

Exercises 1.7

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The properties of area in this set of exercises are to be deduced from the axioms for area stated in the foregoing section.

Exercise 1

Prove that each of the following sets is measurable and has zero area:

⊙ Exercise 1a

A set consisting of a single point.

Proof. Let S be a set consisting of a single point. By definition of a Point, S is a rectangle in which all vertices coincide. By ¶ Choice of Scale, S is measurable with area its width times its height. The width and height of S is trivially zero. Therefore $a(S) = (0)(0) = 0$. □

⊙ Exercise 1b

A set consisting of a finite number of points in a plane.

Proof. Define predicate $P(n)$ as "A set consisting of n points in a plane is measurable with area 0". We use induction to prove $P(n)$ holds for all $n > 0$.

Base Case Consider a set S consisting of a single point in a plane. By ⊙ Exercise 1a, S is measurable with area 0. Thus $P(1)$ holds.

Induction Step Assume induction hypothesis $P(k)$ holds for some $k > 0$. Let S_{k+1} be a set consisting of $k + 1$ points in a plane. Pick an arbitrary point of S_{k+1} . Denote the set containing just this point as T . Denote the remaining set of points as S_k . By construction, $S_{k+1} = S_k \cup T$. By the induction hypothesis, S_k is measurable with area 0. By ⊙ Exercise 1a, T is measurable with area 0. By the ¶ Additive Property, $S_k \cup T$ is measurable, $S_k \cap T$ is measurable, and

$$\begin{aligned} a(S_{k+1}) &= a(S_k \cup T) \\ &= a(S_k) + a(T) - a(S_k \cap T) \\ &= 0 + 0 - a(S_k \cap T). \end{aligned} \tag{1}$$

There are two cases to consider:

Case 1 $S_k \cap T = \emptyset$. Then it trivially follows that $a(S_k \cap T) = 0$.

Case 2 $S_k \cap T \neq \emptyset$. Since T consists of a single point, $S_k \cap T = T$. By [Exercise 1a](#), $a(S_k \cap T) = a(T) = 0$.

In both cases, [\(1\)](#) evaluates to 0, implying $P(k + 1)$ as expected.

Conclusion By mathematical induction, it follows for all $n > 0$, $P(n)$ is true. \square

[Exercise 1c](#)

The union of a finite collection of line segments in a plane.

Proof. Define predicate $P(n)$ as "A set consisting of n line segments in a plane is measurable with area 0". We use induction to prove $P(n)$ holds for all $n > 0$.

Base Case Consider a set S consisting of a single line segment in a plane. By definition of a Line Segment, S is a rectangle in which one side has dimension 0. By [Choice of Scale](#), S is measurable with area its width w times its height h . Therefore $a(S) = wh = 0$. Thus $P(1)$ holds.

Induction Step Assume induction hypothesis $P(k)$ holds for some $k > 0$. Let S_{k+1} be a set consisting of $k + 1$ line segments in a plane. Pick an arbitrary line segment of S_{k+1} . Denote the set containing just this line segment as T . Denote the remaining set of line segments as S_k . By construction, $S_{k+1} = S_k \cup T$. By the induction hypothesis, S_k is measurable with area 0. By the base case, T is measurable with area 0. By the [Additive Property](#), $S_k \cup T$ is measurable, $S_k \cap T$ is measurable, and

$$\begin{aligned} a(S_{k+1}) &= a(S_k \cup T) \\ &= a(S_k) + a(T) - a(S_k \cap T) \\ &= 0 + 0 - a(S_k \cap T). \end{aligned} \tag{2}$$

There are two cases to consider:

Case 1 $S_k \cap T = \emptyset$. Then it trivially follows that $a(S_k \cap T) = 0$.

Case 2 $S_k \cap T \neq \emptyset$. Since T consists of a single point, $S_k \cap T = T$. By the base case, $a(S_k \cap T) = a(T) = 0$.

In both cases, [\(2\)](#) evaluates to 0, implying $P(k + 1)$ as expected.

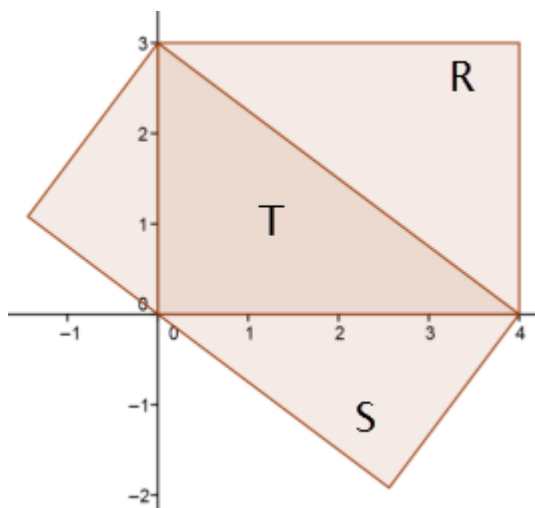
Conclusion By mathematical induction, it follows for all $n > 0$, $P(n)$ is true. \square

! Exercise 2

Every right triangular region is measurable because it can be obtained as the intersection of two rectangles. Prove that every triangular region is measurable and that its area is one half the product of its base and altitude.

Proof. Let T' be a triangular region with base of length a , height of length b , and hypotenuse of length c . Consider the translation and rotation of T' , say T , such that its hypotenuse is entirely within quadrant I and the vertex opposite the hypotenuse is situated at point $(0,0)$.

Let R be a rectangle of width a , height b , and bottom-left corner at $(0,0)$. By construction, R covers all of T . Let S be a rectangle of width c and height $a \sin \theta$, where θ is the acute angle measured from the bottom-right corner of T relative to the x -axis. As an example, consider the image below of triangle T with width 4 and height 3:



By [Choice of Scale](#), both R and S are measurable. By this same axiom, $a(R) = ab$ and $a(S) = ca \sin \theta$. By the [Additive Property](#), $R \cup S$ and $R \cap S$ are both measurable. $a(R \cap S) = a(T)$ and $a(R \cup S)$ can be determined by

noting that R 's construction implies identity $a(R) = 2a(T)$. Therefore

$$\begin{aligned}
 a(T) &= a(R \cap S) \\
 &= a(R) + a(S) - a(R \cup S) \\
 &= ab + ca \sin \theta - a(R \cup S) \\
 &= ab + ca \sin \theta - (ca \sin \theta + \frac{1}{2}a(R)) \\
 &= ab + ca \sin \theta - ca \sin \theta - a(T).
 \end{aligned}$$

Solving for $a(T)$ gives the desired identity:

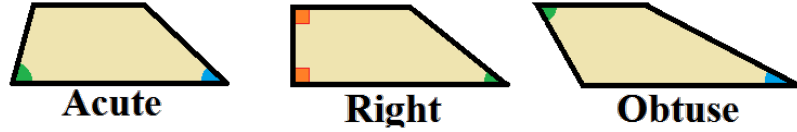
$$a(T) = \frac{1}{2}ab.$$

By [Invariance Under Congruence](#), $a(T') = a(T)$, concluding our proof. \square

! Exercise 3

Prove that every trapezoid and every parallelogram is measurable and derive the usual formulas for their areas.

Proof. We begin by proving the formula for a trapezoid. Let S be a trapezoid with height h and bases b_1 and b_2 , $b_1 < b_2$. There are three cases to consider:



Case 1 Suppose S is a right trapezoid. Then S is the union of non-overlapping rectangle R of width b_1 and height h with right triangle T of base $b_2 - b_1$ and height h . By [Choice of Scale](#), R is measurable. By [Exercise 2](#), T is measurable. By the [Additive Property](#), $R \cup T$ and $R \cap T$ are both measurable

and

$$\begin{aligned}
a(S) &= a(R \cup T) \\
&= a(R) + a(T) - a(R \cap T) \\
&= a(R) + a(T) && \text{by construction} \\
&= b_1 h + a(T) && \text{Choice of Scale} \\
&= b_1 h + \frac{1}{2}(b_2 - b_1)h && \text{! Exercise 2} \\
&= \frac{b_1 + b_2}{2}h.
\end{aligned}$$

Case 2 Suppose S is an acute trapezoid. Then S is the union of non-overlapping triangle T and right trapezoid R . Let c denote the length of base T . Then R has longer base edge of length $b_2 - c$. By ! Exercise 2, T is measurable. By Case 1, R is measurable. By the ! Additive Property, $R \cup T$ and $R \cap T$ are both measurable and

$$\begin{aligned}
a(S) &= a(T) + a(R) - a(R \cap T) \\
&= a(T) + a(R) && \text{by construction} \\
&= \frac{1}{2}ch + a(R) && \text{! Exercise 2} \\
&= \frac{1}{2}ch + \frac{b_1 + b_2 - c}{2}h && \text{Case 1} \\
&= \frac{b_1 + b_2}{2}h.
\end{aligned}$$

Case 3 Suppose S is an obtuse trapezoid. Then S is the union of non-overlapping triangle T and right trapezoid R . Let c denote the length of base T . Reflect T vertically to form another right triangle, say T' . Then $T' \cup R$ is an acute trapezoid. By ! Invariance Under Congruence,

$$a(T' \cup R) = a(T \cup R). \quad (3.1)$$

By construction, $T' \cup R$ has height h and bases $b_1 - c$ and $b_2 + c$ meaning

$$\begin{aligned}
a(T \cup R) &= a(T' \cup R) && (3.1) \\
&= \frac{b_1 - c + b_2 + c}{2}h && \text{Case 2} \\
&= \frac{b_1 + b_2}{2}h.
\end{aligned}$$

Conclusion These cases are exhaustive and in agreement with one another. Thus S is measurable and

$$a(S) = \frac{b_1 + b_2}{2}h.$$

Let P be a parallelogram with base b and height h . Then P is the union of non-overlapping triangle T and right trapezoid R . Let c denote the length of base T . Reflect T vertically to form another right triangle, say T' . Then $T' \cup R$ is an acute trapezoid. By [Invariance Under Congruence](#),

$$a(T' \cup R) = a(T \cup R). \quad (3.2)$$

By construction, $T' \cup R$ has height h and bases $b - c$ and $b + c$ meaning

$$\begin{aligned} a(T \cup R) &= a(T' \cup R) & (3.2) \\ &= \frac{b - c + b + c}{2} h & \text{Area of Trapezoid} \\ &= bh. \end{aligned}$$

□

Exercise 4

Let P be a polygon whose vertices are lattice points. The area of P is $I + \frac{1}{2}B - 1$, where I denotes the number of lattice points inside the polygon and B denotes the number on the boundary.

! Exercise 4a

Prove that the formula is valid for rectangles with sides parallel to the coordinate axes.

Proof. Let P be a rectangle with sides parallel to the coordinate axes, with width w , height h , and lattice points for vertices. We assume P has three non-collinear points, ruling out any instances of points or line segments.

By [Choice of Scale](#), P is measurable with area $a(P) = wh$. By construction, P has $I = (w - 1)(h - 1)$ interior lattice points and $B = 2(w + h)$ lattice points on its boundary. The following shows the lattice point area formula is in agreement with the expected result:

$$\begin{aligned} I + \frac{1}{2}B - 1 &= (w - 1)(h - 1) + \frac{1}{2}[2(w + h)] - 1 \\ &= (wh - w - h + 1) + \frac{1}{2}[2(w + h)] - 1 \\ &= (wh - w - h + 1) + (w + h) - 1 \\ &= wh. \end{aligned}$$

□

! Exercise 4b

Prove that the formula is valid for right triangles and parallelograms.

Proof. Let P be a right triangle with width $w > 0$, height $h > 0$, and lattice points for vertices. Let T be the triangle P translated, rotated, and reflected such that its vertices are $(0,0)$, $(0,w)$, and (w,h) . Let I_T and B_T be the number of interior and boundary points of T respectively. Let H_L denote the number of lattice points on T 's hypotenuse.

Let R be the overlapping rectangle of width w and height h , situated with bottom-left corner at $(0,0)$. Let I_R and B_R be the number of interior and boundary points of R respectively.

By construction, T shares two sides with R . Therefore

$$B_T = \frac{1}{2}B_R - 1 + H_L. \quad (3)$$

Likewise,

$$I_T = \frac{1}{2}(I_R - H_L + 2). \quad (4)$$

The following shows the lattice point area formula is in agreement with the expected result:

$$\begin{aligned} I_T + \frac{1}{2}B_T - 1 &= \frac{1}{2}(I_R - H_L + 2) + \frac{1}{2}B_T - 1 & (4) \\ &= \frac{1}{2}[I_R - H_L + 2 + B_T - 2] \\ &= \frac{1}{2}[I_R - H_L + B_T] \\ &= \frac{1}{2}\left[I_R - H_L + \frac{1}{2}B_R - 1 + H_L\right] & (3) \\ &= \frac{1}{2}\left[I_R + \frac{1}{2}B_R - 1\right] \\ &= \frac{1}{2}\left[(w-1)(h-1) + \frac{1}{2}(2(w+h)) - 1\right] & \text{! Exercise 4a} \\ &= \frac{1}{2}[(w-1)(h-1) + w + h - 1] \\ &= \frac{1}{2}[wh - w - h + 1 + w + h - 1] \\ &= \frac{wh}{2}. \end{aligned}$$

We do not prove this formula is valid for parallelograms here. Instead, refer to ! Exercise 4c below.

□

! Exercise 4c

Use induction on the number of edges to construct a proof for general polygons.

Proof. Define predicate $P(n)$ as "An n -polygon with vertices on lattice points has area $I + \frac{1}{2}B - 1$." We use induction to prove $P(n)$ holds for all $n \geq 3$.

Base Case A 3-polygon is a triangle. By ! Exercise 4b, the lattice point area formula holds. Thus $P(3)$ holds.

Induction Step Assume induction hypothesis $P(k)$ holds for some $k \geq 3$. Let P be a $(k+1)$ -polygon with vertices on lattice points. Such a polygon is equivalent to the union of a k -polygon S with a triangle T . That is, $P = S \cup T$.

Let I_P be the number of interior lattice points of P . Let B_P be the number of boundary lattice points of P . Similarly, let I_S , I_T , B_S , and B_T be the number of interior and boundary lattice points of S and T . Let c denote the number of boundary points shared between S and T .

By our induction hypothesis, $a(S) = I_S + \frac{1}{2}B_S - 1$. By our base case, $a(T) = I_T + \frac{1}{2}B_T - 1$. By construction, it follows:

$$\begin{aligned} I_P &= I_S + I_T + c - 2 \\ B_P &= B_S + B_T - (c - 2) - c \\ &= B_S + B_T - 2c + 2. \end{aligned}$$

Applying the lattice point area formula to P yields the following:

$$\begin{aligned} &I_P + \frac{1}{2}B_P - 1 \\ &= (I_S + I_T + c - 2) + \frac{1}{2}(B_S + B_T - 2c + 2) - 1 \\ &= I_S + I_T + c - 2 + \frac{1}{2}B_S + \frac{1}{2}B_T - c + 1 - 1 \\ &= (I_S + \frac{1}{2}B_S - 1) + (I_T + \frac{1}{2}B_T - 1) \\ &= a(S) + (I_T + \frac{1}{2}B_T - 1) && \text{induction hypothesis} \\ &= a(S) + a(T). && \text{base case} \end{aligned}$$

By the Additive Property, $S \cup T$ is measurable, $S \cap T$ is measurable, and

$$\begin{aligned} a(P) &= a(S \cup T) \\ &= a(S) + a(T) - a(S \cap T) \\ &= a(S) + a(T). && \text{by construction} \end{aligned}$$

This shows the lattice point area formula is in agreement with our axiomatic definition of area. Thus $P(k+1)$ holds.

Conclusion By mathematical induction, it follows for all $n \geq 3$, $P(n)$ is true. \square

! Exercise 5

Prove that a triangle whose vertices are lattice points cannot be equilateral.

[*Hint:* Assume there is such a triangle and compute its area in two ways, using Exercises 2 and 4.]

Proof. Proceed by contradiction. Let T be an equilateral triangle whose vertices are lattice points. Assume each side of T has length a . Then T has height $h = (a\sqrt{3})/2$. By ! Exercise 2,

$$a(T) = \frac{1}{2}ah = \frac{a^2\sqrt{3}}{4}. \quad (5.1)$$

Let I and B denote the number of interior and boundary lattice points of T respectively. By Exercise 4,

$$a(T) = I + \frac{1}{2}B - 1. \quad (5.2)$$

But (5.1) is irrational whereas (5.2) is not. This is a contradiction. Thus, there is *no* equilateral triangle whose vertices are lattice points. \square

! Exercise 6

Let $A = \{1, 2, 3, 4, 5\}$, and let \mathcal{M} denote the class of all subsets of A . (There are 32 altogether, counting A itself and the empty set \emptyset .) For each set S in \mathcal{M} , let $n(S)$ denote the number of distinct elements in S . If $S = \{1, 2, 3, 4\}$ and $T = \{3, 4, 5\}$, compute $n(S \cup T)$, $n(S \cap T)$, $n(S - T)$, and $n(T - S)$. Prove that the set function n satisfies the first three axioms for area.

Proof. Let $S = \{1, 2, 3, 4\}$ and $T = \{3, 4, 5\}$. Then

$$\begin{aligned}
n(S \cup T) &= n(\{1, 2, 3, 4\} \cup \{3, 4, 5\}) \\
&= n(\{1, 2, 3, 4, 5\}) \\
&= 5. \\
n(S \cap T) &= n(\{1, 2, 3, 4\} \cap \{3, 4, 5\}) \\
&= n(\{3, 4\}) \\
&= 2. \\
n(S - T) &= n(\{1, 2, 3, 4\} - \{3, 4, 5\}) \\
&= n(\{1, 2\}) \\
&= 2. \\
n(T - S) &= n(\{3, 4, 5\} - \{1, 2, 3, 4\}) \\
&= n(\{5\}) \\
&= 1.
\end{aligned}$$

We now prove n satisfies the first three axioms for area.

Nonnegative Property n returns the length of some member of \mathcal{M} . By hypothesis, the smallest possible input to n is \emptyset . Since $n(\emptyset) = 0$, it follows $n(S) \geq 0$ for all $S \in \mathcal{M}$.

Additive Property Let S and T be members of \mathcal{M} . It trivially follows that both $S \cup T$ and $S \cap T$ are in \mathcal{M} . Consider the value of $n(S \cup T)$. There are two cases to consider:

Case 1 Suppose $S \cap T = \emptyset$. That is, there is no common element shared between S and T . Thus

$$\begin{aligned}
n(S \cup T) &= n(S) + n(T) \\
&= n(S) + n(T) - 0 \\
&= n(S) + n(T) - n(S \cap T).
\end{aligned}$$

Case 2 Suppose $S \cap T \neq \emptyset$. Then $n(S) + n(T)$ counts each element of $S \cap T$ twice. Therefore $n(S \cup T) = n(S) + n(T) - n(S \cap T)$.

Conclusion These cases are exhaustive and in agreement with one another. Thus $n(S \cup T) = n(S) + n(T) - n(S \cap T)$.

Difference Property Suppose $S, T \in \mathcal{M}$ such that $S \subseteq T$. That is, every member of S is a member of T . By definition, $T - S$ consists of members in T but not in S . Thus $n(T - S) = n(T) - n(S)$.

□