x) = \lambda_0 (2x_1^2 - x_2^2) + \lambda_1 (x_1^2 x_2 - x_2^3)$$
  
 $\partial L/\partial x_1 = 4\lambda_0 x_1 + 2\lambda_1 x_1 \stackrel{\text{def}}{=} 0 \implies x_1 (2\lambda_0 + \lambda_1) = 0 : \boxed{eq1}$   
 $\partial L/\partial x_2 = -2\lambda_0 x_2 + \lambda_1 x_1^2 - 3\lambda_1 x_2^2 \stackrel{\text{def}}{=} 0 : \boxed{eq2}$ 

· Complimentanty:

$$\boxed{c1}: \lambda_1 \left( \chi_1^2 \chi_2 - \chi_2^3 \right) = 0 \implies \text{either } \lambda_1 = 0 \text{ or } \left( \chi_1^2 \chi_2 - \chi_2^3 = 0 \right)$$

· CASE 1: λ₀=0, λ, ≠0

• Eg1) becomes: 
$$\chi_1(\lambda_1)=0 \Rightarrow \chi_1=0$$
 since we specified  $\lambda_1\neq 0$  in case.

• From [C1], 
$$\lambda_1 \neq 0 \implies \chi_1^2 \chi_2 - \chi_2^3 = 0$$
 with  $\chi_1 = 0 \implies \chi_2^3 = 0 \implies \chi_2 = 0$ 

Thus, the cardidate  $[X_1=0, X_2=0]$  with  $\lambda_0=0$  and  $\lambda_1>0$  (to avoid violation of non-triviality condition) is violble per the optimality conditions. From our logiz at the beginning of the problem, we know this is the best solution we can achieve. Let's be thorough and check the remaining Cases.

· CASE 2: \(\lambda\_0 = 1, \lambda\_1 = 0\)

. [c] satisfied since 
$$\lambda_1=0$$
. Eq. 1 becomes  $\chi_1(2\lambda_0)=0 \to \chi_1(2)=0 \Longrightarrow \chi_1=0$ 

Once again, 
$$(x_1=0, x_2=0)$$
 is a viable cardidate, this time with  $\lambda_0=1$  and  $\lambda_1=0$ .

- · CASE 3: λ0=1, λ, > 0
- To satisfy [c1] we have  $x_1^2 x_2 x_2^3 = 0 \rightarrow (x_1^2)(x_2) = (x_2)(x_2^2) \Rightarrow x_1^2 = x_2^2$  $\Rightarrow$  either  $x_1 = x_2$  or  $-x_1 = x_2$  (equivalent to  $x_1 = -x_2$ ).
- · When X1=X2: [2] IS now X, (2+), )=0 = X1 must =0 = X2=0
  - Does [2] hold?  $-2x_2 + \lambda_1 x_1^2 3\lambda_1 x_2^2 = 0 \rightarrow -2x_2 2\lambda_1 x_2^2 = 0$ 
    - $\rightarrow$  -2( $\times_2$   $\lambda_1 \times_2^2$ )=0  $\rightarrow$   $\times_2$  (1- $\lambda_1 \times_2$ )=0, which works when  $\times_2$ =0, which we're explored. So let's explore 1- $\lambda_1 \times_2$ =0  $\rightarrow$   $\times_2$ = $\frac{1}{\lambda_1}$  and  $\frac{1}{\lambda_1}$ = $\frac{1}{\lambda_2}$
  - In this case: objective =  $2\left(\frac{1}{\lambda_1}\right)^2 \left(\frac{1}{\lambda_1}\right)^2 = \frac{1}{\lambda_1}$ . Since  $\lambda_1 > 0$  we can take  $\lambda_1 \to \infty$   $\lambda_1 = 0$ .

Thus, the candidate  $(x_1 = V_{\lambda_1}, x_2 = V_{\lambda_1})$ , with  $\lambda_0 = 1$  and  $\lambda_1 \to \infty$  is a viable candidate, but never quite achieves  $f(x_1 = 0)$ , just  $f(x) \to 0$  as  $\lambda_1 \to \infty$ .

- When  $-x_1=x_2$ :  $\lfloor eq2 \rfloor$  becomes  $-2x_2+\lambda_1(-x_1)^2-3\lambda_1(x_2^2)=0$   $\rightarrow -2x_2+\lambda_1x_2^2-3\lambda_1(x_2^2)=0 \rightarrow -2x_2-2\lambda_1x_2^2$ , which we run into above. Thus, there are no new  $\lambda$ , values to explore.
  - Conduston: we have explored all possible non-trivial cases, and found condidates  $(x_1=0, x_2=0)$  and  $(x_1=x_2=1/h)$  as  $h_1\to\infty$ . The best objective value from these is  $f(x_1)=0$ . Since this isn't a CP, we can't guarantee (0,0) to be a/the global optimizer. However, we know it's a viable candidate according to optimality conditions (both when  $h_0=0$ ,  $h_1>0$  and when  $h_0=1$  and  $h_1=0$ ). Also, based on our reasoning at the beginning of the problem, it is very likely a global optimizer (though rigorous proof would need to hold to guarantee this formally.)

7. Solve the traveling salesman problem using MATLAB's genetic algorithm solver ga. Use the MTZ formulation below.

$$\begin{split} & \text{minimize} & & \sum_{i,j} d_{i,j} x_{i,j} \\ & \text{subject to} & & \sum_{i} x_{i,j} = 1, \forall j \\ & & \sum_{j} x_{i,j} = 1, \forall i \\ & & u_1 = 1 \\ & & 2 \leq u_i \leq n, \forall i \neq 1 \\ & & u_i - u_j + 1 \leq (n-1)(1-x_{i,j}), \forall i \neq 1, \forall j \neq 1 \\ & & x_{i,j} \in \{0,1\} \text{ and } u_i \in \mathbb{R} \end{split}$$

Generate random x and y locations for the cities using random numbers as

- x = 20\*rand(1,n);
- y = 20\*rand(1,n);

See how large you can make n and still solve the problem. Recall that you could solve about 40 cities using the integer programming approach.

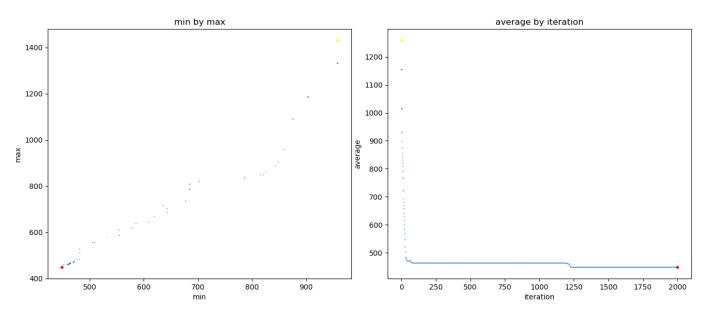


Figure 13: Here are plots (made in Python using matplotlib) showing the progression of the GA, e.g. it "learning" over iterations for 25 randomly-generated cities. This is telling us that the final solution achieved by the GA has a path length of about 450 which it achieved after roughly the  $1250^{th}$  iteration mark (see the right plot). We can see the actual cities and the plot of the final solution in the image directly below.

# Solution for 25 Randomly Created Cities. Total distance: 448.139 units 100 22 20 3 17 40 20 20 40 60 80 100

Figure 14: This plot (made in Python using matplotlib) shows the final path decided upon by the genetic algorithm for 25 cities in the traveling salesman problem. I used the DEAP (Distributed Evolutionary Algorithms in Python) package to do the solving. We can see that this isn't quite an ideal solution; at least from what I saw in the integer programming solutions, we'd expect the most efficient path to be "totally hollow", whereas this is certainly not hollow. I think if it went from 15 over to 12, then kept 12 to 24, and so on up to 10 and then had 10 connect to 4 it would be a more "ideal-looking" solution, but this certainly isn't too far off.

x coordinate of the city

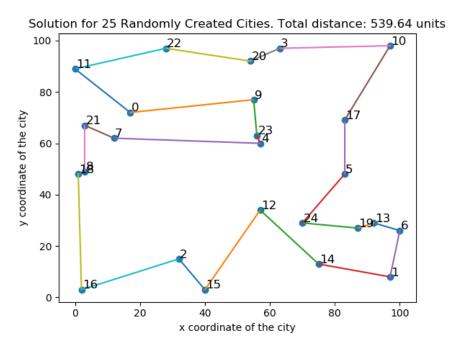


Figure 15: For comparison, I ran the script with a different seed and obtained this plot of cities and ideal route as obtained by the genetic algorithm. Here we can see that this <u>looks</u> like a more ideal route, as it is perfectly hollow and traces the outer edge to connect all the cities, <u>but it is not an/the ideal route!</u> Comparing it to the plot above: we can see that all the cities are in the same place, so it is the same map. But there is a big difference between the 448-unit-length path in Figure 14 and the 539-unit-length path in this plot. Looks can be deceiving!

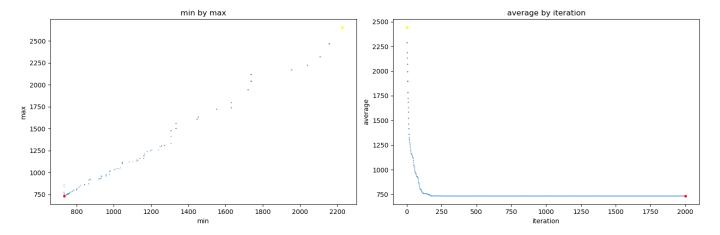


Figure 16: Here are plots (made in Python using matplotlib) showing the progression of the GA, e.g. it "learning" over iterations for 50 randomly-generated cities. This is telling us that the final solution achieved by the GA has a path length of about 750 which it achieved after only about 250 iterations (see the right plot). Surprisingly, this is almost half as many iterations as it took the 25-city version to converge; I honestly don't have a great idea as to why this is. We can see the actual cities and the plot of the final solution in the image directly below.

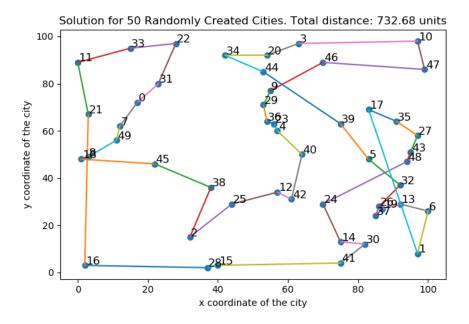


Figure 17: This plot (made in Python using matplotlib) shows the final path decided upon by the genetic algorithm for 50 cities in the traveling salesman problem. I used the DEAP (Distributed Evolutionary Algorithms in Python) package to do the solving. As with the 25-city version, we can see that this isn't quite an ideal solution; we'd expect the most efficient path to be "totally hollow", whereas this is certainly not hollow (the path crosses itself several times). All-in-all though, I wouldn't say it's a terrible solution. Certainly not as good as the integer version I got with Gurobi, but probably better than an average human could draw or come up with (especially in a matter of minutes).

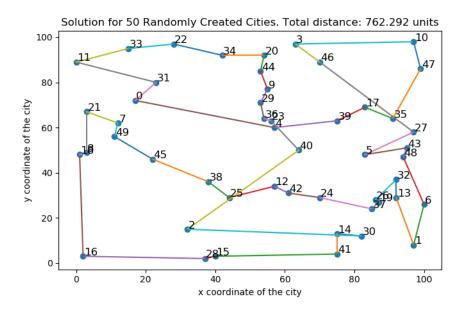


Figure 18: When we did the 25-city version we were able to set a different seed and generate a problem that seemed to yield an ideal solution with the GA solver (although it wasn't actually ideal). I was curious if we could do the same with the 50-city problem. I tried several different seeds, but never could come up with an ideal-looking solution as we did with the 25-city problem. This is an example of another one of the 50-city versions that doesn't quite look ideal when solved with the GA, but it is nice since it converges to roughly the same path length as the route in Figure 17.

Below is my code for this problem.

```
File name : hw7_prob7_TSP_ga.py
Author
           : Jared Hansen
Date created : 11/23/2019
Python version: 3.7.3
DESCRIPTION: The purpose of this script is to solve the TSP (Traveling
        Salesman Problem) using a genetic algorithm.
#==== IMPORT STATEMENTS
import copy
import datetime as dt
import math
import matplotlib.pyplot as plt
import numpy as np
import random
import pylab as pl
from matplotlib import collections as mc
from pulp import * from datetime import datetime
from deap import base, creator, tools
import string
from deap import base, creator, tools
#=== Runner CLASS: uses the DEAP toolbox to set up the GA, track performance
class Runner:
  def __init__(self, toolbox):
    self.toolbox = toolbox
    self.set_params(10, 5, 2)
  def set_params(self, pop_size, iters, num_matings):
    self.iters = iters
    self.pop_size = pop_size
    self.num_matings = num_matings
  def set_fitness(self, population):
     # individual has a shorter path length).
    fitnesses = [
       (individual, self.toolbox.calc_path_len(individual))
       for individual in population
    for individual, fitness in fitnesses:
       individual.fitness.values = (fitness,)
  def get_offspring(self, population):
```

```
n = len(population)
    for _ in range(self.num_matings):
       i1, i2 = np random choice(range(n), size=2, replace=False)
       offspring1, offspring2 = \
          self.toolbox.mate(population[i1], population[i2])
       yield self.toolbox.mutate(offspring1)[0]
       yield self_toolbox_mutate(offspring2)[0]
  @staticmethod
  def return_stats(population, iteration=1):
     fitnesses = [ individual.fitness.values[0] for individual in population ]
       ": iteration.
       'mu': np.mean(fitnesses),
       'std': np.std(fitnesses),
       'max': np.max(fitnesses),
       'min': np.min(fitnesses)
  def Run(self):
     # ("stats" list below) for each iteration, creates offspring, and so on
     population = self.toolbox.population(n=self.pop_size)
     self.set_fitness(population)
    stats = []
    for iteration in list(range(1, self.iters + 1)):
       current population = list(map(self.toolbox.clone, population))
       offspring = list(self.get_offspring(current_population))
       for child in offspring:
          current_population.append(child)
       self_set_fitness(current_population)
       population[:] = self.toolbox.select(current_population, len(population))
       stats.append(
          Runner_return_stats(population, iteration))
     return stats, population
#==== SPECIFYING AND RUNNING THE GA, PLOTTING FINAL RESULTS
creator.create("FitnessMin", base.Fitness, weights=(-1.0,))
creator.create("Individual", list, fitness=creator.FitnessMin)
random.seed(11);
np random seed(121);
INDIVIDUAL_SIZE = NUMBER_OF_CITIES = n = 25
pop_size = 200
num_iters = 1000
num_matings = 50
pop size = 400
num iters = 2000
num_matings = 100
cities = []
for i in range(n):
```

```
cities.append(str(i))
# Create coordinates for each city
random.seed(1)
points = [(random.randint(0,100), random.randint(0,100))] for i in range(n)]
cities_x = np.array(points)[:,0]
cities_y = np.array(points)[:,1]
# Plot the points that we've generated, with cities labeled by number
fig, tsp_img = plt.subplots()
tsp_img.scatter(cities_x, cities_y)
for i, cityNum in enumerate(cities):
  tsp_img.annotate(cityNum, xy=(cities_x[i], cities_y[i]), xytext=(1,1),
              textcoords='offset points',
             fontsize=9)
def calc_dist(city, to_city):
  dist = math.sqrt((points[city][0] - points[to_city][0])**2 +
              (points[city][1] - points[to_city][1])**2)
  return(dist)
# Calculate the distances between each city.
distances = np.zeros((n, n))
for city in range(n):
  for to_city in [i for i in range(n) if not i == city]:
     distances[to_city][city] = distances[city][to_city] = calc_dist(city, to_city)
toolbox = base. Toolbox()
toolbox.register("indices", random.sample, range(INDIVIDUAL_SIZE), INDIVIDUAL_SIZE)
toolbox.register("individual", tools.initlterate, creator.Individual, toolbox.indices)
toolbox register("population", tools.initRepeat, list, toolbox.individual)
# Define the calc_path_len function to calculate the length of the individual path
# e.g. how long to travel along the specified cities path.
def calc_path_len(individual):
  summation = 0
  start = individual[0]
  for i in range(1, len(individual)):
     end = individual[i]
     summation += distances[start][end]
     start = end
  return summation
toolbox.register("calc path len", calc path len)
toolbox.register("mate", tools.cxOrdered)
toolbox_register("mutate", tools_mutShuffleIndexes, indpb=0.01)
toolbox_register("select", tools_selTournament, tournsize=10)
a = Runner(toolbox)
a.set_params(pop_size, num_iters, num_matings)
stats, population = a.Run()
plt.figure(figsize=(15,5))
plt.subplot(1,2,1)
_ = plt.scatter([s['min'] for s in stats], [s['max'] for s in stats], marker='.', s=[(s['std'] + 1) / 20 for s in stats])
_ = plt.title('min by max')
_{-} = plt.xlabel('min')
_{-} = plt.ylabel('max')
_ = plt.plot(stats[0]['min'], stats[0]['max'], marker='.', color='yellow')
_ = plt.plot(stats[-1]['min'], stats[-1]['max'], marker='.', color='red')
plt.subplot(1,2,2)
_= plt.scatter([s['i'] for s in stats], [s['mu'] for s in stats], marker='.', s=[(s['std'] + 1) / 20 for s in stats])
  plt.title('average by iteration')
```

```
plt.xlabel('iteration')
    plt.ylabel('average')
_ = plt.plot(stats[0]['i'], stats[0]['mu'], marker='.', color='yellow')
_ = plt.plot(stats[-1]['i'], stats[-1]['mu'], marker='.', color='red')
plt.tight_layout()
plt.show()
# See how long the best route took.
fitnesses = sorted([
  (i, toolbox.calc_path_len(individual))
  for i, individual in enumerate(population)
], key=lambda x: x[1])
best_fit = np.round(fitnesses[:1][0][1], 3)
calc_path_len(population[0])
best_path = list(population[0])
fig, tsp_img = plt.subplots()
tsp_img.scatter(cities_x, cities_y)
for i, cityNum in enumerate(cities):
  tsp_img annotate(cityNum, xy=(cities_x[i], cities_y[i]), xytext=(1,1),
              textcoords='offset points',
              fontsize=12)
for i in range(len(best_path)):
  if(i < len(best_path)-1):</pre>
     plt.plot([cities_x[best_path[i]], cities_x[best_path[i+1]]],
           [cities_y[best_path[i]], cities_y[best_path[i+1]]])
  else:
     plt.plot([cities_x[best_path[-1]], cities_x[best_path[0]]],
           [cities_y[best_path[-1]], cities_y[best_path[0]]])
plt.title('Solution for ' + str(NUMBER_OF_CITIES) + ' Randomly Created Cities. Total distance: ' + str(best_fit) + ' units')
#plt.title('Solution for ' + str(n) + 'part-B Cities. Total distance: ' + str(total_dist) + 'units')
plt.xlabel('x coordinate of the city')
plt.ylabel('y coordinate of the city')
plt.show()
```

8. The powered descent phase of a planar planetary landing is commonly posed as the following optimization problem:

$$\begin{array}{ll} \text{minimize} & \sum_i ||u_i|| \\ \text{subject to} & x_{i+1} = Ax_i + Bu_i + g, \forall i = 1, \dots, N \\ & 0 < \rho_1 \leq ||u_i|| \leq \rho_2, \forall i = 1, \dots, N \\ & x_1 = a, x_{N+1} = b \end{array}$$

The decision variables are  $x_i \in \mathbb{R}^4 (i = 1, ..., N + 1)$  and  $u_i \in \mathbb{R}^2 (i = 1, ..., N)$ . The x's represent the lander's state vector: horizontal range, vertical attitude, horizontal velocity, and vertical velocity. The u's represent the lander's control vector coming from the thruster in the horizontal and vertical directions.

The objective minimizes the amount of fuel consumed in the descent. The discrete difference equation approximates the differential equations describing the dynamics. There are no nonlinear terms (such as drag) because the thrust force between lower and upper bounds, i.e. the thruster cannot be turned off and it cannot provide infinitely large thrusts. Lastly, the lander is starting at a point  $a \in \mathbb{R}^4$  and landing at a point on the surface  $b \in \mathbb{R}^4$ .

```
A = [1,0,0.1303,0; 0,1,0,0.1303; 0,0,1,0; 0,0,0,1];
B = [0.0085,0; 0,0.0085; 0.1303,0; 0,0.1303];
g = [0;-0.0832;0;-1.2779];
N = 499;
a = [1000;1500;-25;0];
b = [0;0;0;0];
rho1 = 4;
rho2 = 12;
```

Solve the problem using fmincon, ga, and patternsearch. Plot the state and control trajectories as a function of time.

- Before going into the (incorrect) results that I was able to get for this problem I'm going to describe the things that I did to try and solve it (mostly to give an idea of effort, as opposed to my actual output, which is poor.)
- One thing I spent quite a lot of time on was finding a good Python optimizer. My competence with Matlab is relegated to straightforward use of built-in functions like linprog, fminunc, fmincon, etc. Thus it made sense for me to devote some time to finding a good solver to use in Python (since I know Python fairly well).
- The one I landed on, and tried for quite awhile, is called GEKKO. It is built and maintained by a group at BYU. You can find it here if interested: <a href="https://apmonitor.com/wiki/index.php/Main/GekkoPythonOptimization">https://apmonitor.com/wiki/index.php/Main/GekkoPythonOptimization</a> It seemed like a good candidate since it can handle non-linear, non-convex problems like this one.
- The drawback of using this solver is all of the specific syntax used for setting up problems. What I mainly tried to do was use a couple of their examples. See the bottom two examples "Model Predictive Control" and "Optimization of Multiple Linked Phases" at this site <a href="https://gekko.readthedocs.io/en/latest/examples.html">https://gekko.readthedocs.io/en/latest/examples.html</a>. I was never actually able to get code that ran and gave some kind of intelligible output, hence I've omitted that code from my submission.
- One thing I (kind of) got to work was using the CVXPY package in Python. Credit where credit is due, Ronak explained to me (after Dr. Harris explained to him) how this problem can be made convex by using an additional variable γ. Then I was able to use this package to try and solve the problem. (Again credit to Ronak, he helped me quite a lot in writing this code.)
- Problems with convex formulation solution: as can be seen in the figures below, the solution to which I've arrived is incorrect. This is most easily discerned in the plot of range VS altitude. Many of the value (for both range and altitude) are negative, which is nonsensical; this would mean that the aircraft is somehow below the planet's surface. I tried several things to fix this: modifying the objective (to take the norm instead of just sum of  $u_i$  values), and especially modifying the constraints. The most obvious way to fix this is to introduce constraints that specify that the first two elements of  $x_i \in \mathbb{R}^4$  be values  $\geq 0$ . I tried this (commented-out in the code below), but didn't have any luck. Since I'd already spent quite awhile on this problem and problem 7 (probably 6-7 hours between the two of them) I had to call it good here in order to work on homework for other classes.
- Plots below show the (incorrect) solution to which I arrived.

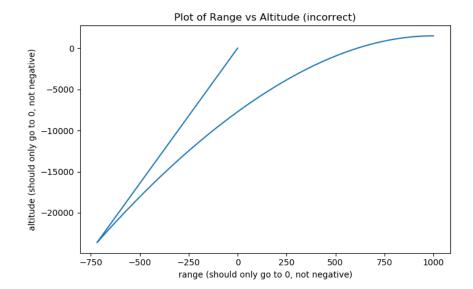


Figure 19: This plot, as mentioned above several times, shows that our solution is incorrect. Something in the formulation doesn't prevent negative values for range or altitude, which leads the solver to determine that going very negative back to (0,0) is the optimal solution. If we're looking hard for any correctness in the plot, at least it starts at the right point (1000,1500) and ends at the right point (0,0). (See later bullet points above for more detailed discussion on how I tried to fix this.)

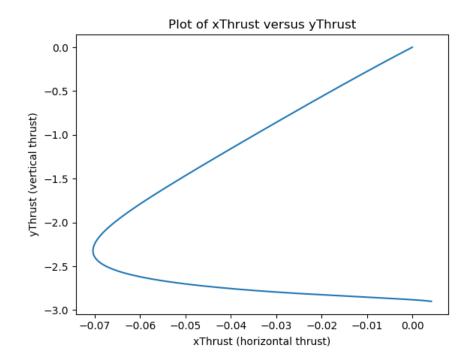


Figure 20: This plot shows horizontal thrust (xThrust) versus (yThrust). It also seem to be pretty incorrect, as we'd expect to see only positive values for yThrust, e.g. propelling the aircraft upward against gravitational forces such that the craft can land with a vertical velocity of 0 when it reaches (0,0).

### Below is my code for this problem.

```
File name : hw7_prob8.py
Author
            : Jared Hansen
Date created : 11/26/2019
Python version: 3.7.3
DESCRIPTION: The purpose of this script is to solve the planetary landing
        problem posed in question 8 in homework 7. Something in the
       formulation is off, since the plot obviously is wrong. It
        satisfies the constraints of x_1 = a, x_{N+1} = b, and starts at
       the correct x and u, but it somehow manages to go negative (e.g.
       crashing far into the ground before coming back up)
#==== IMPORT STATEMENTS
import cvxpy as cp
import numpy as np
import matplotlib.pyplot as plt
import numpy.linalg as npla
#==== DEFINING CONSTANTS, GIVEN INFORMATION
n = 4
m = 2
T = 500
A = np.matrix([[1, 0, 0.1303, 0],
        [0, 1, 0, 0.1303],
        [0, 0, 1, 0],
[0, 0, 0, 1]])
B = np.matrix([[0.0085, 0],
        [0 , 0.0085],
        [0.1303, 0 ],
        [0 , 0.1303]])
g = np.array([0, -0.0832, 0, -1.2779])#.reshape(4,1)
a = np.array([1000, 1500, -25, 0])#.reshape(4,1)
b = np.array([0, 0, 0, 0])#.reshape(4,1)
rho1 = 4;
rho2 = 12:
# Defining variables for the convex solver. We introduce a new variable, gamma,
# to make the problem convex.
x = cp.Variable((n, T))
u = cp.Variable((m, T))
gamma = cp.Variable((1,T))
cost = 0
constr = []
for t in range(T-2):
  cost += cp.sum((gamma[:,t])**2)
  constr += [x[:,t+1] == A@x[:,t] + B@u[:,t] + g,
         cp.norm(u[:,t], 'inf') <= gamma[:,t],</pre>
         gamma[:,t] \leftarrow rho2,
```

```
-gamma[:,t] <= -rho1,
  # This is where the problem lies: there is nothing telling the algorithm
  # that we can't have negative thrust or negative position values. When I
  # try and implement these constraints the solver says that the objective
  # value of the solution is infinite, so this doesn't work with this solver.
  constr += [-x[:,t] \le -np.array([0,0,0,0]),
          -u[:,t] \le np.array([0,0])
constr += [x[:,499] == b, x[:,0] == a]
problem = cp.Problem(cp.Minimize(cost), constr)
problem.solve(solver=cp.ECOS)
#==== PLOTTING (INCORRECT) RESULTS
x_mat = x.value
u_mat = u.value
x_range = x_mat[0,:]
x_altitude = x_mat[1,:]
u_xThrust = u_mat[0,:]
u_yThrust = u_mat[1,:]
# Plot of range versus altitude
plt.plot(x_range, x_altitude)
plt.title("Plot of Range vs Altitude (incorrect)")
plt xlabel('range (should only go to 0, not negative)')
plt.ylabel('altitude (should only go to 0, not negative)')
plt.plot(u_xThrust, u_yThrust)
plt.title('Plot of xThrust versus yThrust')
plt.xlabel('xThrust (horizontal thrust)')
plt.ylabel('yThrust (vertical thrust)')
```