Dynamic Automatic Differentiation of GPU Broadcast Kernels

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Abstract

We show how forward-mode (FM) automatic differentiation can be employed within reverse-mode (RM) computations to dynamically differentiate broadcast operations in a GPUfriendly manner.

Our technique fully exploits the broadcast Jacobian's inherent sparsity structure, and unlike a pure reverse-mode approach, this mixed-mode approach does not require a backwards pass over the broadcasted operation's subgraph, obviating the need for several reverse-mode-specific programmability restrictions on user-authored broadcast operations.

Most notably, this approach allows broadcast fusion in primal code despite the presence of data-dependent control flow. We discuss an experiment in which a Julia implementation of our technique outperformed pure reverse-mode TensorFlow and Julia implementations for differentiating through broadcast operations within an HM-LSTM cell-update calculation.

Broadcasting $p:\mathbb{R}^3 \to \mathbb{R}^2$ over a matrix **A**, scalar α , and vector **a**:

$$p.(\mathbf{A}, \alpha, \mathbf{a}) = \begin{pmatrix} \begin{bmatrix} p(A_{11}, \alpha, a_1)_1 & \dots & p(A_{1m}, \alpha, a_1)_1 \\ \vdots & \ddots & \vdots \\ p(A_{n1}, \alpha, a_n)_1 & \dots & p(A_{nm}, \alpha, a_n)_1 \end{bmatrix}, \begin{bmatrix} p(A_{11}, \alpha, a_1)_2 & \dots & p(A_{1m}, \alpha, a_1)_2 \\ \vdots & \ddots & \vdots \\ p(A_{n1}, \alpha, a_n)_2 & \dots & p(A_{nm}, \alpha, a_n)_2 \end{bmatrix} \end{pmatrix}$$

Reverse-mode AD for an operation containing a broadcast:

Definition	Forward (Primal)	Reverse (Adjoint)
$h(x, \mathbf{y}) = g(\mathbf{f}(x, \mathbf{y}))$ $\mathbf{f}(x, \mathbf{y}) = b.(x, \mathbf{y})$ $b: \mathbb{R}^2 \to \mathbb{R}$ $g: \mathbb{R}^N \to \mathbb{R}$ $x \in \mathbb{R}, \ \mathbf{y} \in \mathbb{R}^N$	$\mathbf{w}_1 = \mathbf{f}(x, \mathbf{y})$ $w_2 = g(\mathbf{w}_1)$	$\overline{w}_2 = 1 \text{ (seed)}$ $\overline{\mathbf{w}}_1 = \overline{w}_2 \frac{\partial w_2}{\partial \mathbf{w}_1}$ $\frac{\partial h}{\partial x} = \overline{\mathbf{w}}_1 \cdot \frac{\partial \mathbf{f}}{\partial x}$ $\frac{\partial h}{\partial \mathbf{y}} = \overline{\mathbf{w}}_1 \cdot \times \frac{\partial f_i}{\partial y_i}$

Our technique: replace b broadcast in forward pass with FM AD'd version to get elementwise derivatives without needing to RM b

$$\left(\frac{\partial \mathbf{f}}{\partial x}, \frac{\partial f_i}{\partial y_i}\right) = \mathbf{D}(b).(x, \mathbf{y})$$

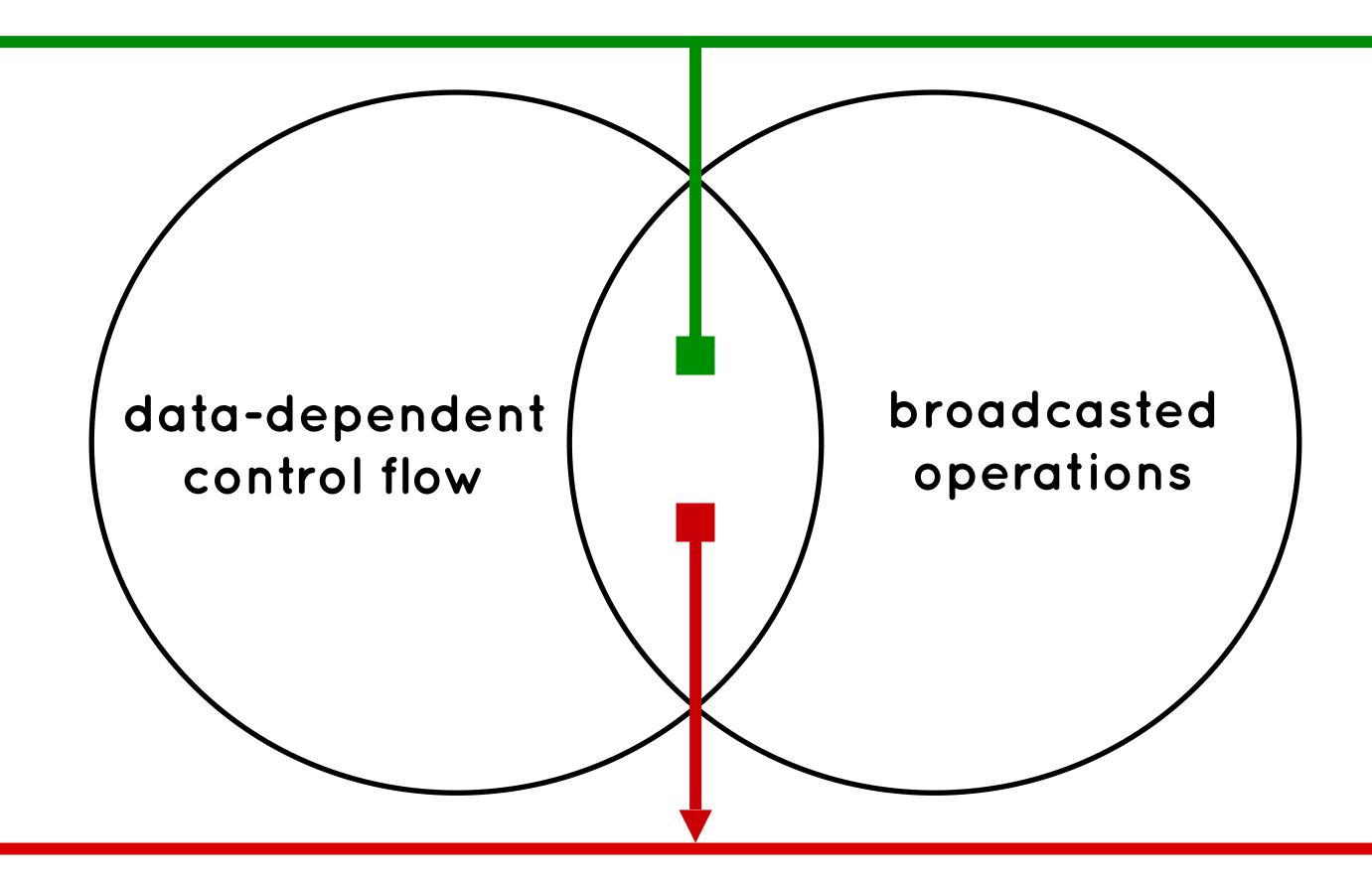
An example where this technique pays off: HM-LSTM cell update!

$$\mathbf{c}_t^\ell = \begin{cases} \sigma.(\mathbf{f}_t^\ell) \times \mathbf{c}_{t-1}^\ell + \sigma.(\mathbf{i}_t^\ell) \times \tanh.(\mathbf{g}_t^\ell) & \text{if } z_{t-1}^\ell = 0, z_t^{\ell-1} = 1 \text{ (UPDATE)} & \leftarrow \text{expensive branch} \\ \mathbf{c}_{t-1}^\ell & \text{if } z_{t-1}^\ell = 0, z_t^{\ell-1} = 0 \text{ (COPY)} & \leftarrow \text{cheap branch} \\ \sigma.(\mathbf{i}_t^\ell) \times \tanh.(\mathbf{g}_t^\ell) & \text{if } z_t^{\ell-1} = 1 \text{ (FLUSH)} \end{cases}$$

TensorFlow def cell_update(z, zb, c, f, i, g):

- i = tf.sigmoid(i)
- g = tf.tanh(g)
- f = tf.sigmoid(f) flush_control = tf.equal(z, tf.constant(1., dtype=tf.float32)) flush = tf.multiply(i, g)
- copy_control = tf.equal(zb, tf.constant(0., dtype=tf.float32)) copy = tf.identity(c)
- update = tf.add(tf.multiply(f, c), tf.multiply(i, g)) return tf.where(flush_control, flush, tf.where(copy_control, copy, update))
- Native Julia
- # can be broadcasted at callsite, e.g. cell_update.(z, zb, c, f, i, g) function cell_update(z, zb, c, f, i, g) return sigm(i) * tanh(g)
- elseif zb == 0.0f0 return c
- return sigm(f) * c + sigm(i) * tanh(g)
- Reverse (TensorFlow) compute time $[\mu s]$ Reverse (Julia) 10^{3} Forward (Julia) 1,024 1,024 5122,048 5125122.0481,024 2.048 GTX Titan (Kepler) Tesla V100 (Volta) Tesla P100 (Pascal)
- FM beats out RM 2x-4x; moreso when controlling for RM implementation
- Newer GPU architectures mitigate negative effects of warp divergence, and thus benefit greatly from FM-enabled fusion

Fusion for these cases: easy to perform + greatly beneficial!

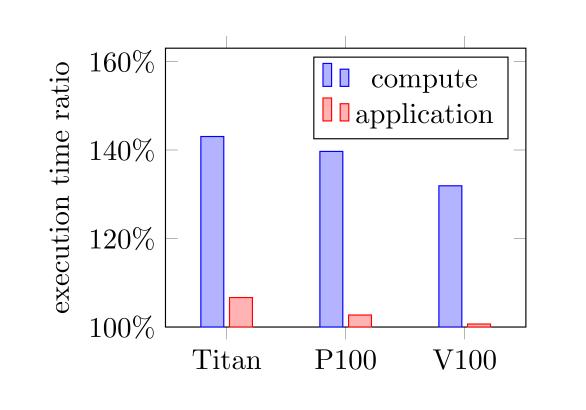


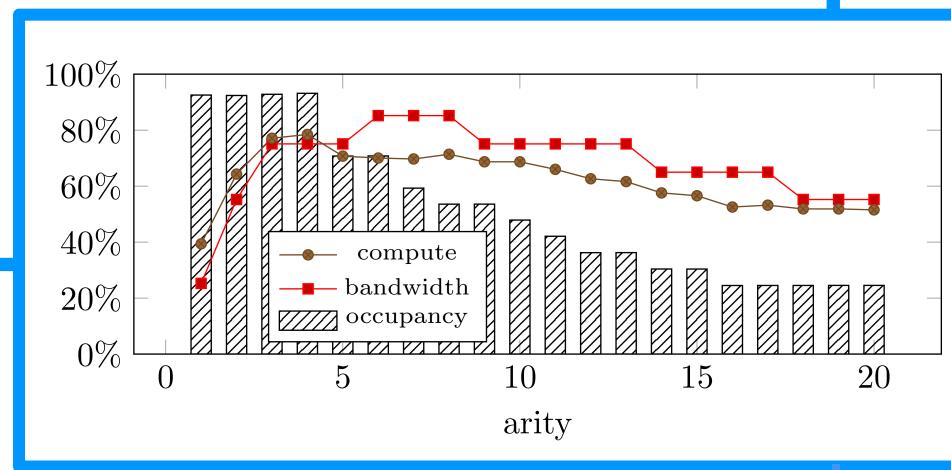
fusion → massive amount of data-dependent scalar operations reverse-mode + dynamic scalar ops → fine-grained dynamic allocations

massive amount of fine-grained allocations → BAD FOR GPU!

Solution: <u>REVERSE-MODE AD</u> the overall computation, but FORWARD-MODE AD the broadcasted operation

- (FM mul-add count)/(RM mul-add count) ~ N/M. For low-arity scalar ops (e.g. broadcasted ops), though, FM performance often matches RM performance even when N > M by making good use of cache bandwidth and exploiting instruction-level parallelism.
- Unlike RM, FM does not require a reversible representation of the operation. Regardless of RM implementation (tape, graph, closure, continuation, etc.), handling data-dependent control flow in target code generally requires dynamic allocation for RM.
- Thus, **FM** enables data-dependent control within broadcasted operations. This capability renders more operations fusable and allows users to express kinds of scalar control-flow previously disallowed in target code, increasing model programmability.
- performance ratio for executing with totally warp-uniform control inputs vs. totally warp-divergent control inputs
- Application overhead due to warp divergence on V100: < 1%





- Hardware utilization drops as register pressure rises
- Need heuristics to tune "chunk size" w.r.t. arity; start by capping at 10.

This work is already used by **Flux/Zygote** on CPUs, GPUs, and **TPUs**. We're planning a new tool called Capstan for mixed-mode AD. Want more code, math, and the bibliography? See the paper!